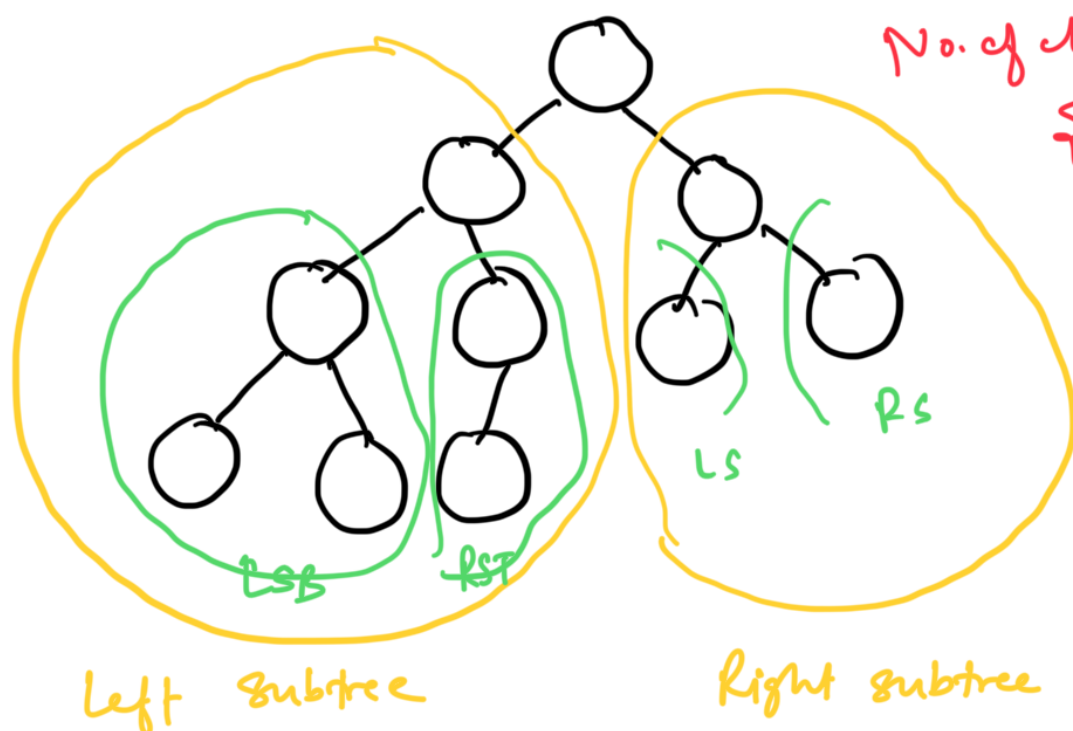


## Binary Tree -

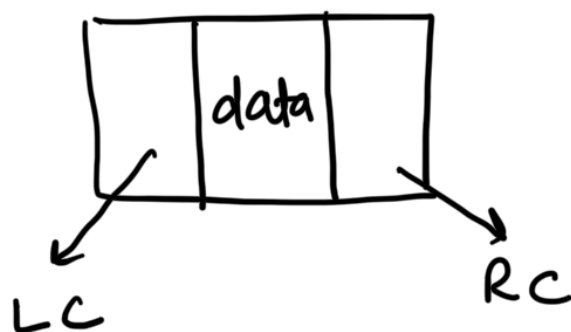
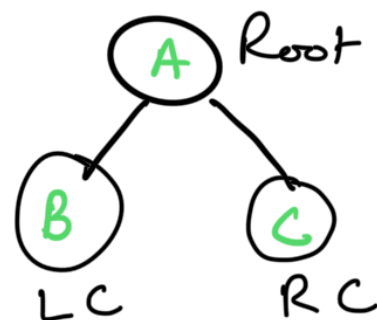
- A BT is a hierarchical structure in which each node has at most two children, referred as left child & right child



### Properties of BT:

#### 1) Basic structure

- Data (value or key)
- LC (Reference pts to LC)
- RC (Reference pts to RC)

