# COMPARISION OF DIFFERENT SORTING ALGORITHMS IN RISCV

SOURABH.T PRUTHVIK KASHYAP S PES1UG20EC339 PES1UG20EC326

#### INTRODUCTION

1. INSERTION SORT

2. SELECTION SORT

3. BINARY SEARCH

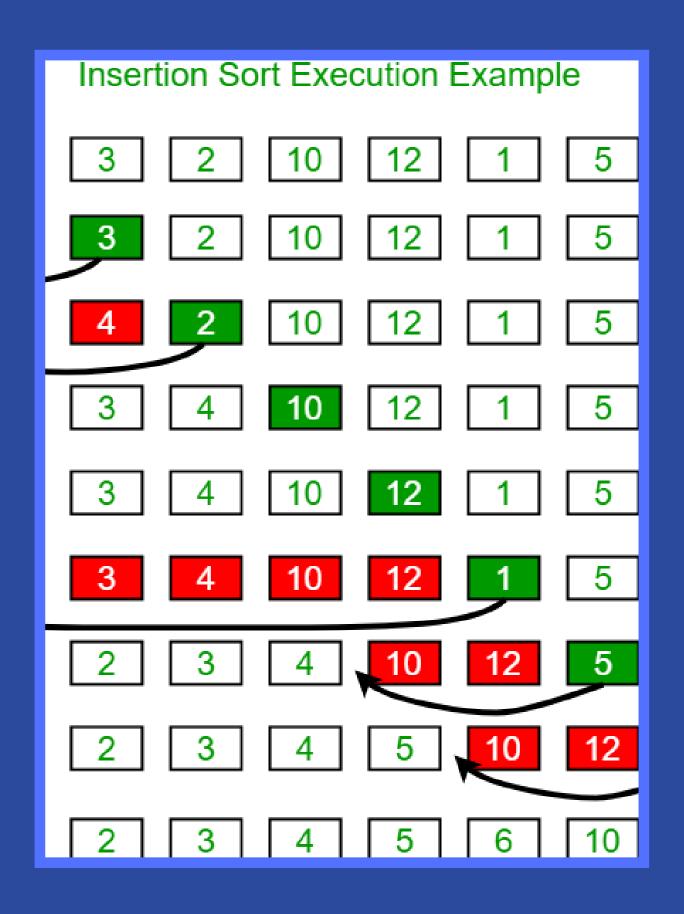
4. QUICK SORT

5. MERGE SORT

## INSERTION SORT

- THIS ALGORITHM IS ONE OF THE SIMPLEST ALGORITHM WITH SIMPLE IMPLEMENTATION
- ASICALLY, INSERTION SORT IS EFFICIENT FOR SMALL DATA VALUES
- INSERTION SORT IS ADAPTIVE IN NATURE,

   I.E. IT IS APPROPRIATE FOR DATA SETS
   WHICH ARE ALREADY PARTIALLY SORTED
   TIME COMPLEXITY O(N²)



## ALGORITHM

- Iterate from arr[1] to arr[N] over the array.
- Compare the current element (key) to its predecessor.
- If the key element is smaller than its predecessor, compare it to the elements before. Move the greater elements one position up to make space for the swapped element.

#### 2

#### SELECTION SORT



### ALGORITHM

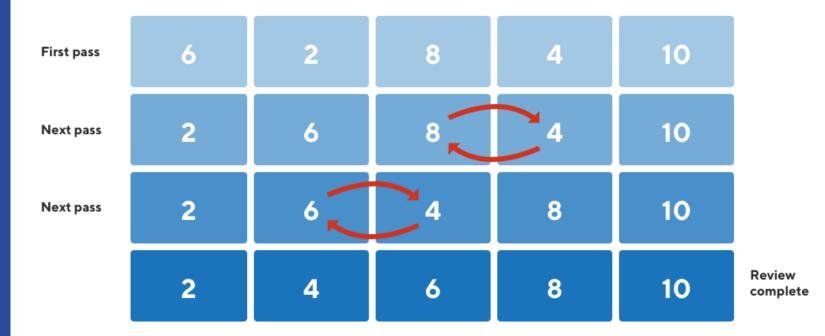
- Initialize minimum value(min\_idx) to location 0.
- Traverse the array to find the minimum element in the array.
- While traversing if any element smaller than min\_idx is found then swap both the values.
- Then, increment min\_idx to point to the next element.
- Repeat until the array is sorted.

