

Chesskell: Embedding a Two-Player Game in Haskell's type system

3rd Year Project Progress Report

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1 Introduction

In 2020, video games are more popular than ever. In the US alone, an ESA report¹ estimates that there are more than 214 million individuals who play games. Considering this, it's surprising how many games are released with major bugs in their software—some of which end up being so notable that news and footage of them appear on mainstream media².

As programming languages have evolved, many have begun to address more errors at compile time. Features similar optional types have been added to languages such as Java and C#, and languages like Rust have pioneered ways of safely handling dynamic allocation through ownership types³. Many compilers now force the developer to handle classes of errors that previously could only be encountered at runtime, such as null pointer exceptions.

However, catching logical errors in imperative languages is almost always done at runtime. Many software systems use runtime features (such as exceptions) to discover and deal with errors and misuse of APIs. Enforcement of invariants and rules is typically done dynamically; if a check fails, an exception is thrown and potentially handled. However, if a programmer forgets to implement such a check, the behaviour is unpredictable. A 2007 study [1] on Java and .NET codebases indicates that exceptions are rarely used to recover from errors, and a 2016 analysis of Java codebases [2] reveals that exceptions are commonly misused in Java.

Recent versions of the *Glasgow Haskell Compiler* (GHC) support programming at the type level, allowing programmers to compute with types in the same way that languages like C or Python compute with values [3], using *type families* [4] [5] that emulate functions at the type-level. These computations run at compile time, before the compiler generates an executable of the source code, allowing programmers to transform logic errors into type errors [6]. The exception misuse described above could be avoided by employing logical invariant checks at the type-level, rather than at runtime.

Since these are relatively recent developments, there are few examples of their usage in complex applications. In this project, we show how to utilise type-level programming features in Haskell in order to model the classic board game Chess in Haskell's type system, ruling out invalid moves at the type-level. A Haskell-embedded Domain-Specific Language (DSL), for describing games of chess, will interact with the type-level model. This Embedded DSL (EDSL) will be modelled on Algebraic Notation, a method of writing down the moves associated with a particular match of chess. We implement the full, official FIDE ruleset for chess.

¹https://www.theesa.com/wp-content/uploads/2020/07/Final-Edited-2020-ESA_Essential_facts.pdf

²<https://www.bbc.co.uk/news/technology-50156033>

³<https://doc.rust-lang.org/book/ch04-01-what-is-ownership.html>

2 Background

In fully dependently-typed languages, such as Idris, there is no clear distinction between types and values. Types’ definitions depend on values, and vice versa; so types and values are interchangeable in usage. Functions can be written to compute with types, as well as values.

Haskell was not designed to be such a language. In Haskell, values have types, and types have kinds. Types of kind `*` (or `Type`) have term-level values; `Int` has kind `*` because it has the inhabitants 0, 1, 2, etc. The `-XDataKinds` extension allows the programmer to define other kinds [3], but these kinds do not have any term-level values.

During compilation of a Haskell program (by GHC), *type erasure* occurs. At runtime, Haskell types have no representation, and so writing runtime computation for type-level constructs is nigh impossible. Much previous academic work has gone into emulating dependent types in Haskell [7], working around type erasure through clever use of GHC extensions. Chesskell uses many of these techniques to blur the line between types and values, allowing compile-time checking of the rules of chess—as we explain below.

2.1 Related Work

There are many publicly available chess-related programs written in Haskell⁴⁵. Many of these are chess engines, which take in a board state and output the move(s) which are strongest, and so therefore perform move checking at the value-level to ensure that the moves that it outputs are valid. Chesskell differs from these in function, in that the end software will not output a list of strong moves; it will simply take in the moves performed, and state whether they are valid chess moves or not. We are not aware of any such type-level chess implementations in Haskell.

There have been allusions to chess at the type-level through solving the N-queens problem in dependently typed languages, such as Idris⁶. The N-queens problem makes use of some chess rules, including the Queen’s attack positions (a straight line in any direction); but as the end goal is not to successfully model a game of chess, it is not a full type-level chess implementation.

An important part of any good chess program is its board representation, since all other parts of the program come from this; move generation, move evaluation, and the entire search space are all defined or influenced by the board representation. A great deal of work has gone into defining memory- or time-efficient chess boards [8] [9], including combinations of multiple representations to yield greater speed [10]. While there is value to be gleaned from examining these representations, Chesskell serves a different purpose; it does not need to search through the valid set of moves to deter-

⁴<https://github.com/mlang/chessIO>

⁵<https://github.com/nionita/Barbarossa>

⁶<https://github.com/ExNexu/nqueens-idris>

mine which are the best, and speed is not its focus. Chesskell’s board representation must be relatively efficient, but it would be naive to expect similar levels of performance from type-level constraint solving computation as from optimised term-level code.