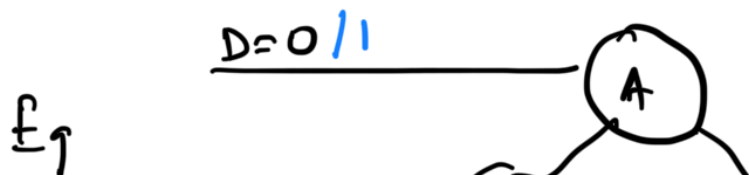


#### 2) Maximum nodes -

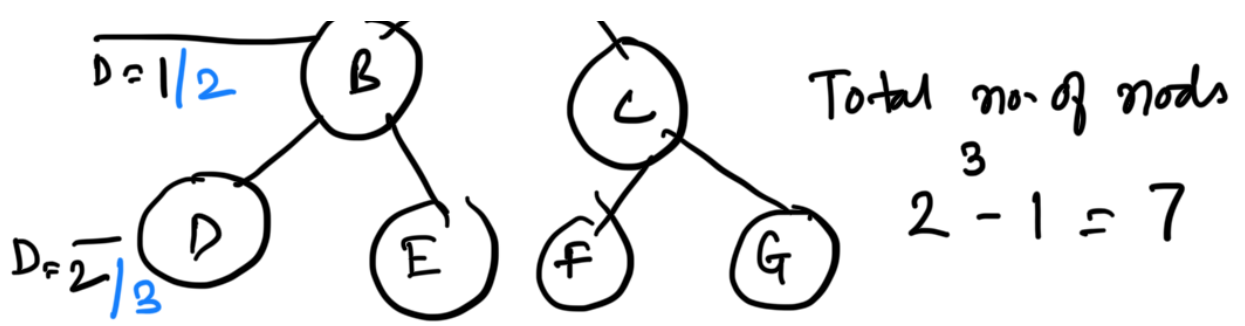
Imp!  
BT

- Max. no. of nodes at level  $L \Rightarrow 2^L$

- Max no. of nodes in a BT of height  $h$  is  $2^h - 1$



Depth = 3



Full BT The max. no. of nodes of level  $i$  of BT is  $2^i, i \geq 0$

$$\text{Max} = 2^i, i \geq 0$$

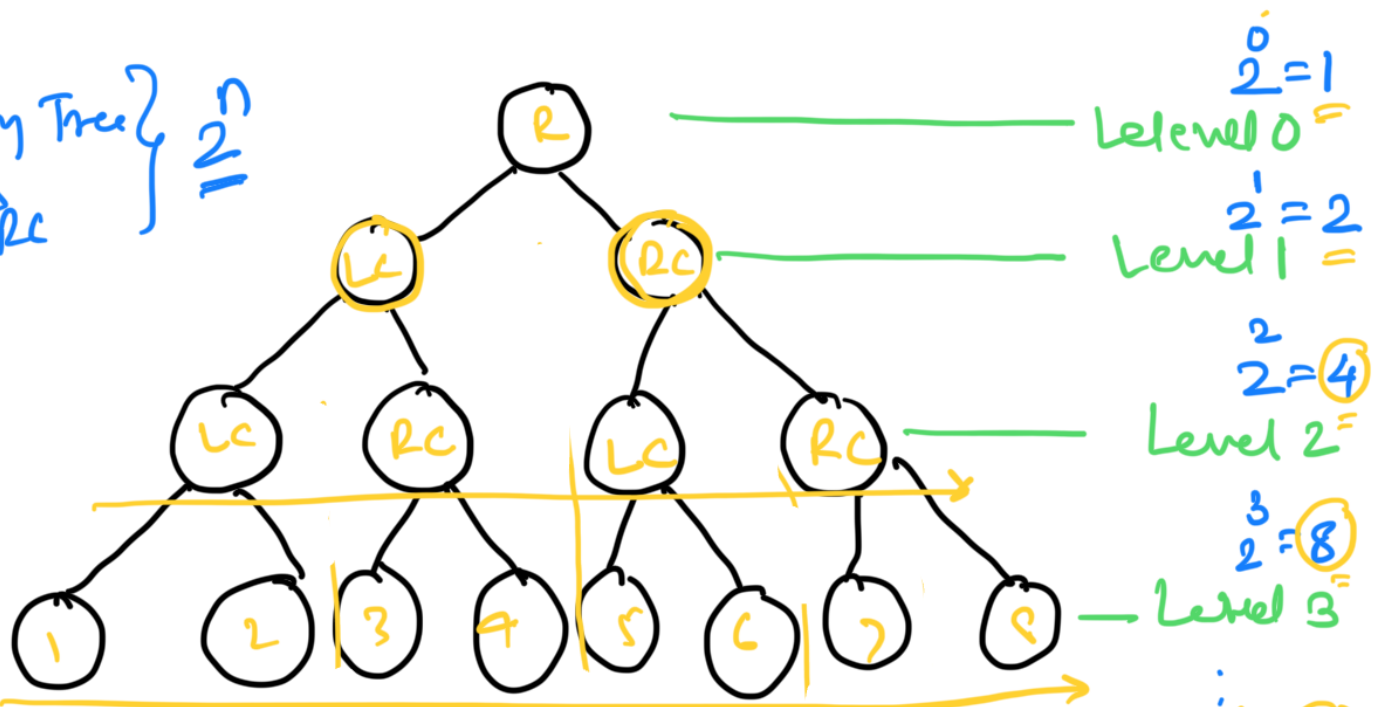
$D=0$

Full BT The max no. of nodes of level  $i$  of BT is  $2^{i-1}, i \geq 1$

$$\text{Max} = 2^{i-1}, i \geq 1$$

$D=1$

Binary Tree }  $2^n$   
LC RC



All 8 nodes present

Full BT

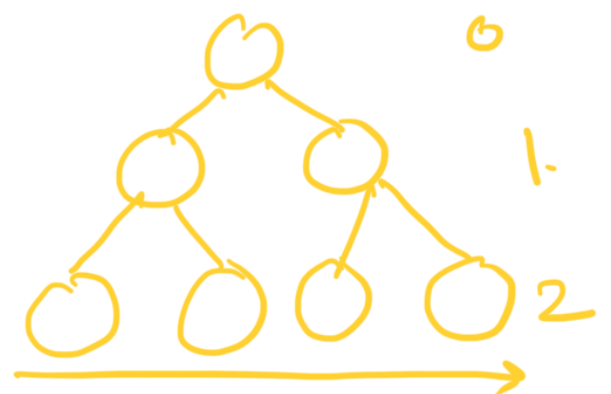
Level 0  $2^0 = 1$   
Level 1  $2^1 = 2$   
Level 2  $2^2 = 4$   
Level 3  $2^3 = 8$   
Level 4  $2^4 = 16$   
Level  $n$   $2^n$

