CodelT

Lesson 2 Objectives:

To gain an understanding of:

- What are priority queues and how we may efficiently implement them
- 2. What are and why we need the heap (binary heap) data structure
- 3. How to analyse and implement heaps
- 4. How to implement heaps with changing keys



Recap: Queues

Queues are FIRST IN FIRST OUT

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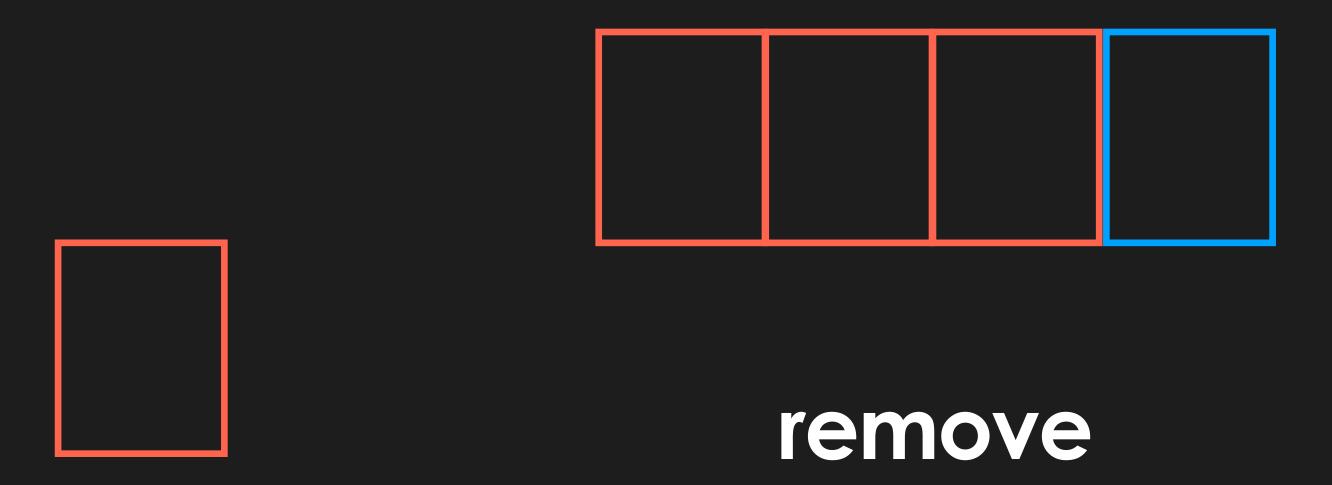
Queues are FIRST IN FIRST OUT



insert



Queues are FIRST IN FIRST OUT





Stock	Estimated Returns
Google	30%
Amazon	20%
Apple	60%
Microsoft	40%
Netflix	25%

Imagine we had the following list of company stocks and their estimated returns on investment (higher is better)

Front
Back

Google	30%
Amazon	20%
Apple	60%
Microsoft	40%
Netflix	25%



Front
Back

Google

Google	30%
Amazon	20%
Apple	60%
Microsoft	40%
Netflix	25%



Front
Back

Google Amazon

Google	30%
Amazon	20%
Apple	60%
Microsoft	40%
Netflix	25%



Front
Back

Apple	Google	Amazon
-------	--------	--------

Google	30%
Amazon	20%
Apple	60%
Microsoft	40%
Netflix	25%



Front
Back

Apple	Microsoft	Google	Amazon
-------	-----------	--------	--------

Google	30%
Amazon	20%
Apple	60%
Microsoft	40%
Netflix	25%



Front				Back
Apple	Microsoft	Google	Netflix	Amazon

Google	30%
Amazon	20%
Apple	60%
Microsoft	40%
Netflix	25%



This is called a **Priority Queue**. We created a queue of stocks with the priority being its returns. How can we create such a structure?

Method 1: Linked Lists

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Insertion: To insert into a PQ linked list, we can simply insert at the front. This takes **O(1)** time



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Insertion: To insert into a PQ linked list, we can simply insert at the front. This takes **O(1)** time

GetMax: To remove the item with highest priority at any time, traverse the list to find the max item using linear search, and remove. This takes **O(N)** time

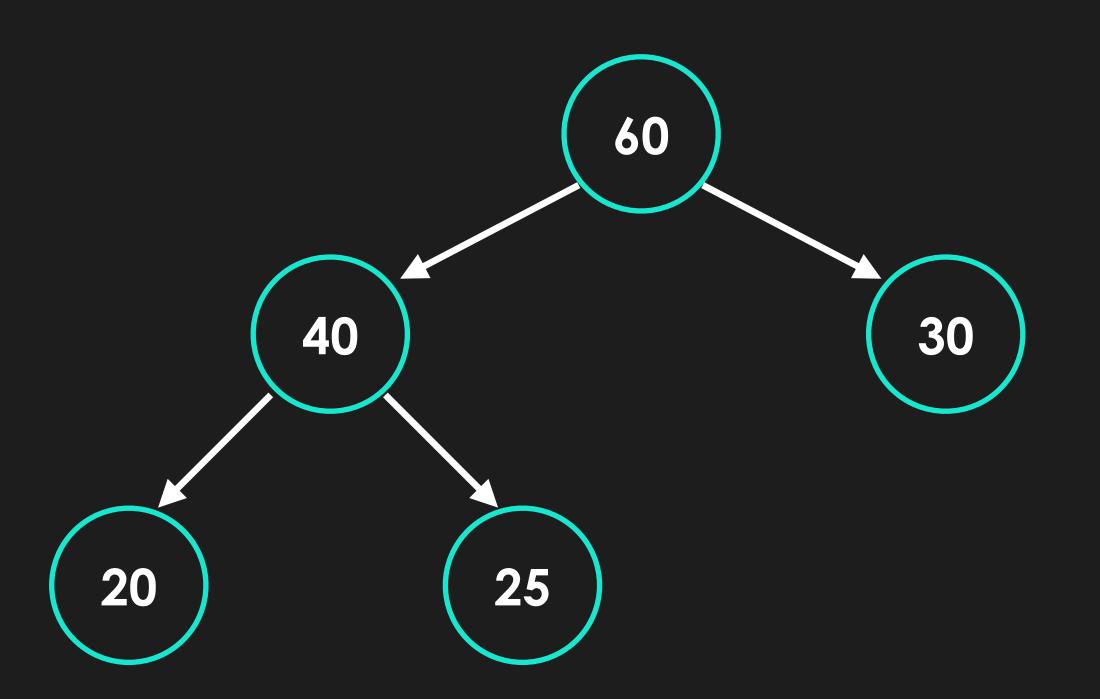


Method 2: Binary Heaps

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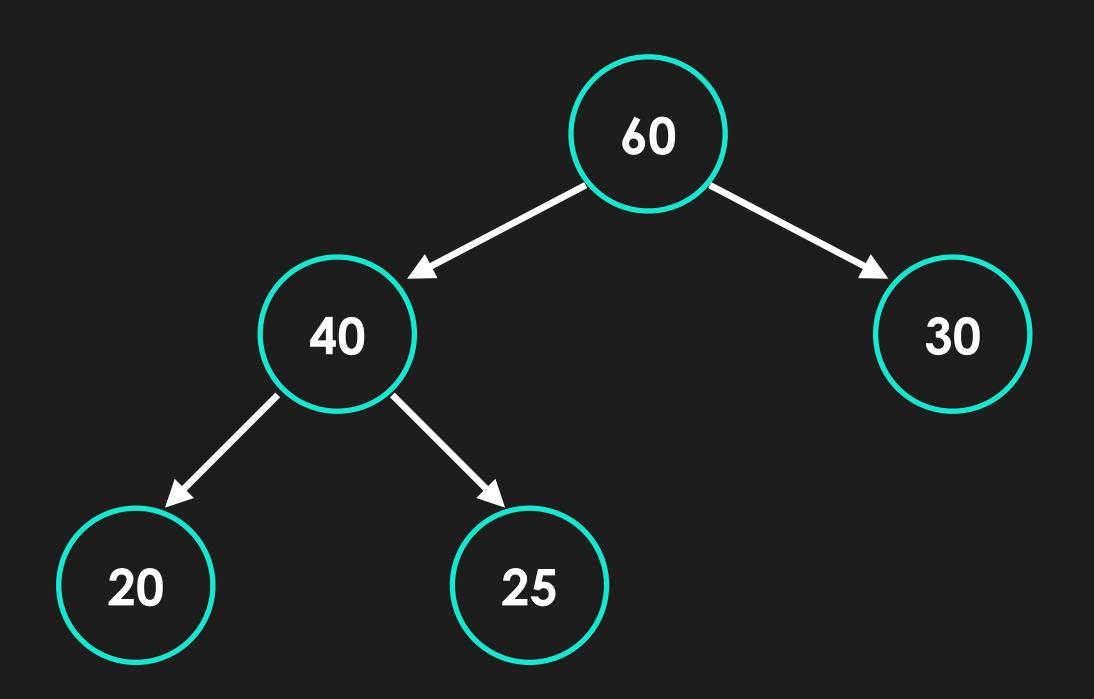
Binary Heaps make use of the binary tree structure to implement the priority queue logic





Google	30%
Amazon	20%
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Netflix	25%

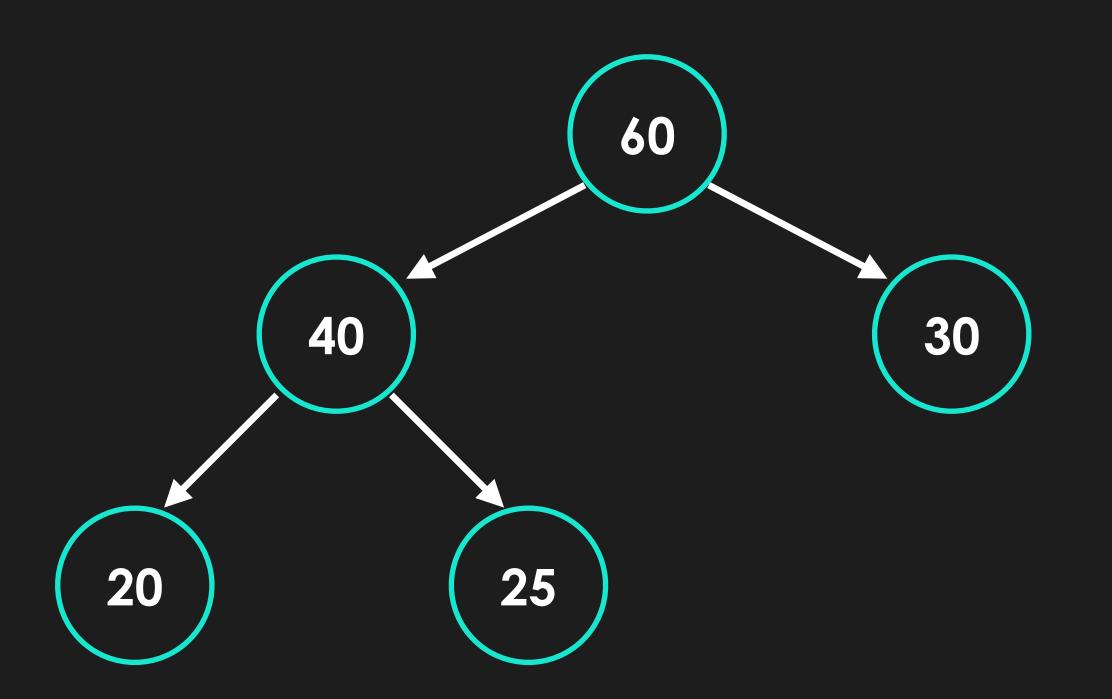




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We will represent each node as it's **priority** (return)



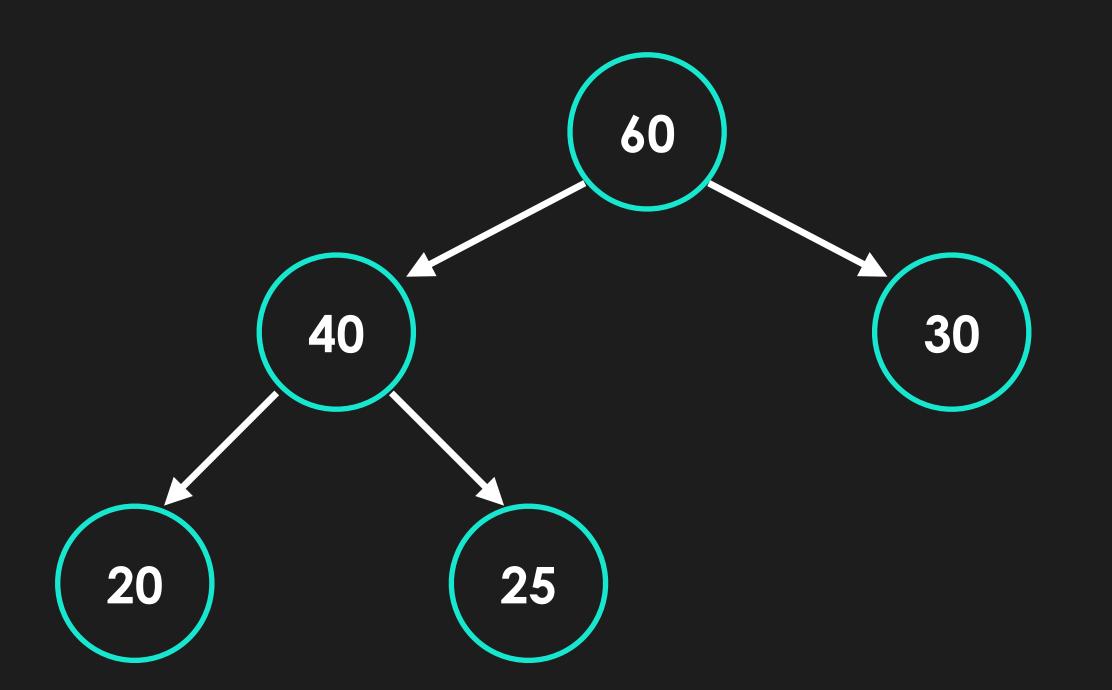


Google	30%
Amazon	20%
Apple	60%
Microsoft	40%
Netflix	25%

Rules for Binary Heap:

- 1. Root node must be entry with max priority
- 2. For each node, every child node below must have lower priority

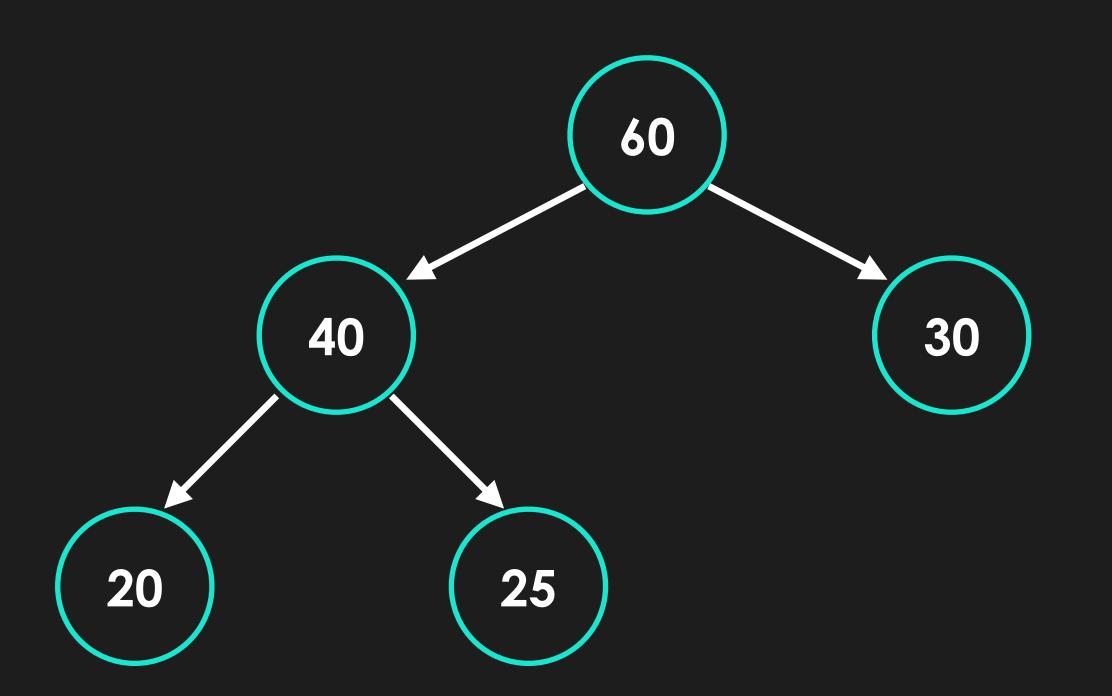




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Amazon	20%
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Microsoft	40%
Netflix	25%

The above is also called a **max heap** because the priority is given to returns of **higher value**





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Amazon	20%
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Microsoft	40%
Netflix	25%

Binary Heaps are a more efficient implementation of priority queues, and we will soon see why!