Despite the apparent lack of work on Chess at the type level specifically, there is work on Haskell-Embedded DSLs in other domains to enforce certain behaviour at compile time. Mezzo [11] is an EDSL for music composition, which checks if the described piece of music conforms to some musical ruleset during compilation of the program. This EDSL is most similar to Chesskell in aim, if not in application domain; performing compile-time checks of rulesets that are commonly checked at runtime.

As another example, BioShake [12] is an EDSL for creating performant bioinformatics computational workflows. The correctness of these workflows is checked during compilation, preventing it from being created if certain errors would occur. For bioinformatics workflows especially, this is ideal since many of these workflows are lengthy. BioShake goes further, however; providing tools to allow parallel execution of these workflows. While it is encouraging to see this (and other EDSLs [13]) focus on (and achieve) high performance, Chesskell will have no such focus. This is primarily because only certain parts of the rule-checking process can be parallelised; much of the move handling and order of rule checks must be done sequentially.

3 Current Progress

3.1 Chess Types Overview

The software has a full representation of a game of chess at the type-level, as we explain below. As the system checks the game move by move, the only required representation is the board at the time of checking; containing all pieces on that board.

The most important types are the board, and the individual pieces on it. Both of these types are composed of other types, making them *composite types*; the board is comprised of pieces, and the pieces contain that piece’s team, name, and position.

3.1.1 Team and PieceName

Both of these types are simple algebraic data types, with all inhabitants defined in code. The team defines whether a piece is black or white, and the piece’s name states which chess piece it is; a pawn, a rook, and so on.

The declarations below, combined with the `-XDataKinds` extension, create two things. The first is a value-level representation of `Team` and `PieceName`, where `Black` is a value of type `Team`. The second is a *type-level* representation, where `'Black` is a type of kind `Team`⁷.

```
data Team = Black | White
```

⁷Note that `'Black` has no term-level values, since it does not have kind `*`

```
data PieceName = Pawn
               | Bishop
               | Knight
               | Rook
               | King
               | Queen
```

3.1.2 Position

The type for holding a piece's position makes use of two more types, one for columns and the other for rows. In chess, columns are labelled with letters and rows are labelled with numbers; "a1" is the top-left of the board, and "h8" is the bottom-right. The column type is another simple algebraic data type with all inhabitants listed. The row type is a type-level implementation of Peano natural numbers.

```
data Column = A | B | C | D | E | F | G | H
data Nat where -- Used for rows
  Z :: Nat
  S :: Nat -> Nat
```

The Position type itself is a *product type*, a type whose cardinality⁸ is equal to the product of the types it takes in. Since Position takes in a type of cardinality 8 (Column) and a type of infinite cardinality (Nat), it has infinite cardinality, even though only 64 of them are valid chess positions.

```
data Position where
  At :: Column -> Nat -> Position
```

3.1.3 The Pieces

Each piece contains information relevant for rule checking: that piece's team, name, and an information type. The information type contains a Nat and a Position, to represent the number of moves that piece has taken, and its current position on the board. Recording the number of moves the piece has taken is important for several rules in chess, including castling and *En Passant*. As such, it is included in the PieceInfo type.

The PieceInfo type was created separate from the plain Piece type so that if any further information was needed, it could be added without breaking existing Piece pattern-match definitions (though no such additional information was required).

```
-- Piece info type, containing general information
data PieceInfo where
  Info :: Nat -> Position -> PieceInfo

-- Piece type, containing the key information for a piece
```

⁸A type's cardinality is the number of inhabitants it has.

```
data Piece where
  MkPiece :: Team -> PieceName -> PieceInfo -> Piece
```

3.1.4 The Board

At the type level, the most common way of storing n elements is with the built-in list type (similar to the value-level). However, a singly-linked list has issues; how can we ensure that the chess board is the appropriate size (an 8x8 grid) without a length check each move? This would take at least 56 additions, since list length is computed recursively; as well as 7 more addition operations to put together the list lengths.