$$BT = O(\log n)$$

BT = Level =  $2^n$  nodes  $\rightarrow$  present

Full BT

$$\text{level} = 2 \rightarrow 2^2 = 4$$



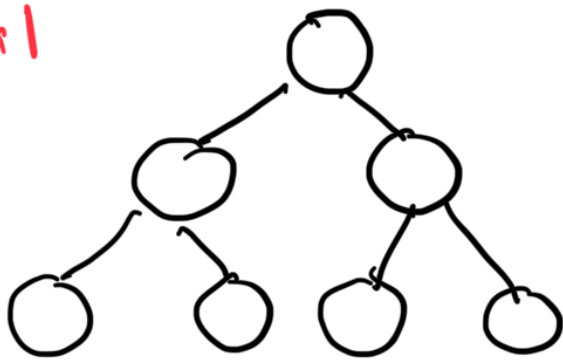
## ② Height of Tree $\rightarrow$

The height of tree is the number of edges in the longest path from the root to leaf node.

The height of tree with  $n$  nodes is at most  $\underline{n-1} \rightarrow \underline{O(\log n)}$  Time Complexity

### Type of BT:

Ex 1

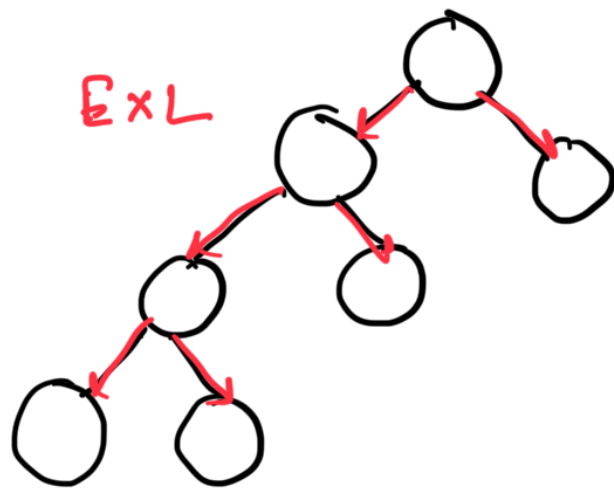


Full BT

1) Full BT  $\rightarrow$  BT with complete levels

- Every node has 0 or 2 children

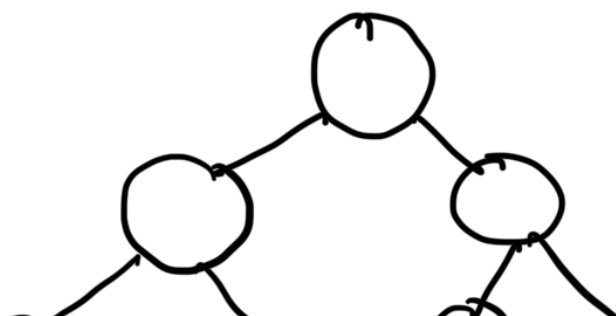
Ex 2



Full BT

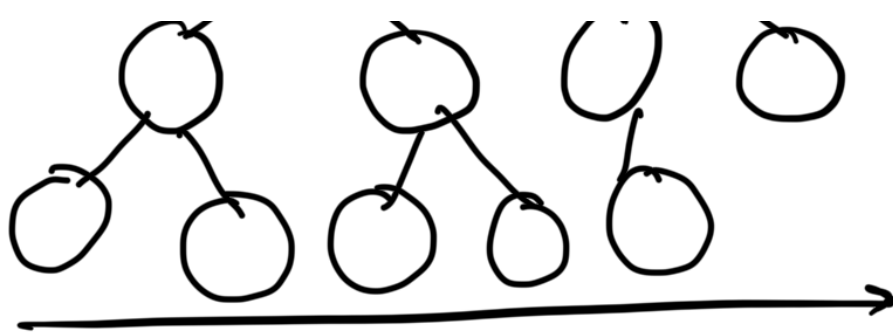
2) Complete BT : - BT in which at every

level, except possibly the last, has to be filled and all nodes are as far as left as possible