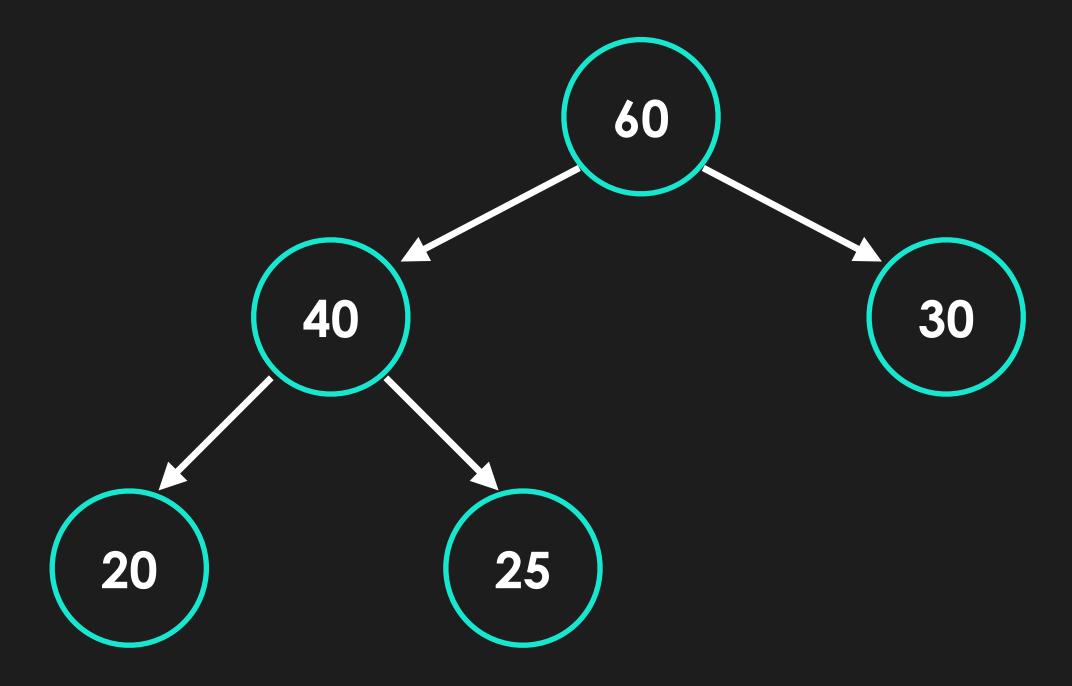


Binary Heap Representation

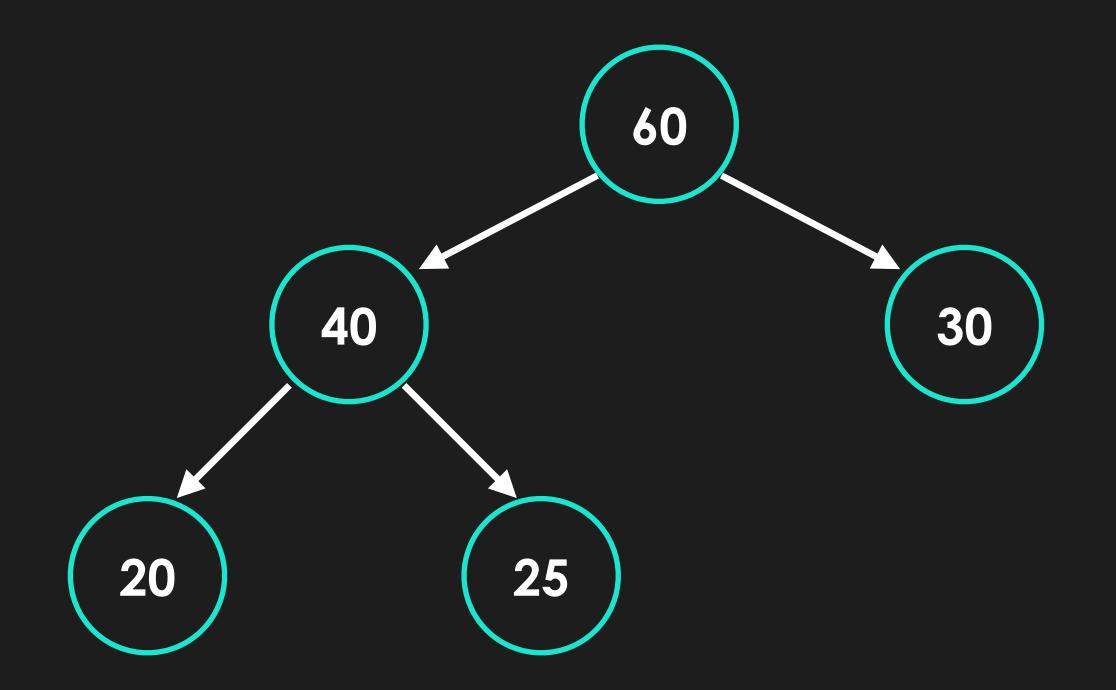
Heapitem Class

```
class HeapItem:
    def __init__(self, key, value):
        self.key = key
        self.value = value
```



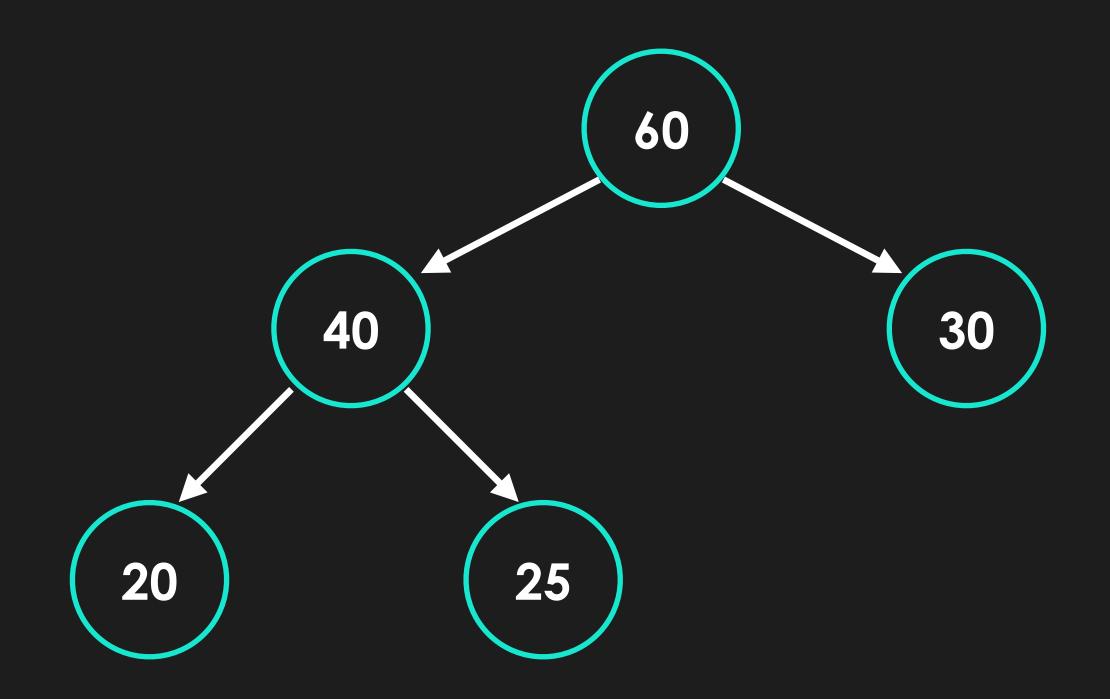


To represent binary heaps, we often use arrays (lists)



We only show the **value** of each Heapltem in the visualisation. However, each element in the heap is really a **Heapltem** class

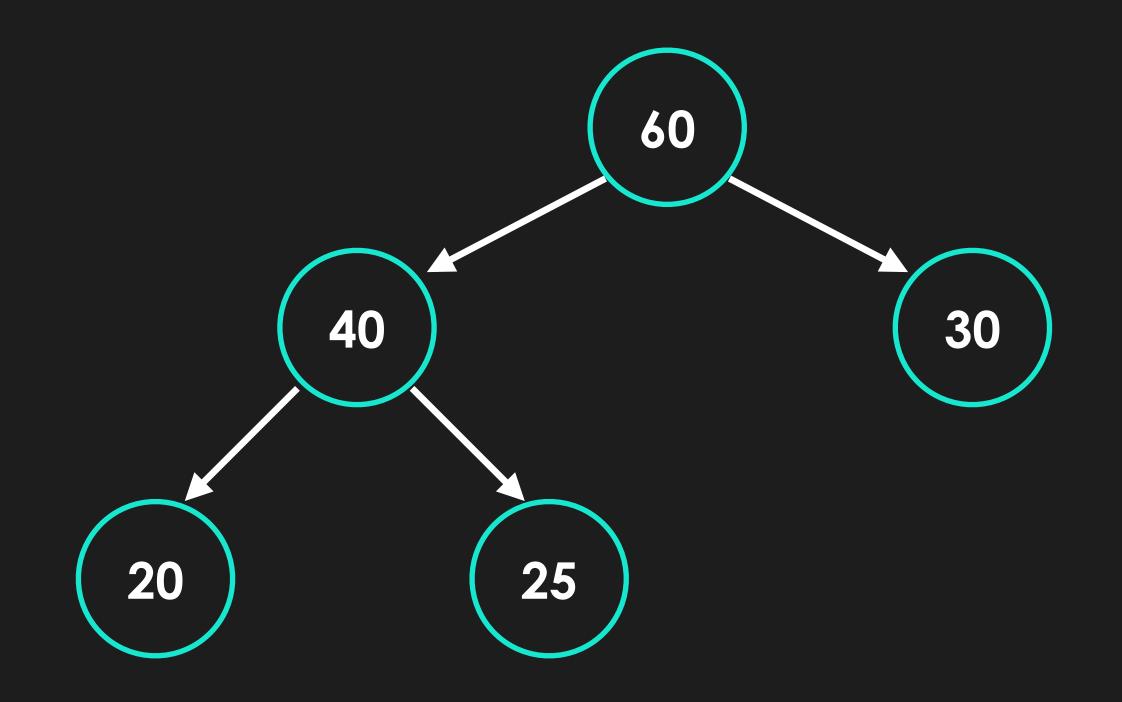


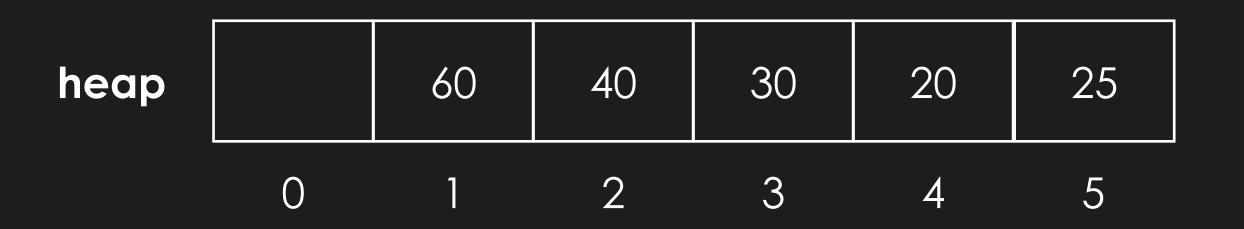


heap 60 40 30 20 25 0 1 2 3 4 5

For a node at index N:

- It's left child node will be at index 2 * N
- It's right child node will be at index (2 * N + 1)





For a node at index N:

- It's **left child node** will be at index 2 * N
- It's right child node will be at index (2 * N + 1)

The **root node** will always hold **index 1** and we leave index 0 of our list empty to make things easier



Implementation of Max Heap

API

```
class MaxHeap:
    def __init__ (self, capacity):
        self.capacity = capacity
        self.heap = [None] * (capacity + 1)
    def sink(self, index):
        pass
    def swim (self, index):
        pass
    def insert(self, index):
        pass
    def getMax(self, index):
        pass
```



Helpers

```
# HELPER FUNCTIONS
def parent (index):
    return index // 2
def left (index):
    return index * 2
def right (index):
    return index * 2 + 1
def swap (array, index1, index2):
    array[index1], array[index2] = array[index2], array[index1]
```



Max Heap Construction



Initialisation

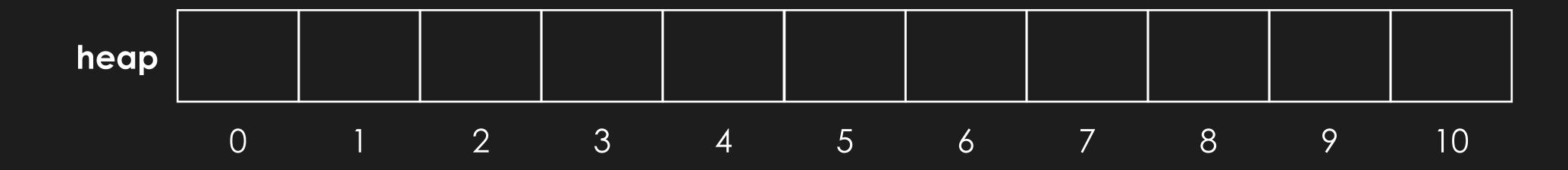
pq = MaxHeap(10)

```
def __init__ (self, maxsize):
    self.maxsize = maxsize
    self.size = 0
    self.heap = [None] * (maxsize + 1)
```



```
pq = MaxHeap(10)
```

```
def __init__ (self, maxsize):
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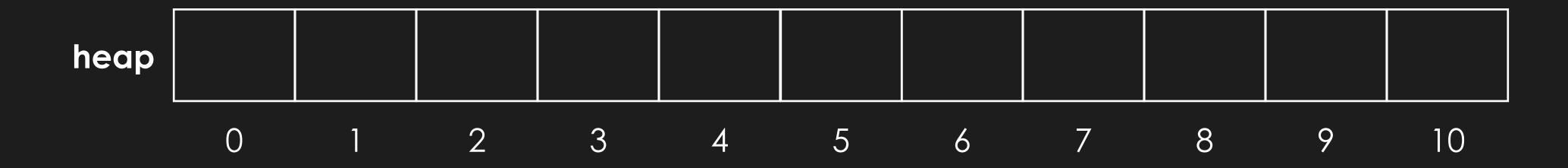


Insertion

STEP 1: Check if capacity is reached

pq.insert(30)