#### 3

#### BUBBLE SORT

- BUBBLE SORT IS THE SIMPLEST SORTING ALGORITHM THAT WORKS BY REPEATEDLY SWAPPING THE ADJACENT ELEMENTS IF THEY ARE IN THE WRONG ORDER.
- THIS ALGORITHM IS NOT SUITABLE FOR LARGE DATA SETS AS ITS AVERAGE AND WORST-CASE TIME COMPLEXITY IS QUITE HIGH.
  - TIME COMPLEXITY  $O(N^2)$

#### **Bubble Sort**



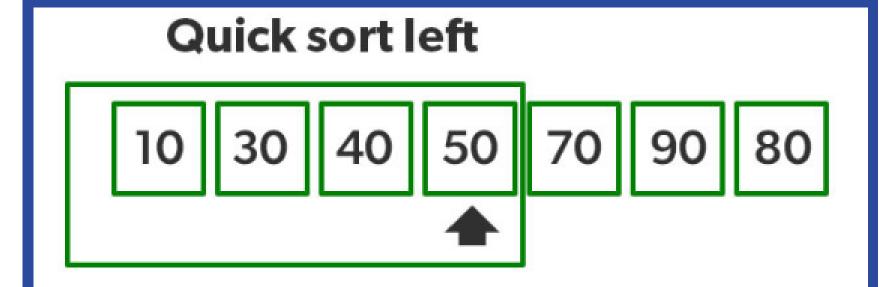
## ALGORITHM

- Run a nested for loop to traverse the input array using two variables i and j, such that 0 ≤ i < n-1 and 0 ≤ j < n-i-</li>
   1
- If arr[j] is greater than arr[j+1] then swap these adjacent elements, else move on
  - Print the sorted array

# 10 80 30 90 40 50 70 Pivot

## QUICK SORT

- QUICKSORT IS A DIVIDE AND CONQUER ALGORITHM.
- IT PICKS AN ELEMENT AS A PIVOT AND PARTITIONS THE GIVEN ARRAY AROUND THE PICKED PIVOT.
- THERE ARE MANY DIFFERENT VERSIONS
   OF QUICKSORT THAT PICK PIVOT IN
   DIFFERENT WAYS.
  - TIMING COMPLEXITY O(N\*LOGN)





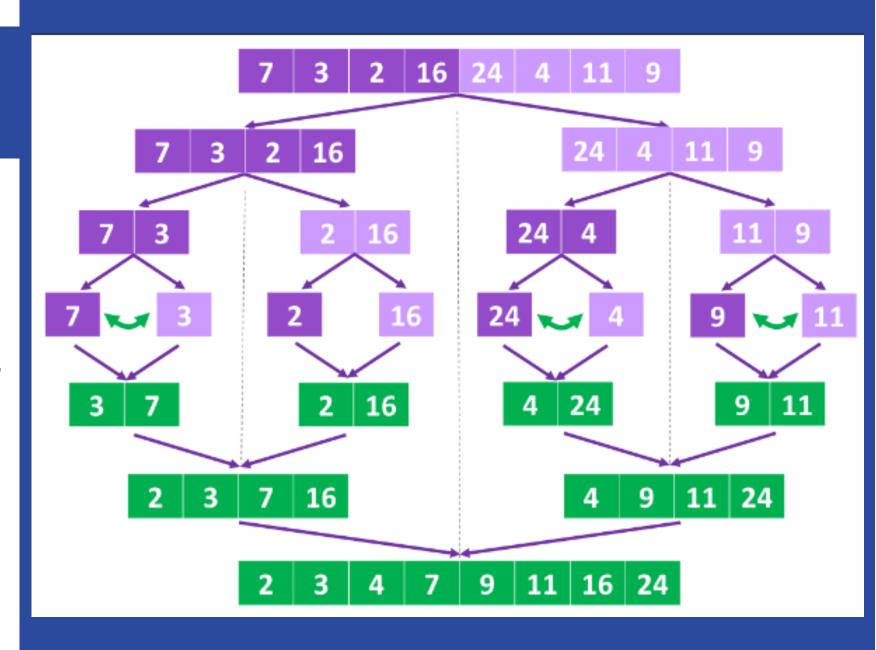
## PARTITION ALGORITHM

- The logic is simple, we start from the leftmost element and keep track of the index of smaller (or equal to) elements as i.
- While traversing, if we find a smaller element, we swap the current element with arr[i]. Otherwise, we ignore the current element.

