

# *Compilers*

## *Design and Implementation*

### *Traditional Optimizations*

Algebraic Simplification, Copy Propagation,  
& Constant Propagation

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# Outline

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- Overview of Control-Flow/Data-Flow Analysis
- Algebraic Simplification
- Copy Propagation
- Constant Propagation

# Available Expressions

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- Domain
  - Set of Expressions
- Data-Flow Direction
  - Forward: Out values computed based on In values
- Data-Flow Functions:
  - $OUT = gen \cup (IN - kill)$
  - $gen = \{ exp \mid exp \text{ is calculated in the Basic Block} \}$
  - $kill = \{ exp \mid \exists \text{a variable } v \in exp \text{ that is defined in the Basic Block} \}$
- Meet Operation
  - $IN = \bigcap OUT$  for all the predecessors of a Basic Block
- Initial values
  - Empty Set

# DU-Chain (Reaching Definitions)

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- Domain
  - Set of definitions
- Data-Flow Direction
  - Forward: Out values computed based on In values
- Data-Flow Transfer Function
  - $\text{OUT} = \text{Gen} \cup (\text{IN} - \text{Kill})$
  - $\text{Gen} = \{ x \mid x \text{ is defined in the Basic Block/Statement} \}$
  - $\text{Kill} = \{ x \mid \text{LHS var. of } x \text{ is redefined in the Basic Block/Statement} \}$
- Meet Operation
  - $\text{IN} = \bigcup \text{OUT}$  for all the predecessors of a Basic Block
- Initial Values
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# Analysis Framework

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- Control-Flow Analysis
  - Identify the Structure of the Program
  - Help build Data-Flow Analysis
- Data-Flow Analysis
  - Framework to find information needed for optimizations
  - So far we found
    - Available Expressions
    - UD and DU chains
  - Now let use the information to do something interesting!!!

# “Optimizations”

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- Each Optimization is very simple
  - Reduces Complexity
- Multiple Optimizations are needed
- Optimizations may need to be applied multiple times

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# Algebraic Simplification

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- Apply our knowledge from algebra, number theory etc. to simplify expressions

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- Example

$$\begin{array}{ll} - a + 0 & \Rightarrow a \\ - a * 1 & \Rightarrow a \\ - a / 1 & \Rightarrow a \\ - a * 0 & \Rightarrow 0 \\ - 0 - a & \Rightarrow -a \\ - a + (-b) & \Rightarrow a - b \\ - -(-a) & \Rightarrow a \end{array}$$

# Algebraic Simplification

---

- Apply our knowledge from algebra, number theory etc. to simplify expressions
- Example
  - $a \wedge \text{true} \Rightarrow a$
  - $a \wedge \text{false} \Rightarrow \text{false}$
  - $a \vee \text{true} \Rightarrow \text{true}$
  - $a \vee \text{false} \Rightarrow a$

# Algebraic Simplification

---

- Apply our knowledge from algebra, number theory etc. to simplify expressions
- Example

$$\begin{array}{ll} - a^2 & \Rightarrow a * a \\ - a * 2 & \Rightarrow a + a \\ - a * 8 & \Rightarrow a \ll 3 \end{array}$$

# Algebraic Simplification Opportunities

---

- In the Code
  - Programmers are lazy to simplify expressions
  - Programs are more readable with full expressions
- After Compiler Expansion
  - Example: Array read A[8][12] will get expanded to
  - $*(A_{\text{base}} + 4*(12 + 8*256))$  which can be simplified
- After other Optimizations

# Usefulness of Algebraic Simplification

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- Reduces the number of instructions
- Uses less expensive instructions
- Enable other optimizations

# Implementation

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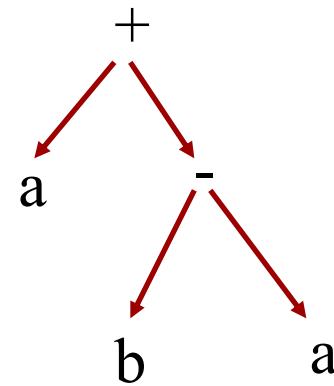
- Not a Data-Flow optimization!
- Find candidates that matches the simplification rules and simplify the expression trees
- Candidates may not be obvious

# Implementation

---

- Not a Data-Flow optimization! Why?
- Find candidates that matches the simplification rules and simplify the expression trees
- Candidates may not be obvious
  - Example

$$a + b - a$$



# Use Knowledge about Operators

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- Commutative Operators
  - $a \text{ op } b = b \text{ op } a$
  -
- Associative Operators
  - $(a \text{ op } b) \text{ op } c = b \text{ op } (a \text{ op } c)$

# Canonical Format

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- Put expression trees into a canonical format
  - Sum of multiplicands
  - Example
$$(a + 3) * (a + 8) * 4 \Rightarrow 4*a*a + 44*a + 96$$

# Effects on the Numerical Stability

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- Some algebraic simplifications may produce incorrect results

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- Example
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# Effects on the Numerical Stability

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- Some algebraic simplifications may produce incorrect results
- Example
  - $(a / b)^0 + c$
  - we can simplify this to  $c$
  - But what about when  $b = 0$ ? should be a exception, ...

# Outline

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- Overview of Control-Flow/Data-Flow Analysis
- Algebraic Simplification
- Copy Propagation
- Constant Propagation

# Copy Propagation

---

- Bypass Multiple Copying
  - propagate a value directly to its use
- Example

$$a = b + c$$

$$d = a$$

$$e = d$$

$$f = d + 2*e + c$$

# Copy Propagation

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- Bypass Multiple Copying
  - propagate a value directly to its use
- Example

$$a = b + c$$

$$\begin{aligned} d &= a \\ e &= d \end{aligned}$$


$$f = d + 2 * e + c$$

# Copy Propagation

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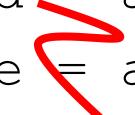
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# Opportunities for Copy Propagation

---

- Exists in User Code
- After other Optimizations
  - Example: Algebraic simplification

# Advantages of Copy Propagation

---

- Leads to further algebraic simplification

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# Advantages of Copy Propagation

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- Leads to further algebraic simplification
- Example

$$a = b + c$$

$$d = a$$

$$e = a$$

$$f = \mathbf{a} + 2*a + c$$

# Advantages of Copy Propagation

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- Leads to further algebraic simplification
- Example

$$a = b + c$$

$$d = a$$

$$e = a$$

$$f = \mathbf{3*a} + c$$

# Advantages of Copy Propagation

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- Reduce Instructions by Eliminating Copy Operations
  - Creates dead code that can be eliminated

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- Example

$$a = b + c$$
$$\textcolor{red}{d = a}$$
$$\textcolor{red}{e = a}$$
$$f = 3*a + c$$

# Advantages of Copy Propagation

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- Reduce Instructions by Eliminating Copy Operations
  - Creates dead code that can be eliminated
- Example

$$a = b + c$$

$$f = 3*a + c$$

# How to Perform Copy Propagation ?

---

- At each RHS Expression  
and for each variable  $v$  used in the RHS expression
  - if the variable  $v$  is defined by a statement of the form  $v = u$
  - replace the variable  $v$  by  $u$
- At each Point of the Program Need to Know
  - The variables that are equal
  - Track equal variables by keeping tuples of the form  $\langle u, v \rangle$  in a set iff  $v = u$  at that point in the program ( $u, v$  are variables)

# How to Perform Copy Propagation ?

---

- An assignment of  $v = u$  is still valid at a given point of the execution if and only if
  - An statement of  $v = u$  occurs in every execution path that reaches the current point
  - The variable  $v$  is not redefined in *any these execution paths* between the assign statement and the current point
  - The variable  $u$  is not redefined in *any these execution paths* between the assign statement and the current point
- A Data-Flow Problem !!!

# Copy Propagation Data-Flow Problem

---

- Domain
  - set of tuples  $\langle v, u \rangle$  representing a statement  $v = u$
- Data-Flow Direction
  - Forward
- Data-Flow Function
  - $OUT = Gen \cup (IN - Kill)$
  - $Gen = \{ \langle v, u \rangle \mid v = u \text{ is the statement} \}$
  - $Kill = \{ \langle v, u \rangle \mid \text{LHS var. of an assignment stmt. is either } v \text{ or } u \}$
- Meet Operation
  - $IN = \bigcap OUT$
- Initial Values
  - Empty Set

# Example

 $b = a$  $c = b + 1$  $d = b$  $b = d + c$  $b = d$

# Example

gen = {  $\langle v, u \rangle \mid v = u$  is the statement }  
kill = {  $\langle v, u \rangle \mid$  LHS var. of the assignment stmt. is either  $v$  or  $u$  }

$b = a$

$c = b + 1$

$d = b$

$b = d + c$

$b = d$

# Example

gen = {  $\langle v, u \rangle$  |  $v = u$  is the statement }  
kill = {  $\langle v, u \rangle$  | LHS var. of the assignment stmt. is either  $v$  or  $u$  }

gen = {  $\langle b, a \rangle$  }

$b = a$

gen = { }

$c = b + 1$

gen = {  $\langle d, b \rangle$  }

$d = b$

gen = { }

$b = d + c$

gen = {  $\langle b, d \rangle$  }

$b = d$

# Example

gen = {  $\langle v, u \rangle$  |  $v = u$  is the statement }  
kill = {  $\langle v, u \rangle$  | LHS var. of the assignment stmt. is either  $v$  or  $u$  }

$b = a$	gen = { $\langle b, a \rangle$ } kill = { $\langle d, b \rangle, \langle b, d \rangle, \langle b, a \rangle, \langle a, b \rangle$ }
$c = b + 1$	gen = { } kill = { $\langle c \rangle$ }
$d = b$	gen = { $\langle d, b \rangle$ } kill = { $\langle b, d \rangle, \langle d, b \rangle$ }
$b = d + c$	gen = { } kill = { $\langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle$ }
$b = d$	gen = { $\langle b, d \rangle$ } kill = { $\langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle$ }

# Example

gen = {  $\langle v,u \rangle \mid v = u$  is the statement }  
kill = {  $\langle v,u \rangle \mid$  LHS var. of the assignment stmt. is either  $v$  or  $u$  }

$b = a$

gen = {  $\langle b,a \rangle$  }  
kill = {  $\langle d,b \rangle, \langle b,d \rangle, \langle b,a \rangle, \langle a,b \rangle$  }

Kills any tuple  
with  $\langle c \rangle$

$c = b + 1$

gen = { }  
kill = {  $\langle c \rangle$  } 

$d = b$

gen = {  $\langle d,b \rangle$  }  
kill = {  $\langle b,d \rangle, \langle d,b \rangle$  }

$b = d + c$

gen = { }  
kill = {  $\langle a,b \rangle, \langle b,a \rangle, \langle d,b \rangle, \langle b,d \rangle$  }

$b = d$

gen = {  $\langle b,d \rangle$  }  
kill = {  $\langle a,b \rangle, \langle b,a \rangle, \langle d,b \rangle, \langle b,d \rangle$  }

# Example

$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$

$b = a$        $\text{gen} = \{ <b,a> \}$

$\text{kill} = \{ <d,b>, <b,d>, <b,a>, <a,b> \}$

$c = b + 1$        $\text{gen} = \{ \ }$

$\text{kill} = \{ <c> \}$

$d = b$        $\text{gen} = \{ <d,b> \}$

$\text{kill} = \{ <b,d>, <d,b> \}$

$b = d + c$        $\text{gen} = \{ \ }$

$\text{kill} = \{ <a,b>, <b,a>, <d,b>, <b,d> \}$

$b = d$        $\text{gen} = \{ <b,d> \}$

$\text{kill} = \{ <a,b>, <b,a>, <d,b>, <b,d> \}$

# Example

$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$

$b = a$

gen = {  $\langle b, a \rangle$  }

IN = { }

kill = {  $\langle d, b \rangle, \langle b, d \rangle, \langle b, a \rangle, \langle a, b \rangle$  }

OUT = {  $\langle b, a \rangle$  }

$c = b + 1$

gen = { }

kill = {  $\langle c \rangle$  }

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gen = {  $\langle d, b \rangle$  }

kill = {  $\langle b, d \rangle, \langle d, b \rangle$  }

$b = d + c$

gen = { }

kill = {  $\langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle$  }

$b = d$

gen = {  $\langle b, d \rangle$  }

kill = {  $\langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle$  }

# Example

$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$

$b = a$	$\text{gen} = \{ \langle b, a \rangle \}$	$\text{IN} = \{ \}$
	$\text{kill} = \{ \langle d, b \rangle, \langle b, d \rangle, \langle b, a \rangle, \langle a, b \rangle \}$	$\text{IN} = \{ \langle b, a \rangle \}$
$c = b + 1$	$\text{gen} = \{ \}$	
	$\text{kill} = \{ \langle c \rangle \}$	$\text{OUT} = \{ \langle b, a \rangle \}$
$d = b$	$\text{gen} = \{ \langle d, b \rangle \}$	
	$\text{kill} = \{ \langle b, d \rangle, \langle d, b \rangle \}$	
$b = d + c$	$\text{gen} = \{ \}$	
	$\text{kill} = \{ \langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle \}$	
$b = d$	$\text{gen} = \{ \langle b, d \rangle \}$	
	$\text{kill} = \{ \langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle \}$	

# Example

$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$

$b = a$

gen = {  $\langle b, a \rangle$  }

IN = { }

kill = {  $\langle d, b \rangle, \langle b, d \rangle, \langle b, a \rangle, \langle a, b \rangle$  }

IN = {  $\langle b, a \rangle$  }

$c = b + 1$

gen = { }

IN = {  $\langle b, a \rangle$  }

kill = {  $\langle c \rangle$  }

$d = b$

gen = {  $\langle d, b \rangle$  }

OUT = {  $\langle b, a \rangle, \langle d, b \rangle$  }

kill = {  $\langle b, d \rangle, \langle d, b \rangle$  }

$b = d + c$

gen = { }

kill = {  $\langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle$  }

$b = d$

gen = {  $\langle b, d \rangle$  }

kill = {  $\langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle$  }

# Example

$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$

$b = a$	$\text{gen} = \{ <b,a> \}$ $\text{kill} = \{ <d,b>, <b,d>, <b,a>, <a,b> \}$	$\text{IN} = \{ \ }$ $\text{IN} = \{ <b,a> \}$
$c = b + 1$	$\text{gen} = \{ \ }$ $\text{kill} = \{ <c> \}$	$\text{IN} = \{ <b,a> \}$
$d = b$	$\text{gen} = \{ <d,b> \}$ $\text{kill} = \{ <b,d>, <d,b> \}$	$\text{IN} = \{ <b,a>, <d,b> \}$
$b = d + c$	$\text{gen} = \{ \ }$ $\text{kill} = \{ <a,b>, <b,a>, <d,b>, <b,d> \}$	$\text{OUT} = \{ \ }$
$b = d$	$\text{gen} = \{ <b,d> \}$ $\text{kill} = \{ <a,b>, <b,a>, <d,b>, <b,d> \}$	

# Example

$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$

$b = a$	$\text{gen} = \{ \langle b, a \rangle \}$ $\text{kill} = \{ \langle d, b \rangle, \langle b, d \rangle, \langle b, a \rangle, \langle a, b \rangle \}$	$\text{IN} = \{ \}$ $\text{IN} = \{ \langle b, a \rangle \}$
$c = b + 1$	$\text{gen} = \{ \}$ $\text{kill} = \{ \langle c \rangle \}$	$\text{IN} = \{ \langle b, a \rangle \}$
$d = b$	$\text{gen} = \{ \langle d, b \rangle \}$ $\text{kill} = \{ \langle b, d \rangle, \langle d, b \rangle \}$	$\text{IN} = \{ \langle b, a \rangle, \langle d, b \rangle \}$
$b = d + c$	$\text{gen} = \{ \}$ $\text{kill} = \{ \langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle \}$	$\text{IN} = \{ \}$
$b = d$	$\text{gen} = \{ \langle b, d \rangle \}$ $\text{kill} = \{ \langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle \}$	$\text{OUT} = \{ \langle b, d \rangle \}$

# Example

$b = a$	gen = { $\langle b, a \rangle$ }	{ }
	kill = { $\langle d, b \rangle, \langle b, d \rangle, \langle b, a \rangle, \langle a, b \rangle$ }	{ $\langle b, a \rangle$ }
$c = b + 1$	gen = { }	
	kill = { $\langle c \rangle$ }	{ $\langle b, a \rangle$ }
$d = b$	gen = { $\langle d, b \rangle$ }	
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$b = d + c$	gen = { }	
	kill = { $\langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle$ }	{ }
$b = d$	gen = { $\langle b, d \rangle$ }	
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$b = d + c$	gen = { }	
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$b = d$	gen = { $\langle b, d \rangle$ }	
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# Example

$b = a$	gen = { $\langle b, a \rangle$ }	{ } $\{ \langle b, a \rangle \}$
	kill = { $\langle d, b \rangle$ , $\langle b, d \rangle$ , $\langle b, a \rangle$ , $\langle a, b \rangle$ }	
$c = b + 1$	gen = { }	
	kill = { $\langle c \rangle$ }	$\{ \langle b, a \rangle \}$
$d = a$	gen = { $\langle d, b \rangle$ }	
	kill = { $\langle b, d \rangle$ , $\langle d, b \rangle$ }	$\{ \langle b, a \rangle, \langle d, b \rangle \}$
$b = d + c$	gen = { }	
	kill = { $\langle a, b \rangle$ , $\langle b, a \rangle$ , $\langle d, b \rangle$ , $\langle b, d \rangle$ }	{ } $\{ \langle b, d \rangle \}$
$b = d$	gen = { $\langle b, d \rangle$ }	
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	kill = { $\langle b, d \rangle, \langle d, b \rangle$ }	{ $\langle b, a \rangle, \langle d, b \rangle$ }
$b = d + c$	gen = { }	
	kill = { $\langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle$ }	{ }
$b = d$	gen = { $\langle b, d \rangle$ }	
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	kill = { $\langle b, d \rangle, \langle d, b \rangle$ }	{ $\langle b, a \rangle, \langle d, b \rangle$ }
$b = b + c$	gen = { }	
	kill = { $\langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle$ }	{ }
$b = d$	gen = { $\langle b, d \rangle$ }	
	kill = { $\langle a, b \rangle, \langle b, a \rangle, \langle d, b \rangle, \langle b, d \rangle$ }	{ $\langle b, d \rangle$ }

# Example

	gen = { <b,a> }	{ }
b = a	kill = { <d,b>, <b,d>, <b,a>, <a,b> }	{ <b,a> }
c = b + 1	gen = { }	
	kill = { <c> }	{ <b,a> }
d = a	gen = { <d,b> }	
	kill = { <b,d>, <d,b> }	{ <b,a>, <d,b> }
b = <b>b</b> + c	gen = { }	
	kill = { <a,b>, <b,a>, <d,b>, <b,d> }	{ }
b = d	gen = { <b,d> }	
	kill = { <a,b>, <b,a>, <d,b>, <b,d> }	{ <b,d> }