```
if self.size >= self.maxsize:
    print('size limit reached')
    return
```

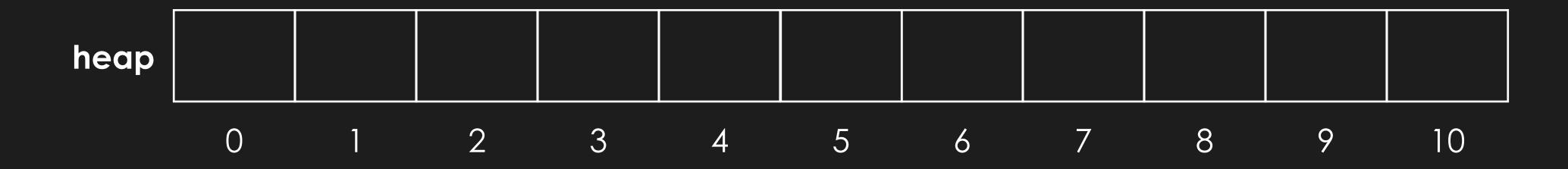




STEP 2: Increment size and insert into heap at index size

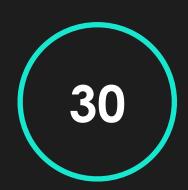
pq.insert("Google", 30)

```
self.size += 1
self.heap[self.size] = HeapItem(newKey,
newValue)
```



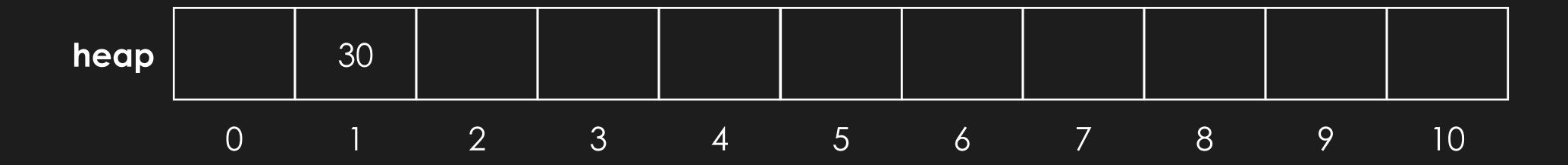


STEP 2: Increment size and insert into heap at index size



pq.insert("Google", 30)

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self.size += 1
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```





STEP 3: While parent node has higher priority, swap with parent node

30

pq.insert("Google", 30)

self.swim(self.size)





STEP 1: Check if capacity is reached

30

pq.insert("Amazon", 20)

```
if self.size >= self.maxsize:
    print('size limit reached')
    return
```



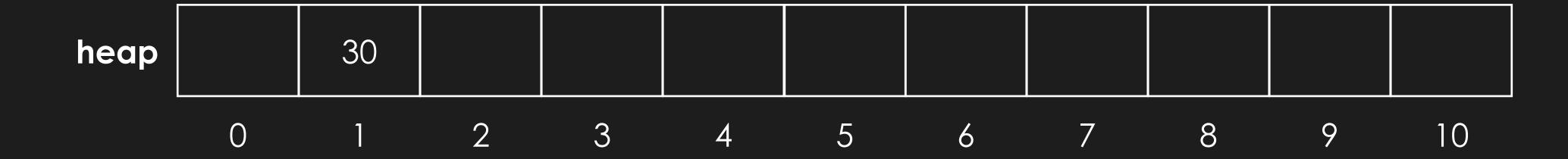


STEP 2: Increment size and insert into heap at index size



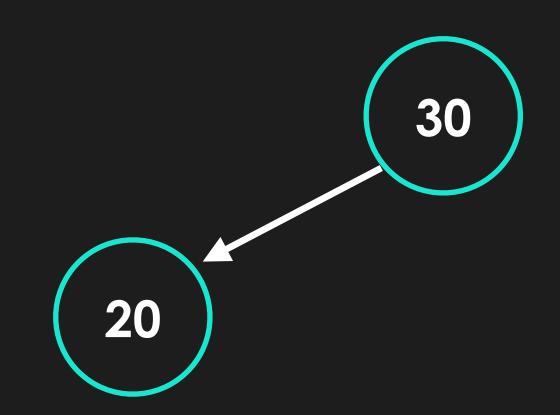
pq.insert("Amazon", 20)

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self.size += 1
self.heap[self.size] = HeapItem(newKey,
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```



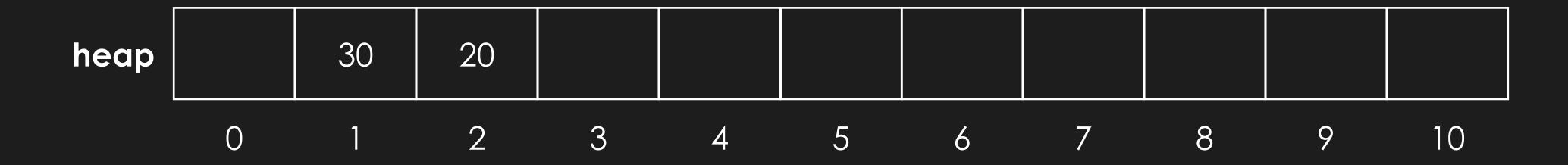


STEP 2: Increment size and insert into heap at index size



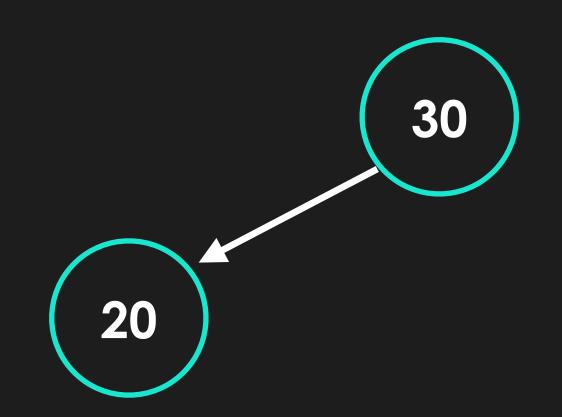
pq.insert("Amazon", 20)

```
self.size += 1
self.heap[self.size] = HeapItem(newKey,
newValue)
```





STEP 3: While parent node has higher priority, swap with parent node



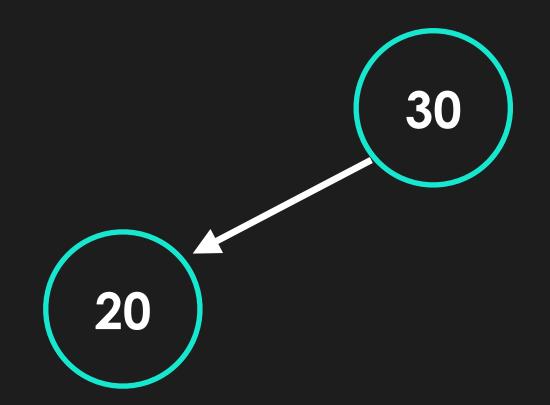
pq.insert("Amazon", 20)

self.swim(self.size)



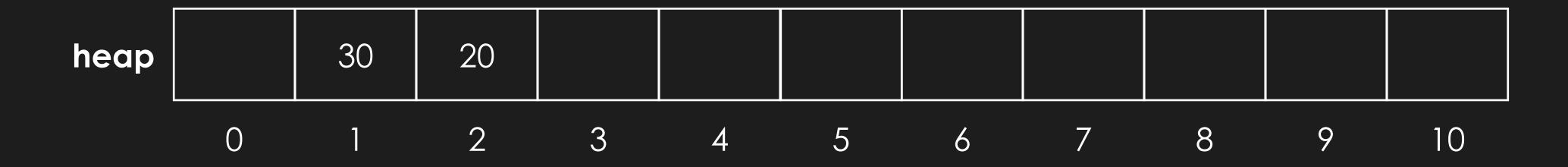


STEP 1: Check if capacity is reached



pq.insert("Apple", 60)

```
if self.size >= self.maxsize:
    print('size limit reached')
    return
```



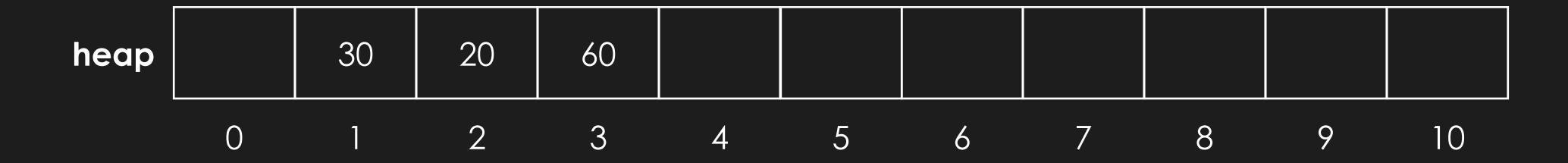


STEP 2: Increment size and insert into heap at index size



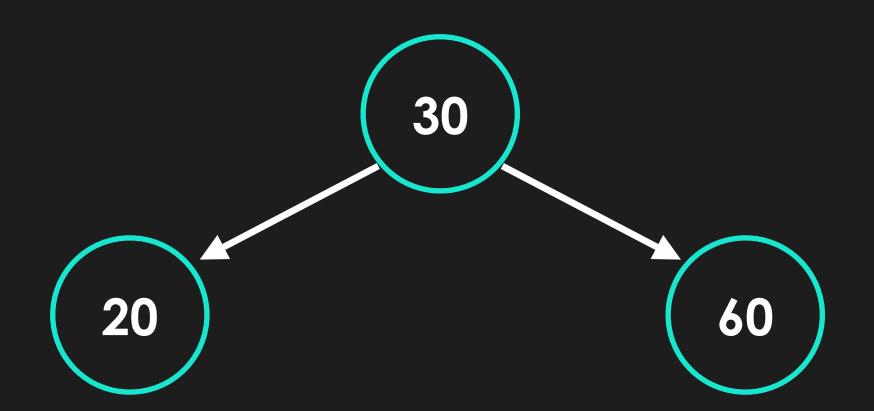
pq.insert("Apple", 60)