Performing such a check each move would be wasteful. Instead, length-indexed vectors are used. `Nat` values specify the length of a vector, resulting in a type error if the vector is any other length. As with most things in Haskell, we use recursive definitions; an empty vector has length 0, and you express a vector of length $(n + 1)$ by pushing an element to the front of a vector of length n .

```
data Vec (n :: Nat) (a :: Type) where
  VEnd    :: Vec Z a
  (:->)   :: a -> Vec n a -> Vec (S n) a
```

For example, a vector of 3 boolean types has the type: `Vec (S (S (S Z))) Bool`. The board can be expressed using these vectors—a vector of 8 vectors of 8 pieces. (We use `Maybe Piece` instead, since a board square does not necessarily contain a piece.)

```
type Eight = (S (S (S (S (S (S (S (S Z)))))))
type Row   = Vec Eight (Maybe Piece)  -- Row of eight Maybe
                                              Pieces
type Board = Vec Eight Row              -- Grid of eight rows
```

Although this is the main board type, it is augmented with a `BoardDecorator`, so named because the intention is similar to the decorator design pattern [14] (with the exception that subclassing and superclassing are not features of Haskell). `BoardDecorator` should be used instead of `Board`, since it provides additional information:

- The last team to move;
- The last position moved to;
- The White and Black king positions, stored as a tuple;
- The number of moves in the game thus far.

Previous versions of the program, to find the king positions, would pass repeatedly over the `Board`. Having their positions available in the decorator saves the system from making these passes. While there is the overhead that the `BoardDecorator` must be updated at each move, the code is much conceptually clearer with the use of the decorator.

```

data BoardDecorator where
  Dec :: Board           -- Chess board
  -> Team                -- Last team moved
  -> Position            -- Last position moved to
  -> (Position, Position) -- King positions
  -> Nat                 -- No. of moves so far
  -> BoardDecorator

```

3.2 Checking Chess Rules

We perform all rule-checking with type families that take in the `BoardDecorator` (among other arguments) and output either a correct result, or a type error. Rule-checking occurs during movement, since the chess game effectively hangs while players do not move.

3.2.1 Type families in Haskell

Although Haskell allows type-level computation via type families, these type families can't be partially applied like normal Haskell functions can. However, Li-yao Xia's work on First Class Families⁹ defines a kind of type-level interpreter, allowing similar behaviour to partial application.

This means that the board is mappable, performing computation on the pieces therein without writing a new function for each map. Incrementing a piece's move count, or changing the name of a piece, is therefore much easier with First Class Families than it would otherwise be with regular type families.

3.2.2 Move Lists

Each piece, depending on its team and name, can move to a certain number of spaces. For instance, a King can move a single space in any direction. The `PieceMoveList` type family formalises this, returning a list of spaces that a piece can move to (given a piece and board decorator).

```

data PieceMoveList :: Piece -> BoardDecorator -> Exp [Position]
-- If it's the pawn's first move (i.e. it has a move count of
-- 0), allow it to move two spaces, otherwise move one
type instance Eval (PieceMoveList (MkPiece team Pawn info)
  boardDec) = Eval (If (Eval ((IsZero . GetMoveCount) info))
  (PawnStartMove (MkPiece team Pawn info) boardDec)
  (PawnPostStart (MkPiece team Pawn info) boardDec))

-- Bishops move diagonally

```

⁹<https://github.com/Lysxia/first-class-families>


```

type instance Eval (PieceMoveList (MkPiece team Bishop info)
    boardDec) = Eval (AllReachableDiag team boardDec (Eval
    (GetPosition info)))

-- Knights move in an L-shape
type instance Eval (PieceMoveList (MkPiece team Knight info)
    boardDec) = Eval (AllReachableGivenList team boardDec (Eval
    (GetAllKnightPositions (Eval (GetPosition info)))))

-- Rooks move in straight lines
type instance Eval (PieceMoveList (MkPiece team Rook info)
    boardDec) = Eval (AllReachableStraightLine team boardDec
    (Eval (GetPosition info)))

-- Queens can move in any direction
type instance Eval (PieceMoveList (MkPiece team Queen info)
    boardDec) = Eval (AllReachableLineAndDiag team boardDec (Eval
    (GetPosition info)))

-- Kings can move a single space in any direction
type instance Eval (PieceMoveList (MkPiece team King info)
    boardDec) = KingMoveList (MkPiece team King info) boardDec

```

There is another function, `PieceAttackList`, which is similar to `PieceMoveList` but gives the list of spaces that a piece can attack. This is different from a piece's move list in a handful of cases (for instance, pieces can attack the opposite team's King, but cannot move to its position and take it).