# Example

 $b = a$  $c = b + 1$  $d = a$  $b = b + c$  $b = d$

# Example

gen = {  $\langle v, u \rangle$  |  $v = u$  is the statement }  
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gen = {  $\langle b, a \rangle$  }

$b = a$

gen = { }

$c = b + 1$

gen = {  $\langle d, a \rangle$  }

$d = a$

gen = { }

$b = b + c$

gen = {  $\langle b, d \rangle$  }

$b = d$

# Example

gen = {  $\langle v, u \rangle$  |  $v = u$  is the statement }  
kill = {  $\langle v, u \rangle$  | LHS var. of an assignment stmt. is either  $v$  or  $u$  }

b = a                    gen = {  $\langle b, a \rangle$  }  
                          kill = {  $\langle d, b \rangle, \langle b, d \rangle, \langle b, a \rangle, \langle a, b \rangle$  }

c = b + 1            gen = { }  
                          kill = {  $\langle c \rangle$  }

d = a                    gen = {  $\langle d, a \rangle$  }  
                          kill = {  $\langle b, d \rangle, \langle d, b \rangle, \langle d, a \rangle, \langle a, d \rangle$  }

b = b + c            gen = { }  
                          kill = {  $\langle b, a \rangle, \langle a, b \rangle, \langle b, d \rangle, \langle d, b \rangle$  }

b = d                    gen = {  $\langle b, d \rangle$  }  
                          kill = {  $\langle b, a \rangle, \langle a, b \rangle, \langle b, d \rangle, \langle d, b \rangle$  }

# Example

$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$

$b = a$        $\text{gen} = \{ \langle b, a \rangle \}$   
 $\text{kill} = \{ \langle d, b \rangle, \langle b, d \rangle, \langle b, a \rangle, \langle a, b \rangle \}$

$c = b + 1$        $\text{gen} = \{ \}$   
 $\text{kill} = \{ \langle c \rangle \}$

$d = a$        $\text{gen} = \{ \langle d, a \rangle \}$   
 $\text{kill} = \{ \langle b, d \rangle, \langle d, b \rangle, \langle d, a \rangle, \langle a, d \rangle \}$

$b = b + c$        $\text{gen} = \{ \}$   
 $\text{kill} = \{ \langle b, a \rangle, \langle a, b \rangle, \langle b, d \rangle, \langle d, b \rangle \}$

$b = d$        $\text{gen} = \{ \langle b, d \rangle \}$   
 $\text{kill} = \{ \langle b, a \rangle, \langle a, b \rangle, \langle b, d \rangle, \langle d, b \rangle \}$

# Example

$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$

$$\text{gen} = \{ \langle b,a \rangle \} \quad \text{IN} = \{ \}$$

$$b = a \quad \text{kill} = \{ \langle d,b \rangle, \langle b,d \rangle, \langle b,a \rangle, \langle a,b \rangle \}$$

$$\text{OUT} = \{ \langle b,a \rangle \}$$

$$c = b + 1 \quad \text{gen} = \{ \}$$

$$\text{kill} = \{ \langle c \rangle \}$$

$$d = a \quad \text{gen} = \{ \langle d,a \rangle \}$$

$$\text{kill} = \{ \langle b,d \rangle, \langle d,b \rangle, \langle d,a \rangle, \langle a,d \rangle \}$$

$$b = b + c \quad \text{gen} = \{ \}$$

$$\text{kill} = \{ \langle b,a \rangle, \langle a,b \rangle, \langle b,d \rangle, \langle d,b \rangle \}$$

$$b = d \quad \text{gen} = \{ \langle b,d \rangle \}$$

$$\text{kill} = \{ \langle b,a \rangle, \langle a,b \rangle, \langle b,d \rangle, \langle d,b \rangle \}$$

# Example

$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$

	$\text{gen} = \{ \langle b,a \rangle \}$	$\text{IN} = \{ \}$
$b = a$	$\text{kill} = \{ \langle d,b \rangle, \langle b,d \rangle, \langle b,a \rangle, \langle a,b \rangle \}$	$\text{IN} = \{ \langle b,a \rangle \}$
$c = b + 1$	$\text{gen} = \{ \}$ $\text{kill} = \{ \}$	$\text{OUT} = \{ \langle b,a \rangle \}$
$d = a$	$\text{gen} = \{ \langle d,a \rangle \}$ $\text{kill} = \{ \langle b,d \rangle, \langle d,b \rangle, \langle d,a \rangle, \langle a,d \rangle \}$	
$b = b + c$	$\text{gen} = \{ \}$ $\text{kill} = \{ \langle b,a \rangle, \langle a,b \rangle, \langle b,d \rangle, \langle d,b \rangle \}$	
$b = d$	$\text{gen} = \{ \langle b,d \rangle \}$ $\text{kill} = \{ \langle b,a \rangle, \langle a,b \rangle, \langle b,d \rangle, \langle d,b \rangle \}$	

# Example

$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$

	$\text{gen} = \{ \langle b,a \rangle \}$	$\text{IN} = \{ \}$
$b = a$	$\text{kill} = \{ \langle d,b \rangle, \langle b,d \rangle, \langle b,a \rangle, \langle a,b \rangle \}$	$\text{IN} = \{ \langle b,a \rangle \}$
$c = b + 1$	$\text{gen} = \{ \}$ $\text{kill} = \{ \}$	$\text{IN} = \{ \langle b,a \rangle \}$
$d = a$	$\text{gen} = \{ \langle d,a \rangle \}$ $\text{kill} = \{ \langle b,d \rangle, \langle d,b \rangle, \langle d,a \rangle, \langle a,d \rangle \}$	$\text{OUT} = \{ \langle d,a \rangle, \langle b,a \rangle \}$
$b = b + c$	$\text{gen} = \{ \}$ $\text{kill} = \{ \langle b,a \rangle, \langle a,b \rangle, \langle b,d \rangle, \langle d,b \rangle \}$	
$b = d$	$\text{gen} = \{ \langle b,d \rangle \}$ $\text{kill} = \{ \langle b,a \rangle, \langle a,b \rangle, \langle b,d \rangle, \langle d,b \rangle \}$	

# Example

$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$

	$\text{gen} = \{ <\!\!b,a\!\!> \}$	$\text{IN} = \{ \ }$
$b = a$	$\text{kill} = \{ <\!\!d,b\!\!>, <\!\!b,d\!\!>, <\!\!b,a\!\!>, <\!\!a,b\!\!> \}$	$\text{IN} = \{ <\!\!b,a\!\!> \}$
$c = b + 1$	$\text{gen} = \{ \ }$ $\text{kill} = \{ \ }$	$\text{IN} = \{ <\!\!b,a\!\!> \}$
$d = a$	$\text{gen} = \{ <\!\!d,a\!\!> \}$ $\text{kill} = \{ <\!\!b,d\!\!>, <\!\!d,b\!\!>, <\!\!d,a\!\!>, <\!\!a,d\!\!> \}$	$\text{IN} = \{ <\!\!d,a\!\!>, <\!\!b,a\!\!> \}$
$b = b + c$	$\text{gen} = \{ \ }$ $\text{kill} = \{ <\!\!b,a\!\!>, <\!\!a,b\!\!>, <\!\!b,d\!\!>, <\!\!d,b\!\!> \}$	$\text{OUT} = \{ <\!\!d,a\!\!> \}$
$b = d$	$\text{gen} = \{ <\!\!b,d\!\!> \}$ $\text{kill} = \{ <\!\!b,a\!\!>, <\!\!a,b\!\!>, <\!\!b,d\!\!>, <\!\!d,b\!\!> \}$	

# Example

$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$

	$\text{gen} = \{ <\!\!b,a\!\!> \}$	$\text{IN} = \{ \ }$
$b = a$	$\text{kill} = \{ <\!\!d,b\!\!>, <\!\!b,d\!\!>, <\!\!b,a\!\!>, <\!\!a,b\!\!> \}$	$\text{IN} = \{ <\!\!b,a\!\!> \}$
$c = b + 1$	$\text{gen} = \{ \ }$ $\text{kill} = \{ \ }$	$\text{IN} = \{ <\!\!b,a\!\!> \}$
$d = a$	$\text{gen} = \{ <\!\!d,a\!\!> \}$ $\text{kill} = \{ <\!\!b,d\!\!>, <\!\!d,b\!\!>, <\!\!d,a\!\!>, <\!\!a,d\!\!> \}$	$\text{IN} = \{ <\!\!d,a\!\!>, <\!\!b,a\!\!> \}$
$b = b + c$	$\text{gen} = \{ \ }$ $\text{kill} = \{ <\!\!b,a\!\!>, <\!\!a,b\!\!>, <\!\!b,d\!\!>, <\!\!d,b\!\!> \}$	$\text{IN} = \{ <\!\!d,a\!\!> \}$
$b = d$	$\text{gen} = \{ <\!\!b,d\!\!> \}$ $\text{kill} = \{ <\!\!b,a\!\!>, <\!\!a,b\!\!>, <\!\!b,d\!\!>, <\!\!d,b\!\!> \}$	$\text{OUT} = \{ <\!\!d,a\!\!>, <\!\!b,d\!\!> \}$

# Example

	gen = { <b,a> }	{ }
b = a	kill = { <d,b>, <b,d>, <b,a>, <a,b> }	{ <b,a> }
c = b + 1	gen = { }	
	kill = { }	{ <b,a> }
d = a	gen = { <d,a> }	
	kill = { <b,d>, <d,b>, <d,a>, <a,d> }	{ <d,a>, <b,a> }
b = b + c	gen = { }	
	kill = { <b,a>, <a,b>, <b,d>, <d,b> }	{ <d,a> }
b = d	gen = { <b,d> }	
	kill = { <b,a>, <a,b>, <b,d>, <d,b> }	{ <d,a>, <b,d> }

# Example

	gen = { <b,a> }	{ }
b = a	kill = { <d,b>, <b,d>, <b,a>, <a,b> }	{ <b,a> }
c = b + 1	gen = { }	
	kill = { }	{ <b,a> }
d = a	gen = { <d,a> }	
	kill = { <b,d>, <d,b>, <d,a>, <a,d> }	{ <d,a>, <b,a> }
b = <b>b</b> + c	gen = { }	
	kill = { <b,a>, <a,b>, <b,d>, <d,b> }	{ <d,a> }
b = d	gen = { <b,d> }	
	kill = { <b,a>, <a,b>, <b,d>, <d,b> }	{ <d,a>, <b,d> }

# Example

	gen = { <b,a> }	{ }
b = a	kill = { <d,b>, <b,d>, <b,a>, <a,b> }	{ <b,a> }
c = b + 1	gen = { }	
	kill = { }	{ <b,a> }
d = a	gen = { <d,a> }	
	kill = { <b,d>, <d,b>, <d,a>, <a,d> }	{ <d,a>, <b,a> }
b = a + c	gen = { }	
	kill = { <b,a>, <a,b>, <b,d>, <d,b> }	{ <d,a> }
b = d	gen = { <b,d> }	
	kill = { <b,a>, <a,b>, <b,d>, <d,b> }	{ <d,a>, <b,d> }

# Example

	gen = { <b,a> }	{ }
b = a	kill = { <d,b>, <b,d>, <b,a>, <a,b> }	{ <b,a> }
c = b + 1	gen = { }	
	kill = { }	{ <b,a> }
d = a	gen = { <d,a> }	
	kill = { <b,d>, <d,b>, <d,a>, <a,d> }	{ <d,a>, <b,a> }
b = a + c	gen = { }	
	kill = { <b,a>, <a,b>, <b,d>, <d,b> }	{ <d,a> }
b = d	gen = { <b,d> }	
	kill = { <b,a>, <a,b>, <b,d>, <d,b> }	{ <d,a>, <b,d> }

# Example

	gen = { <b,a> }	{ }
b = a	kill = { <d,b>, <b,d>, <b,a>, <a,b> }	{ <b,a> }
c = b + 1	gen = { }	
	kill = { }	{ <b,a> }
d = a	gen = { <d,a> }	
	kill = { <b,d>, <d,b>, <d,a>, <a,d> }	{ <d,a>, <b,a> }
b = a + c	gen = { }	
	kill = { <b,a>, <a,b>, <b,d>, <d,b> }	{ <d,a> }
b = a	gen = { <b,d> }	
	kill = { <b,a>, <a,b>, <b,d>, <d,b> }	{ <d,a>, <b,d> }

# Example

	gen = { <b,a> }	{ }
b = a	kill = { <d,b>, <b,d>, <b,a>, <a,b> }	{ <b,a> }
c = b + 1	gen = { }	
	kill = { }	{ <b,a> }
d = a	gen = { <d,a> }	
	kill = { <b,d>, <d,b>, <d,a>, <a,d> }	{ <d,a>, <b,a> }
b = a + c	gen = { }	
	kill = { <b,a>, <a,b>, <b,d>, <d,b> }	{ <d,a> }
b = a	gen = { <b,d> }	
	kill = { <b,a>, <a,b>, <b,d>, <d,b> }	{ <d,a>, <b,d> }

# Example

 $b = a$  $c = b + 1$  $d = a$  $b = a + c$  $b = a$ 

ARE WE DONE?  
YES!!

# Example

 $b = a$  $c = b + 1$  $d = a$  $b = a + c$  $b = a$ 

Can we do other Transformations?

# Example

$$\cancel{b = a}$$

$$c = b + 1$$

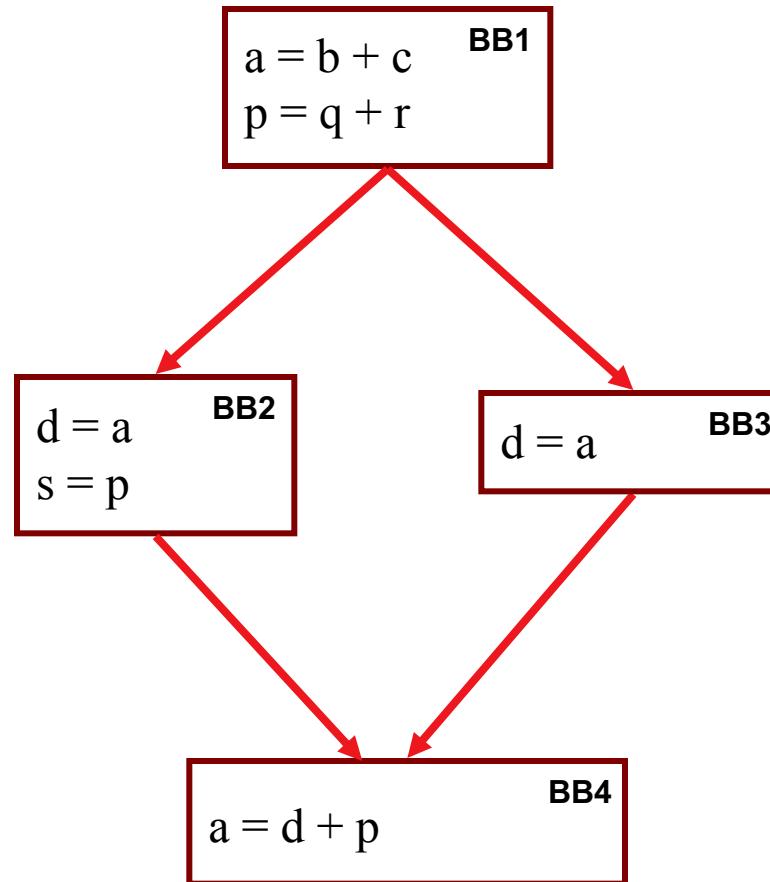
$$\cancel{d = a} \quad ?$$

$$b = a + c$$

$$b = a$$

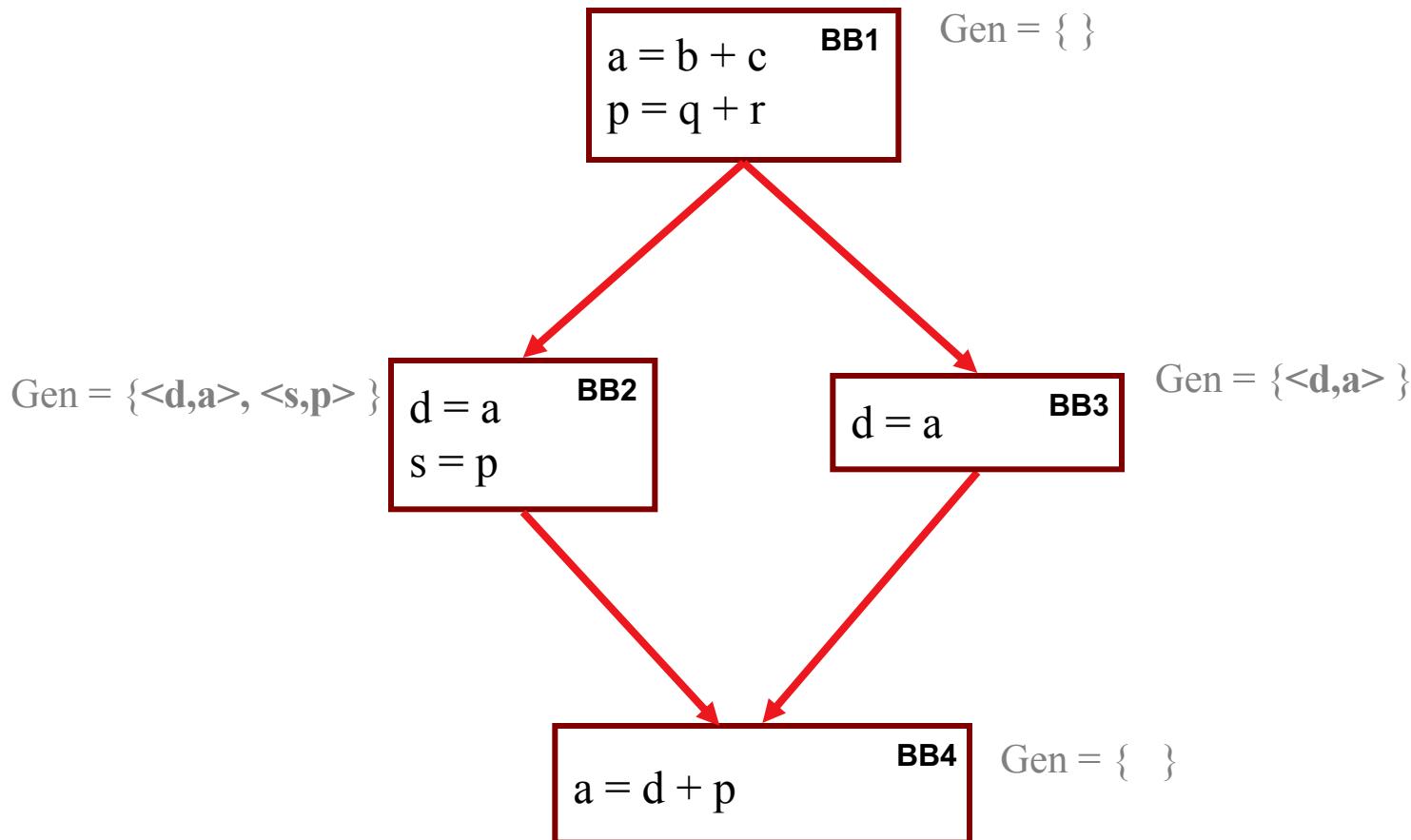
Can we do other Transformations?  
**YES!!**