```
self.size += 1
self.heap[self.size] = HeapItem(newKey,
newValue)
```





STEP 3: While parent node has higher priority, swap with parent node



pq.insert("Apple", 60)

self.swim(self.size)





swim

greater

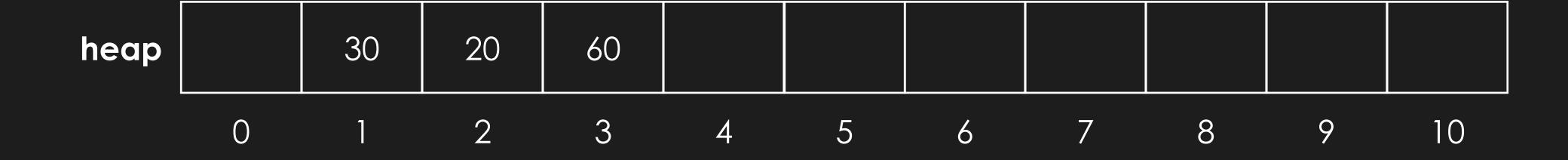
```
def greater(self, index1, index2):
   if self.heap[index1].value > self.heap[index2].value:
      return True
   else:
      return False
```

We define a method **greater** which returns **True** if HeapItem at **index1** has a greater value (higher priority) then HeapItem at **index2**





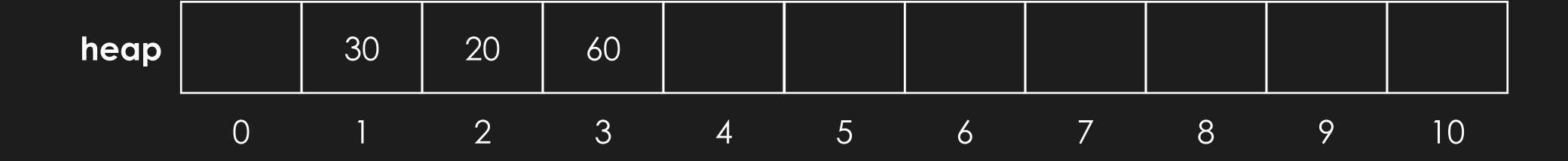
```
while (index != 1 and self.greater(index,
parent(index))):
    swap(self.heap, index, parent(index))
    index = parent(index)
```







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while (index != 1 and self.greater(index,
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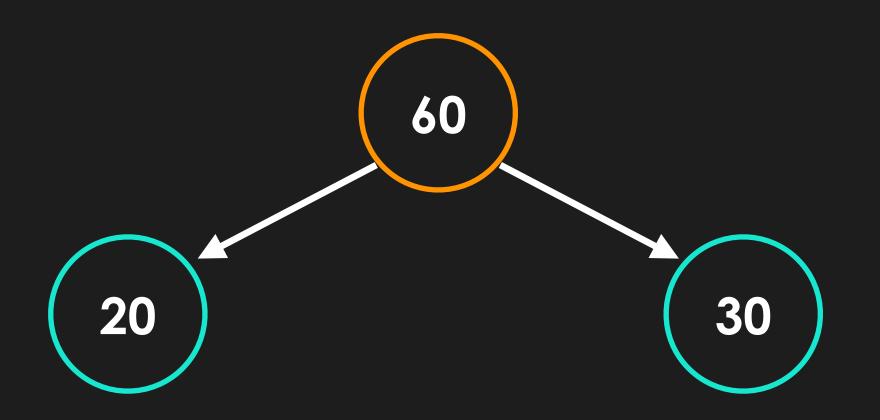




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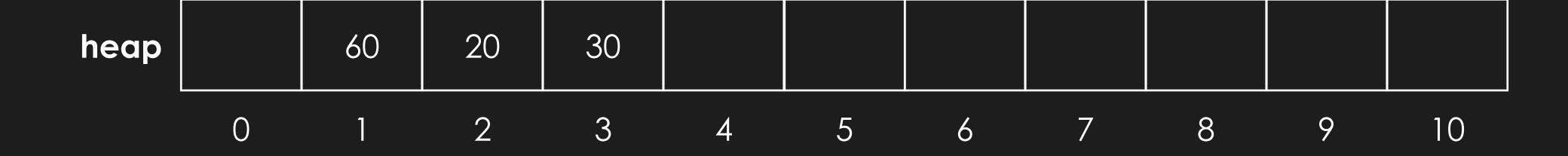
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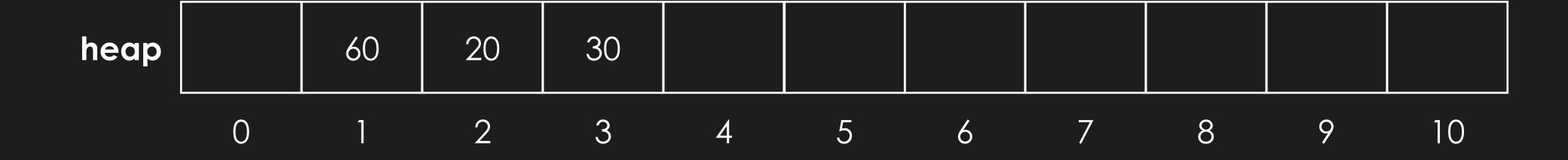
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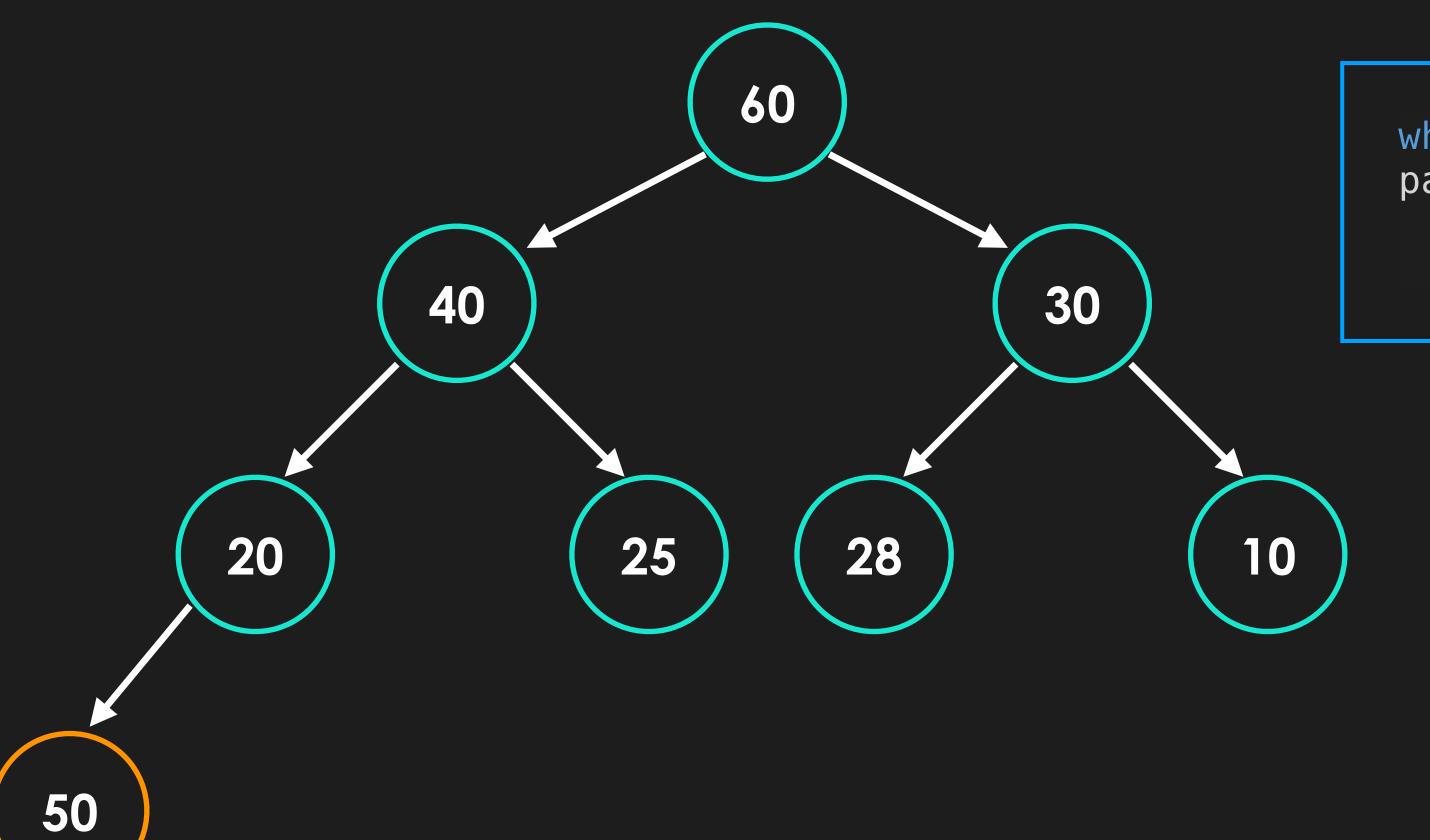
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while (index != 1 and self.greater(index,