Root  
LC  
RC

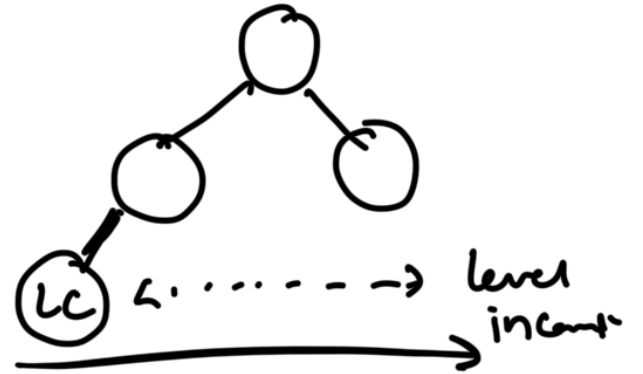
$\{0, 1, 2\}$



Complete BT

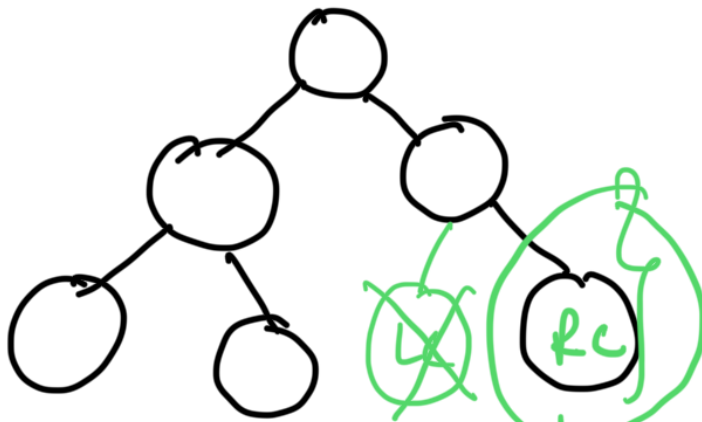
↓  
LC

Ex 2



Complete BT

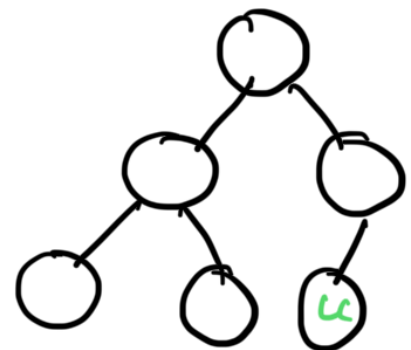
Ex 3



X Complete BT

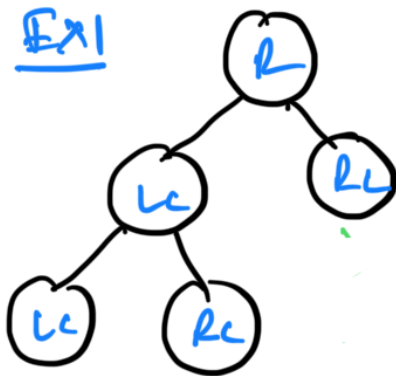
Incomplete BT

Ex 4



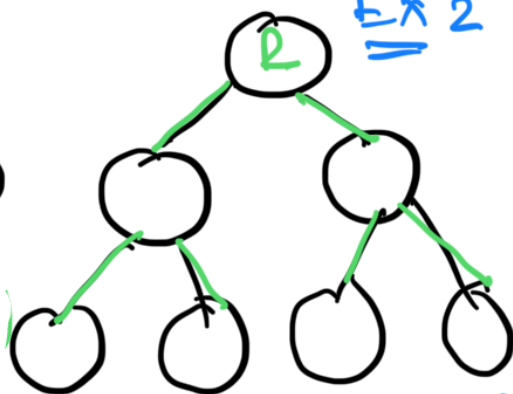
Complete BT

Ex 1

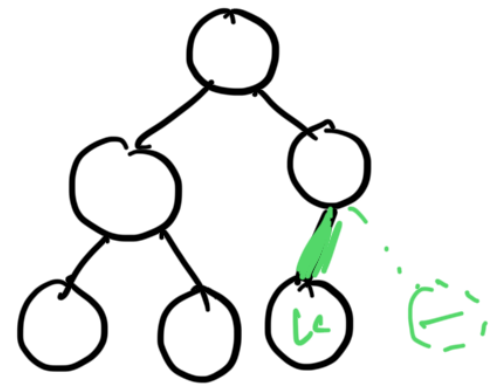


BT  
Full BT  
CBT ✓

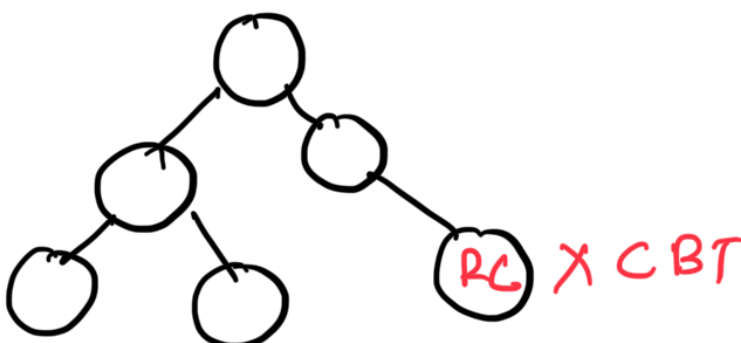
Ex 2



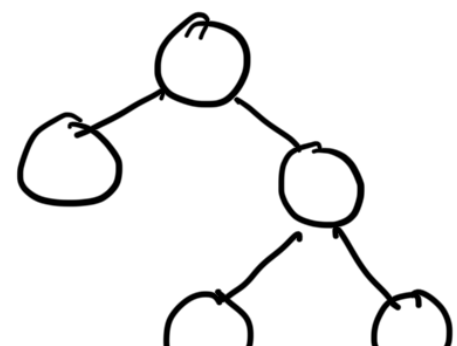
BT  
FBT  
CBT  
Perfect BT