# Another Example



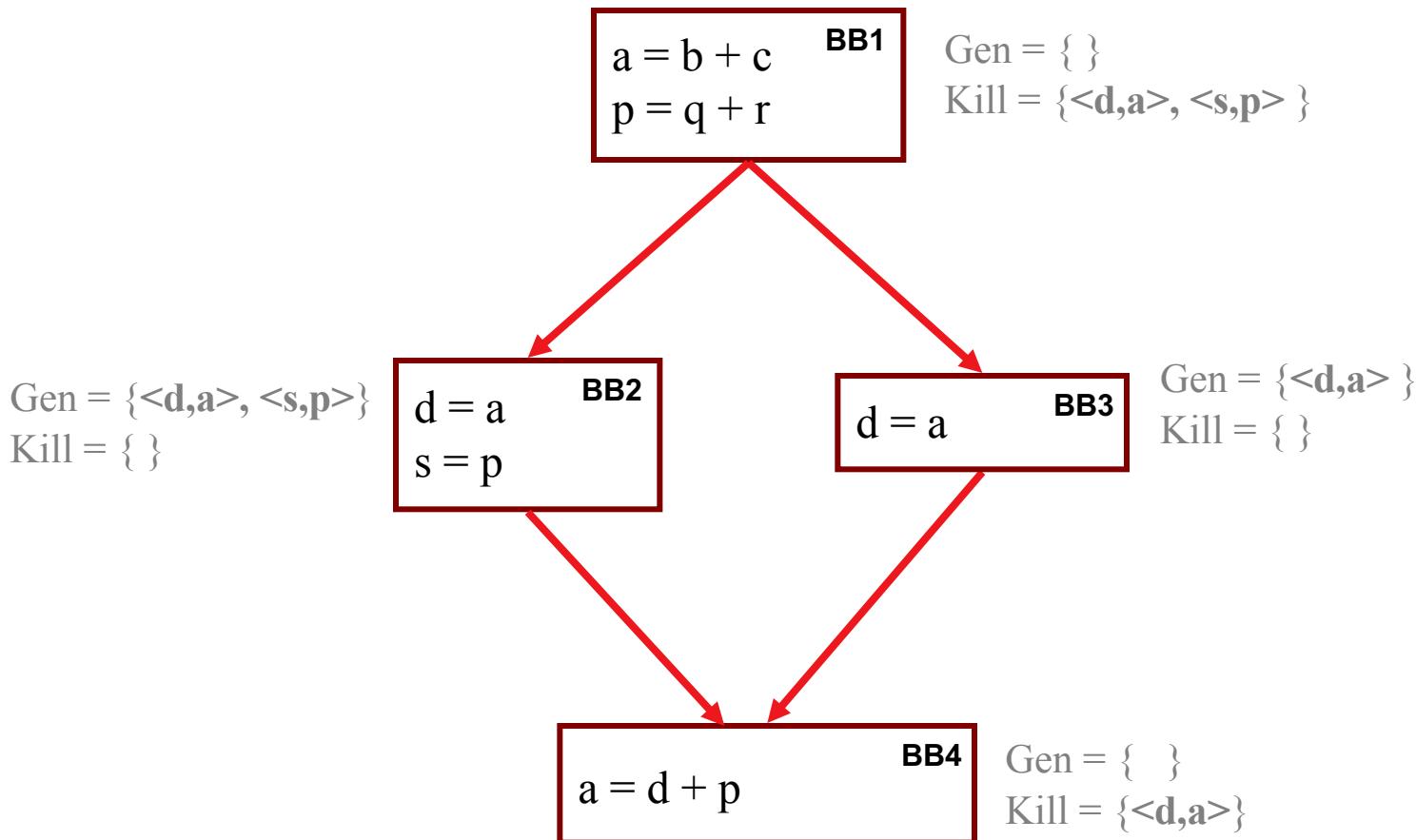
# Another Example

gen = {  $\langle v, u \rangle \mid v = u$  is the statement }  
kill = {  $\langle v, u \rangle \mid$  LHS var. of an assignment stmt. is either  $v$  or  $u$  }



# Another Example

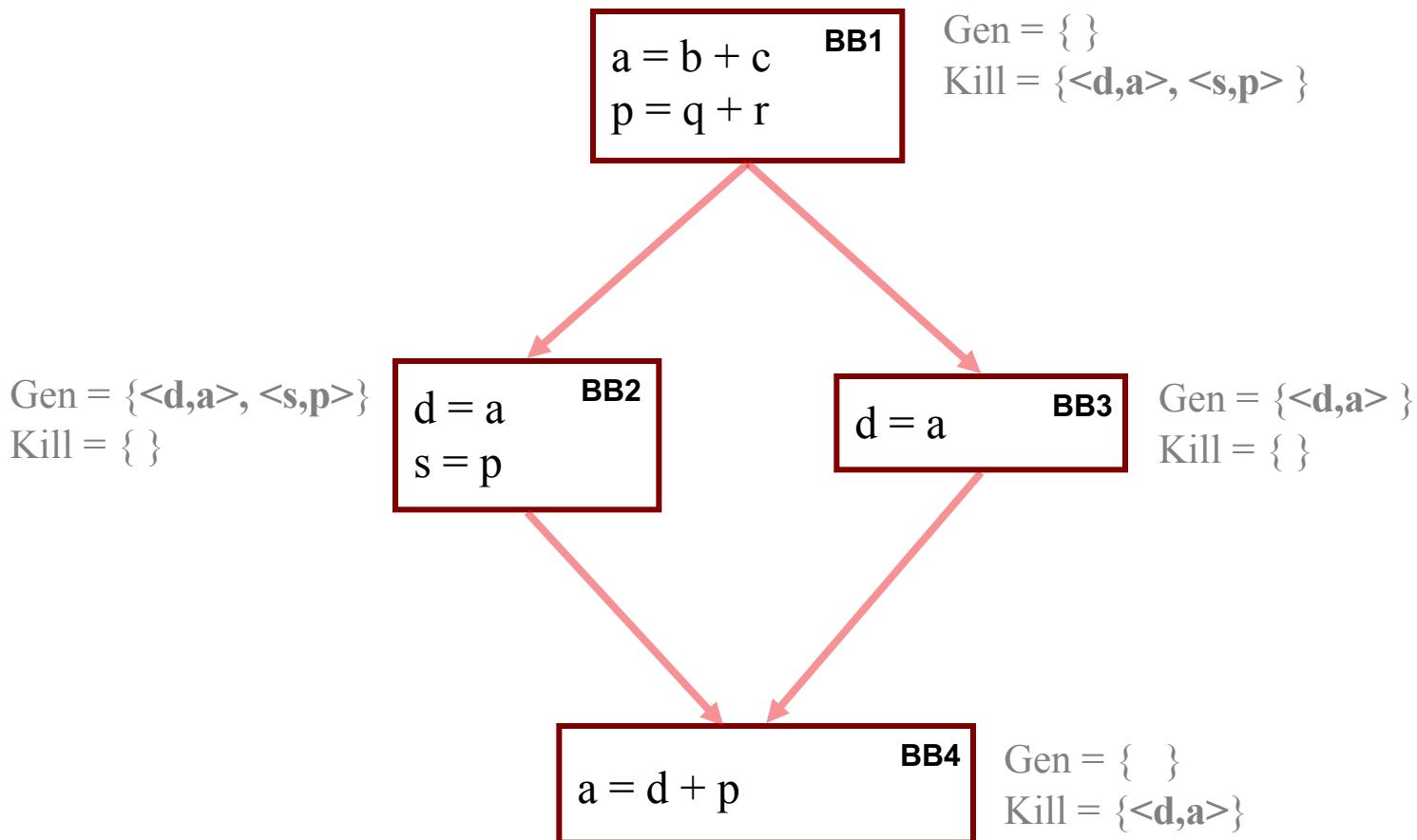
gen = {  $\langle v, u \rangle \mid v = u$  is the statement }  
kill = {  $\langle v, u \rangle \mid$  LHS var. of an assignment stmt. is either  $v$  or  $u$  }



# Another Example

$$\text{IN} = \cap \text{OUT}$$

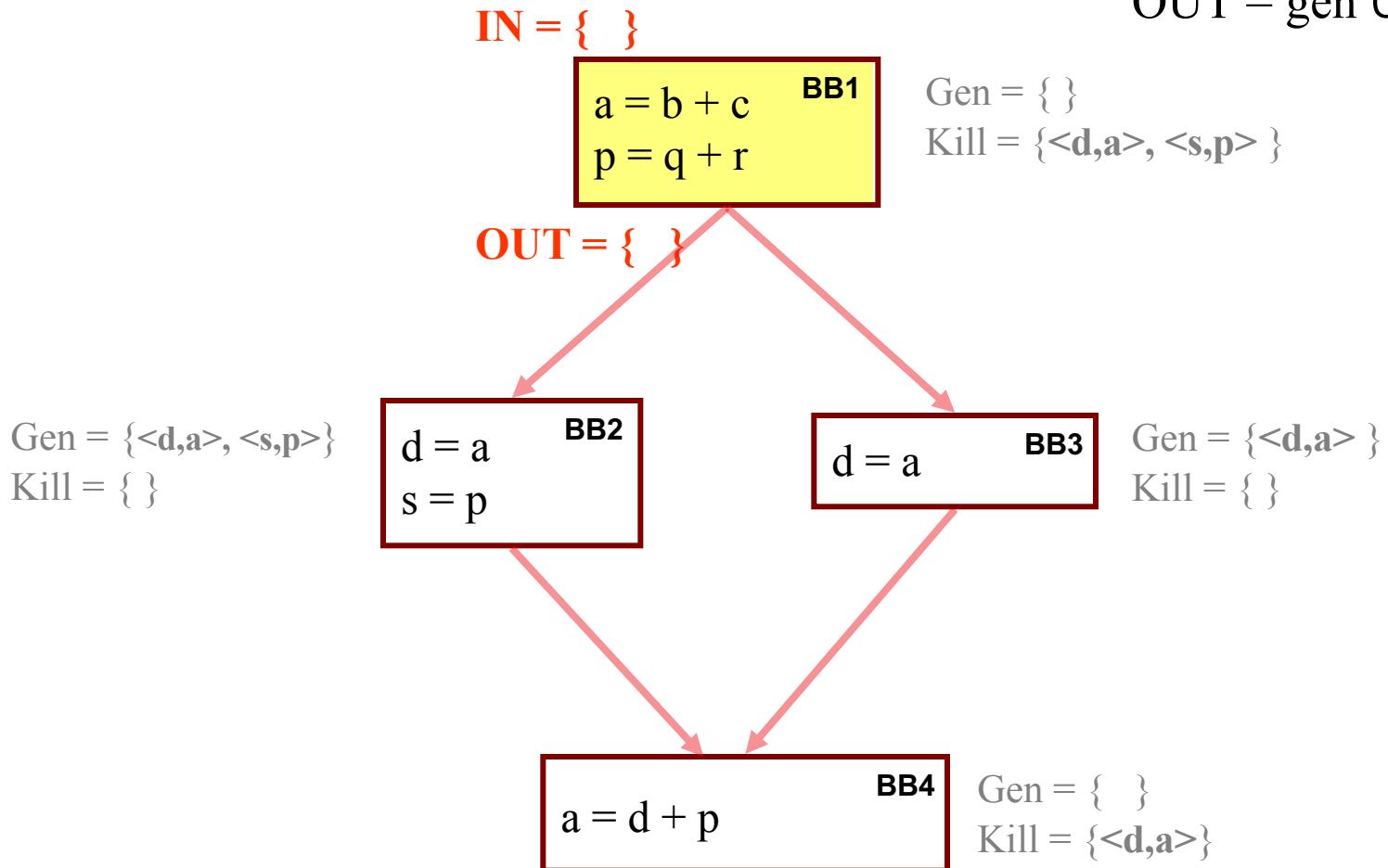
$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$



# Another Example

$$\text{IN} = \cap \text{OUT}$$

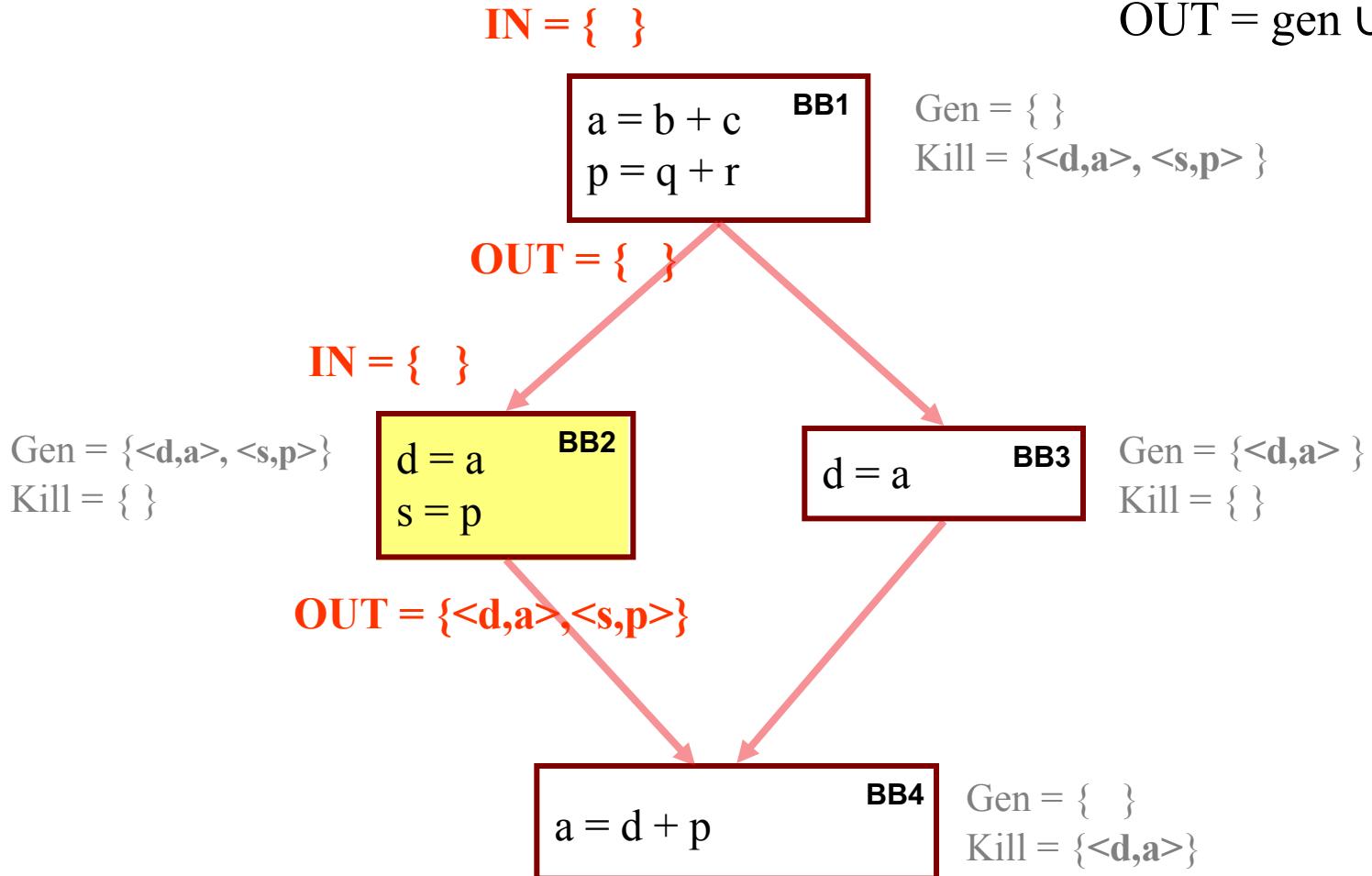
$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$



# Another Example

$$\text{IN} = \cap \text{OUT}$$

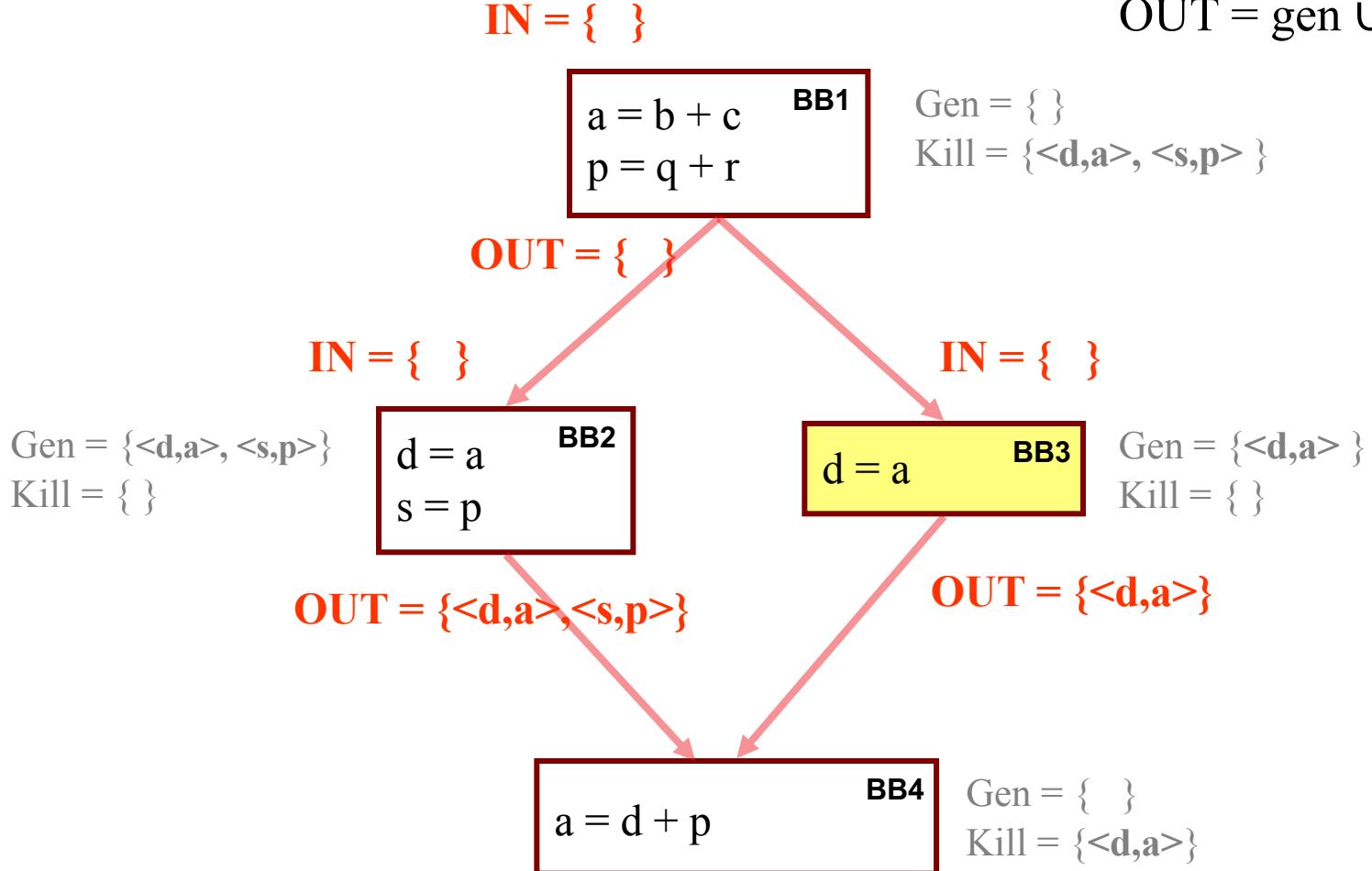
$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$



