# 操作系统实验报告Prj2-Part1

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# 主要设计思路

# 进程调度

# PCB包含的信息

主要包含以下信息: (代码注释中)

```
typedef struct pcb
   //寄存器(用作恢复现场)
   regs_context_t kernel_context;
   regs_context_t user_context;
   //栈顶位置
   uint32_t kernel_stack_top;
   uint32_t user_stack_top;
   //上一个进程和下一个进程指针,用作进程队列
   void *prev;
   void *next;
   //pid
   pid_t pid;
   //区分这是用户进程还是系统进程
   task_type_t type;
   //进程状态:准备、阻塞、退出
   task status t status;
   //光标位置
   int cursor_x;
   int cursor_y;
} pcb_t;
```

# 如何启动一个task

### 如何获得task的入口地址

将 \$ra 寄存器设置为task的入口地址。因为在调度结束的时候,会执行 jr \$ra 的指令。

### 启动时需要设置哪些寄存器

需要设置 ra , sp 寄存器。

## 进程之间调度的流程

- 1.一个进程使用 do schedule 函数,将控制权交给操作系统进行调度。
- 2.操作系统将进程现场进行保存(包括寄存器,CPO特殊寄存器数据),并根据进程的状态决定进程进入阻塞队列、准备队列还是退出进程。
- 3.从准备队列中选取下一个进程。
- 4.对选取的进程进行现场恢复。(包括寄存器, CPO特殊寄存器数据)

# 锁

# spin-lock和mutual-lock的区别

自旋锁的工作原理是:一旦无法获得锁,就会一直请求获得锁,直到获得锁为止。

互斥锁的工作原理是:如果没有获得锁,就会将其放入阻塞队列中,当锁解锁的时候,会将其从阻塞队

列中拿出,放入准备队列中。

## 锁的处理流程

1.在锁的初始化阶段,将锁的状态设定为 UNBLOCKED 。 2.程序请求锁之后,将指定id的锁的状态设置为 LOCKED 。 3.如果在请求锁的过程中,发现已经被锁,则将程序状态设定为 TASK\_STATUS ,并调用 do\_schedule 函数,将其放入阻塞队列。

4.程序在使用锁完毕之后,会将锁进行解锁,并将相应阻塞队列中的进程设定为 TASK\_READY ,并放入准备队列中。

# 源码设计(基本包含所有代码)

### 进程调度

### 保护现场

```
.macro SAVE_CONTEXT offset
   #
       $k0(26)
                  current running
                                     # Current running is a global pointer.
   lw
             k0, current_running
             $1, OFFSET REG1(k0)
   SW
   SW
             $2, OFFSET_REG2(k0)
             $3, OFFSET_REG3(k0)
   SW
             $4, OFFSET_REG4(k0)
   SW
   SW
             $5, OFFSET REG5(k0)
              $6, OFFSET_REG6(k0)
   SW
              $7, OFFSET_REG7(k0)
              $8, OFFSET REG8(k0)
   SW
             $9, OFFSET_REG9(k0)
   SW
              $10, OFFSET_REG10(k0)
   SW
             $11, OFFSET_REG11(k0)
              $12, OFFSET REG12(k0)
   SW
              $13, OFFSET_REG13(k0)
   SW
              $14, OFFSET REG14(k0)
   SW
              $15, OFFSET_REG15(k0)
   SW
              $16, OFFSET_REG16(k0)
   SW
             $17, OFFSET_REG17(k0)
   SW
             $18, OFFSET_REG18(k0)
             $19, OFFSET_REG19(k0)
   SW
             $20, OFFSET REG20(k0)
             $21, OFFSET_REG21(k0)
```

```
$22, OFFSET_REG22(k0)
   SW
   SW
              $23, OFFSET_REG23(k0)
              $24, OFFSET_REG24(k0)
   SW
              $25, OFFSET_REG25(k0)
   SW
   # no
            26
   # no
            27
              $28, OFFSET_REG28(k0)
   SW
              $29, OFFSET_REG29(k0)
   SW
              $30, OFFSET_REG30(k0)
   SW
              $31, OFFSET_REG31(k0)
            k1, CP0_STATUS
   mfc0
   nop
            k1, OFFSET_STATUS(k0)
   SW
            k1, CP0_STATUS
   mfc0
   nop
            k1, OFFSET_STATUS(k0)
   SW
   mfhi
            k1
   nop
            k1, OFFSET_HI(k0)
   SW
            k1
   mflo
   nop
            k1, OFFSET_LO(k0)
   SW
   mfc0
            k1, CP0 BADVADDR
   nop
            k1, OFFSET_BADVADDR(k0)
   SW
   mfc0
            k1, CP0_CAUSE
   nop
   SW
            k1, OFFSET_CAUSE(k0)
   mfc0
            k1, CP0 EPC
   nop
            k1, OFFSET_EPC(k0)
   SW
.endm
```

