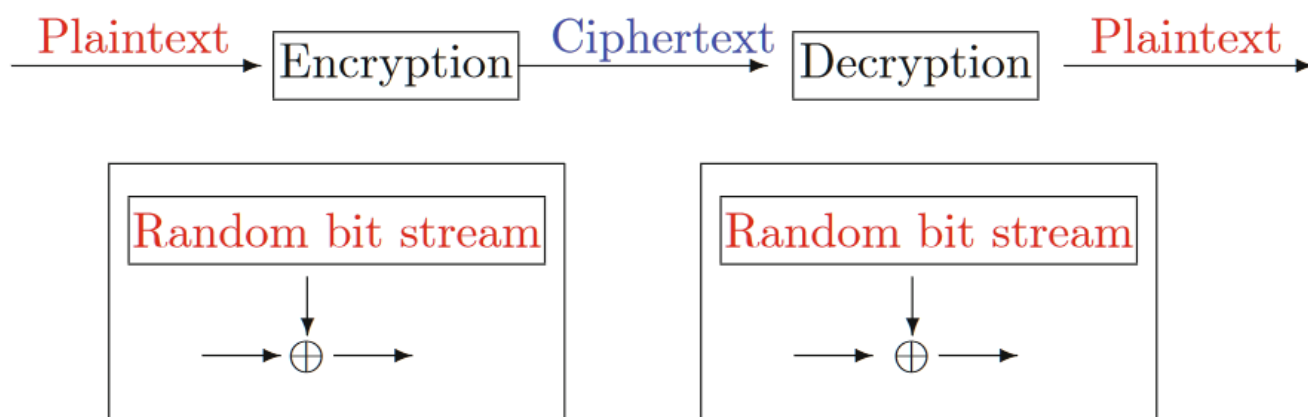


# Stream Ciphers



## Stream ciphers

The idea is to make our ciphertext as indistinguishable from random as possible. Since the message is fixed we want to generate the key in a smart way to make it as random as possible. We are using pseudorandom generators to generate the key and  $\oplus$ -ing it with the message like in the OTP setting.

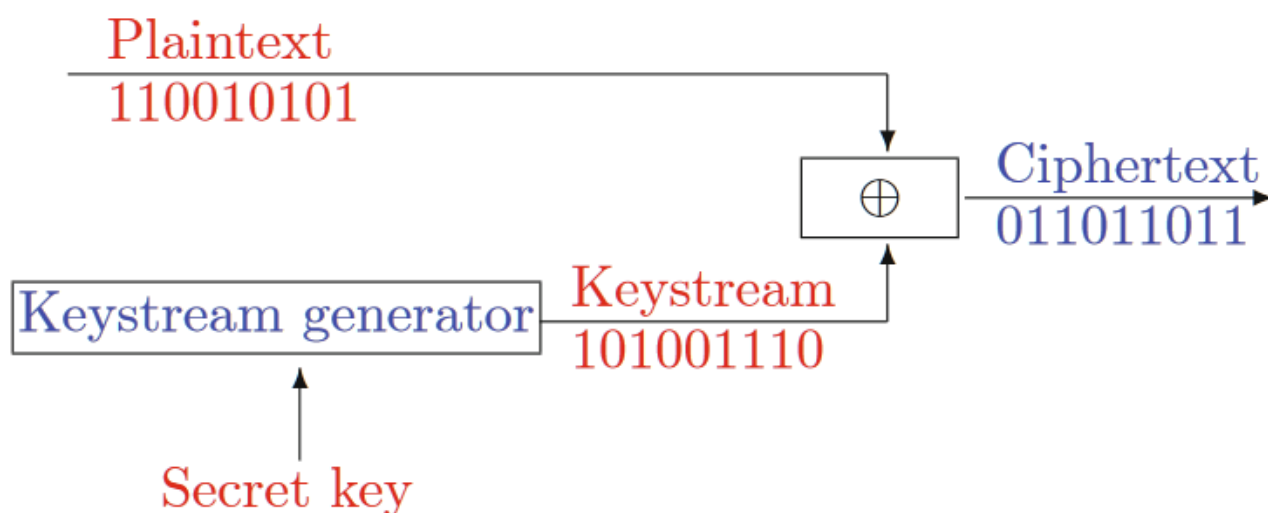


FIGURE 10.2. Stream ciphers

We use a seed to generate a key (Pseudorandom generator). We put a 128 seed into a function  $G$  that outputs a 2048 number.

$$k = G(s)$$

where  $s$  = seed

Let  $(E, D)$  be a cipher.

- $c = E(k, m) = E(G(s), m) = G(s) \oplus m$

### Remark

- Stream ciphers cannot have perfect secrecy since  $|\mathcal{M}| < |\mathcal{K}|$ ! We need a new definition for security.

PRG must be **unpredictable**. A PRG is predictable if given the first  $n - 1$  bits we can predict the next bits

### PRG security

We want the PRG to be indistinguishable from the true random distribution.

Let  $A(x)$  be a **statistical test**:

- 0 if the output is not random
- 1 if the output is random
- Examples
  - $|nr(0) - nr(1)| \leq 10 \cdot \sqrt{(n)}$  (nr of 0 - nr of 1)
  - $nr(00) \leq 10 \cdot \sqrt{(n)}$
  - longest sequence of 0

Let  $G : K \rightarrow \{0, 1\}^n$  and define **Advantage** as:

$$\text{Adv}(A, G) = |Pr[A(G(k)) = 1] - Pr[A(k) = 1]| \in [0, 1]$$

- If  $\text{Adv} \rightarrow 1 \Rightarrow A$  can distinguish from random
- If  $\text{Adv} \rightarrow 0 \Rightarrow A$  can't distinguish from random

**Example:**

Suppose

- $msb(G(k)) = 1$  for  $2/3$  of  $k \in K$
- $A(x) = 1 \iff msb(x) = 1$

Then

- $\text{Adv}(A, G) = |Pr[A(G(k)) = 1] - Pr[A(k) = 1]| = |2/3 - 1/2| = 1/6$

**PRG security -- Definition**

A PRG is secure if for all efficient statistical tests  $A$  the  $\text{Adv}(A, G)$  is negligible

**Theorem**

If the generator  $G$  is secure  $\Rightarrow$  a PRG based on it is unpredictable

**Theorem**

An unpredictable PRG is secure (Then the  $G$  is secure)

**Theorem**

$G : K \rightarrow \{0, 1\}^n$  is a secure PRG  $\Rightarrow$  the Stream cipher is semantically secure

## More Resources

---

- [https://en.wikipedia.org/wiki/Stream\\_cipher](https://en.wikipedia.org/wiki/Stream_cipher) - wiki page
- <https://www.youtube.com/watch?v=rAFNmO-4CIA> - Another short explanation
- <https://www.youtube.com/watch?v=W39KqX0ZTbU> - Another long explanation
- <http://netlab.cs.ucla.edu/wiki/files/shannon1949.pdf> - Shannon security paper