# Another Example

$$\text{IN} = \cap \text{OUT}$$

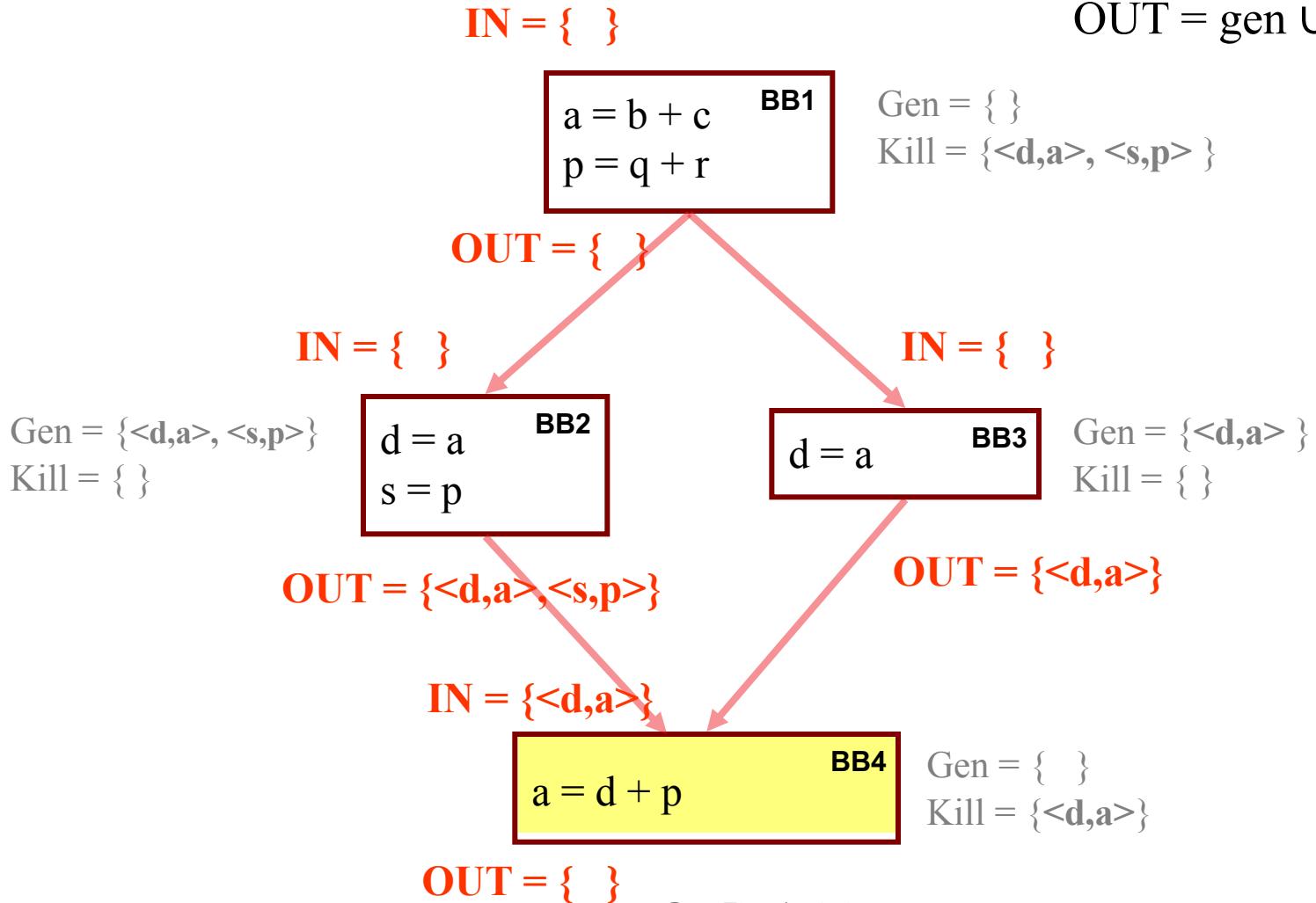
$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$



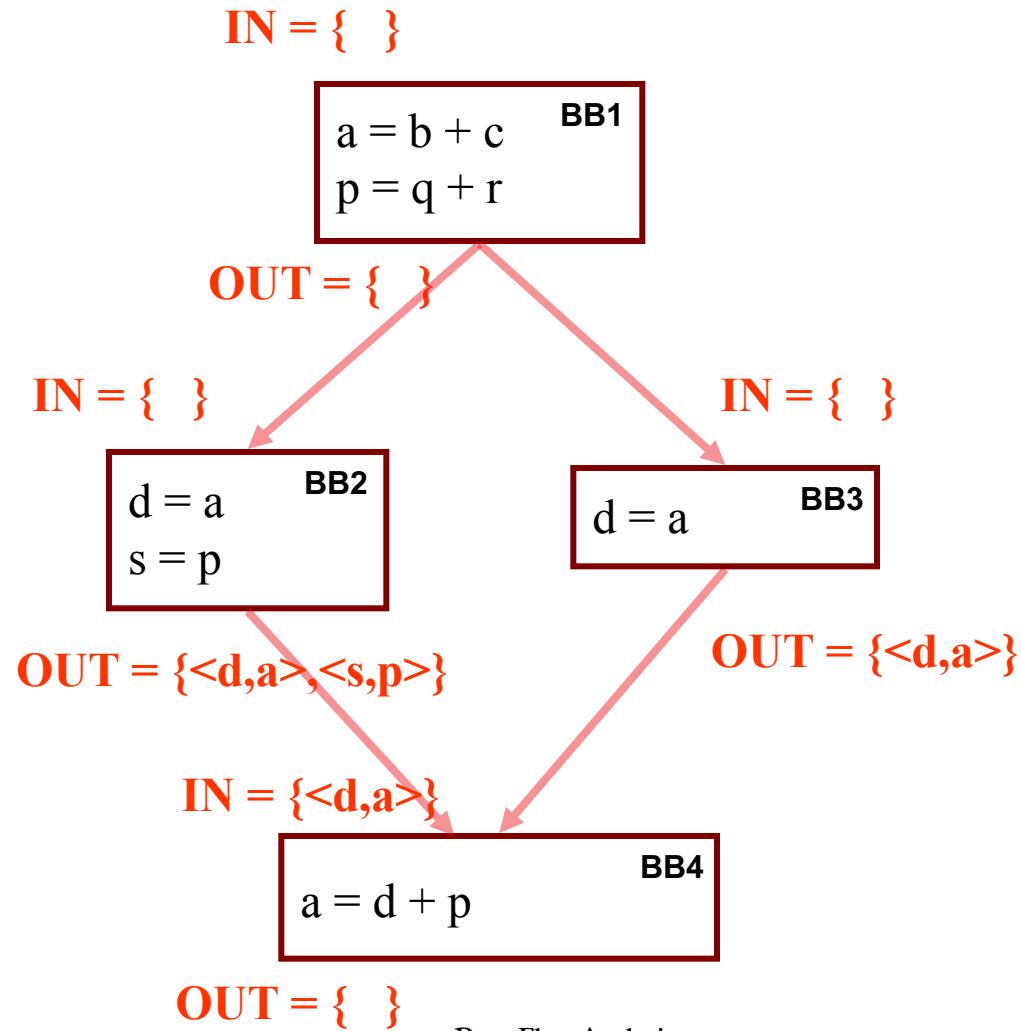
# Another Example

$$\text{IN} = \cap \text{OUT}$$

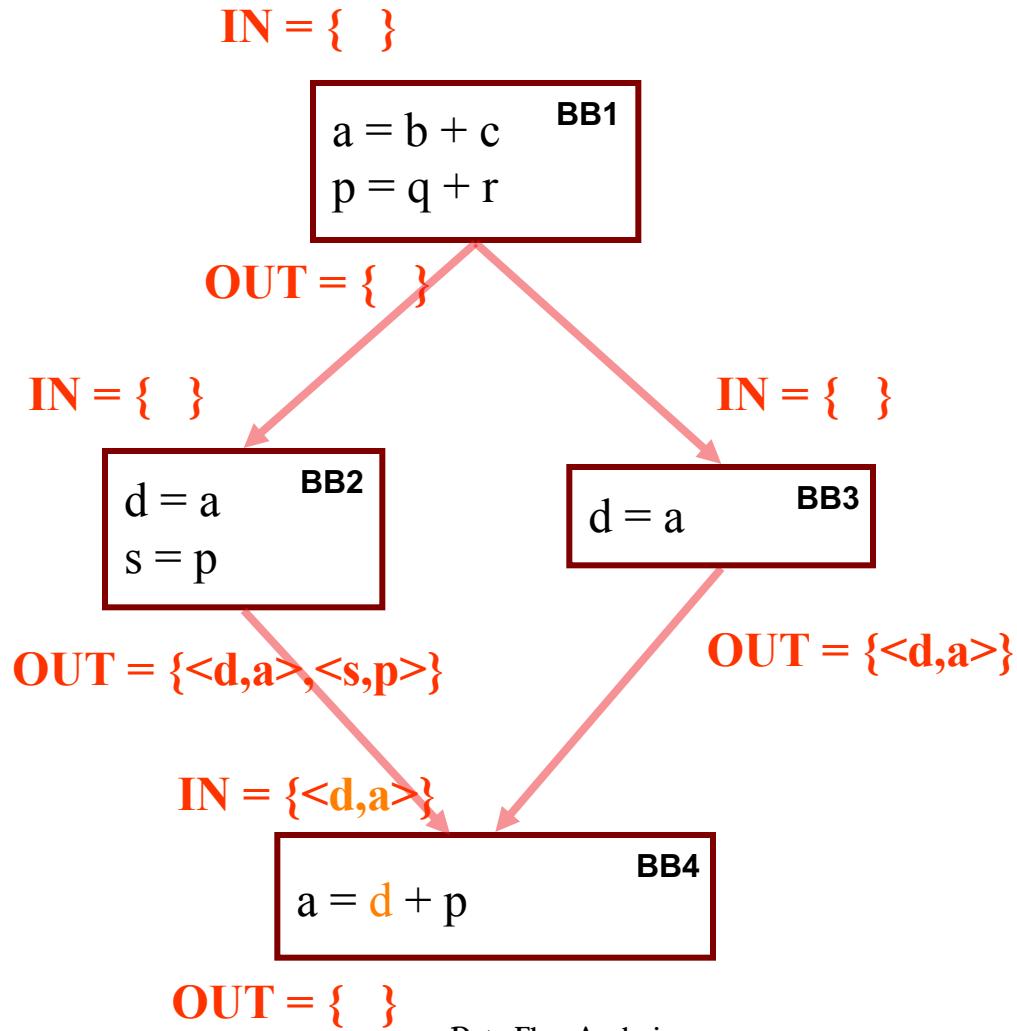
$$\text{OUT} = \text{gen} \cup (\text{IN} - \text{kill})$$



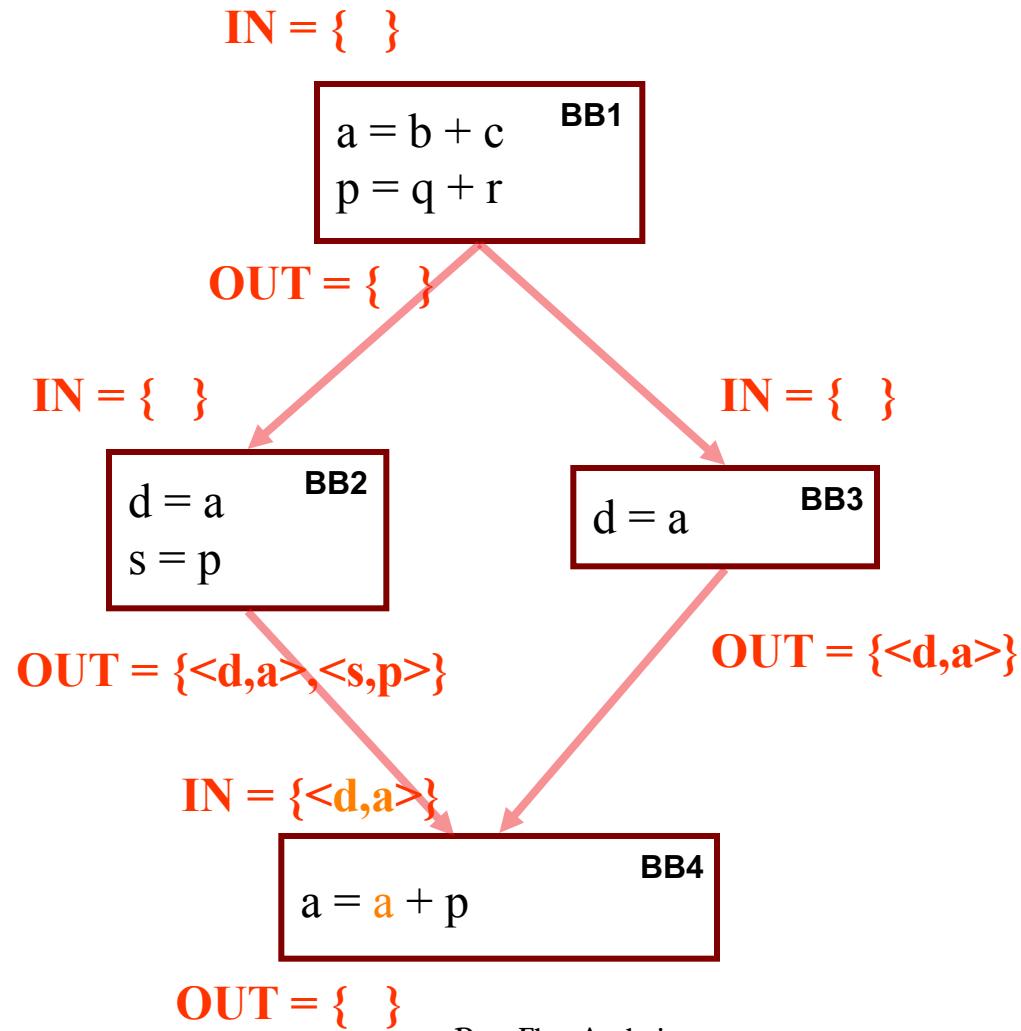
# Another Example



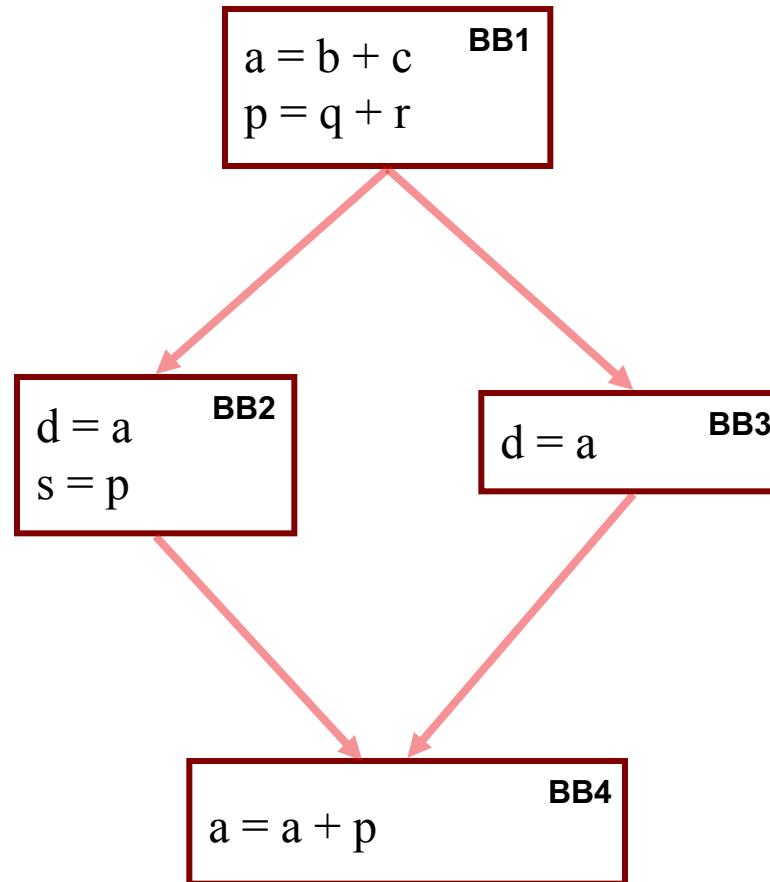
# Another Example



# Another Example



# Another Example



# Outline

---

- Overview of Control-Flow Analysis
- Algebraic Simplification
- Copy Propagation
- Constant Propagation

# Constant Propagation

---

- Use Constant Values
  - Use the known constant of a variable

# Constant Propagation

---

- Use Constant Values
  - Use the known constant of a variable
- Example

a = 43

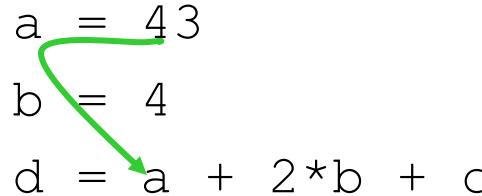
b = 4

d = a + 2\*b + c

# Constant Propagation

- Use Constant Values
  - Use the known constant of a variable
- Example

a = 43  
b = 4  
d = a + 2\*b + c



# Constant Propagation

---

- Use Constant Values
  - Use the known constant of a variable
- Example

$$a = 43$$
$$b = 4$$
$$d = 43 + 2 * b + c$$


# Constant Propagation

---

- Use Constant Values
  - Use the known constant of a variable
- Example

$$a = 43$$
$$b = 4$$
$$d = a + 2^*b + c$$

# Constant Propagation

---

- Use Constant Values
  - Use the known constant of a variable
- Example

$$a = 43$$
$$b = 4$$
$$d = 43 + 2 * 4 + c$$


# Constant Propagation Opportunities

---

- User-Defined Constants
  - Same Constants propagating from many different paths
  - Symbolic Constants defined as variables
- Constants Known to the Compiler
  - data sizes, stack offsets
- Constants Available after Other Optimizations
  - Algebraic Simplification
  - Copy propagation

# Advantages of Constant Propagation

---

- Simplification of the Program

# Advantages of Constant Propagation

---

- Simplification of the Program
- Example

a = 43

b = 4

d = 43 + 2\*4 + c

# Advantages of Constant Propagation

---

- Simplification of the Program
- Example

a = 43

b = 4

d = **43** + **2\*4** + c

# Advantages of Constant Propagation

---

- Simplification of the Program
- Example

a = 43

b = 4

d = **51** + c

# Advantages of Constant Propagation

---

- Enabling further Optimizations

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

a = 4

b = 8

d = 2\*a - b + c

e = c + d

# Advantages of Constant Propagation

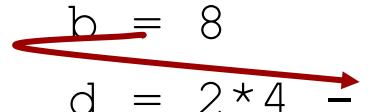
- Enabling further Optimizations
- Example

$$\begin{array}{l} a \equiv 4 \\ b = 8 \\ d = 2 * a - b + c \\ e = c + d \end{array}$$

A red curved arrow starts at the assignment statement  $a \equiv 4$  and points to the term  $2 * a$  in the expression  $d = 2 * a - b + c$ .