#### 调度函数

```
void scheduler(void)
{
    if(current_running->status == TASK_RUNNING) {
        current_running->status = TASK_READY;
        queue_push(&ready_queue, current_running);
    } else if(current_running->status == TASK_BLOCKED) {
        queue_push(&block_queue, current_running);
    }
    current_running = queue_dequeue(&ready_queue);
    current_running->status = TASK_RUNNING;
}
```

### 恢复现场

```
.macro RESTORE_CONTEXT offset
     # $k0(26)
                    current_running
   lw
           k0, current_running
                                   # Current running is a global pointer.
            $1, OFFSET REG1(k0)
   lw
   lw
             $2, OFFSET_REG2(k0)
             $3, OFFSET REG3(k0)
   lw
   lw
             $4, OFFSET_REG4(k0)
             $5, OFFSET_REG5(k0)
            $6, OFFSET_REG6(k0)
   lw
            $7, OFFSET REG7(k0)
   lw
            $8, OFFSET_REG8(k0)
   lw
            $9, OFFSET REG9(k0)
            $10, OFFSET_REG10(k0)
   lw
   lw
            $11, OFFSET_REG11(k0)
   lw
            $12, OFFSET_REG12(k0)
            $13, OFFSET_REG13(k0)
   lw
   lw
            $14, OFFSET_REG14(k0)
            $15, OFFSET_REG15(k0)
   lw
            $16, OFFSET REG16(k0)
   lw
            $17, OFFSET_REG17(k0)
   lw
            $18, OFFSET_REG18(k0)
   lw
            $19, OFFSET_REG19(k0)
            $20, OFFSET REG20(k0)
            $21, OFFSET_REG21(k0)
   lw
            $22, OFFSET_REG22(k0)
            $23, OFFSET_REG23(k0)
   ٦w
   lw
            $24, OFFSET REG24(k0)
   lw
            $25, OFFSET_REG25(k0)
           26
   # no
   # no
           27
            $28, OFFSET_REG28(k0)
   lw
            $29, OFFSET_REG29(k0)
   lw
            $30, OFFSET_REG30(k0)
   lw
   lw
            $31, OFFSET REG31(k0)
   lw
           k1, OFFSET STATUS(k0)
           k1, CP0_STATUS
   mtc0
   nop
   lw
           k1, OFFSET_STATUS(k0)
           k1, CP0_STATUS
   mtc0
   nop
           k1, OFFSET_HI(k0)
   lw
   mthi
           k1
   nop
   lw
           k1, OFFSET_LO(k0)
   mtlo
   nop
           k1, OFFSET_BADVADDR(k0)
   lw
   mtc0
           k1, CP0 BADVADDR
   nop
           k1, OFFSET_CAUSE(k0)
   lw
   mtc0
           k1, CP0 CAUSE
   nop
```

lw

k1, OFFSET\_EPC(k0)

```
mtc0 k1, CP0_EPC
nop
.endm
```

### 阻塞与释放

```
void do_block()
{
    // block the current_running task into the queue
    current_running->status = TASK_BLOCKED;
    // push to blockqueue
    queue_push(&block_queue, current_running);
    do_scheduler();
}

void do_unblock_one()
{
    // unblock the head task from the queue
    pcb_t *unblock_one;
    unblock_one = queue_dequeue(&block_queue);
    unblock_one->status = TASK_READY;
    queue_push(&ready_queue, unblock_one);
}
```