BT  
FBT X {0, 2}  
CBT {0, 1, 2}  
↓  
LC



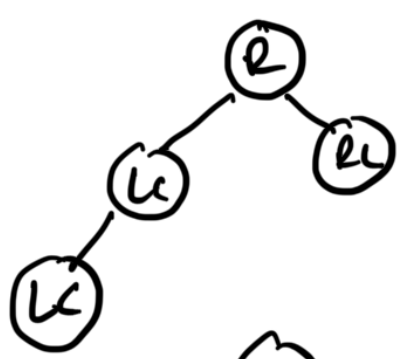
X CBT



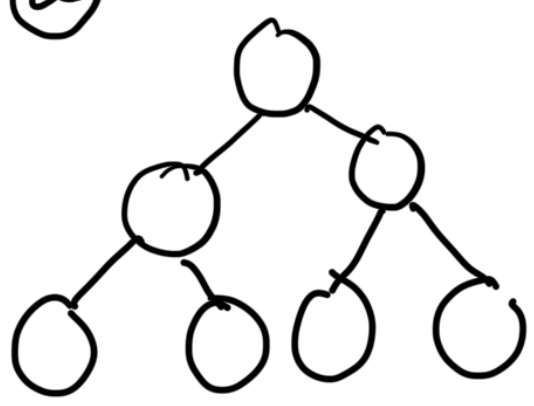
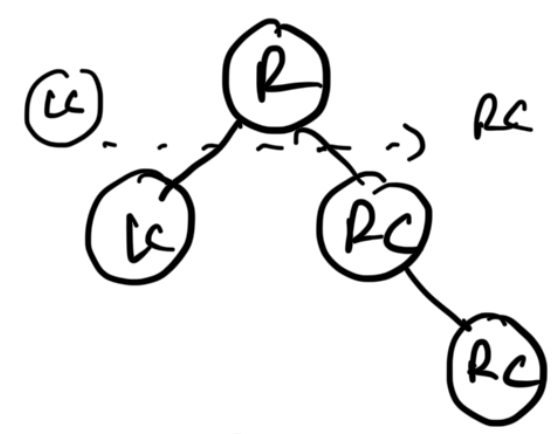
BT  
 FBT  $\times \{0, 2\}$   
 CBT  $\times \{0, 1, 2\}$   
 ↓  
RC  $\times$

ICBT

BT  
 FBT  $\{0, 2\}$   
 CBT  $\times \{0, 1, 2\}$   
 (left  $\times$  first preference)

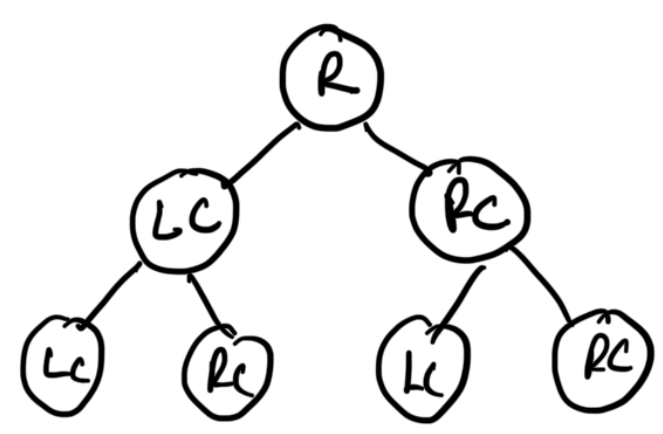


BT  
 CBT

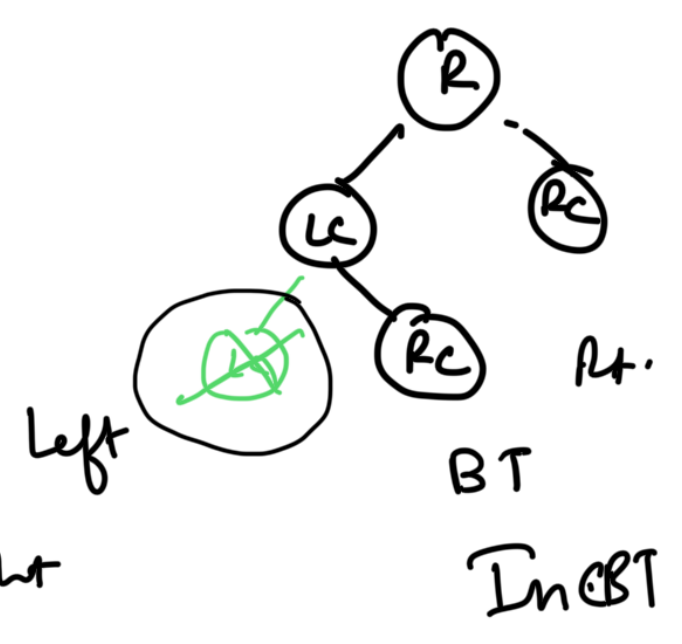


BT  
 CBT  
 Full BT  
 Perfect BT

BT  
 InCBT  
 (LC is missing)