# Advantages of Constant Propagation

- Enabling further Optimizations
- Example

$$a = 4$$
$$\text{b} = 8$$
$$d = 2 * 4 - b + c$$
$$e = c + d$$


# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

$$a = 4$$

$$b = 8$$

$$d = \mathbf{2*4} - 8 + c$$

$$e = c + d$$

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

$$a = 4$$

$$b = 8$$

$$d = \mathbf{8} - 8 + c$$

$$e = c + d$$

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

$$a = 4$$

$$b = 8$$

$$d = \mathbf{8} - \mathbf{8} + c$$

$$e = c + d$$

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

$$a = 4$$

$$b = 8$$

$$d = \mathbf{0} + c$$

$$e = c + d$$

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

a = 4

b = 8

d = **0 + c**

e = c + d

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

a = 4

b = 8

d = **c**

e = c + d

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

a = 4

b = 8

d = c

e = c + d

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

a = 4

b = 8

**d = c**

e = c + **d**

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

a = 4

b = 8

**d = c**

e = c + **c**

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

a = 4

b = 8

d = c

e = c + c

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

a = 4

b = 8

d = c

e = **c + c**

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

a = 4

b = 8

d = c

e = **2\*c**

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

a = 4

b = 8

**d = c**

e = 2\*c

# Advantages of Constant Propagation

---

- Enabling further Optimizations
- Example

a = 4

b = 8

e = 2 \* c

# How to Perform Constant Propagation

---

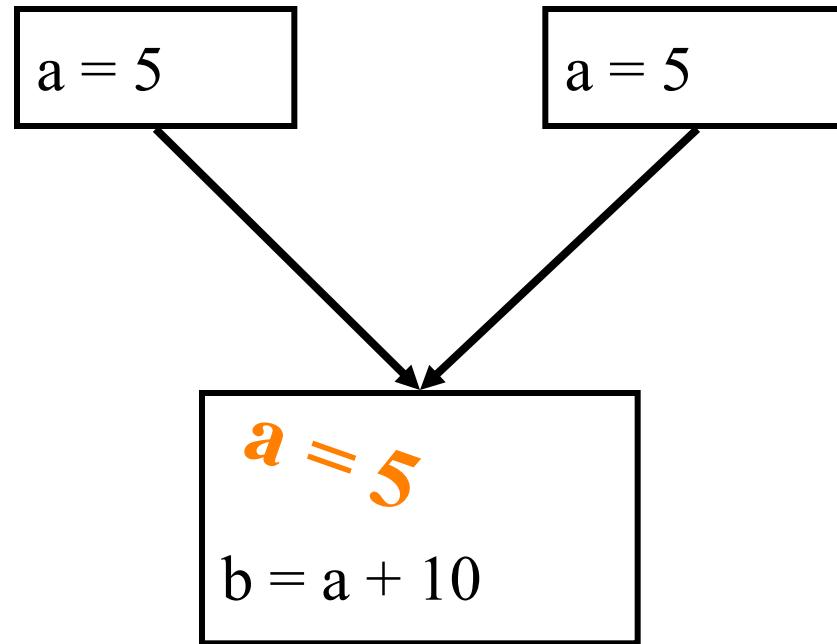
- At each RHS Expression
  - For each variable  $v$  used in the RHS
    - If the variable  $v$  is a known constant  $k$
    - Replace the variable  $v$  by  $k$
- At Each Point of the Program need to know:
  - For each variable  $v$ , if  $v$  is a constant,
  - If so, what the specific constant value is

# How to Perform Constant Propagation

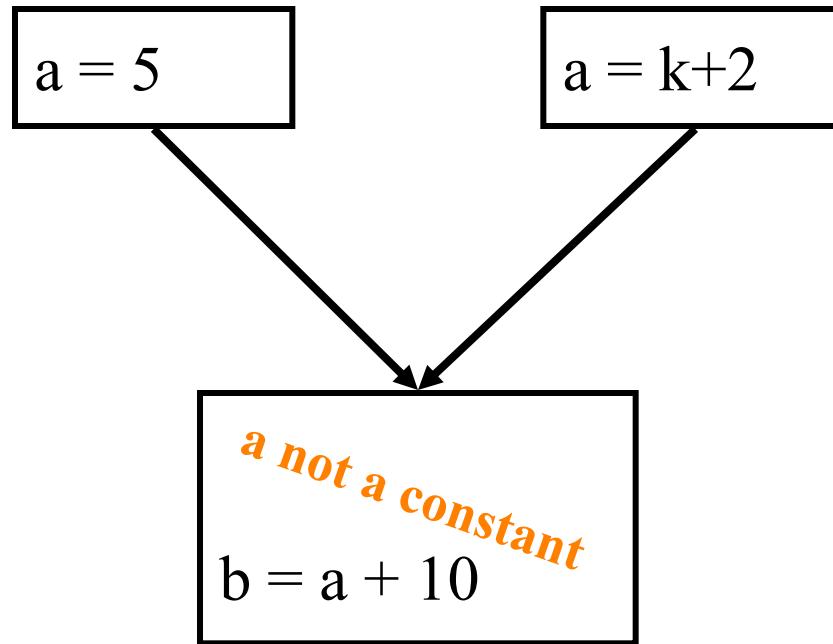
---

- A variable  $v$  is the constant  $k$  at a point of the execution if and only if
  - The current statement is  $v = k$
  - or
  - *Every path reaching the current point* has  $k$  assigned to  $v$
- A Data-Flow Problem !!!

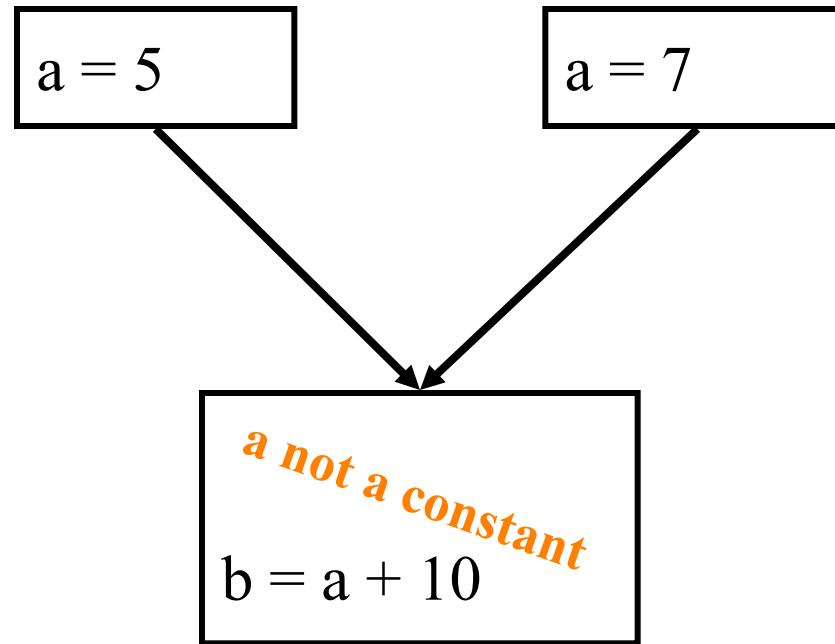
# Values from Two Paths



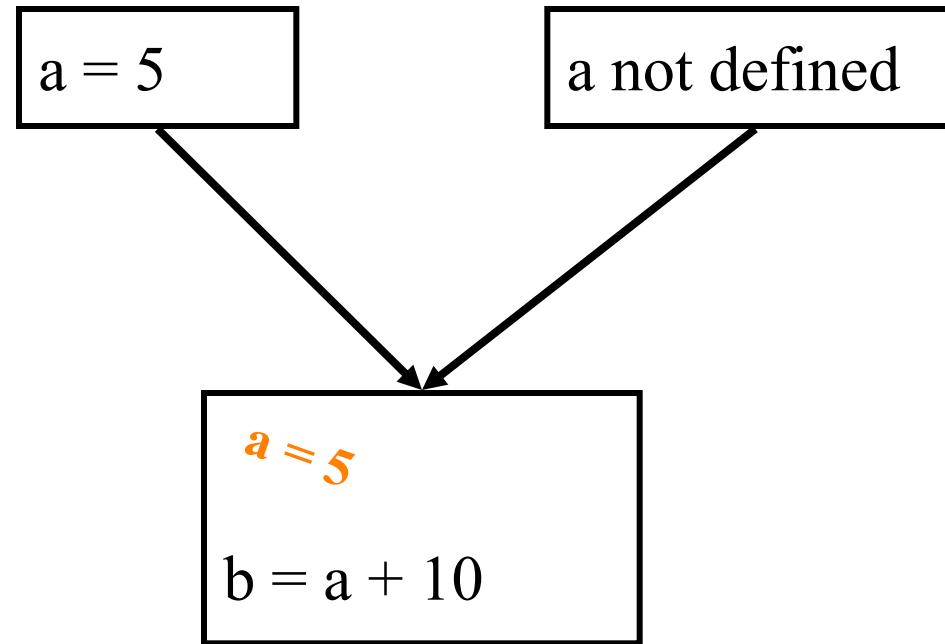
# Values from Two Paths



# Values from Two Paths

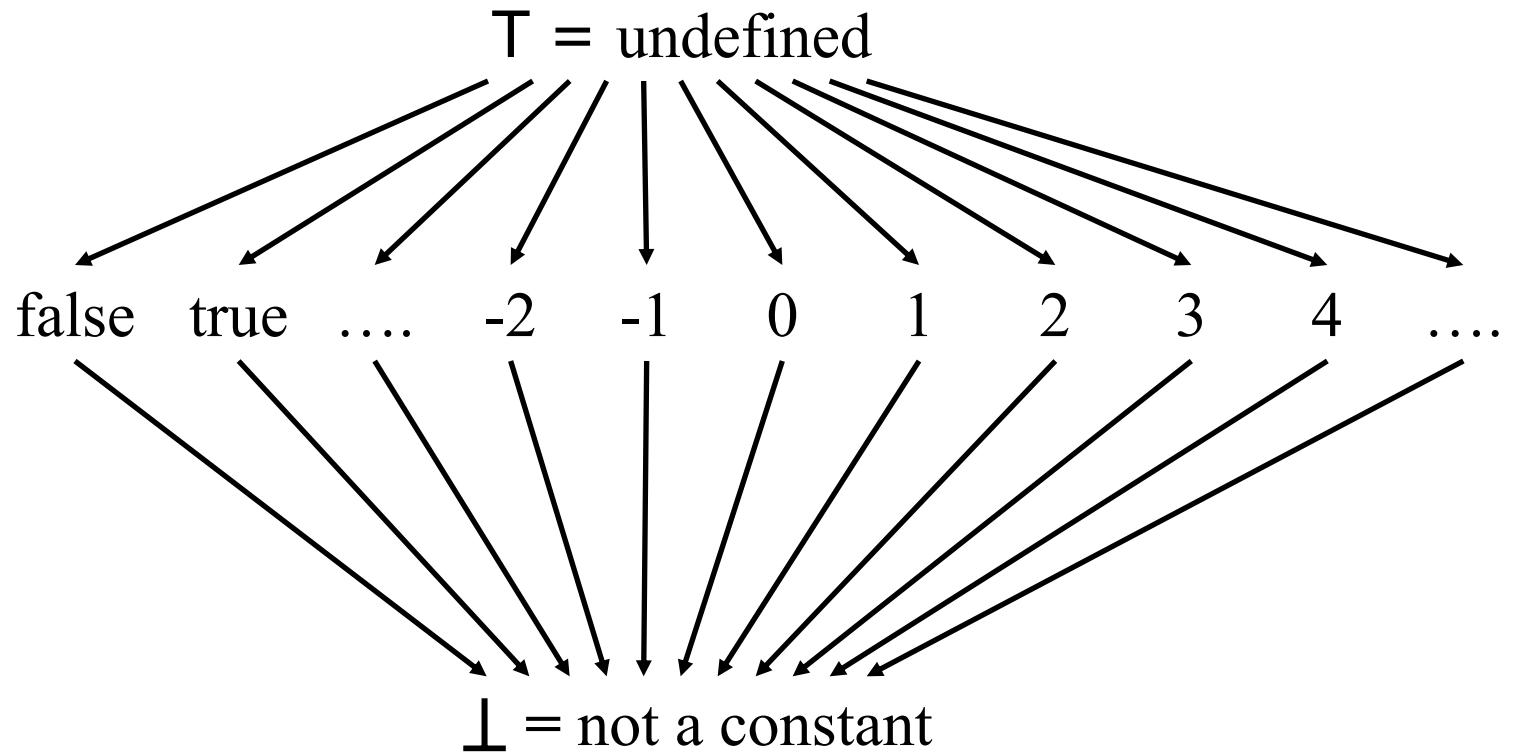


# Values from Two Paths

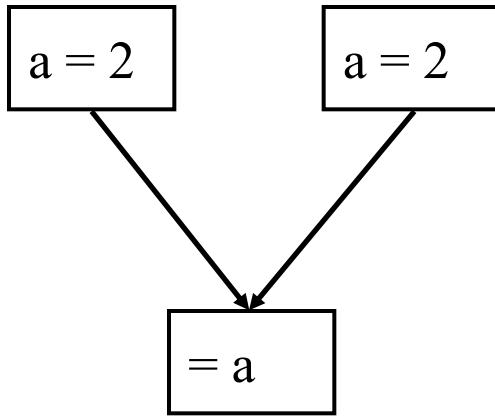


- Brain-dead program, uses an uninitialized value
- High level semantics the compiler don't understand makes this a correct program.