### 初始化进程

```
static void init_pcb()
{
    int i;
    int j;
    current_running = &pcb[0];
    process_id = 1;
    pcb[0].pid = 1;
    pcb[0].status = TASK_EXITED;
    pcb[0].cursor_x = 0;
    pcb[0].cursor_y = 0;
    for(i = 1; i \leftarrow 31; i++){
        pcb[0].kernel_context.regs[i] = 0;
    }
    pcb[0].kernel_context.regs[29] = 0xa0f01000;
    pcb[0].kernel_context.regs[31] = 0xa0800200;
    pcb[0].kernel_context.cp0_status = 0;
    pcb[0].kernel_context.hi = 0;
    pcb[0].kernel_context.lo = 0;
    pcb[0].kernel_context.cp0_badvaddr = 0;
    pcb[0].kernel_context.cp0_cause = 0;
    pcb[0].kernel_context.cp0_epc = 0;
    for(j = 0; j < num\_sched1\_tasks; j++){
        pcb[j + 1].pid = 2 + j;
        pcb[j + 1].status = TASK_READY;
        pcb[j + 1].cursor_x = 0;
        pcb[j + 1].cursor_y = 0;
        for(i = 1; i <= 31; i++){
            pcb[j + 1].kernel_context.regs[i] = 0;
        pcb[j + 1].kernel\_context.regs[29] = 0xa0f000000 - j * 0x1000; //sp
```

```
pcb[j + 1].kernel_context.regs[31] = sched1_tasks[j]->entry_point;
        pcb[j + 1].kernel_context.cp0_status = 0;
        pcb[j + 1].kernel_context.hi = 0;
        pcb[j + 1].kernel_context.lo = 0;
        pcb[j + 1].kernel_context.cp0_badvaddr = 0;
        pcb[j + 1].kernel_context.cp0_cause = 0;
           pcb[j + 1].kernel_context.cp0_epc = 0;
        queue_push(&ready_queue, &pcb[j + 1]);
    for(j = num_sched1_tasks; j < num_lock_tasks + num_sched1_tasks; j++){</pre>
        pcb[j + 1].pid = 2 + j;
        pcb[j + 1].status = TASK_READY;
        pcb[j + 1].cursor_x = 0;
        pcb[j + 1].cursor_y = 0;
        for(i = 1; i \leftarrow 31; i++){
            pcb[j + 1].kernel_context.regs[i] = 0;
        pcb[j + 1].kernel_context.regs[29] = 0xa0f00000 - j * 0x1000; //sp
        pcb[j + 1].kernel_context.regs[31] = lock_tasks[j - num_sched1_tasks]->entry_point;
        pcb[j + 1].kernel_context.cp0_status = 0;
        pcb[j + 1].kernel_context.hi = 0;
        pcb[j + 1].kernel_context.lo = 0;
        pcb[j + 1].kernel_context.cp0_badvaddr = 0;
        pcb[j + 1].kernel_context.cp0_cause = 0;
           pcb[j + 1].kernel_context.cp0_epc = 0;
        queue_push(&ready_queue, &pcb[j + 1]);
   }
}
```

## 锁

#### 锁的初始化

```
void do_mutex_lock_init(mutex_lock_t *lock)
{
    // if the current lock already be locked
    if(lock->status == LOCKED) {
        // set it to block and wait
        do_block(&block_queue[lock->id]);
    }
    lock->status = LOCKED;
}
```

### 锁的获取

```
void do_mutex_lock_acquire(mutex_lock_t *lock)
{
    if(lock->status == LOCKED)
    {
        do_block(&block_queue[lock->id]);
    }
    lock->status = LOCKED;
}
```

```
void do_mutex_lock_release(mutex_lock_t *lock)
{
    if(!queue_is_empty(&block_queue[lock->id])) {
        do_unblock_one(&block_queue[lock->id]);
        lock->status = LOCKED;
    }
    else
        lock->status = UNLOCKED;
}
```