left right

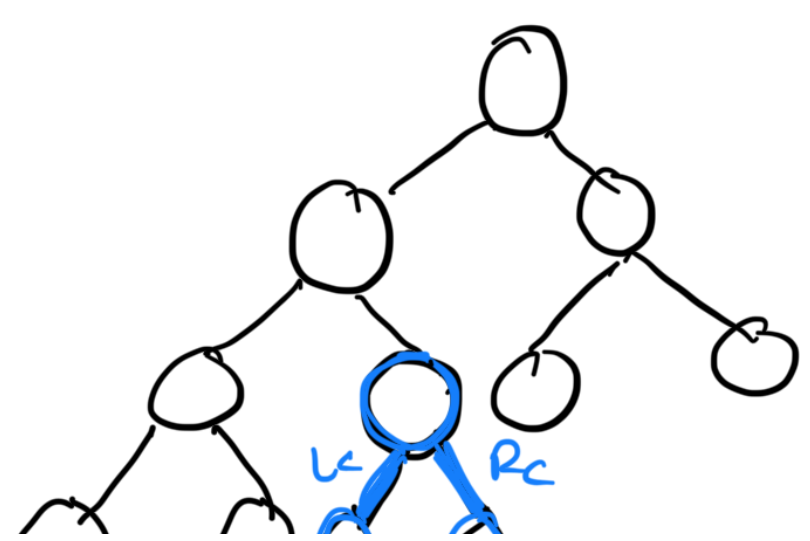


BT  
 InCBT

preference to LC  $\rightarrow$  CBT ✓

RC  $\rightarrow$  } In CBT  
 LC Missing }

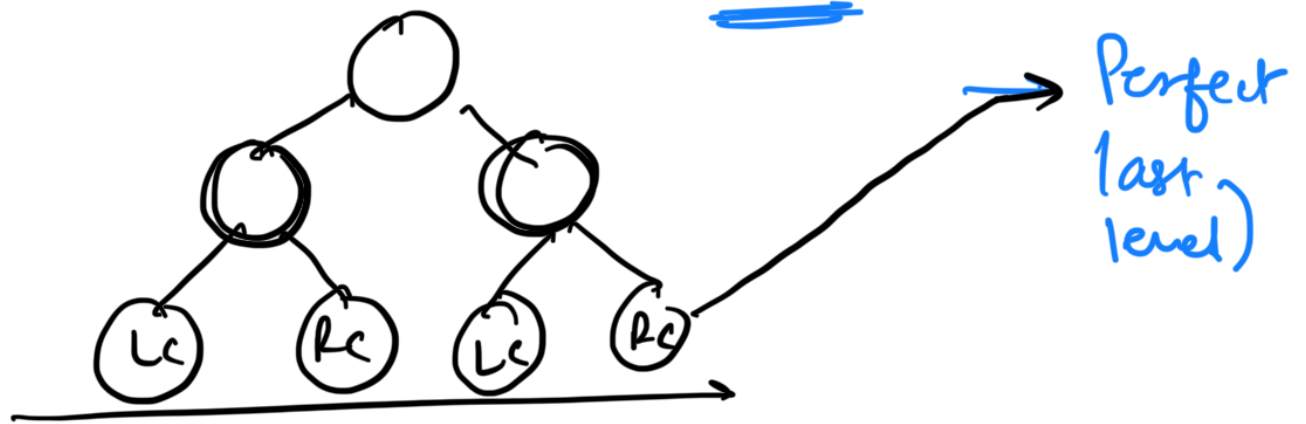
Ex



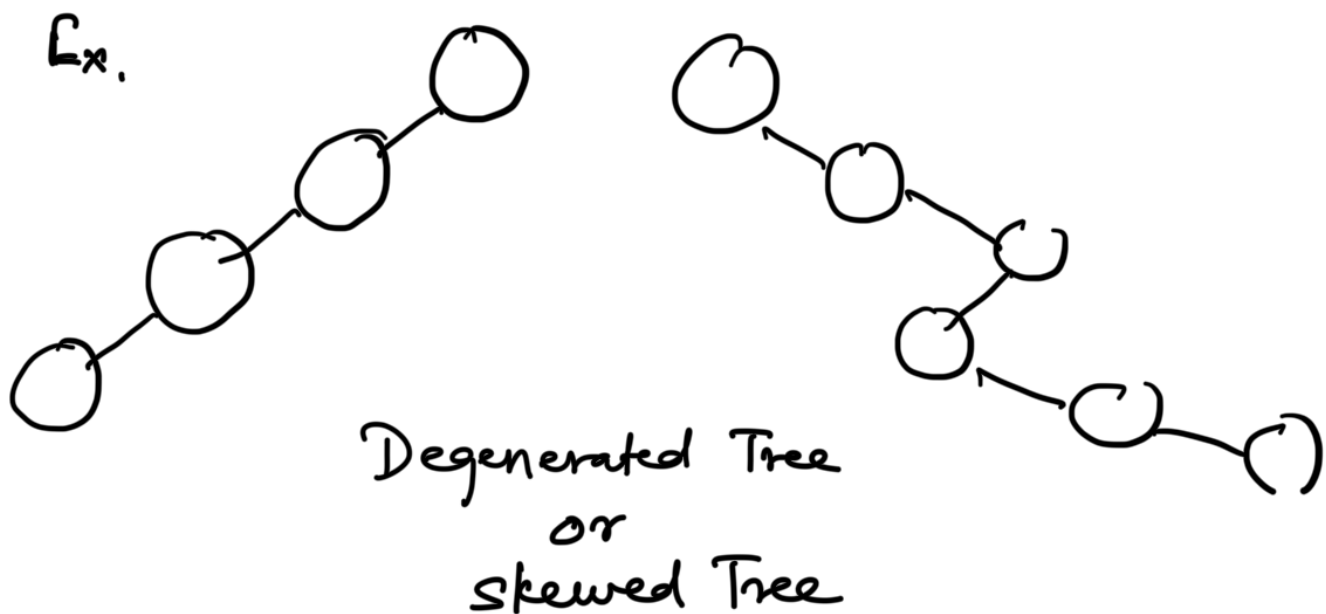


○ ○ ○ ○ +  $\{0, 2\}$  → Full BT

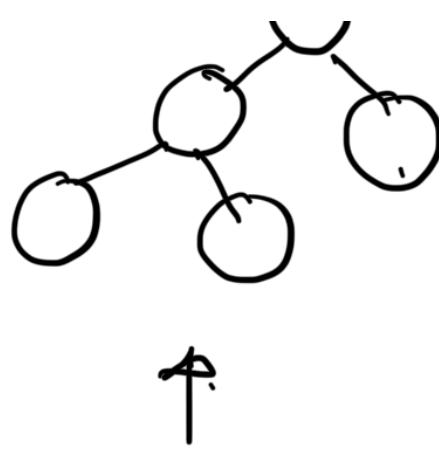
$\{0, 1, 2\}$  → CBT



1. Full BT → Every node has 0 or 2 children
2. Perfect BT → All internal nodes have 2 children.  
and all leaf nodes are at same level
3. Complete BT - All levels are completely filled  