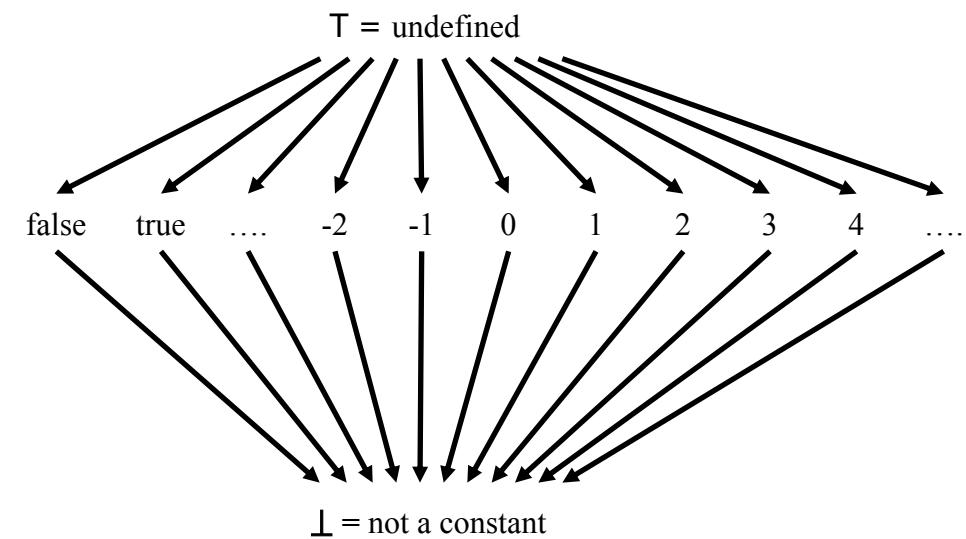
# Lattice for Constant Propagation



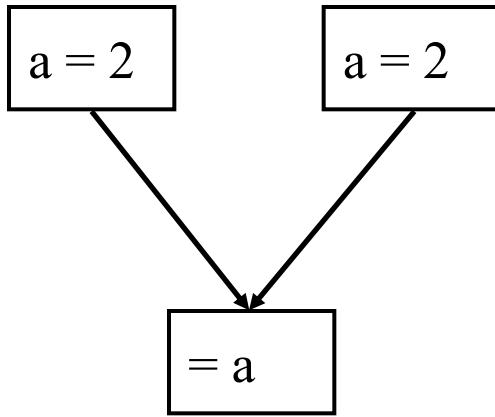
# Meet Operations on the Lattice



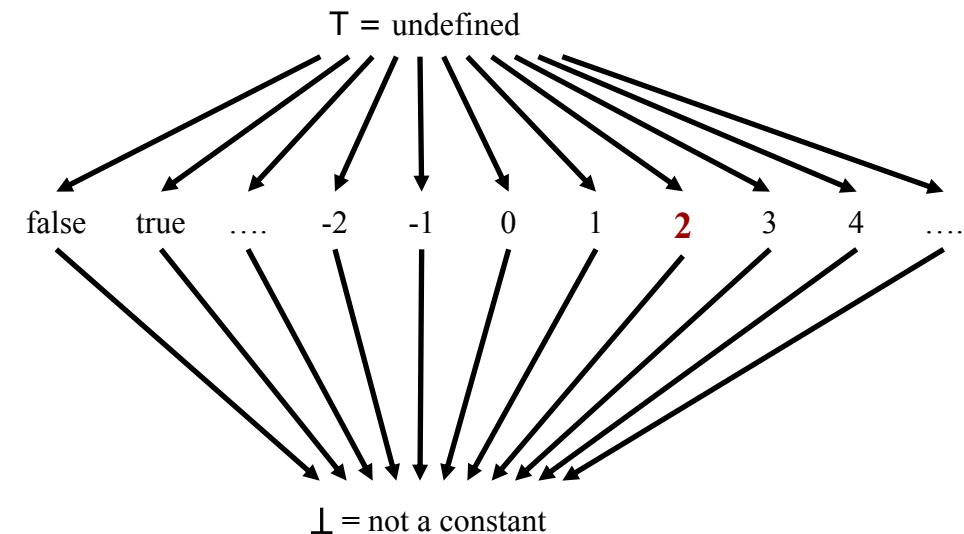
$$2 \wedge 2 =$$



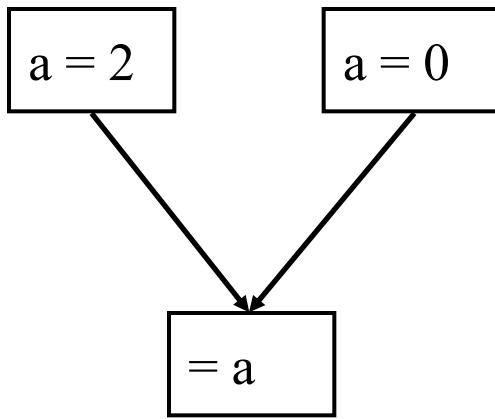
# Meet Operations on the Lattice



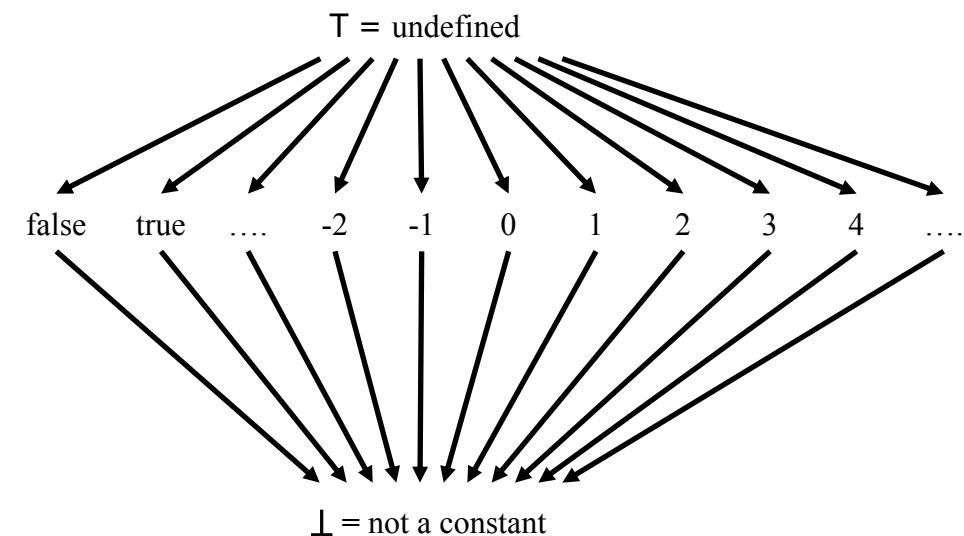
$$2 \wedge 2 = 2$$



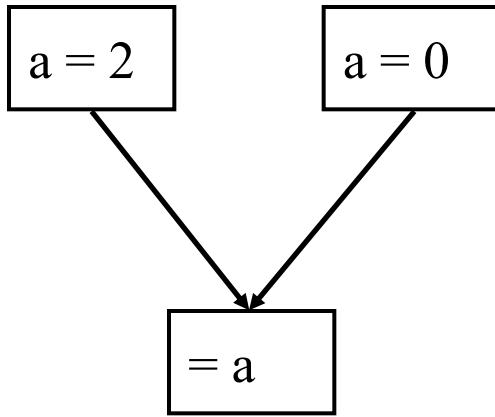
# Meet Operations on the Lattice



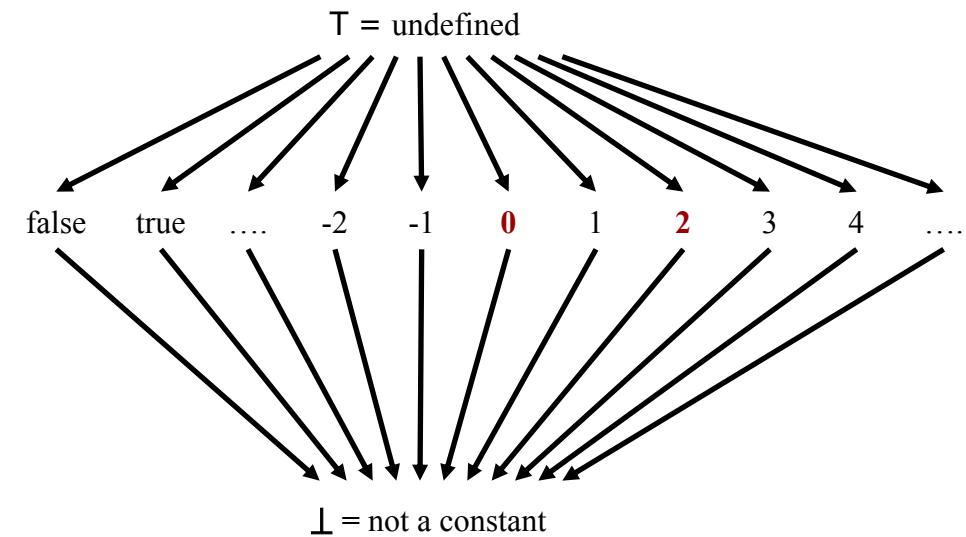
$$2 \wedge 0 =$$



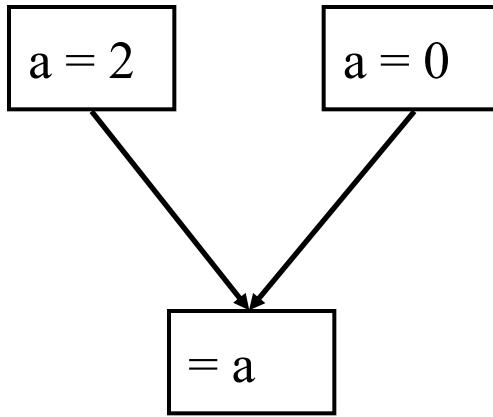
# Meet Operations on the Lattice



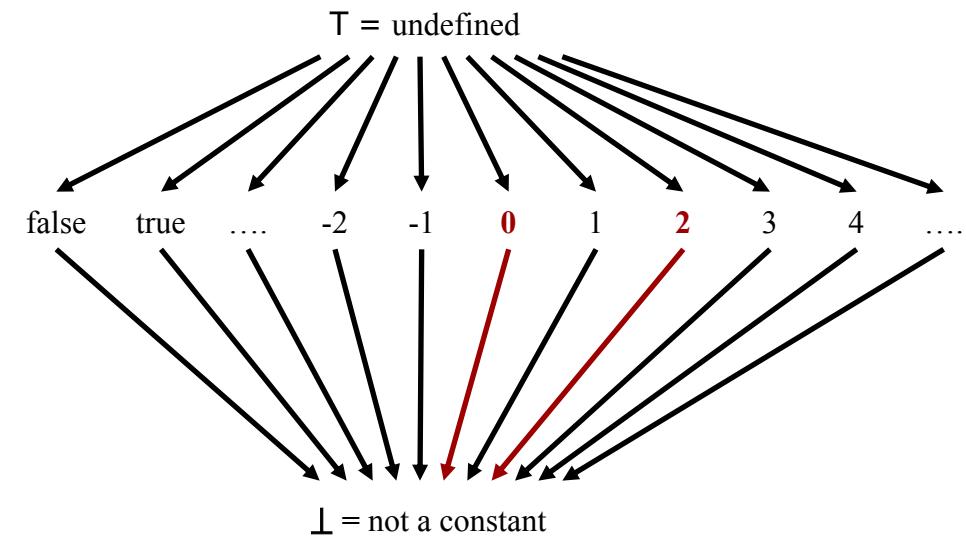
$$2 \wedge 0 =$$



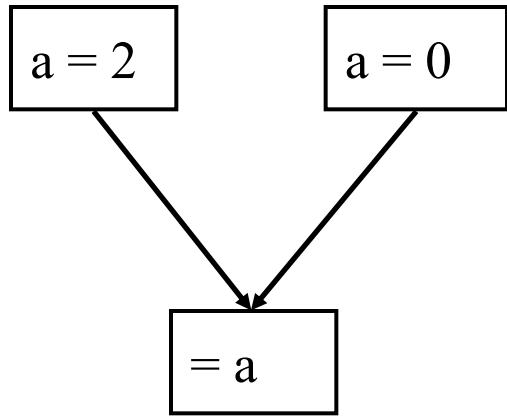
# Meet Operations on the Lattice



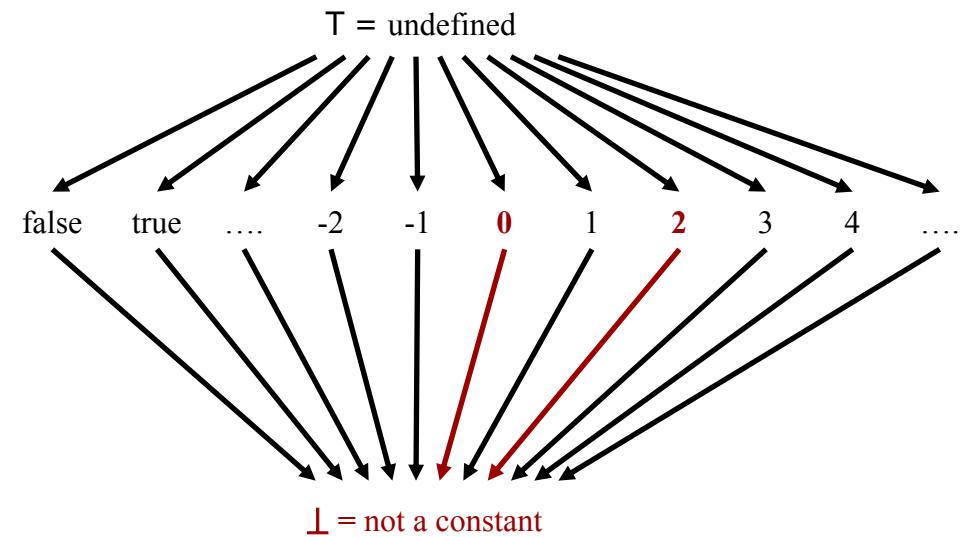
$$2 \wedge 0 =$$



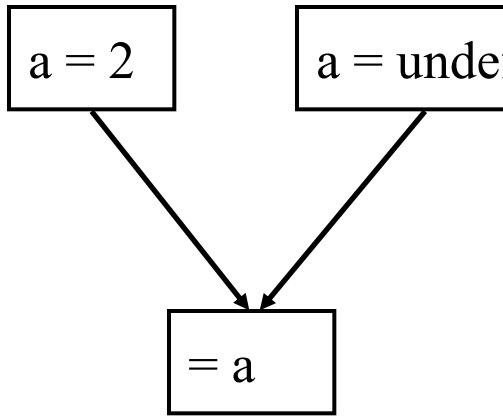
# Meet Operations on the Lattice



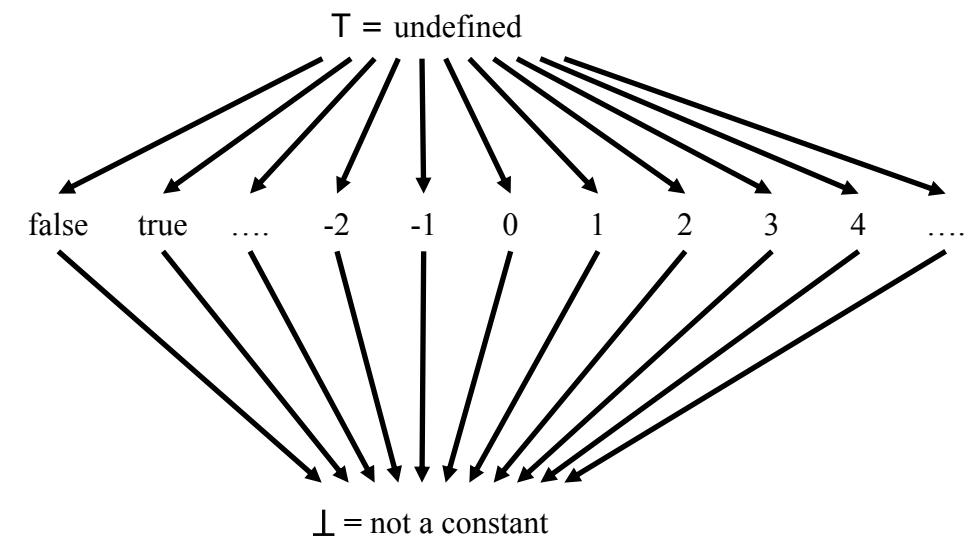
$$2 \wedge 0 = \text{not a constant}$$



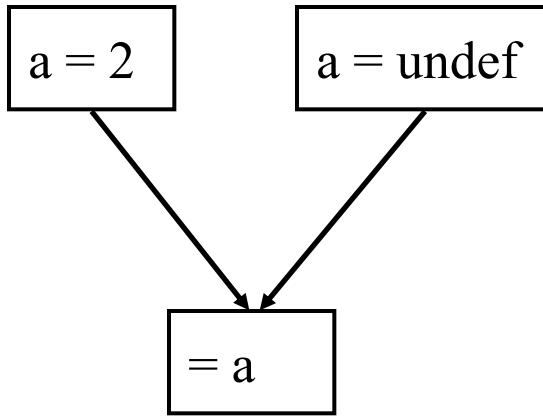
# Meet Operations on the Lattice



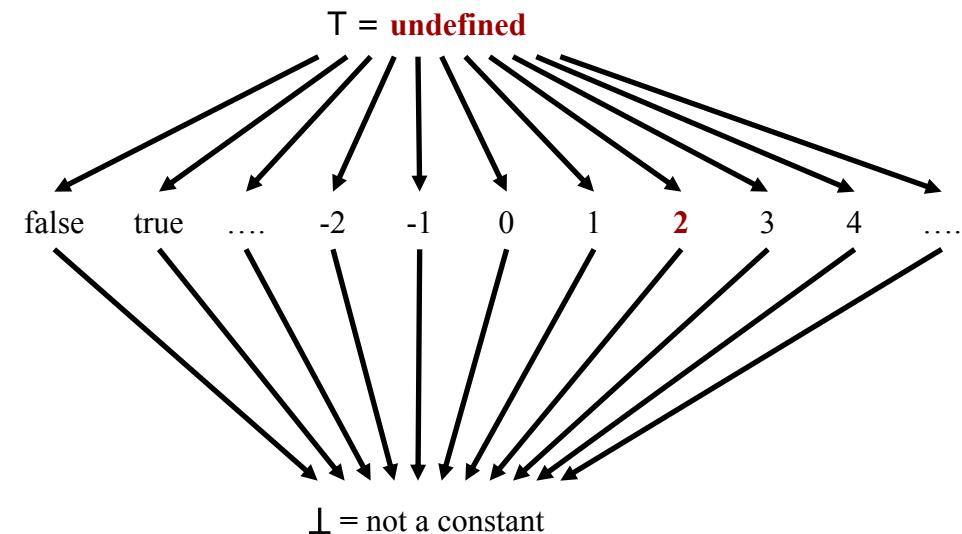
$$2 \wedge \text{undef} =$$



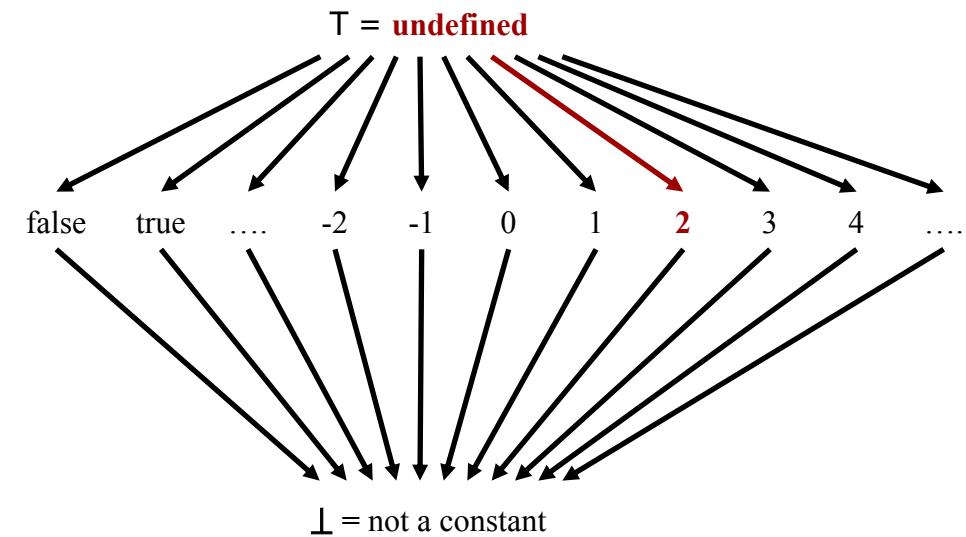
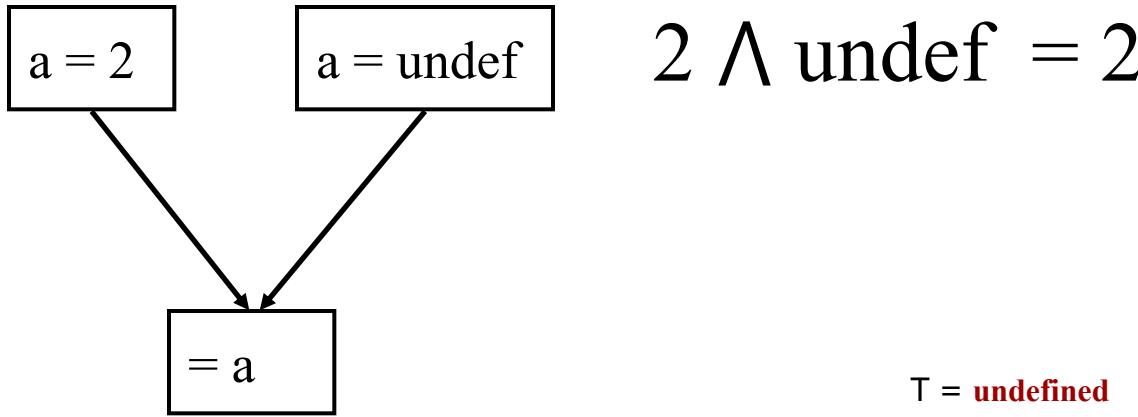
# Meet Operations on the Lattice



$$2 \wedge \text{undef} =$$



# Meet Operations on the Lattice



# Constant Propagation Data-Flow Problem

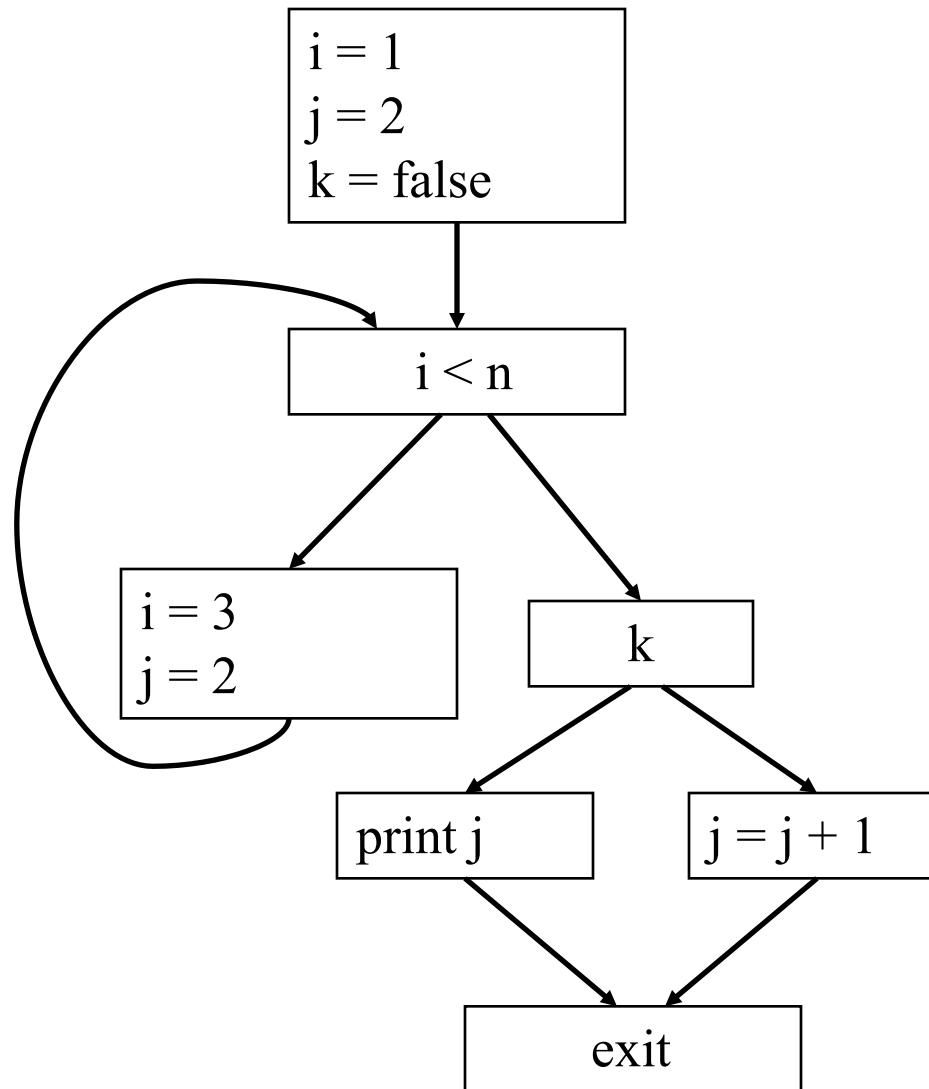
---

- Domain
  - For each variable a lattice  $L_v$
- Data-Flow Direction
  - Forward
- Data-Flow Function
  - $OUT = gen \wedge (IN \vee prsv)$