#### 5

#### MERGE SORT

- THE MERGE SORT ALGORITHM IS A SORTING ALGORITHM THAT IS BASED ON THE DIVIDE AND CONQUER PARADIGM.
- IN THIS ALGORITHM, THE ARRAY IS INITIALLY DIVIDED INTO TWO EQUAL HALVES AND THEN THEY ARE COMBINED IN A SORTED MANNER.
  - TIMING COMPLEXITY O(N\*LOGN)

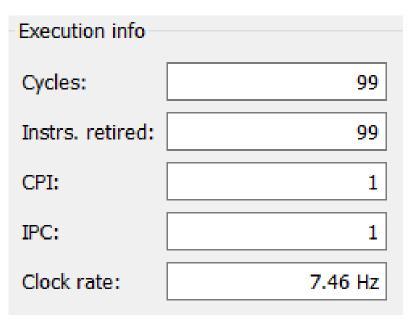


## ALGORITHM

- It as a recursive algorithm continuously splits the array in half until it cannot be further divided.
- If the array has multiple elements, split the array into halves and recursively invoke the merge sort on each of the halves.
- Finally, when both halves are sorted, the merge operation is applied.
- Merge operation is the process of taking two smaller sorted arrays and combining them to eventually make a larger one.

#### COMPARISION - SINGLE CYCLE

#### 1.INSERTION SORT



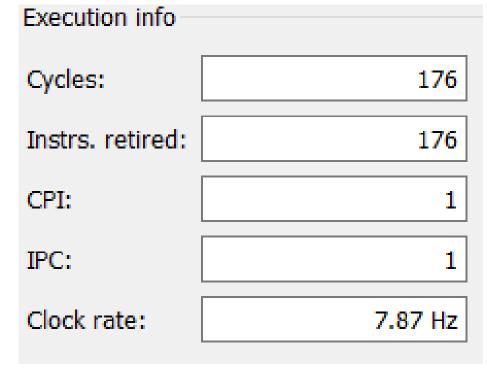
#### 2.SELECTION SORT

Execution info	
Cycles:	127
Instrs. retired:	127
CPI:	1
IPC:	1
Clock rate:	10.75 Hz

#### 3. BUBBLE SORT

Execution info	
Cycles:	174
Instrs. retired:	174
CPI:	1
IPC:	1
Clock rate:	50.58 Hz

#### **4.QUICK SORT**

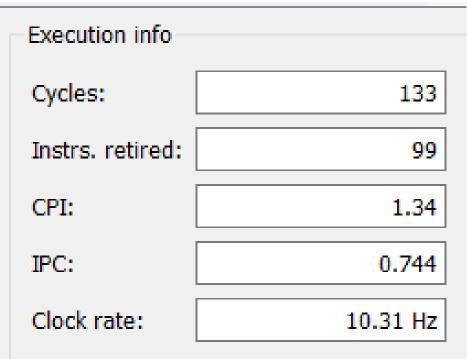


#### **5.MERGE SORT**

Execution info	
Cycles:	263
Instrs. retired:	263
CPI:	1
IPC:	1
Clock rate:	9.35 Hz

## COMPARISION - 5 CYCLE(S)

#### 1.INSERTION SORT



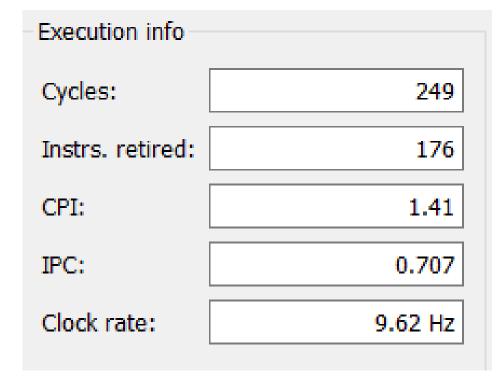
#### 2.SELECTION SORT

Execution info	
Cycles:	197
Instrs. retired:	127
CPI:	1.55
IPC:	0.645
Clock rate:	9.26 Hz

#### **3.BUBBLE SORT**

Execution info	
Cycles:	242
Instrs. retired:	174
CPI:	1.39
IPC:	0.719
Clock rate:	124.61 Hz

#### 4.QUICK SORT



#### **5.MERGE SORT**

Execution info	
Cycles:	373
Instrs. retired:	263
CPI:	1.42
IPC:	0.705
Clock rate:	271.74 Hz

## BUBBLE VS QUICK SORT FOR LARGE DATA SET

#### **BUBBLE SORT**

- Compares the current element with the next on n-1 time.
- Relatively less efficient.
- Timing complexity O(n2)

Execution info	
Cycles:	760
Instrs. retired:	610
CPI:	1.25
IPC:	0.803
Clock rate:	983.12 Hz

#### **QUICK SORT**

- Compares the pivot with the rest of the elements n-1 times
- Relatively more efficient.
- Timing complexity O(n\*logn)

Execution info	
Cycles:	1206
Instrs. retired:	839
CPI:	1.44
IPC:	0.696
Clock rate:	3.36 KHz

## THANK YOU