3.2.3 Moving the Pieces

The core movement function, aptly named `Move`, performs a series of checks on the board state, moves the piece, and then performs further checks. The actual movement code is omitted here for brevity; but all moves except for pawn promotion are complete.

```

-- MoveNoChecks moves the piece depending on its' type,
-- handling any relevant side effects
data Move :: Position -> Position -> BoardDecorator -> Exp
    BoardDecorator
type instance Eval (Move fromPos toPos boardDec) = Eval ((
    CheckNoCheck (GetMovingTeam boardDec) . -- Post-move checks
    MoveNoChecks fromPos toPos .             -- Move the piece
    NotTakingKingCheck toPos .               -- Pre-move checks
    CanMoveCheck fromPos toPos .
    NotTakingOwnTeamCheck toPos .
    NotSamePosCheck fromPos toPos .
    NotLastToMoveCheck fromPos .

```

```
TeamCheck fromPos) boardDec)
```

3.3 EDSL

3.3.1 Proxies and Singletons

Type-level information is not naturally available to Haskell’s term-level facilities, due to type erasure. As such, there are various methods to simulate computation with types at runtime.

Proxy types are parameterised with kind-polymorphic type variables [3]. A `Proxy` value can have type `Proxy Int` or `Proxy 'Black`, for example. Proxies are used to make runtime board decorator information available at the type level in Chesskell.

Singleton types mimic dependent types [7]; each type has a single, unambiguous value, and vice versa. We use this value to pass around type information at the term-level. Chesskell uses singletons to perform computation with positions and piece names.

3.3.2 EDSL Implementation

A Continuation Passing Style [15] scheme forms the foundation for the EDSL, with inspiration taken from Dmitriy Szamozvancev’s Flat Builders pattern [16]. The core idea is value transformation through a series of continuation function applications, until the final continuation function returns a value.

```
-- Starts with a value of type t, gets a continuation of type
   (t -> m), and applies it to the value
type Spec t = forall m. (t -> m) -> m

-- Converts a value of type s into a value of type t, passing
   that t-value into a continuation
type Conv s t = s -> Spec t

-- Ends the continuation stream
type Term t r = t -> r
```

Type-level computations within the EDSL are achieved through *type applications* [17], which allow the programmer to specify the exact type of a kind-polymorphic type variable. Using a combination of singletons, proxies, kind signatures, and type applications, the value-level Haskell code for the EDSL can have specific type variables and involve type family application. Essentially, the term-level EDSL can involve type-level rule checking.

The chess game starts with a `Proxy` value, with its type containing a `BoardDecorator` type. Continuations are applied, transforming that value, until the chess game ends (or a rule is broken). The core continuations are named after the pieces, such as `pawn`

and king. Each of them takes in an `SPosition`, a singleton version of `Position`. We define a new datatype, `MoveArgs`, in order to simplify the process of passing information between the continuations.

```
-- StartDec is a board decorator holding all the information
for the start of a chess game
chess :: Spec (Proxy StartDec)
chess cont = cont (Proxy @StartDec)

data MoveArgs where
    MA :: BoardDecorator -> Position -> PieceName -> Position
    -> MoveArgs

-- Passes a created MoveArgs type constructor to the
continuation
pawn :: Proxy (b :: BoardDecorator) -> SPosition fromPos ->
    Spec (Proxy (MA b fromPos 'Pawn'))
pawn (dec :: Proxy b) (from :: SPosition fromPos) cont = cont
    (Proxy @(MA b fromPos Pawn))
```

The next continuation, `to`, takes in another `SPosition`, performs the move computation, puts the resulting board decorator into a `Proxy` type, and passes that `Proxy` into the continuation given.

```
-- Moves a piece from one position to another
to :: Proxy (MA (b :: BoardDecorator) (fromPos :: Position) (n
    :: PieceName)) -> SPosition toPos -> Spec (Proxy (Eval
    (MoveWithStateCheck n fromPos toPos b)))
to (args :: Proxy (MA (b :: BoardDecorator) (fromPos ::
    Position) (n :: PieceName))) (to' :: SPosition toPos) cont =
    cont (Proxy @(Eval (MoveWithStateCheck n fromPos toPos b)))
```

The final relevant definition is of `end`, which ends the chess game (and the continuation stream).

```
end :: Term (Proxy (b :: BoardDecorator)) (Proxy (b ::
    BoardDecorator))
end = id
```