# Constant Propagation Data-Flow Problem

- Domain
  - For each variable a lattice  $L_v$
- Data-Flow Direction
  - Forward
- Data-Flow Function
  - $OUT = gen \vee (IN \wedge prsv)$
  - $gen = \left\{ x_v \mid \begin{array}{ll} T & \text{if } v \text{ is not LHS} \\ \perp & \text{otherwise} \end{array} \right. \text{ if } v \text{ is the LHS \& RHS value is a const.} \right\}$
  - $prsv = \left\{ x_v \mid \begin{array}{ll} T & \text{if } v \text{ is the LHS} \\ \perp & \text{if } v \text{ is not the LHS} \end{array} \right\}$
- **Obs:** For the  $prsv$  set, if a variable is not the LHS, we get  $\perp$  and the value of  $L_v$  is  $\perp$ ; For the  $gen$  set, if a variable is not the LHS, we get  $T$  and the value of  $L_v$  does not change.

# Example



# Example

```
i = 1  
j = 2  
k = false
```

$$\text{gen} = \{ \quad i:1, \quad j:2, \quad k:\text{false}, \quad n:T \}$$

$$\text{prsv} = \{ \quad i:T, \quad j:T, \quad k:T, \quad n: \perp \}$$

$$\text{gen} = \left\{ x_v \mid \begin{array}{ll} T & \text{if } v \text{ is not LHS} \\ x_v = \text{value} & \text{if } v \text{ is the LHS \& RHS value is a const.} \\ \perp & \text{otherwise} \end{array} \right\}$$
$$\text{prsv} = \left\{ x_v \mid \begin{array}{ll} T & \text{if } v \text{ is the LHS} \\ \perp & \text{if } v \text{ is not the LHS} \end{array} \right\}$$

# Example

```
j = j + 1
```

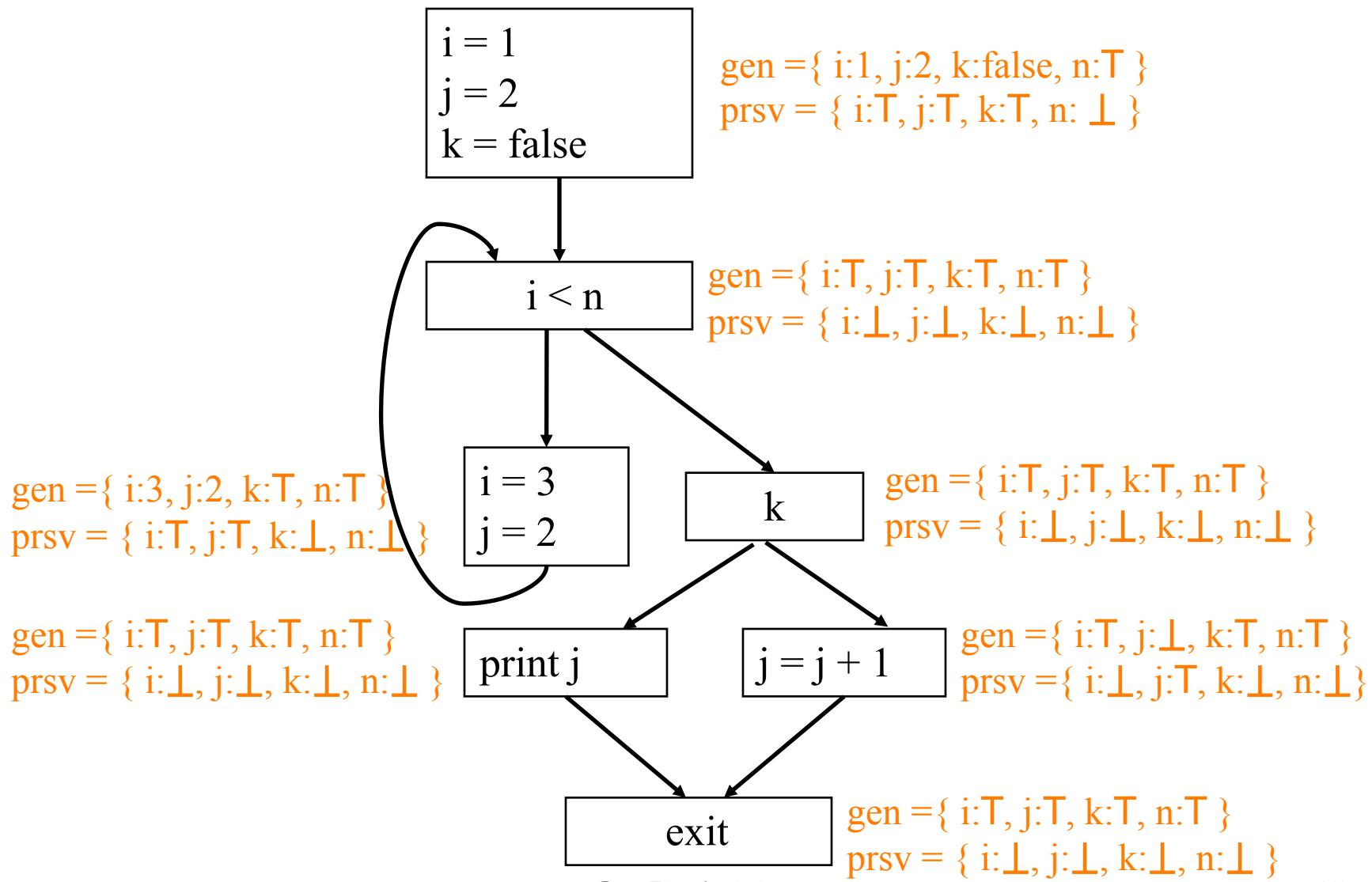
$$\text{gen} = \{ \quad i : T, \quad j : \perp, \quad k : T, \quad n : T \}$$

$$\text{prsv} = \{ \quad i : \perp, \quad j : T, \quad k : \perp, \quad n : \perp \}$$

$$\text{gen} = \left\{ x_v \mid \begin{array}{ll} T & \text{if } v \text{ is not LHS} \\ x_v = \text{value} & \text{if } v \text{ is the LHS \& RHS value is a const.} \\ \perp & \text{otherwise} \end{array} \right\}$$

$$\text{prsv} = \left\{ x_v \mid \begin{array}{ll} T & \text{if } v \text{ is the LHS} \\ \perp & \text{if } v \text{ is not the LHS} \end{array} \right\}$$

# Example



# Example

$\text{IN} = \{ i:\top, j:\top, k:\top, n:\top \}$

$i = 1$   
 $j = 2$   
 $k = \text{false}$

$\text{gen} = \{ i:1, j:2, k:\text{false}, n:\top \}$   
 $\text{prsv} = \{ i:\top, j:\top, k:\top, n:\perp \}$

$i < n$

$\text{gen} = \{ i:\top, j:\top, k:\top, n:\top \}$   
 $\text{prsv} = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$i = 3$   
 $j = 2$

$k$

$\text{gen} = \{ i:\top, j:\top, k:\top, n:\top \}$   
 $\text{prsv} = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$\text{gen} = \{ i:3, j:2, k:\top, n:\top \}$   
 $\text{prsv} = \{ i:\top, j:\top, k:\perp, n:\perp \}$

$\text{gen} = \{ i:\top, j:\top, k:\top, n:\top \}$   
 $\text{prsv} = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$\text{print } j$

$j = j + 1$

$\text{gen} = \{ i:\top, j:\perp, k:\top, n:\top \}$   
 $\text{prsv} = \{ i:\perp, j:\top, k:\perp, n:\perp \}$

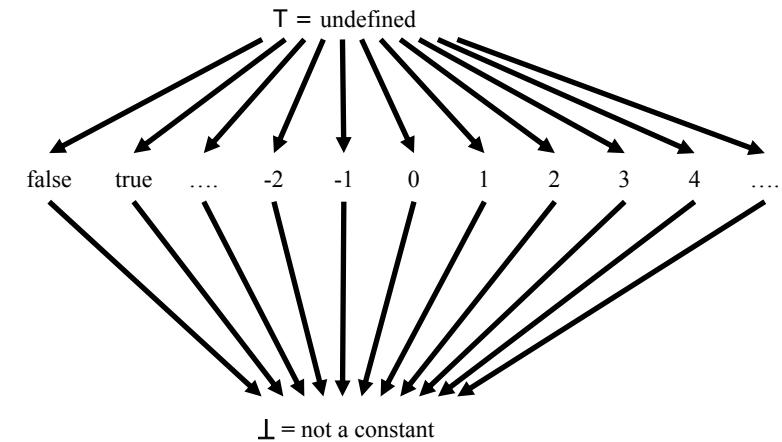
$\text{exit}$

$\text{gen} = \{ i:\top, j:\top, k:\top, n:\top \}$   
 $\text{prsv} = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

# Example

---

```
i = 1
j = 2
k = false
```



$$\text{gen} = \{ i:1, j:2, k:\text{false}, n:T \}$$

$$\text{prsv} = \{ i:T, j:T, k:T, n:\perp \}$$

$$\text{IN} = \{ i:T, j:T, k:T, n:T \}$$

$$\text{OUT} = \text{gen} \wedge (\text{IN} \vee \text{prsv})$$

$$\text{OUT} = \{ i:1, j:2, k:\text{false}, n:T \}$$

# Example

$\text{IN} = \{ i:T, j:T, k:T, n:T \}$

```
i = 1
j = 2
k = false
```

$\text{gen} = \{ i:1, j:2, k:\text{false}, n:T \}$   
 $\text{prsv} = \{ i:T, j:T, k:T, n:\perp \}$

$\text{OUT} = \{ i:1, j:2, k:\text{false}, n:T \}$

```
i < n
```

$\text{gen} = \{ i:T, j:T, k:T, n:T \}$   
 $\text{prsv} = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$\text{gen} = \{ i:3, j:2, k:T, n:T \}$   
 $\text{prsv} = \{ i:T, j:T, k:\perp, n:\perp \}$

```
i = 3
j = 2
```

```
k
```

$\text{gen} = \{ i:T, j:T, k:T, n:T \}$   
 $\text{prsv} = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$\text{gen} = \{ i:T, j:T, k:T, n:T \}$   
 $\text{prsv} = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

```
print j
```

```
j = j + 1
```

$\text{gen} = \{ i:T, j:\perp, k:T, n:T \}$   
 $\text{prsv} = \{ i:\perp, j:T, k:\perp, n:\perp \}$

```
exit
```

$\text{gen} = \{ i:T, j:T, k:T, n:T \}$   
 $\text{prsv} = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

# Example

IN = { i:T, j:T, k:T, n:T }

```
i = 1
j = 2
k = false
```

gen = { i:1, j:2, k:false, n:T }  
 prsv = { i:T, j:T, k:T, n: $\perp$  }

OUT = { i:1, j:2, k:false, n:T }

```
i < n
```

gen = { i:T, j:T, k:T, n:T }  
 prsv = { i: $\perp$ , j: $\perp$ , k: $\perp$ , n: $\perp$  }

gen = { i:3, j:2, k:T, n:T }  
 prsv = { i:T, j:T, k: $\perp$ , n: $\perp$  }

```
i = 3
j = 2
```

```
k
```

gen = { i:T, j:T, k:T, n:T }  
 prsv = { i: $\perp$ , j: $\perp$ , k: $\perp$ , n: $\perp$  }

OUT = { i:T, j:T, k:T, n:T }

gen = { i:T, j:T, k:T, n:T }  
 prsv = { i: $\perp$ , j: $\perp$ , k: $\perp$ , n: $\perp$  }

```
print j
```

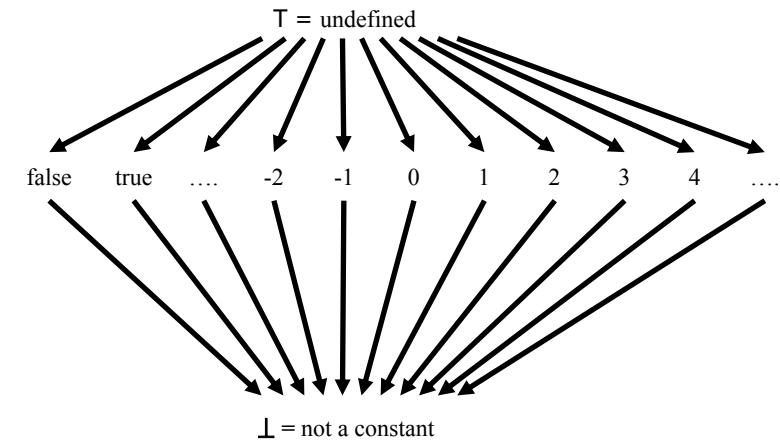
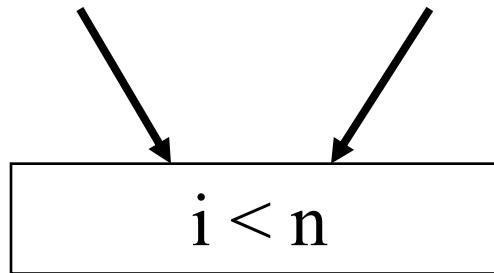
```
j = j + 1
```

gen = { i:T, j: $\perp$ , k:T, n:T }  
 prsv = { i: $\perp$ , j:T, k: $\perp$ , n: $\perp$  }

```
exit
```

gen = { i:T, j:T, k:T, n:T }  
 prsv = { i: $\perp$ , j: $\perp$ , k: $\perp$ , n: $\perp$  }

# Example



$\text{out1} = \{ i:T, j:T, k:T, n:T \}$

$\text{out2} = \{ i:1, j:2, k:\text{false}, n:T \}$

$\text{IN} = \text{out1} \wedge \text{out2}$

$\text{IN} = \{ i:1, j:2, k:\text{false}, n:T \}$

# Example

$IN = \{ i:T, j:T, k:T, n:T \}$

```
i = 1
j = 2
k = false
```

$gen = \{ i:1, j:2, k:false, n:T \}$   
 $prsv = \{ i:T, j:T, k:T, n:\perp \}$

$OUT = \{ i:1, j:2, k:false, n:T \}$

$IN = \{ i:1, j:2, k:false, n:T \}$

$i < n$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$OUT = \{ i:1, j:2, k:false, n:T \}$

$gen = \{ i:3, j:2, k:T, n:T \}$

$prsv = \{ i:T, j:T, k:\perp, n:\perp \}$

```
i = 3
j = 2
```

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$OUT = \{ i:T, j:T, k:T, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

print j

$j = j + 1$

$gen = \{ i:T, j:\perp, k:T, n:T \}$

$prsv = \{ i:\perp, j:T, k:\perp, n:\perp \}$

exit

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

# Example

$IN = \{ i:T, j:T, k:T, n:T \}$

```
i = 1
j = 2
k = false
```

$gen = \{ i:1, j:2, k:false, n:T \}$   
 $prsv = \{ i:T, j:T, k:T, n:\perp \}$

$OUT = \{ i:1, j:2, k:false, n:T \}$

$IN = \{ i:1, j:2, k:false, n:T \}$

$i < n$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$OUT = \{ i:1, j:2, k:false, n:T \}$

$IN = \{ i:1, j:2, k:false, n:T \}$

$gen = \{ i:3, j:2, k:T, n:T \}$

$prsv = \{ i:T, j:T, k:\perp, n:\perp \}$

$OUT = \{ i:3, j:2, k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$i = 3$   
 $j = 2$

$k$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$print\ j$

$j = j + 1$

$gen = \{ i:T, j:\perp, k:T, n:T \}$

$prsv = \{ i:\perp, j:T, k:\perp, n:\perp \}$

$exit$

$gen = \{ i:T, j:T, k:T, n:T \}$   
 $prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

# Example

$IN = \{ i:T, j:T, k:T, n:T \}$

```
i = 1
j = 2
k = false
```

$gen = \{ i:1, j:2, k:false, n:T \}$   
 $prsv = \{ i:T, j:T, k:T, n:\perp \}$

$OUT = \{ i:1, j:2, k:false, n:T \}$

$IN = \{ i:1, j:2, k:false, n:T \}$

$i < n$

$gen = \{ i:T, j:T, k:T, n:T \}$   
 $prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$OUT = \{ i:1, j:2, k:false, n:T \}$

$IN = \{ i:1, j:2, k:false, n:T \}$

$gen = \{ i:3, j:2, k:T, n:T \}$   
 $prsv = \{ i:T, j:T, k:\perp, n:\perp \}$

$OUT = \{ i: \perp, j:2, k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$i = 3$   
 $j = 2$