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```





swim example 2

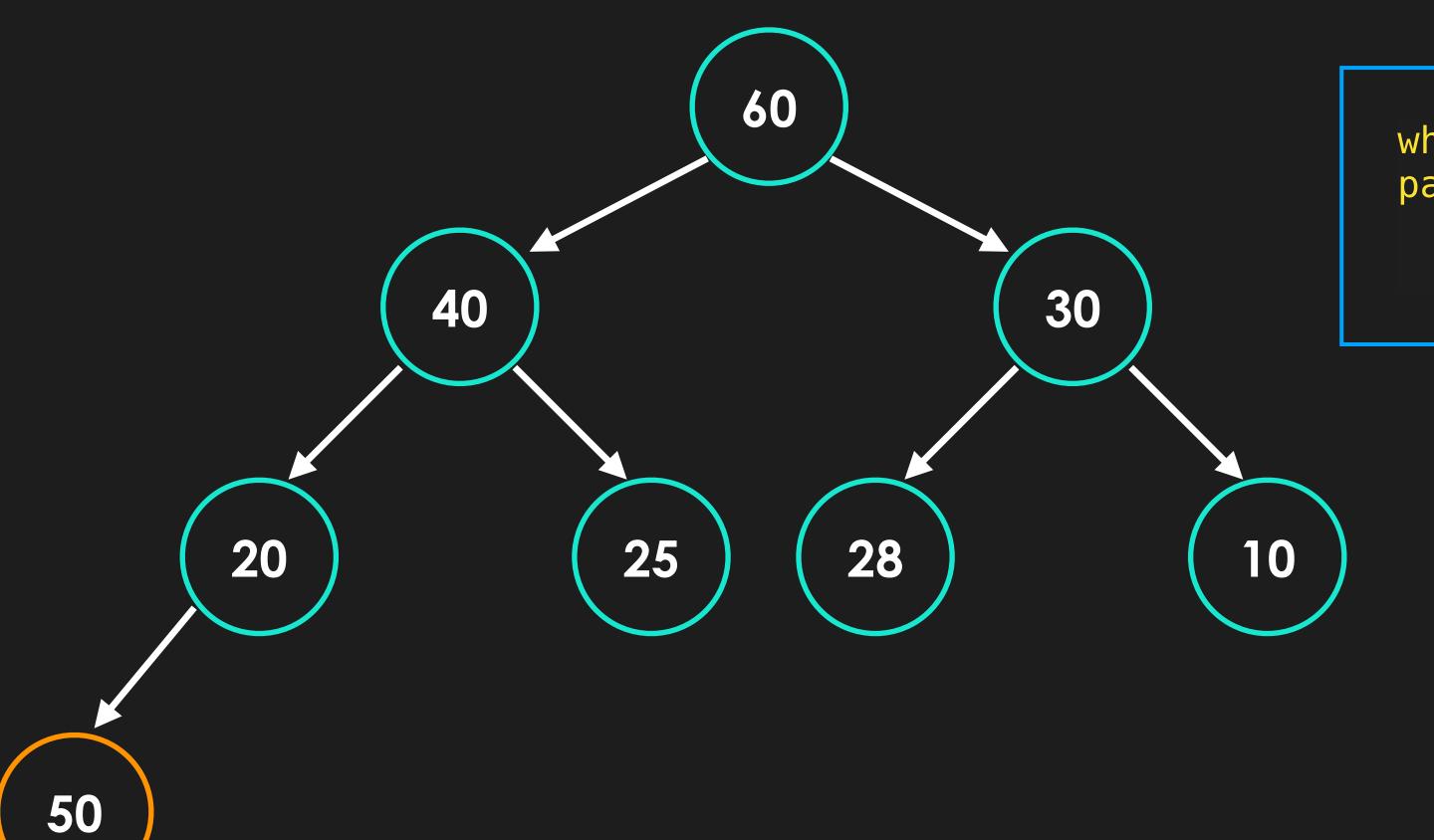




while (index != 1 and self.greater(index,
parent(index))):
 swap(self.heap, index, parent(index))
 index = parent(index)

heap	60	40	30	20	25	28	10	50		
		2							9	10

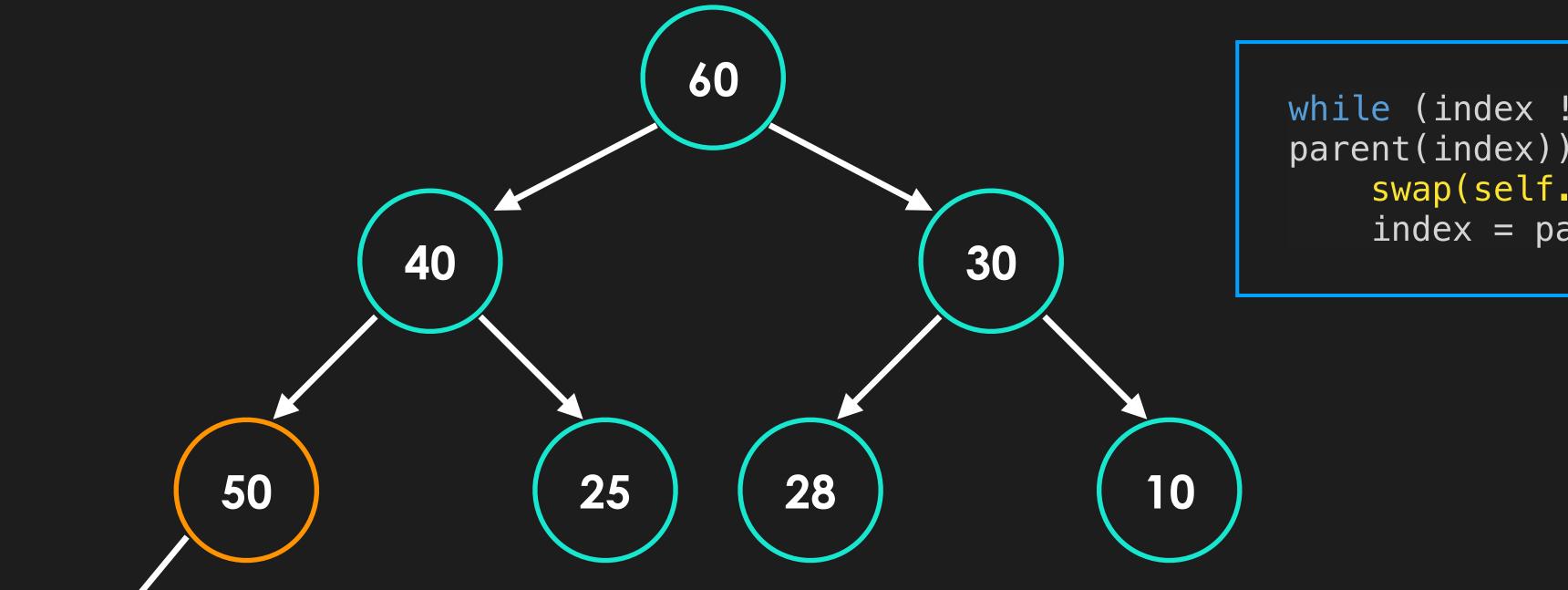




while (index != 1 and self.greater(index,
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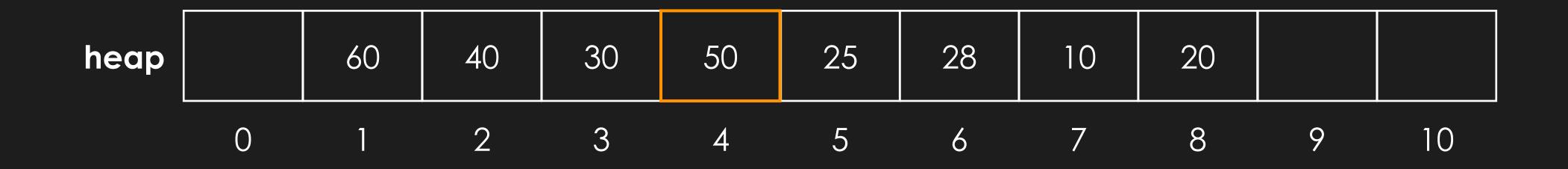




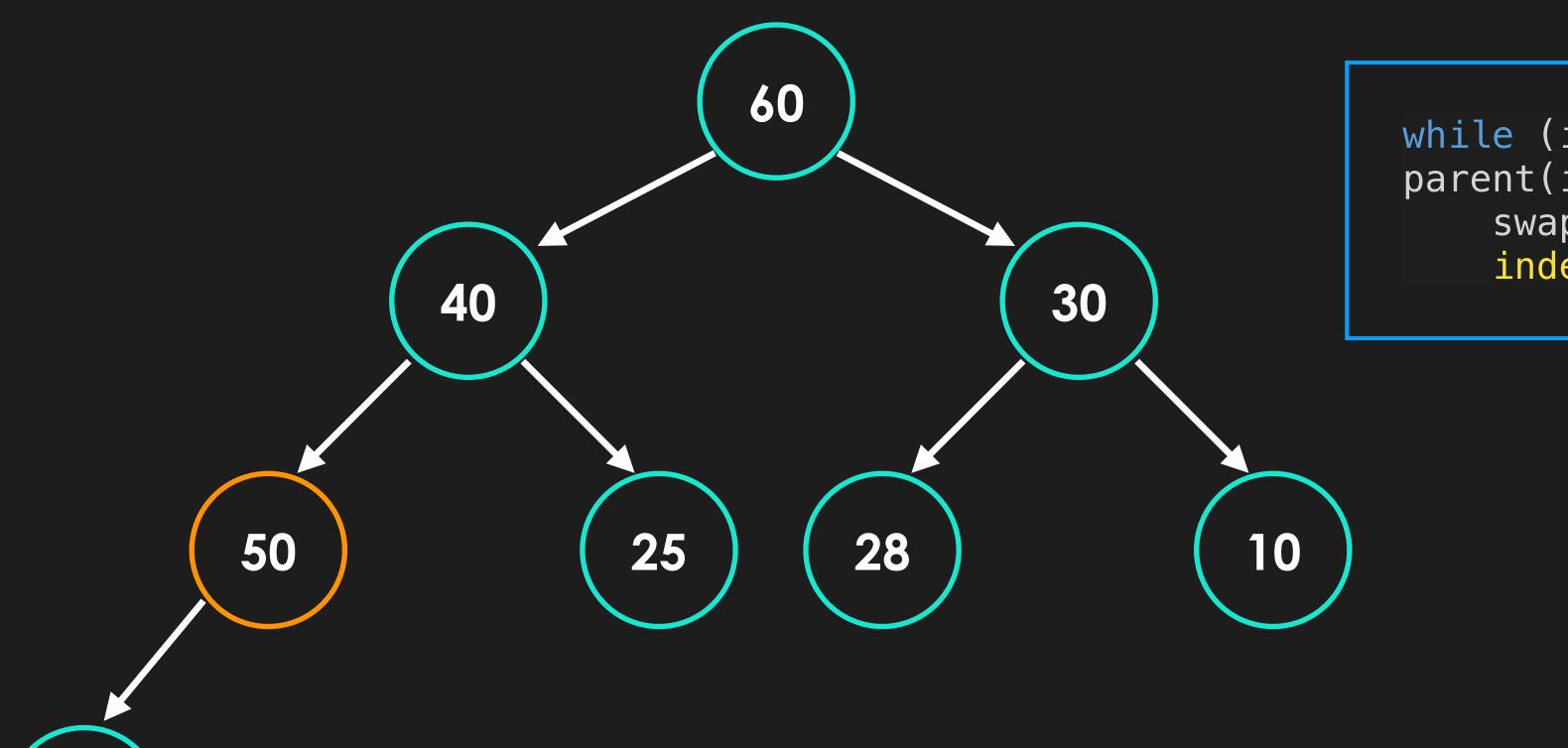


20

while (index != 1 and self.greater(index,
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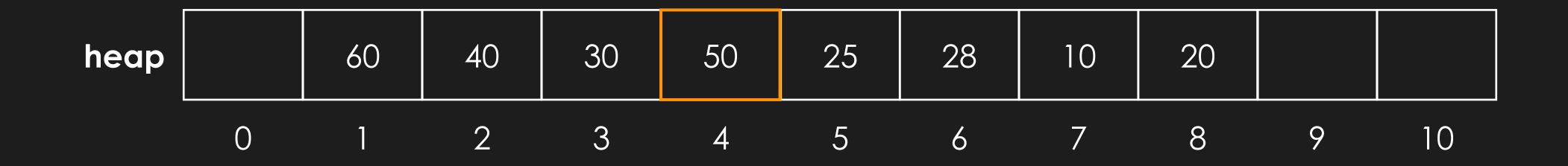




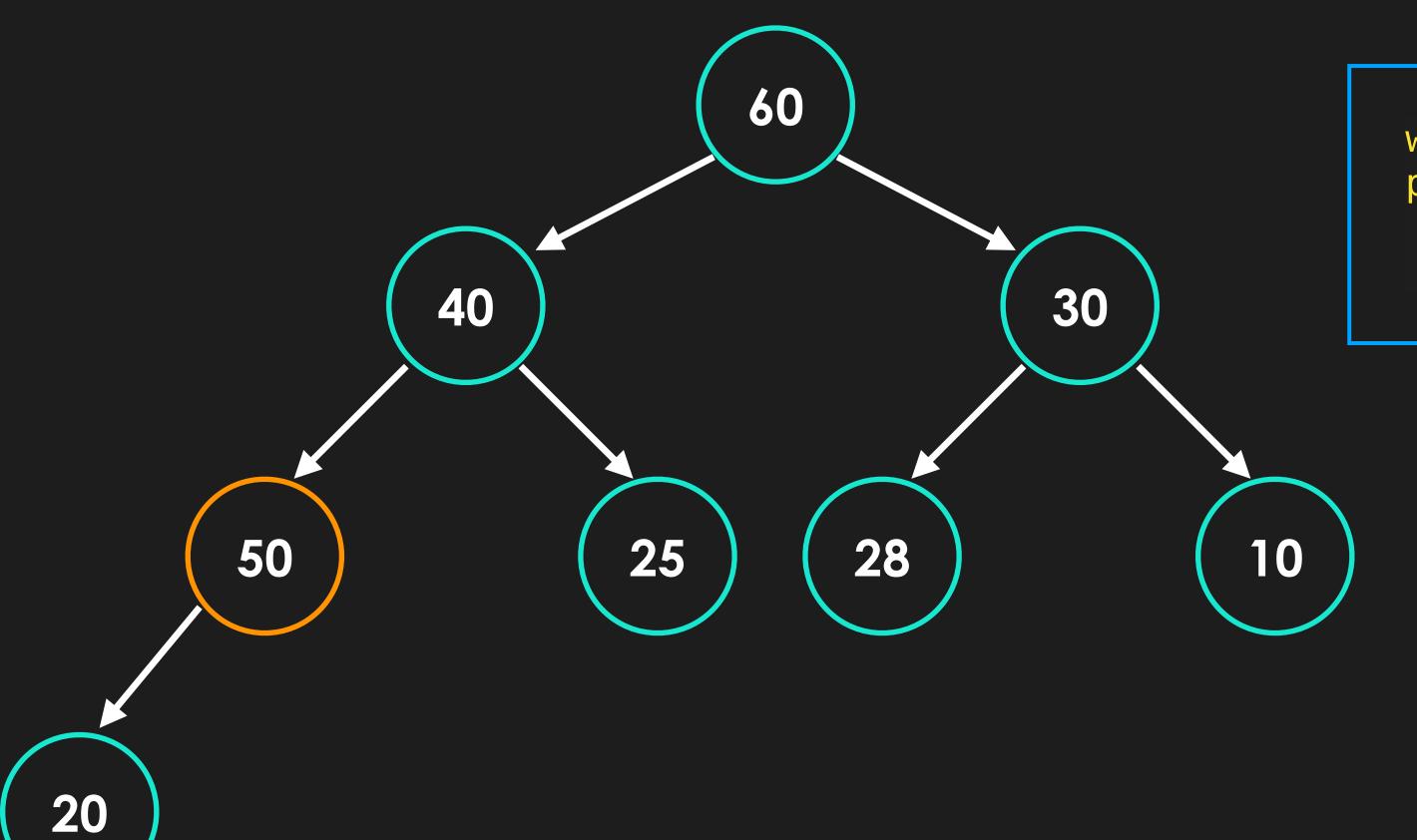


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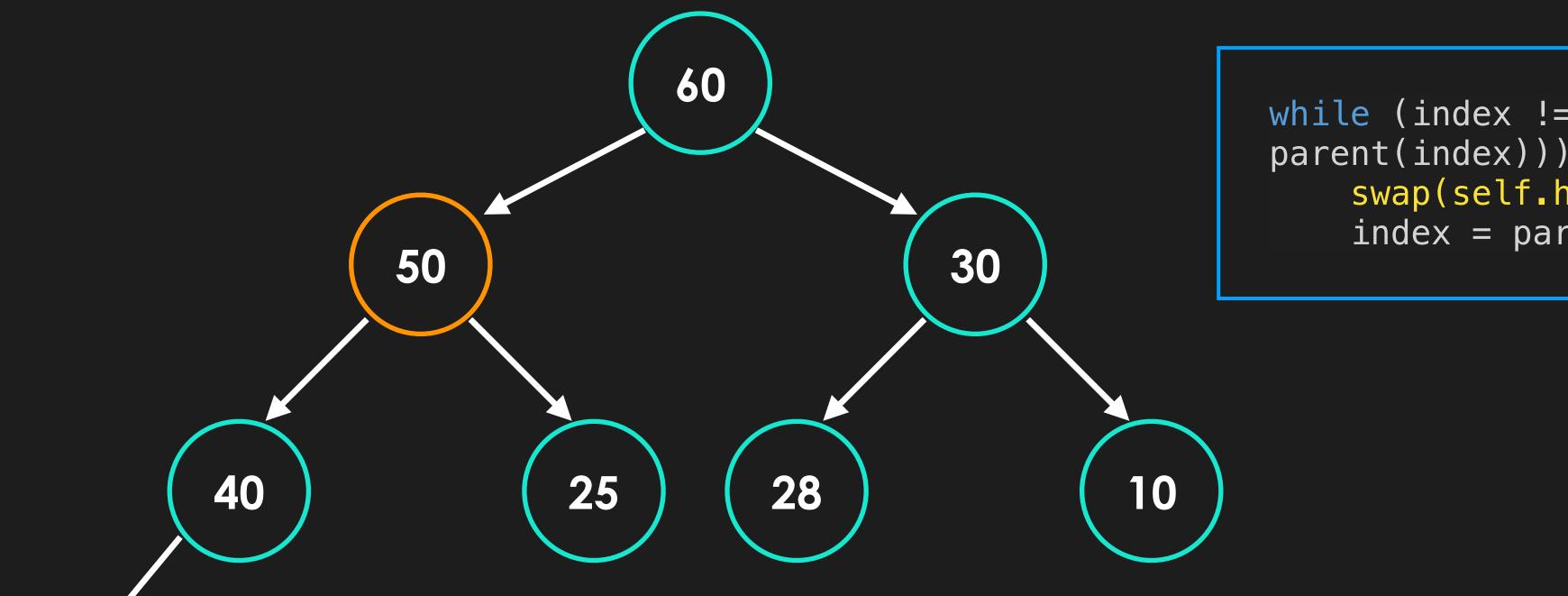




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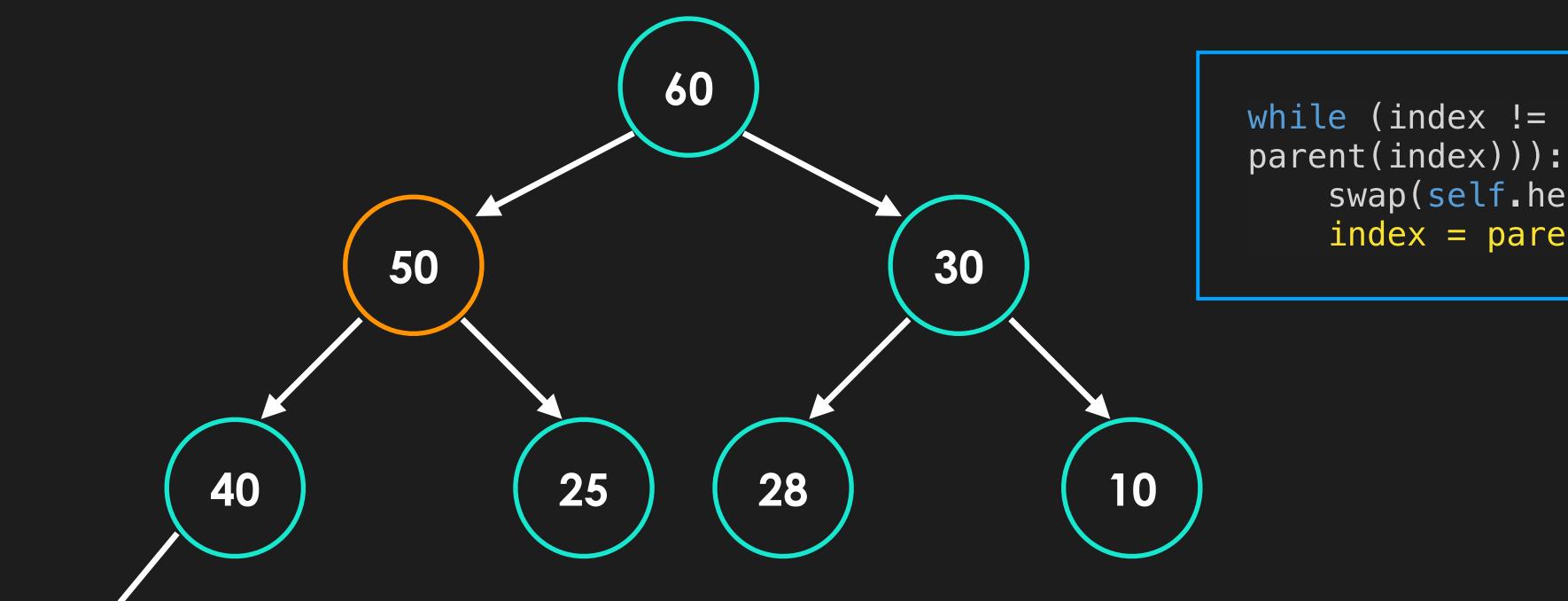


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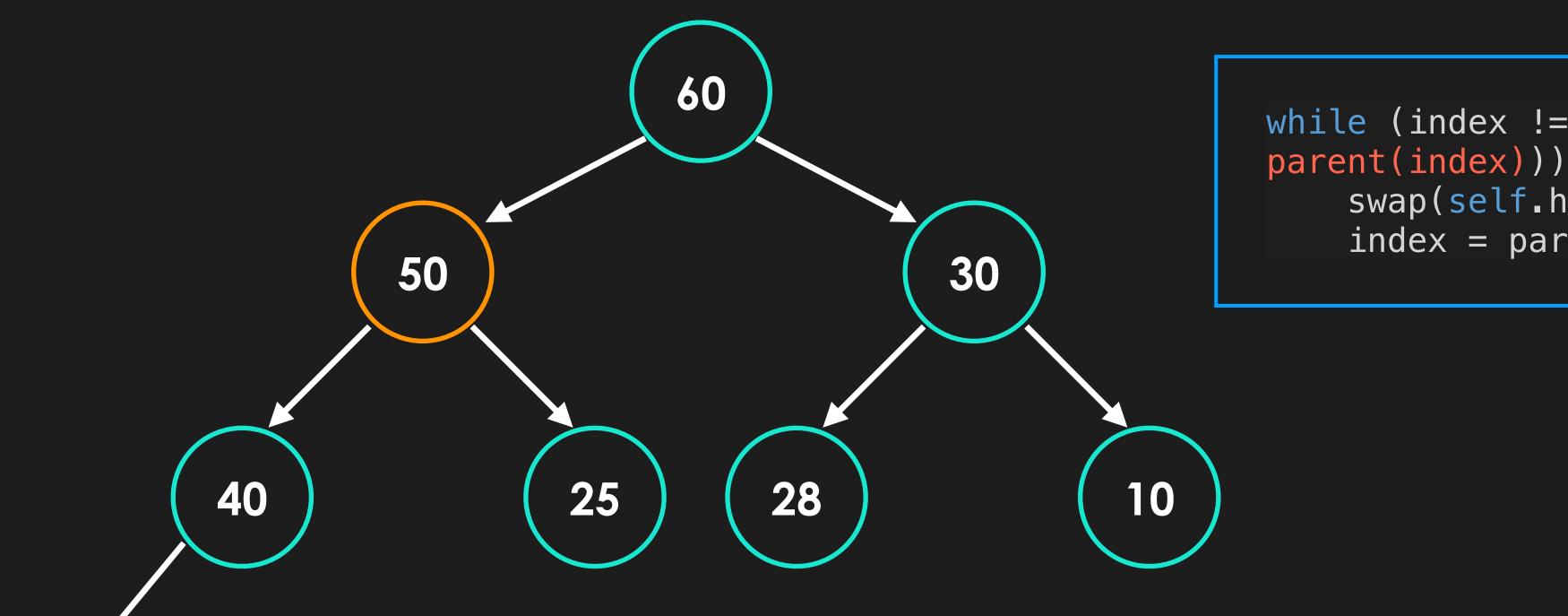


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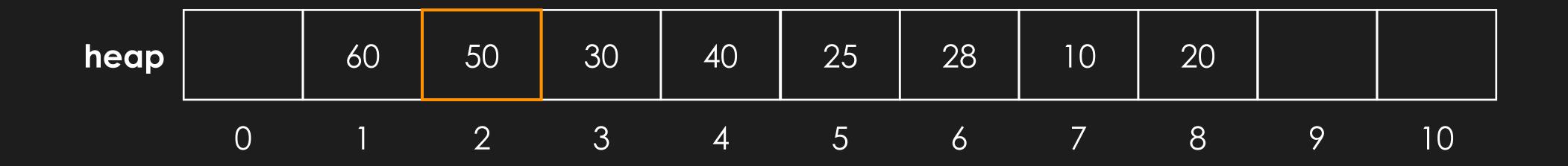






20

while (index != 1 and self.greater(index,
 parent(index))):
 swap(self.heap, index, parent(index))
 index = parent(index)





MaxHeap Insertion

