

Show that the following Hoare triple is satisfied under partial correctness:

$$\{T\} \text{ if } (i < j) \{m := i\} \text{ else } \{m := j\} \{(m \leq i) \wedge (m \leq j)\}$$

Proof. We first prove the following Hoare triples

$$\{T \wedge (i < j)\} m := i \{(m \leq i) \wedge (m \leq j)\} \quad (1)$$

$$\{T \wedge \neg(i < j)\} m := j \{(m \leq i) \wedge (m \leq j)\} \quad (2)$$

starting with (1). We have

$$\{(i \leq i) \wedge (i \leq j)\} m := i \{(m \leq i) \wedge (m \leq j)\} \quad (\text{Assignment})$$

$$i \leq i \iff T \quad (\text{Logic})$$

$$T \wedge (i \leq j) \iff (i \leq i) \wedge (i \leq j) \quad (\text{Logic})$$

$$i \leq j \iff T \wedge (i \leq j) \quad (\text{Logic})$$

$$\{i \leq j\} m := i \{(m \leq i) \wedge (m \leq j)\} \quad (\text{PreConEq})$$

$$i < j \implies i \leq j \quad (\text{Logic})$$

$$T \wedge (i < j) \iff i < j \quad (\text{Logic})$$

$$\{T \wedge (i < j)\} m := i \{(m \leq i) \wedge (m \leq j)\} \quad (\text{PreConStr})$$

For (2), we have

$$\{(j \leq i) \wedge (j \leq j)\} m := j \{(m \leq i) \wedge (m \leq j)\} \quad (\text{Assignment})$$

$$j \leq j \iff T \quad (\text{Logic})$$

$$(j \leq i) \wedge T \iff (j \leq i) \wedge (j \leq j) \quad (\text{Logic})$$

$$j \leq i \iff (j \leq i) \wedge T \quad (\text{Logic})$$

$$j \leq i \iff \neg(i < j) \quad (\text{Logic})$$

$$T \wedge \neg(i < j) \iff \neg(i < j) \quad (\text{Logic})$$

$$\{T \wedge \neg(i < j)\} m := j \{(m \leq i) \wedge (m \leq j)\} \quad (\text{PreConEq})$$

Hence we have

$$\frac{\begin{array}{l} \{T \wedge (i < j)\} m := i \{(m \leq i) \wedge (m \leq j)\} \\ \{T \wedge \neg(i < j)\} m := j \{(m \leq i) \wedge (m \leq j)\} \end{array}}{\{T\} \text{ if } (i < j) \{m := i\} \text{ else } \{m := j\} \{(m \leq i) \wedge (m \leq j)\}}$$

by Conditionals, as desired. \square

Find an invariant for the following Hoare triple and prove partial correctness.

$\{x = 0 \wedge y = 1 \wedge z = 1 \wedge n \geq 1\} \text{ while } (z < n) \{y := x + y; x := y - x; z := z + 1\} \quad \{y = f(n)\}$