#### DAA HOMEWORK 3

1. function 
$$x = f(n)$$

$$x = 1$$
;  
for  $i = 1$ :n  
for  $j = 1$ :n

$$x = x + 1$$
;

# **Step-by-step analysis:**

# 1. Initialization:

x = 1; This operation takes constant time O(1).

# 2. Outer loop:

for i = 1:n loop runs n times.

# 3. Inner loop:

for j = 1:n loop also runs n times for each iteration of the outer loop.

### 4. Inside the inner loop:

The operation x = x + 1; is executed once when the inner loop runs, which takes O(1) time.

#### **Total number of iterations:**

The total number of iterations for the x = x + 1; for 2 loops. Since both loops run n times,

Total iterations = 
$$\sum_{i=0}^{n} \sum_{j=1}^{n} 1 = \sum_{i=1}^{n} n = n * n = n^2$$

### **Runtime:**

- The initialization step takes constant time: O(1).
- The nested loops iterate n<sup>2</sup> times.

Thus, the total runtime T(n) of the function is:

$$T(n)=O(1)+O(n^2)=O(n^2)$$

# **Conclusion:**

The runtime of the algorithm is  $O(n^2)$ .

3.

Since the runtime of the algorithm is  $T(n) = (n^2)$  which is quadratic, from the above analysis

#### Big-O:

The function is  $O(n^2)$  since the curve is a quadratic, which represents the upper bound.

### Big-Omega:

The function is  $\Omega$  (n<sup>2</sup>) since data is at quadratic growth rate, representing the lower bound.

# Big-Theta:

Since both Big-O and Big-Omega are  $n^2$ , the Big-Theta is also  $\Theta\left(n^2\right)$ 

## 4. Refer plot.png image for the graph

$$n = 5$$
, we can set  $n = 5$ .

n 0 is chosen because, after this point, the data follows the expected polynomial curve. This is where the measured time starts to follow the polynomial equation of  $n^2$ 

5. If I modified the function to be:

```
x = f(n)

x = 1;

y = 1;

for i = 1:n

for j = 1:n

x = x + 1;

y = i + j;
```

Yes, adding y = i + j; increases the number of iterations within the innermost loop, thereby increasing the actual time taken per iteration. However, this additional operation is still O(1), so the time complexity remains O( $n^2$ )

6.

- **Time Complexity**: The asymptotic runtime remains  $O(n^2)$  since adding a constant-time operation does not change the dominant term in the complexity analysis.
- **Measured Execution Time**: The actual execution time will increase slightly due to the extra operation, but the corresponding bounds (Big-O, Big-Omega, Big-Theta) remain the same.

The analysis remains accurate, and only the constant factor changes slightly.