except possibly the last, which is filled  
from left to right (LC → priority)
4. Degenerated Tree (Skewed tree) - Every parent  
node has only one child (left or right)  
making the tree behaving like a linked list.



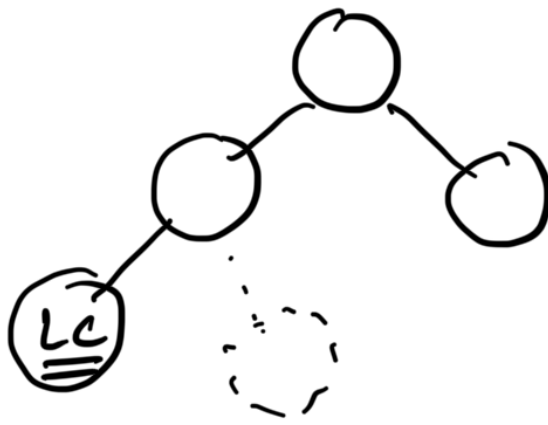
<u>Strict BT</u> - Strictly BT $\{0, 2\}$ Ex.	one single child is
--	------------------------



BT  
Full BT  
CBT  $\rightarrow$  Left  
PBTX  
ICBTX

not allowed

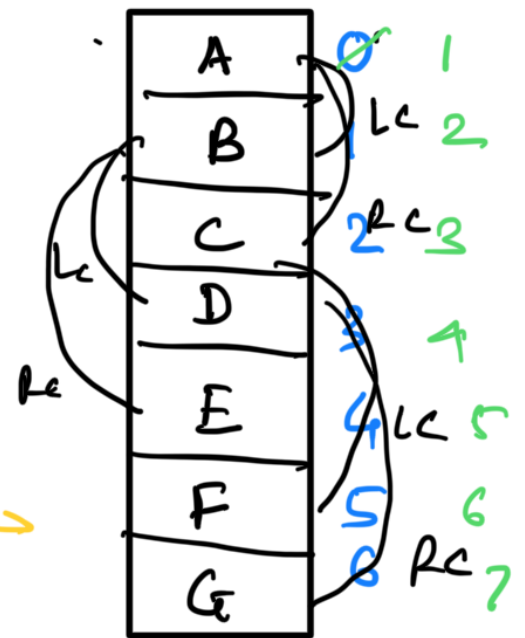
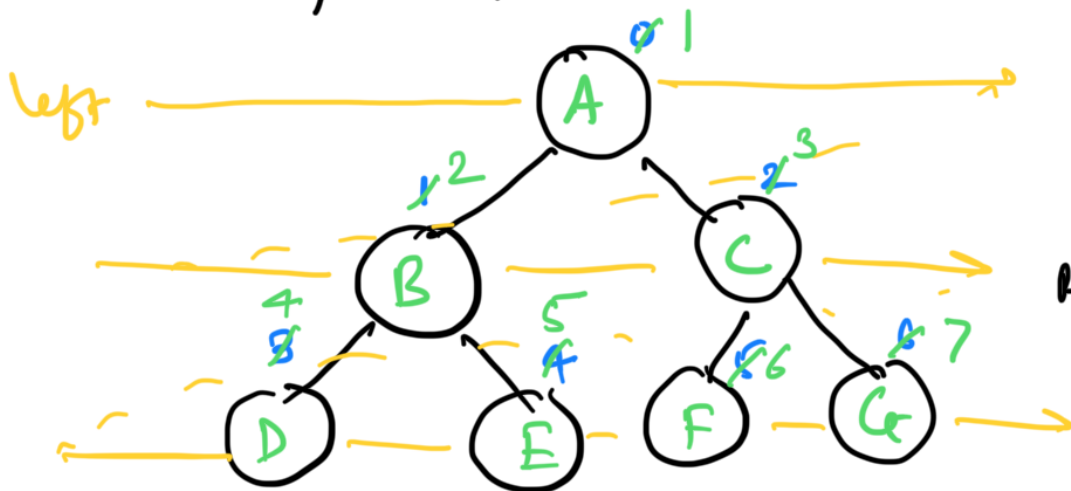
ACBT  $\rightarrow$  Almost Complete BT