```
def swim (self, index):
   while (index != 1 and self.greater(index, parent(index))):
        swap(self.heap, index, parent(index))
        index = parent(index)
    return
def insert(self, newKey, newValue):
    if self.size >= self.maxsize:
        print('size limit reached')
        return
    self.size += 1
    self.heap[self.size] = HeapItem(newKey, newValue)
    self.swim(self.size)
```



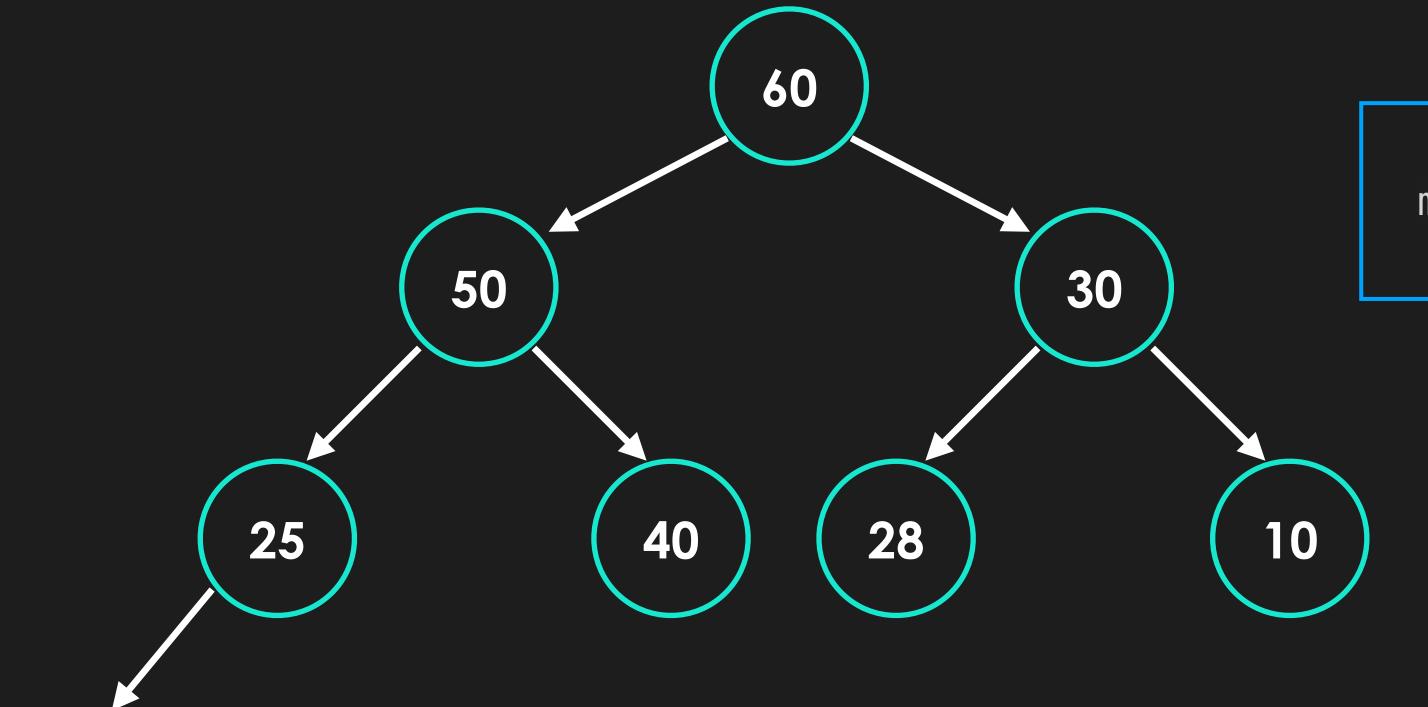
Max Extraction



Max Extraction

- Remove max element
- Ensure Heap order is maintained

getMax()



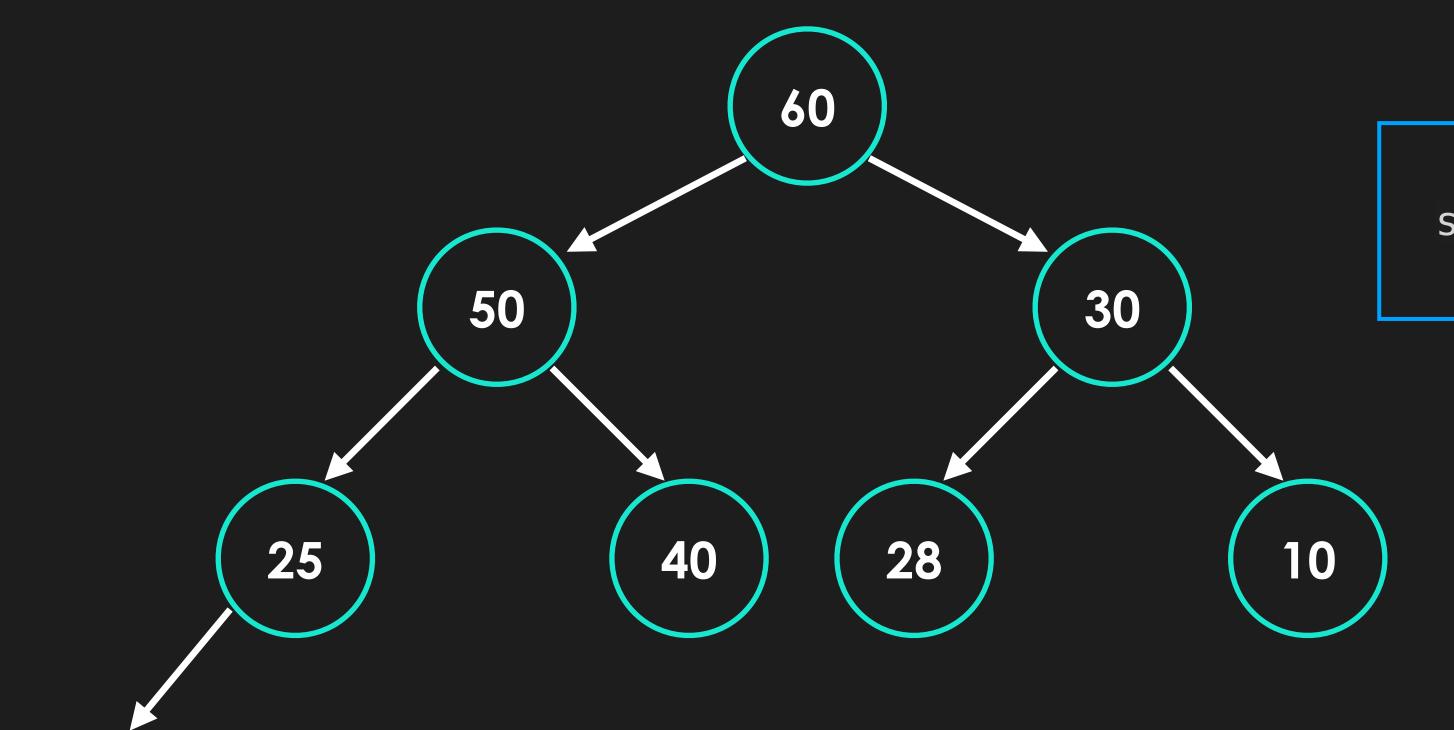
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maxItem = self.heap[1]

STEP 1: Save max element into a variable







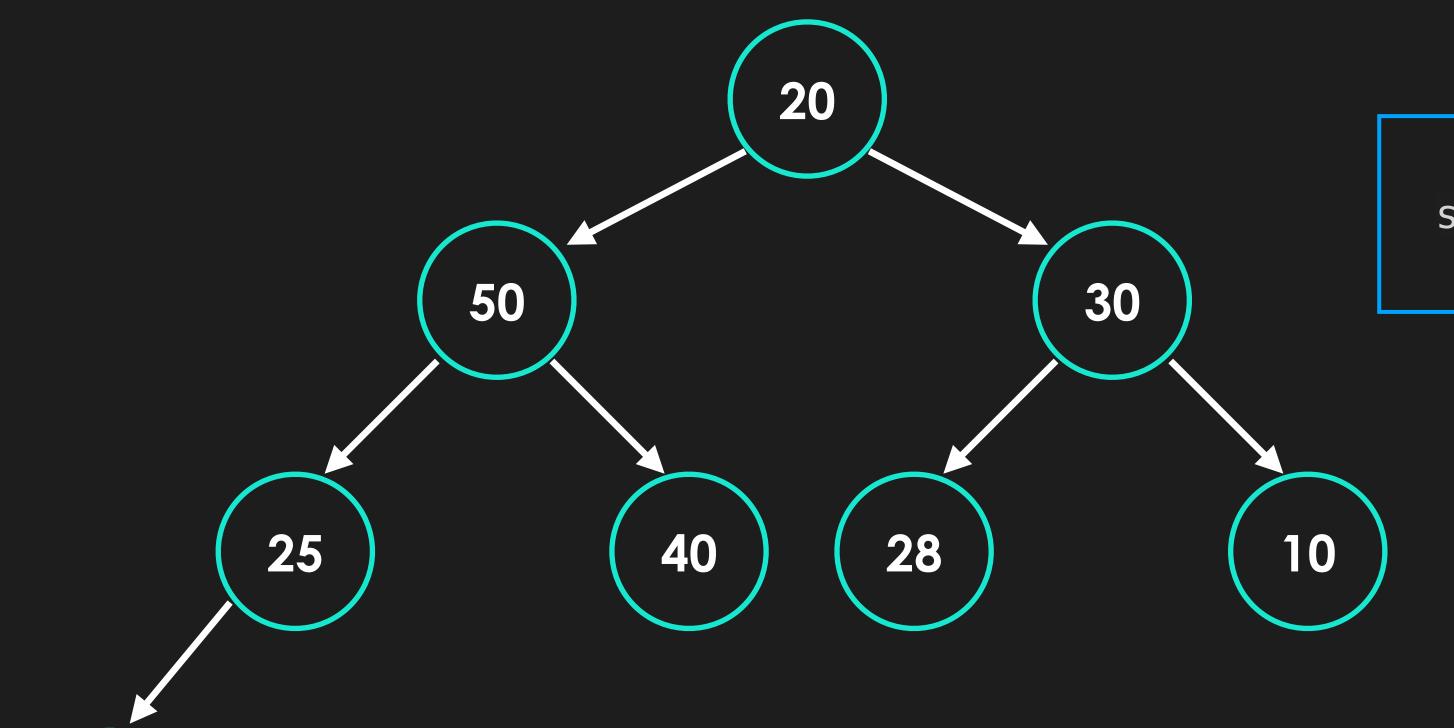
20

swap(self.heap, 1, self.size)

STEP 2: Swap first and last element

heap		60	50	30	25	40	28	10	20		
	O	1	2	3	4	5	6	7	8	9	10





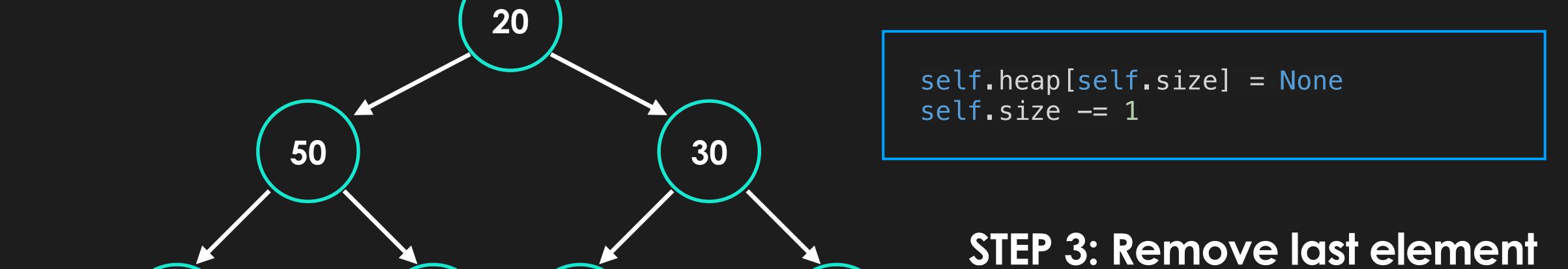
60

swap(self.heap, 1, self.size)

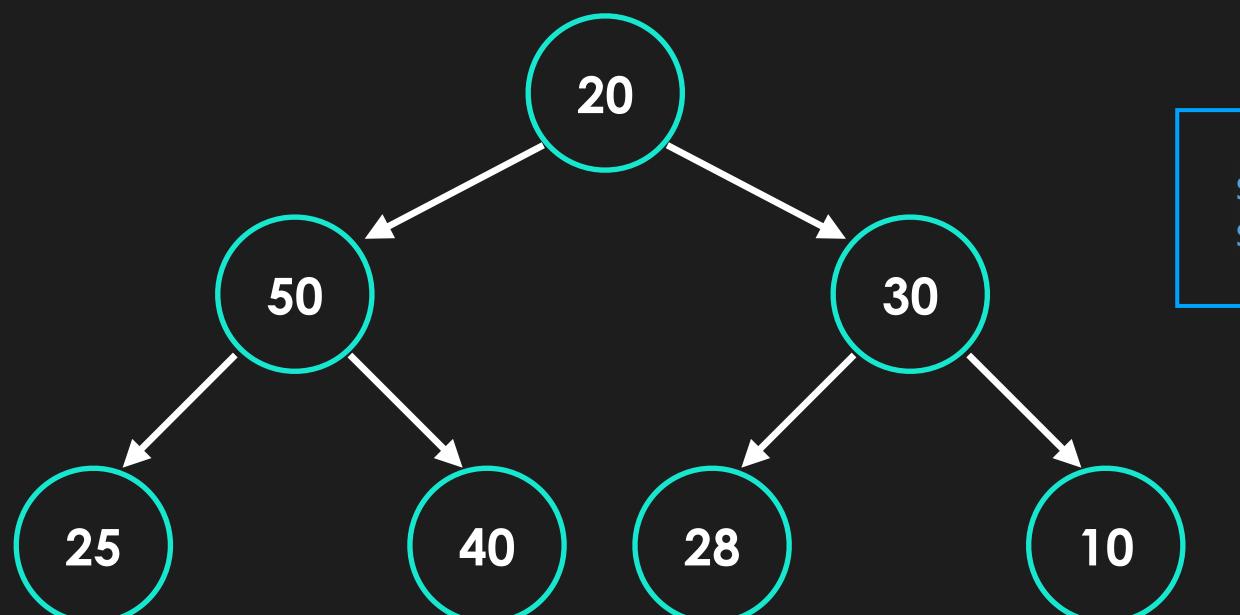
STEP 2: Swap first and last element









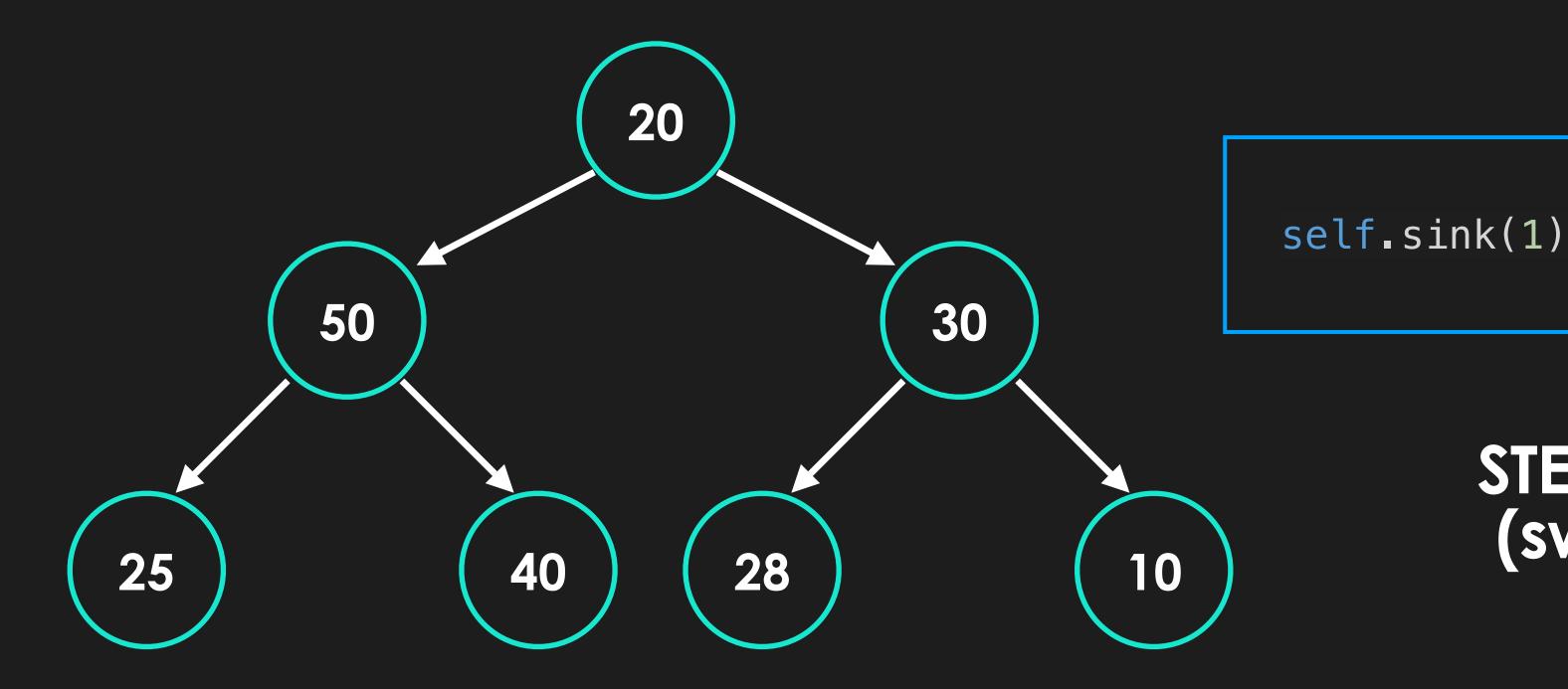


self.heap[self.size] = None
self.size -= 1

STEP 3: Remove last element

heap		20	50	30	25	40	28	10			
	0	1	2	3	4	5	6	7	8	9	10



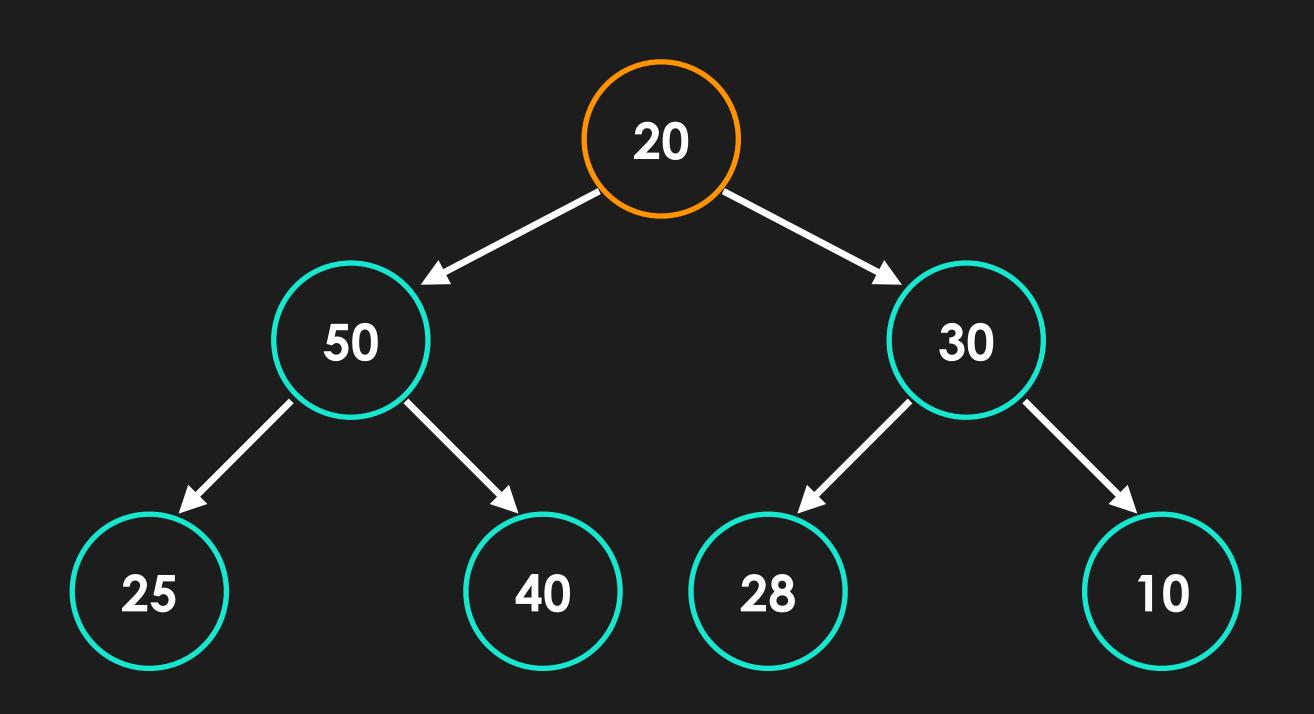


STEP 4: Sink first element, (swap with child node if priority is lower)





sink