$k$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$print\ j$

$j = j + 1$

$gen = \{ i:T, j:\perp, k:T, n:T \}$

$prsv = \{ i:\perp, j:T, k:\perp, n:\perp \}$

$exit$

$gen = \{ i:T, j:T, k:T, n:T \}$   
 $prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

# Example

$IN = \{ i:T, j:T, k:T, n:T \}$

```
i = 1
j = 2
k = false
```

$gen = \{ i:1, j:2, k:false, n:T \}$   
 $prsv = \{ i:T, j:T, k:T, n:\perp \}$

$OUT = \{ i:1, j:2, k:false, n:T \}$

$IN = \{ i:\perp, j:2, k:false, n:T \}$

$i < n$

$gen = \{ i:T, j:T, k:T, n:T \}$   
 $prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$OUT = \{ i:\perp, j:2, k:false, n:T \}$

$IN = \{ i:\perp, j:2, k:false, n:T \}$

$gen = \{ i:3, j:2, k:T, n:T \}$

$prsv = \{ i:T, j:T, k:\perp, n:\perp \}$

```
i = 3
j = 2
```

```
k
```

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$OUT = \{ i:\perp, j:2, k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

print j

$j = j + 1$

$gen = \{ i:T, j:\perp, k:T, n:T \}$

$prsv = \{ i:\perp, j:T, k:\perp, n:\perp \}$

exit

$gen = \{ i:T, j:T, k:T, n:T \}$   
 $prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

# Example

$IN = \{ i:T, j:T, k:T, n:T \}$

```
i = 1
j = 2
k = false
```

$gen = \{ i:1, j:2, k:false, n:T \}$   
 $prsv = \{ i:T, j:T, k:T, n:\perp \}$

$OUT = \{ i:1, j:2, k:false, n:T \}$

```
i < n
```

$gen = \{ i:T, j:T, k:T, n:T \}$   
 $prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$OUT = \{ i:\perp, j:2, k:false, n:T \}$

$IN = \{ i:\perp, j:2, k:false, n:T \}$

$gen = \{ i:3, j:2, k:T, n:T \}$

$prsv = \{ i:T, j:T, k:\perp, n:\perp \}$

$OUT = \{ i:\perp, j:2, k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

```
i = 3
j = 2
```

k

$IN = \{ i:\perp, j:2, k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$OUT = \{ i:\perp, j:2, k:false, n:T \}$

$gen = \{ i:T, j:\perp, k:T, n:T \}$

$prsv = \{ i:\perp, j:T, k:\perp, n:\perp \}$

print j

j = j + 1

exit

$gen = \{ i:T, j:T, k:T, n:T \}$   
 $prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

# Example

$IN = \{ i:T, j:T, k:T, n:T \}$

```
i = 1
j = 2
k = false
```

$gen = \{ i:1, j:2, k:false, n:T \}$   
 $prsv = \{ i:T, j:T, k:T, n:\perp \}$

$OUT = \{ i:1, j:2, k:false, n:T \}$

```
i < n
```

$gen = \{ i:T, j:T, k:T, n:T \}$   
 $prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$OUT = \{ i:\perp, j:2, k:false, n:T \}$

$IN = \{ i:\perp, j:2, k:false, n:T \}$

$gen = \{ i:3, j:2, k:T, n:T \}$

$prsv = \{ i:T, j:T, k:\perp, n:\perp \}$

$OUT = \{ i:\perp, j:2, k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

```
i = 3
j = 2
```

```
k
```

$IN = \{ i:\perp, j:2, k:false, n:T \}$   
 $gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$OUT = \{ i:\perp, j:2, k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

print j

$j = j + 1$

$OUT = \{ i:\perp, j:2, k:false, n:T \}$

exit

$gen = \{ i:T, j:T, k:T, n:T \}$   
 $prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

# Example

$IN = \{ i:T, j:T, k:T, n:T \}$

```
i = 1
j = 2
k = false
```

$gen = \{ i:1, j:2, k:false, n:T \}$   
 $prsv = \{ i:T, j:T, k:T, n:\perp \}$

$OUT = \{ i:1, j:2, k:false, n:T \}$

```
IN = {i: $\perp$ , j:2, k:false, n:T}
i < n
```

$gen = \{ i:T, j:T, k:T, n:T \}$   
 $prsv = \{ i: $\perp$ , j: $\perp$ , k: $\perp$ , n: $\perp \}$$

$OUT = \{ i: $\perp$ , j:2, k:false, n:T \}$

$IN = \{ i: $\perp$ , j:2, k:false, n:T \}$

$gen = \{ i:3, j:2, k:T, n:T \}$

$prsv = \{ i:T, j:T, k: $\perp$ , n: $\perp \}$$

$OUT = \{ i: $\perp$ , j:2, k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i: $\perp$ , j: $\perp$ , k: $\perp$ , n: $\perp \}$$

```
i = 3
j = 2
```

```
k
```

$IN = \{ i: $\perp$ , j:2, k:false, n:T \}$   
 $gen = \{ i:T, j:T, k:T, n:T \}$   
 $prsv = \{ i: $\perp$ , j: $\perp$ , k: $\perp$ , n: $\perp \}$$

$OUT = \{ i: $\perp$ , j:2, k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i: $\perp$ , j: $\perp$ , k: $\perp$ , n: $\perp \}$$

$IN = \{ i: $\perp$ , j:2, k:false, n:T \}$

$OUT = \{ i: $\perp$ , j:2, k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i: $\perp$ , j: $\perp$ , k: $\perp$ , n: $\perp \}$$

$IN = \{ i: $\perp$ , j:2, k:false, n:T \}$

$OUT = \{ i: $\perp$ , j:2, k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i: $\perp$ , j: $\perp$ , k: $\perp$ , n: $\perp \}$$

$IN = \{ i: $\perp$ , j:2, k:false, n:T \}$

$OUT = \{ i: $\perp$ , j: $\perp$ , k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i: $\perp$ , j: $\perp$ , k: $\perp$ , n: $\perp \}$$

Data-Flow Analysis

# Example

$IN = \{ i:T, j:T, k:T, n:T \}$

```
i = 1
j = 2
k = false
```

$gen = \{ i:1, j:2, k:false, n:T \}$   
 $prsv = \{ i:T, j:T, k:T, n:\perp \}$

$OUT = \{ i:1, j:2, k:false, n:T \}$

```
i < n
```

$gen = \{ i:T, j:T, k:T, n:T \}$   
 $prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$OUT = \{ i:\perp, j:2, k:false, n:T \}$

$IN = \{ i:\perp, j:2, k:false, n:T \}$

$gen = \{ i:3, j:2, k:T, n:T \}$

$prsv = \{ i:T, j:T, k:\perp, n:\perp \}$

$OUT = \{ i:\perp, j:2, k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

```
i = 3
j = 2
```

```
k
```

$IN = \{ i:\perp, j:2, k:false, n:T \}$

$gen = \{ i:T, j:T, k:T, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

$OUT = \{ i:\perp, j:2, k:false, n:T \}$

$gen = \{ i:\perp, j:2, k:false, n:T \}$

$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

```
print j
```

```
j = j + 1
```

$OUT = \{ i:\perp, j:2, k:false, n:T \}$

$OUT = \{ i:\perp, j:\perp, k:false, n:T \}$

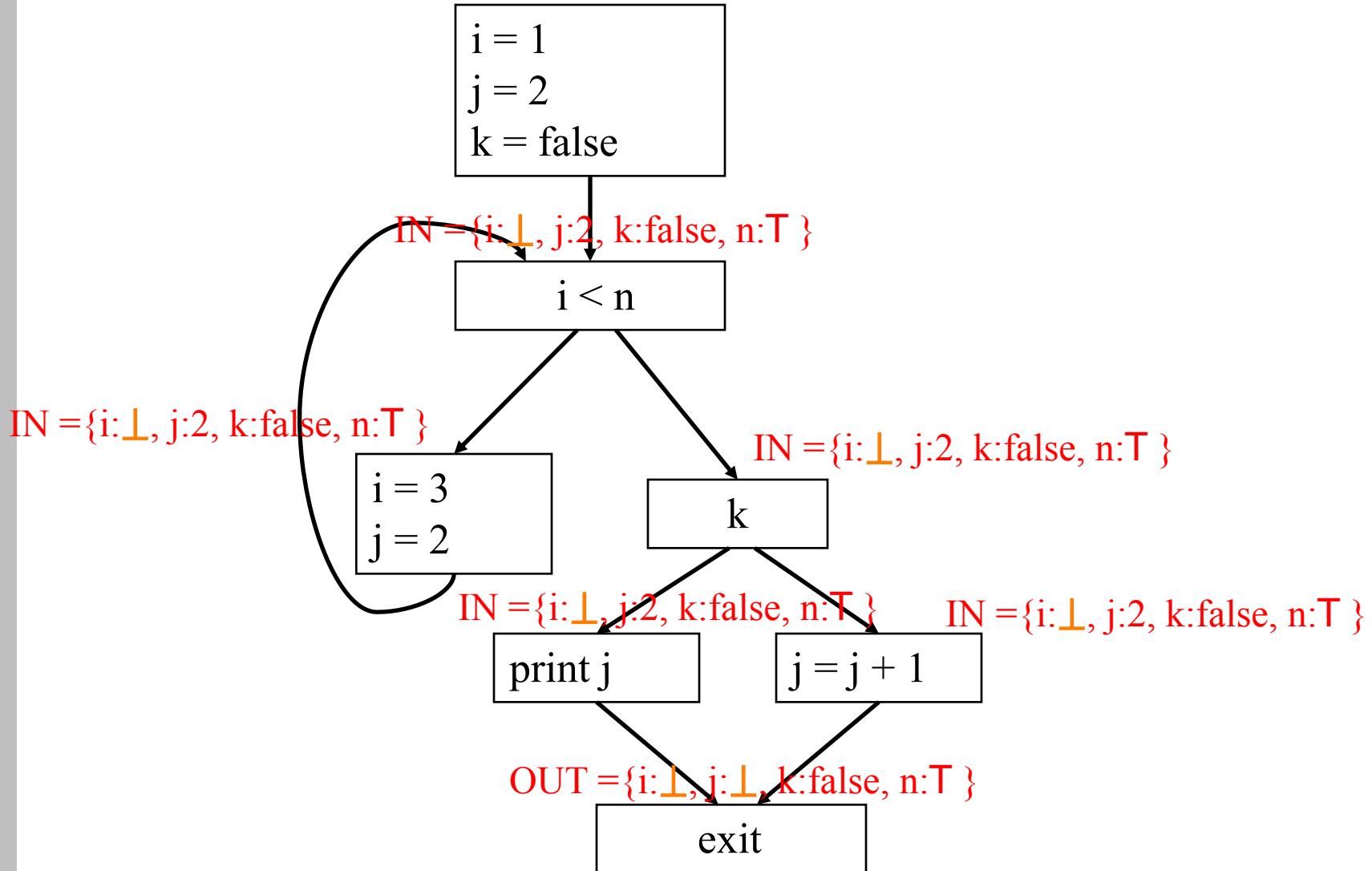
exit

$gen = \{ i:T, j:T, k:T, n:T \}$

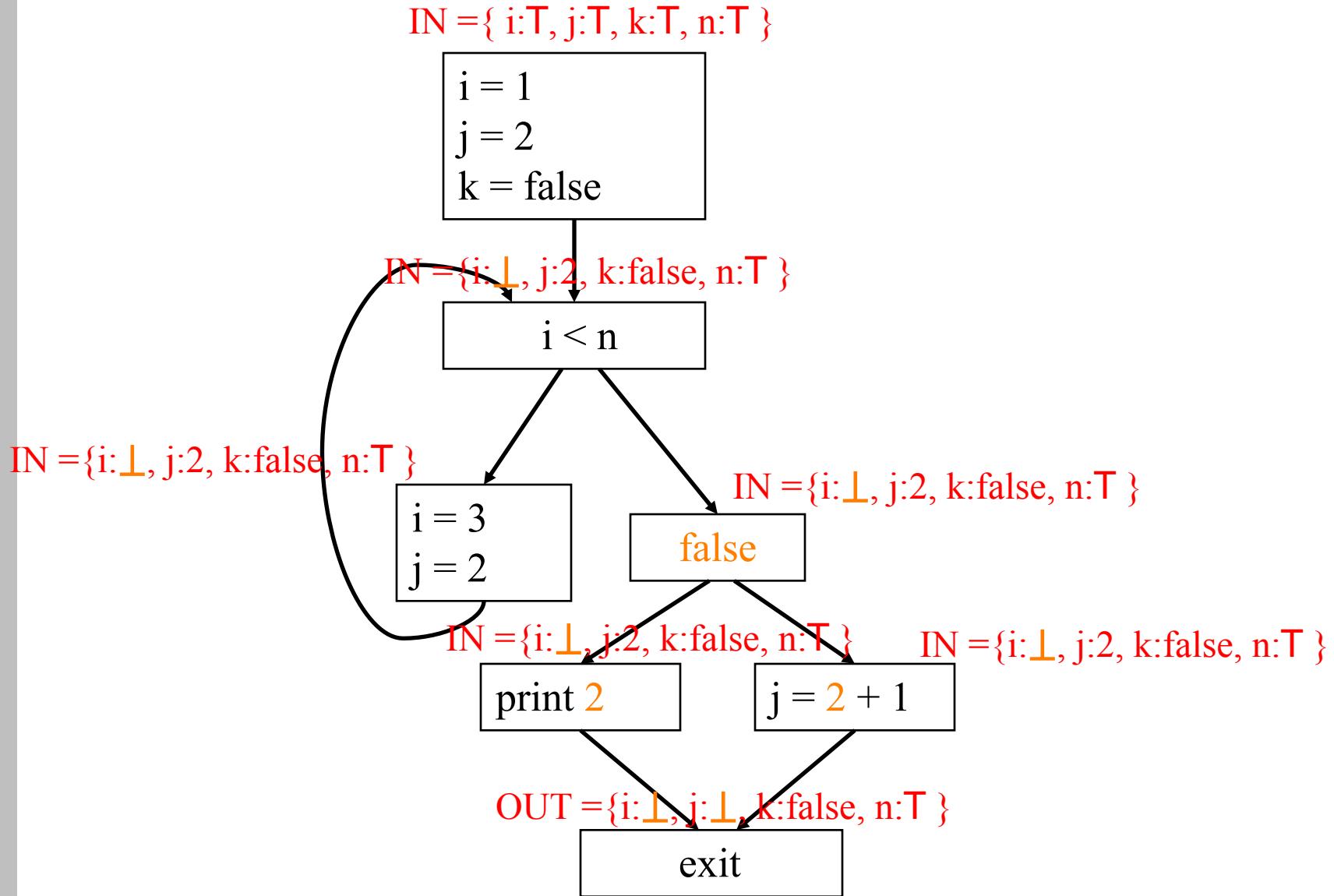
$prsv = \{ i:\perp, j:\perp, k:\perp, n:\perp \}$

# Example

IN = { i:T, j:T, k:T, n:T }

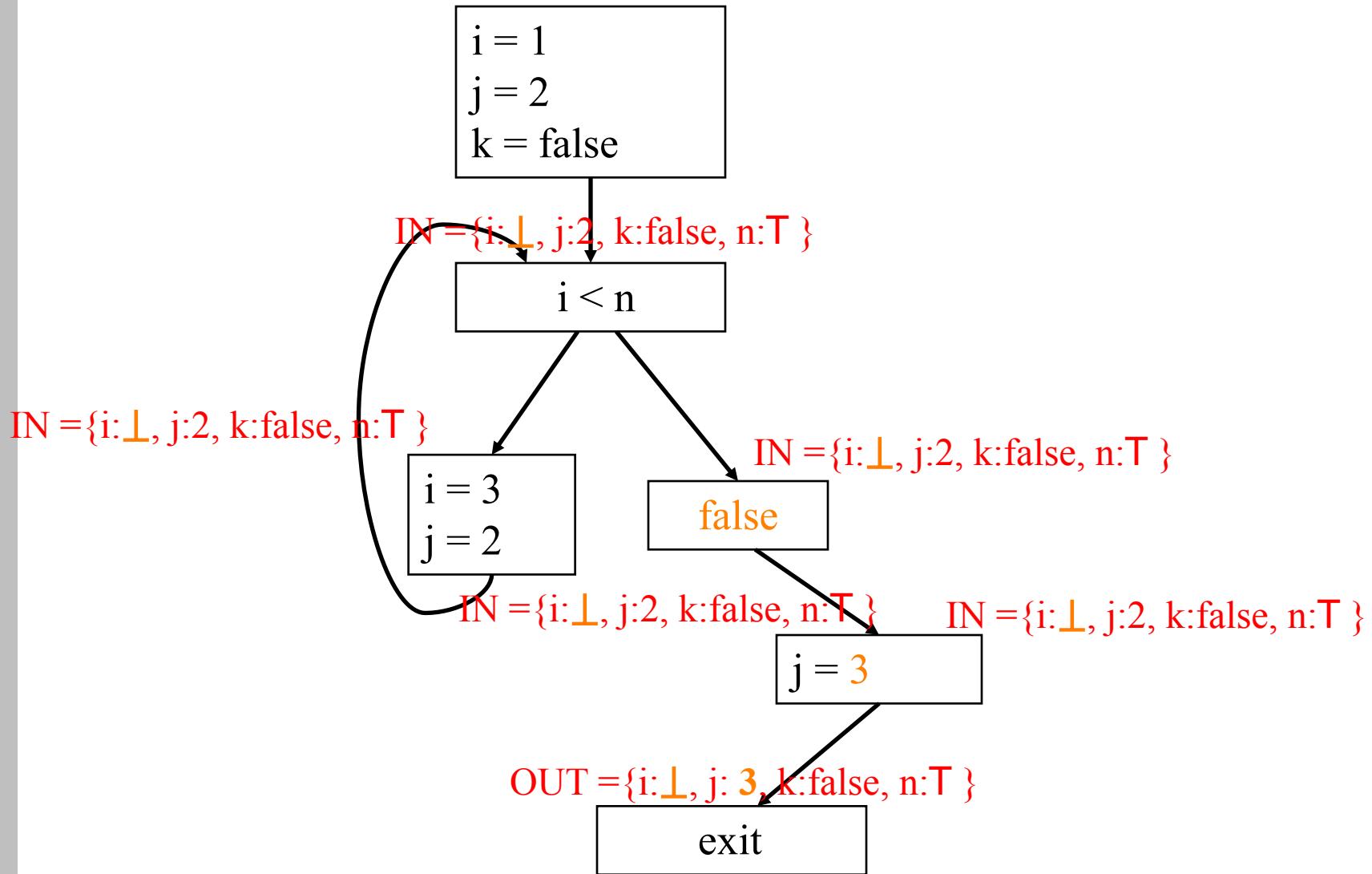


# Example



# Example

IN = { i:T, j:T, k:false, n:T }



# Summary

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- Overview of Control-Flow Analysis
- Algebraic Simplification
- Copy Propagation
- Constant Propagation