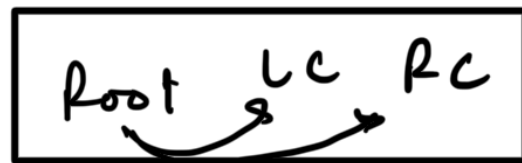
BT  
CBT  
ACBT

BT Representation in 2 ways -

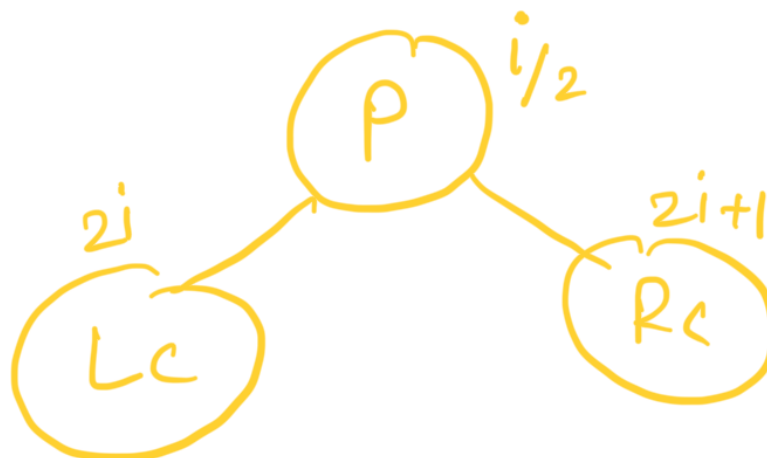
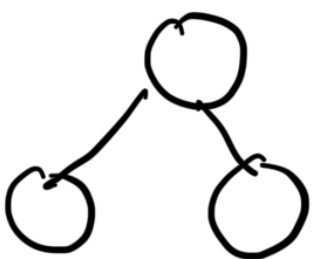
1. Array Implementation



$B \rightarrow 2i$   
 $\rightarrow 2(1) = 2$

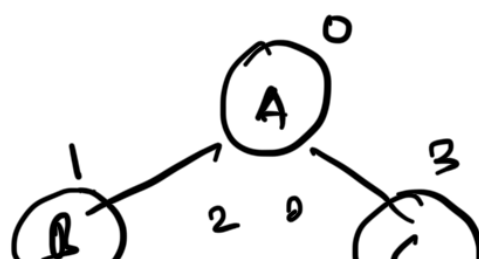


Array  
Implementation

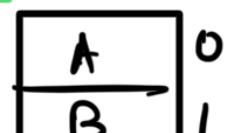


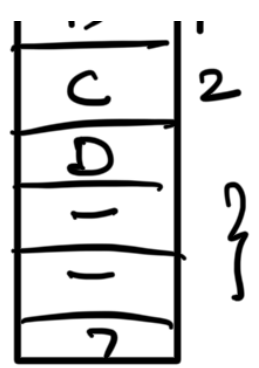
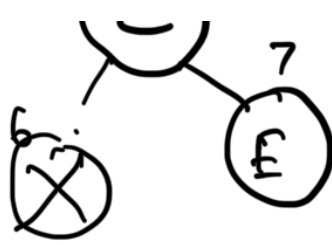
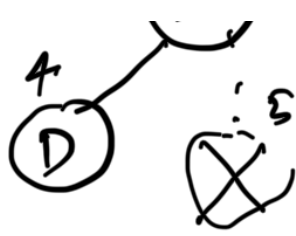
Parent =  $\frac{i}{2}$   
  
 $LC = 2i$   
 $RC = 2i+1$

Ex



Index calculate





Disadvantages of Array Implementation } Memory wastage

## 2. Dynamic Implementation using DLL