```
while (2 * index <= self.size):
    j = left(index)

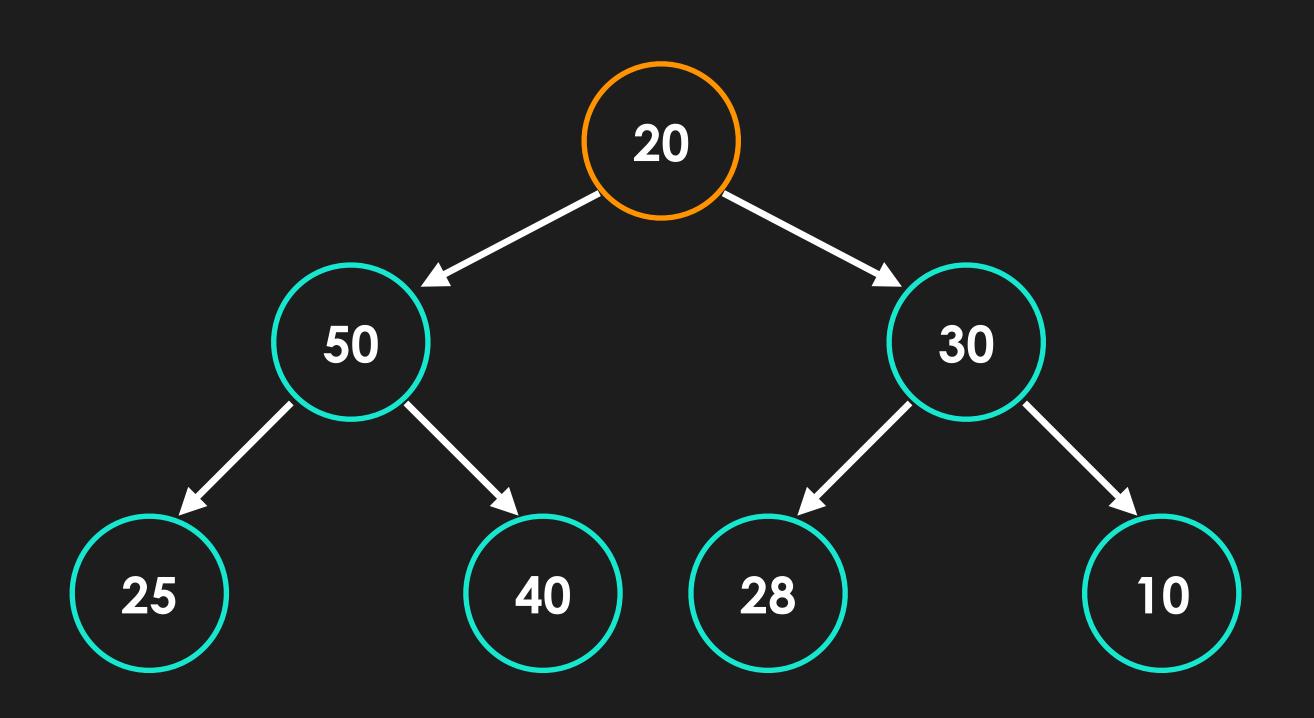
if (j < self.size and
    self.greater(right(index), j)):
        j = right(index)