Using the above continuations, we can lay out a chess game, move by move. Consider a short game, involving a 3-move checkmate by White. The game can be laid out as below in the EDSL, and will compile successfully:

```
-- Describes a game with a checkmate by White in 3 moves
game = chess
    pawn _e2 to _e4      -- White begins
    pawn _f7 to _f5      -- Black move
    queen _d1 to _f3
    pawn _g7 to _g5
```

```

        queen _f3 to _h5 -- Checkmate!
    end

```

However, should Black attempt to move after checkmate, an error will be logged, since no moves by Black will be valid:

```

-- Below results in the following type error:
-- * The Black King is in check after a Black move. This is
  not allowed.
-- * When checking the inferred type
--   game :: Data.Proxy.Proxy (TypeError ...)
game = chess
    pawn _e2 to _e4 -- White begins
    pawn _f7 to _f5 -- Black move
    queen _d1 to _f3
    pawn _g7 to _g5
    queen _f3 to _h5 -- Checkmate!
    pawn _g5 to _g4 -- Invalid move by black - Black
    King is in checkmate
end

```

Or, should White attempt an impossible move in the middle of the game, such as moving a Queen through another piece, a different type error will occur:

```

-- Below results in the following type error:
-- * There is no valid move from D1 to D3.
-- The Queen at D1 can move to: E2, F3, G4, H5
-- * When checking the inferred type
--   game :: Data.Proxy.Proxy (...)
game = chess
    pawn _e2 to _e4 -- White begins
    pawn _f7 to _f5 -- Black move
    queen _d1 to _d3 -- White Queen moves through a
    Pawn - not allowed
    pawn _g7 to _g5
    queen _f3 to _h5 -- Checkmate!
end

```

3.3.3 EDSL Board Creation

There are additional statements to help create an initial chess board for easy testing. A short example is below:

```

-- A simple example board created for testing castling code
example
    = create -- create chess board
        put _Wh _R at _a8 -- Set a white rook at a8

```

```

    put _Wh _R at _h1
    put _Bl _P at _h7  -- Set a black rook at h7
    lastTeam _Bl      -- Set black as last team moved
startMoves           -- Begin chess game
    rook _a8 to _a2
    pawn _h7 to _h6
    rook _h1 to _h2
end

```

3.3.4 Pawn Promotion

The EDSL function `becomes` handles Pawn promotion, transforming a Pawn into another piece after it has moved to the opposite side of the board. Since the whole chess game is expressed in continuation-passing style, `becomes` simply edits a Pawn's `PieceName` in place.

```

game = chess
    -- ...
    pawn _a7 to _a8 becomes _Queen
    -- ...
end

```

We can combine this with the board creation functionality to create a simple game which compiles:

```

game = create
    put _Wh _P at _a7
    startMoves
    pawn _a7 to _a8 becomes _Queen
end

```

However, the below code will also compile, when it should not, since according to the FIDE 2018 Laws of Chess¹⁰, promotion is not optional, and must occur.

```

game = create
    put _Wh _P at _a7
    startMoves
    pawn _a7 to _a8  -- No promotion - should have type
error
end

```

3.3.5 EDSL Issues

GHC had issues compiling early versions of the EDSL, due to the initial definition of `chess`; the programmer did not manually create `StartDec`, but pieced it together through a series of type family applications:

¹⁰<https://handbook.fide.com/chapter/E012018>

```
-- Create StartDec with expensive type family computations
type StartDec = MakeDecorator (...)
```

This early version of `StartDec` was used to set up the chess game, with attempts at two definitions of `chess`; one using a type application, and the other using a type signature.

```
-- Early definition of chess with type application
chess :: Spec (Proxy StartDec)
chess cont = cont (Proxy @StartDec)

-- Version of chess with type signature
chess :: Spec (Proxy StartDec)
chess cont = cont (Proxy :: Proxy StartDec)
```

The version of `chess` that used a type application would compile without issues, but would cause lengthy pauses at runtime when used. Conversely, when compiling the type signature version, GHC would consume high amounts of memory, and would either crash or hang for multiple minutes. This behaviour is unusual, so we filed a GHC bug report¹¹. As a work-around until the bug is fixed, `StartDec` is manually written out in full in the source Haskell file, allowing compilation and usage of either definition with no issues.

Additionally, adding the move number to the type error¹² causes a spike in compile time; not as excessive as above, as GHC does not crash, but still doubling or tripling compile time. The move number will only be re-added to the error if this problem can be fixed or worked around, since it would greatly slow the testing of a large number of games.

3.4 Testing

The test suite is a mixture of HSpec behaviour-driven tests [18] and module declarations, where the HSpec tests cover many of the declared type families (such as `Move` and `PieceMoveList`).

The original plan was to have all tests be HSpec tests; many of these HSpec tests make use of the type-level testing facilities laid out in other works [16]. However, running the test suite on EDSL statements, as well as some of the expensive type families (such as those involving castling) causes excessive memory usage, exceeding 27 GB in some cases.

Work on a solution is ongoing, but for now, we express EDSL tests as high-level declarations in a Haskell module. The tests are run by seeing whether these declarations compile or not (and if not, what error they throw). This will hopefully only be temporary, but it should suffice as a rudimentary (if slow) method of integration testing.

¹¹<https://gitlab.haskell.org/ghc/ghc/-/issues/18902>

¹²e.g. `Error at move 1: Pawn cannot move to...`

4 Next Steps

As the project's main feature set is nearly complete, the remaining time will be spent testing, evaluating, and improving the project. Full-length chess games will be expressed in the EDSL, to determine if the program compiles correct chess games. Then, we will introduce subtle errors to those games, to test that incorrect games fail to compile. Of course, the code itself is only part of the project; work on the dissertation should begin early next term/year.

4.1 Chess Rules

At the moment, the only incomplete chess rule is *mandatory promotion*. Promotion in the EDSL is optional, meaning a Pawn can reach the opposite end of the board and remain a Pawn. The current implementation of promotion does not involve the `Move` function, but since pawn promotion must occur by the end of its' move, the `Move` function must include this functionality.

4.2 EDSL

The only major addition to the EDSL is mandatory promotion, as we explain above. However, a few additional helper functions or syntax changes may be helpful during development, and so the EDSL is subject to change.

4.3 Performance

The current memory usage and compile time of the EDSL is acceptable, but only short (<10 moves) chess games have been described and tested thus far. Memory usage and compile time are likely to rise both with the length of the chess game, and the number of games tested. Memory usage optimisations may become necessary as testing continues.

4.4 Revised Timetable/Plan

4.4.1 Weeks 9-10: 30th November to 13th December

These two weeks will be spent implementing mandatory promotion, and performing integration/system testing. A curated data set of chess games will be expressed in the EDSL and put through the GHC compiler. Some of these games will be valid, and some will be invalid. Only valid games should compile, and invalid games should fail to compile with an appropriate type error.

4.4.2 Weeks 11-14: 14th December 2020 to 10th January 2021

This is the first section of allocated empty space. If the test suite is posing any issues (with regards to compile time and memory usage) then this additional month should help with solving them. But should the testing be going well, with promotion complete, then this extra time will be spent either starting early on the dissertation, or adding extensions to the project to further explore the modelling of games at the type-level.

4.4.3 Weeks 15-18: 11th January to 7th February

Writing the dissertation is planned for this stretch of time; it will be planned out, section by section, while gathering any relevant graphs, figures, and citations.

The module CS324 Computer Graphics has a coursework due on the 20th January; as such, planning the dissertation could take longer than expected. However, since a month is set aside for just planning, delays are unlikely.

4.4.4 Weeks 19-22: 8th February to 7th March

Once the previous detailed planning stage is complete, writing shall begin; with a detailed enough plan, this section should not take longer than a month. We will complete an initial draft by the 7th of March.

The project itself will be evaluated during this period. We will examine the cleaned and completed code for any subtle bugs. Work on the project presentation will also start in earnest, writing out a script.

4.4.5 Weeks 23-24: 8th March to 21st March

These two weeks are more empty space; set aside to act as a buffer for delays in dissertation writing. The project presentation occurs on the 19th March; so this empty space will ideally be spent completing the presentation and rehearsing it.

4.4.6 Weeks 25-29: 22nd March to 25th April

Drafting, re-drafting, and refining the dissertation (with the help of the supervisor) will take place during this time. This is the final stretch, and will be spent ensuring that the final piece of writing is as good as it can be.

Revision for examinations is key during this period; however, since the dissertation should be complete and this period is for evaluation and not the main bulk of writing, there should be ample time.

5 Risks

The biggest risk going into the project was that it was unachievable; however, that risk has not materialised—the implementation is almost complete. The largest project-related risks left are the memory usage and compile time issues.

There are, of course, other risks not directly associated with the project; in 2020, there is a large-scale pandemic going on, which could pose considerable personal risk. However, by following Government and University guidelines, the chances of contracting COVID-19 will be kept to a minimum, and sudden illness should not interrupt the project.

6 References

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