if self.greater(index, j):
        break

else:
        swap(self.heap, j, index)
        index = j</pre>
```

heap	20	50	30	40	25	28	10		
							7	9	10





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while (2 * index <= self.size):
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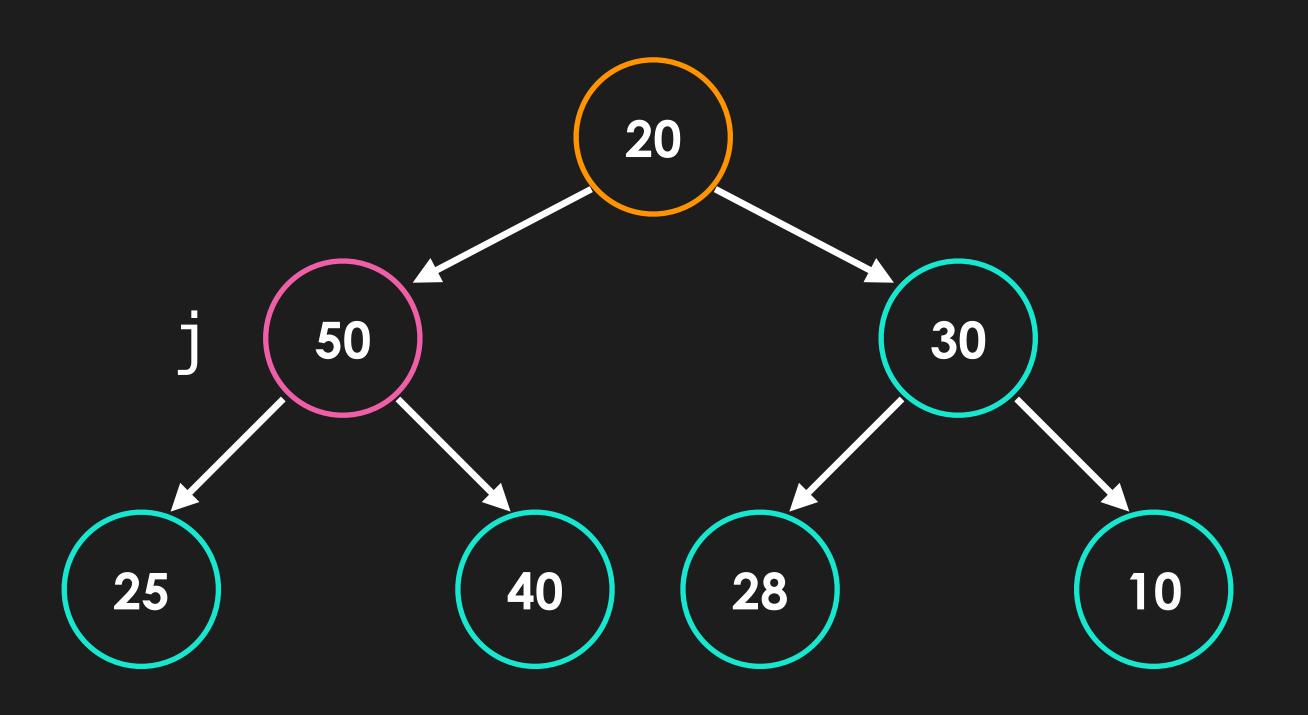
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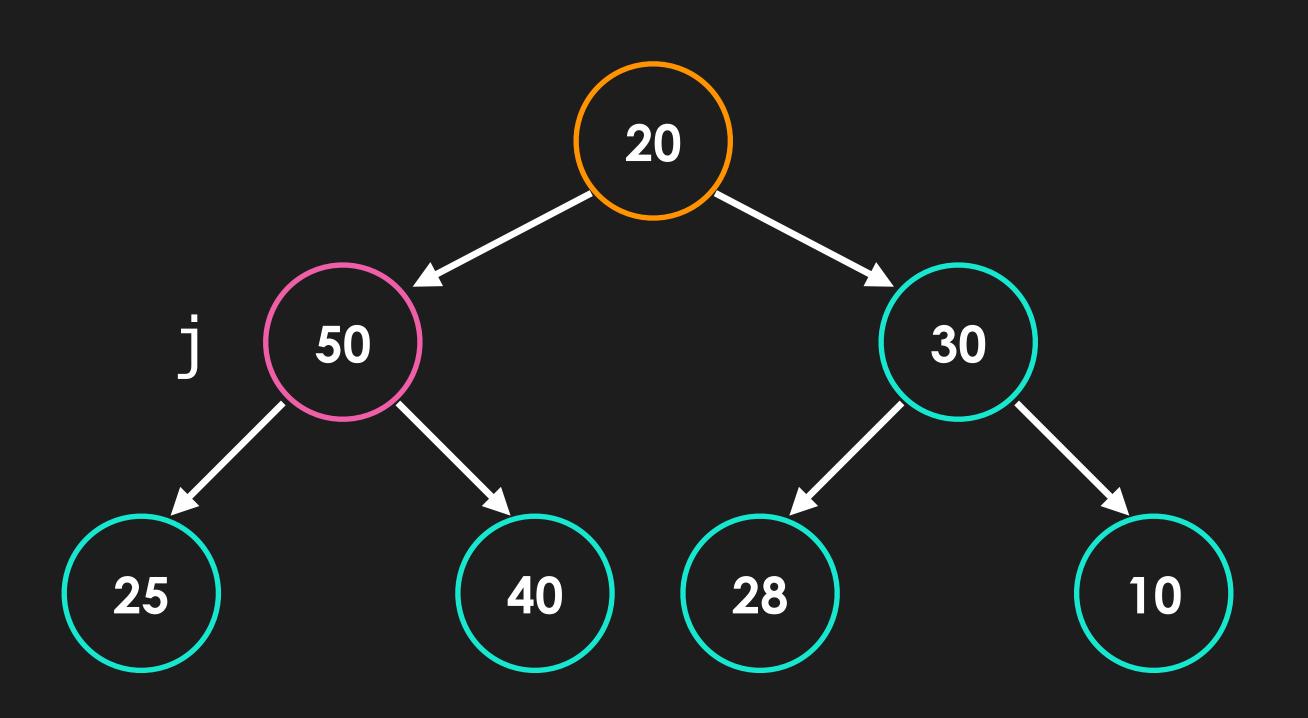
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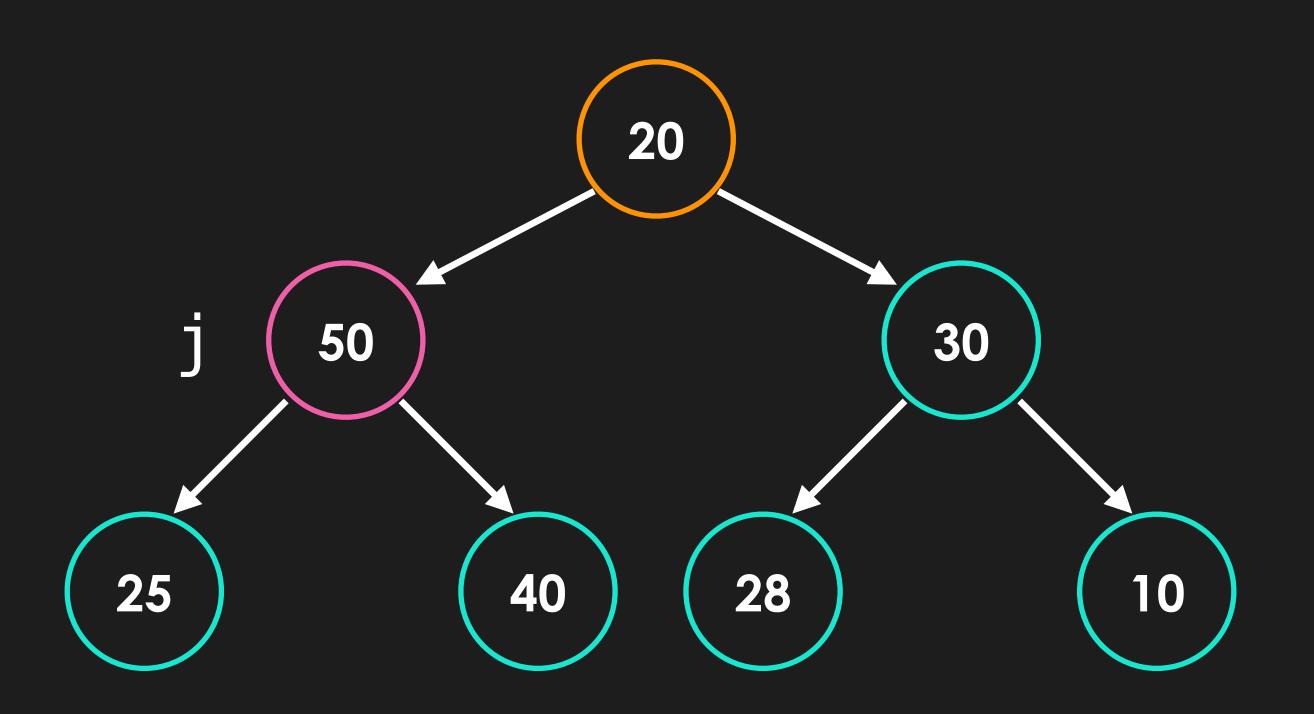
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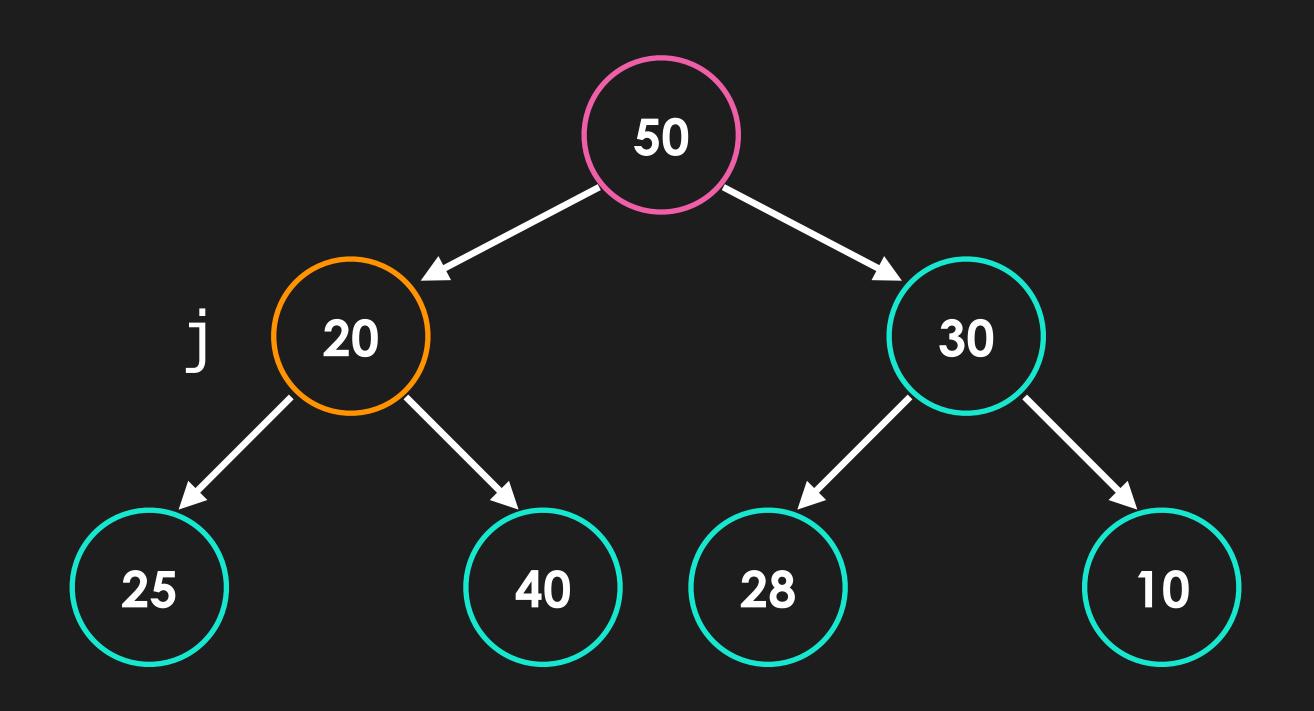
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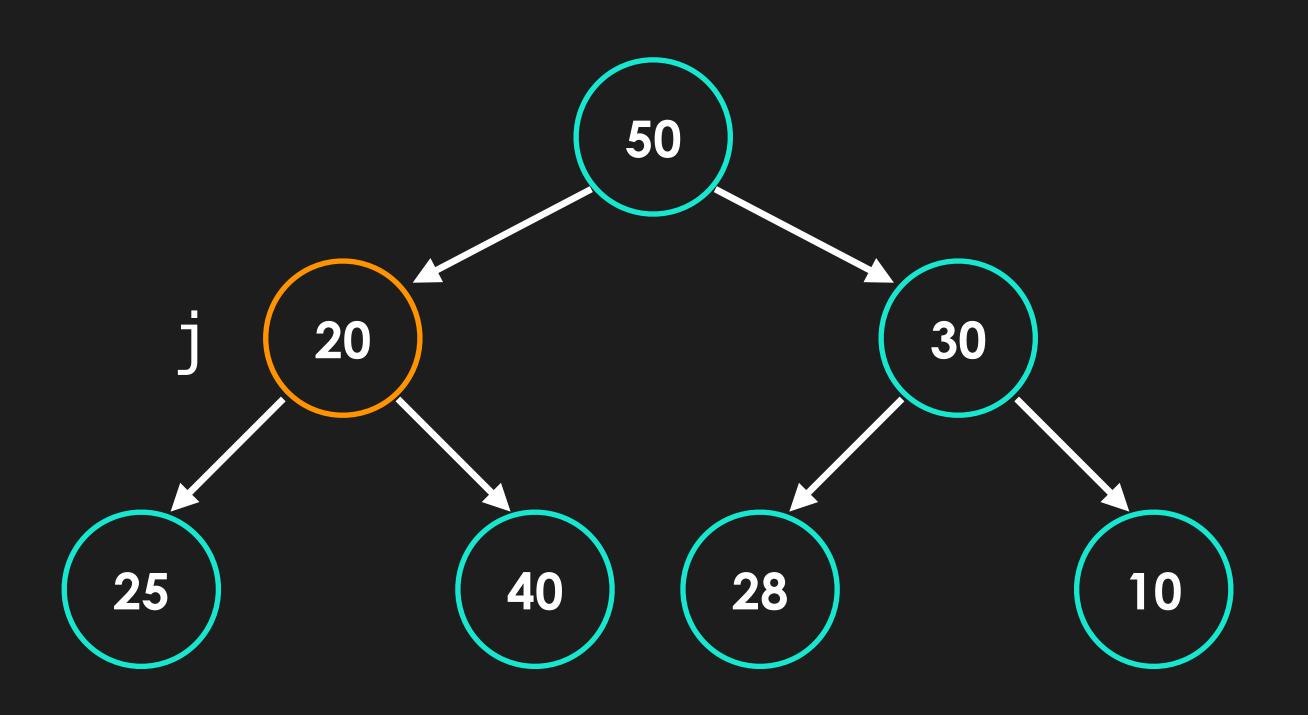
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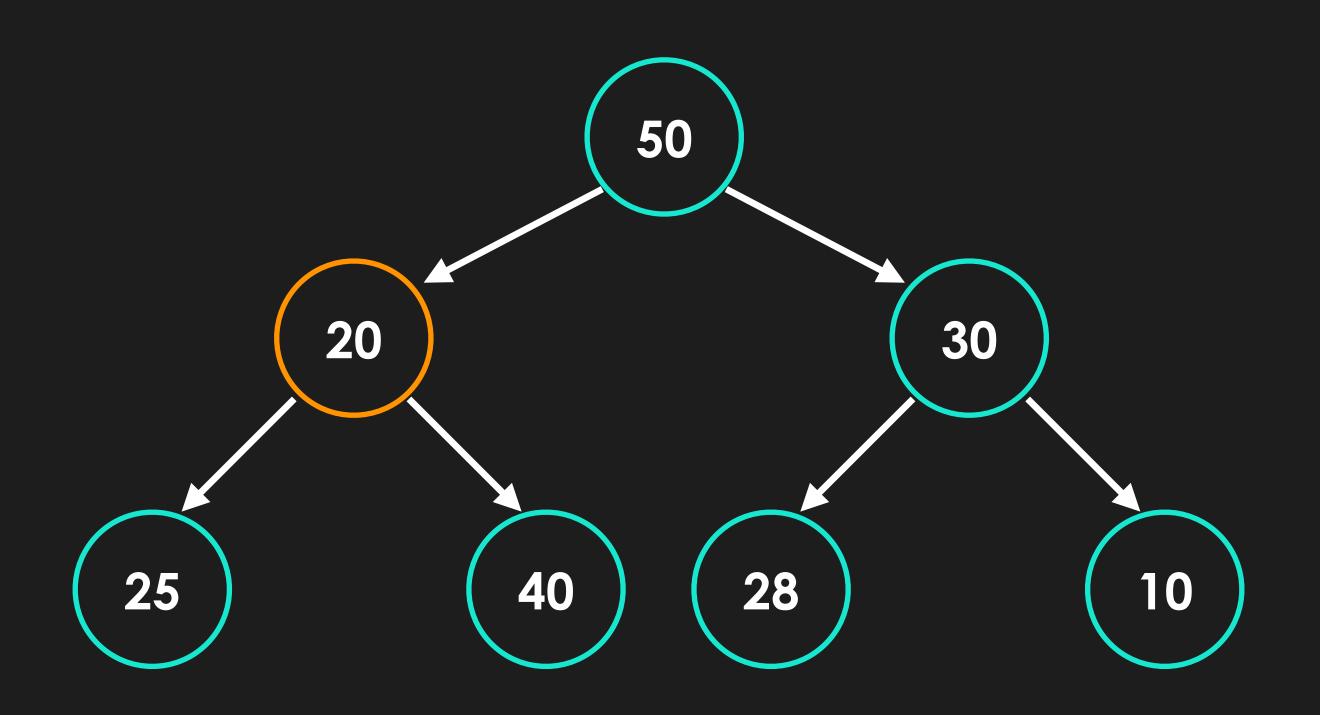
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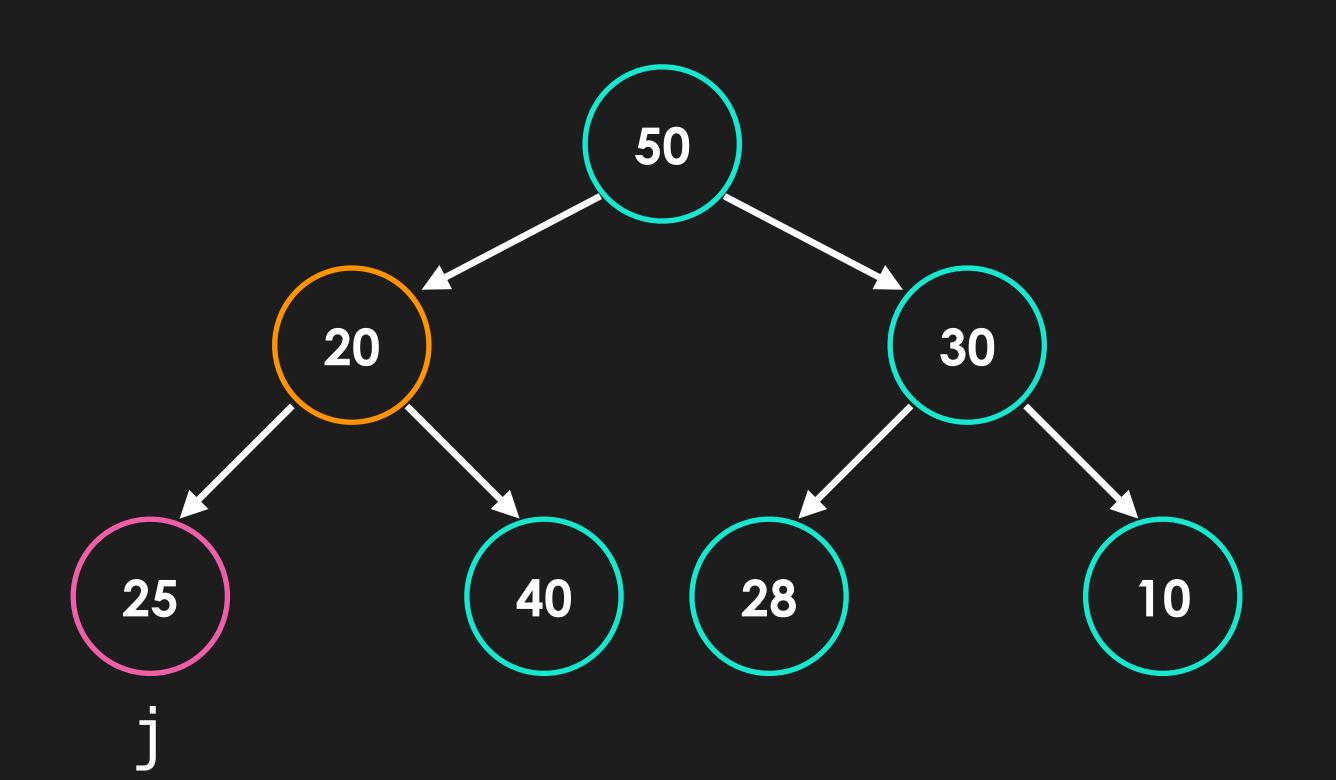
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```



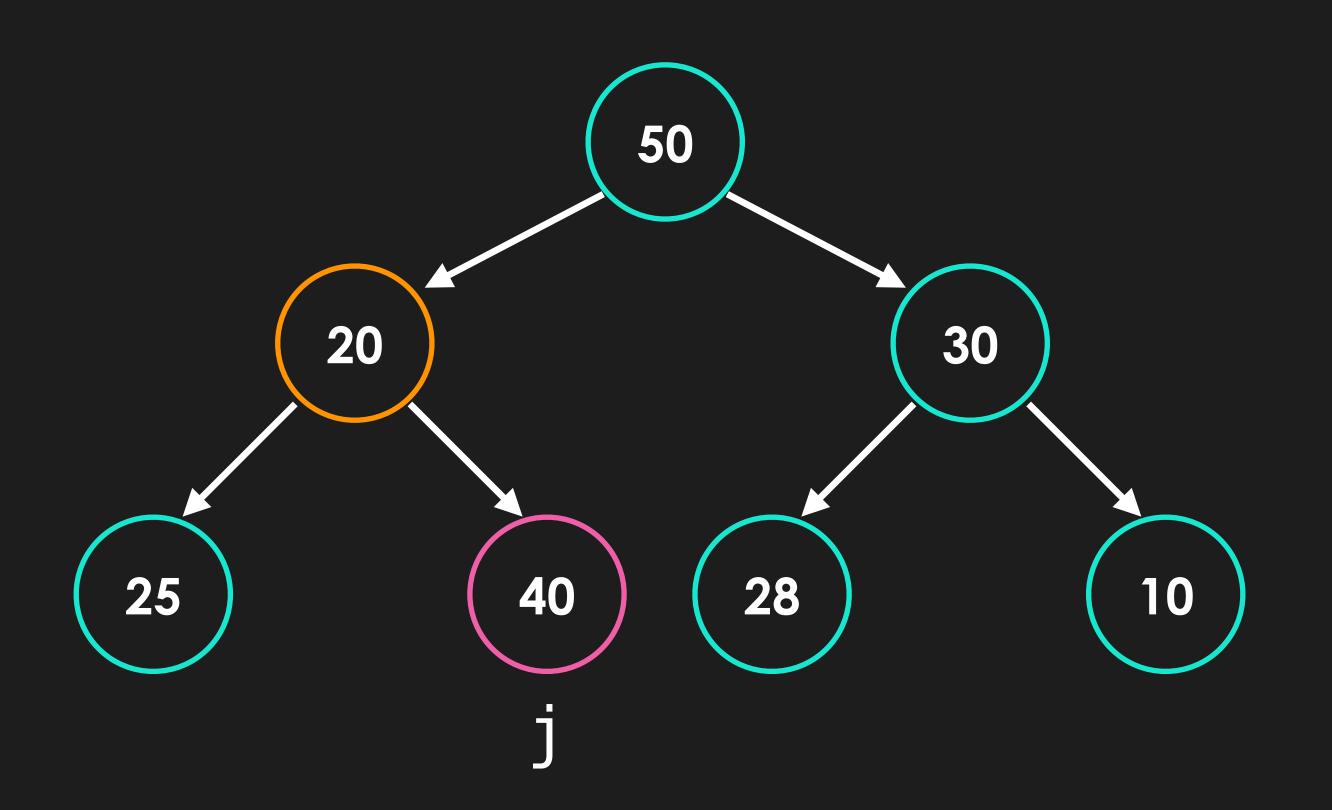




```
while (2 * index <= self.size):</pre>
    j = left(index)
    if (j < self.size and</pre>
    self.greater(right(index), j)):
        j = right(index)
    if self.greater(index, j):
        break
    else:
        swap(self.heap, j, index)
        index = j
```



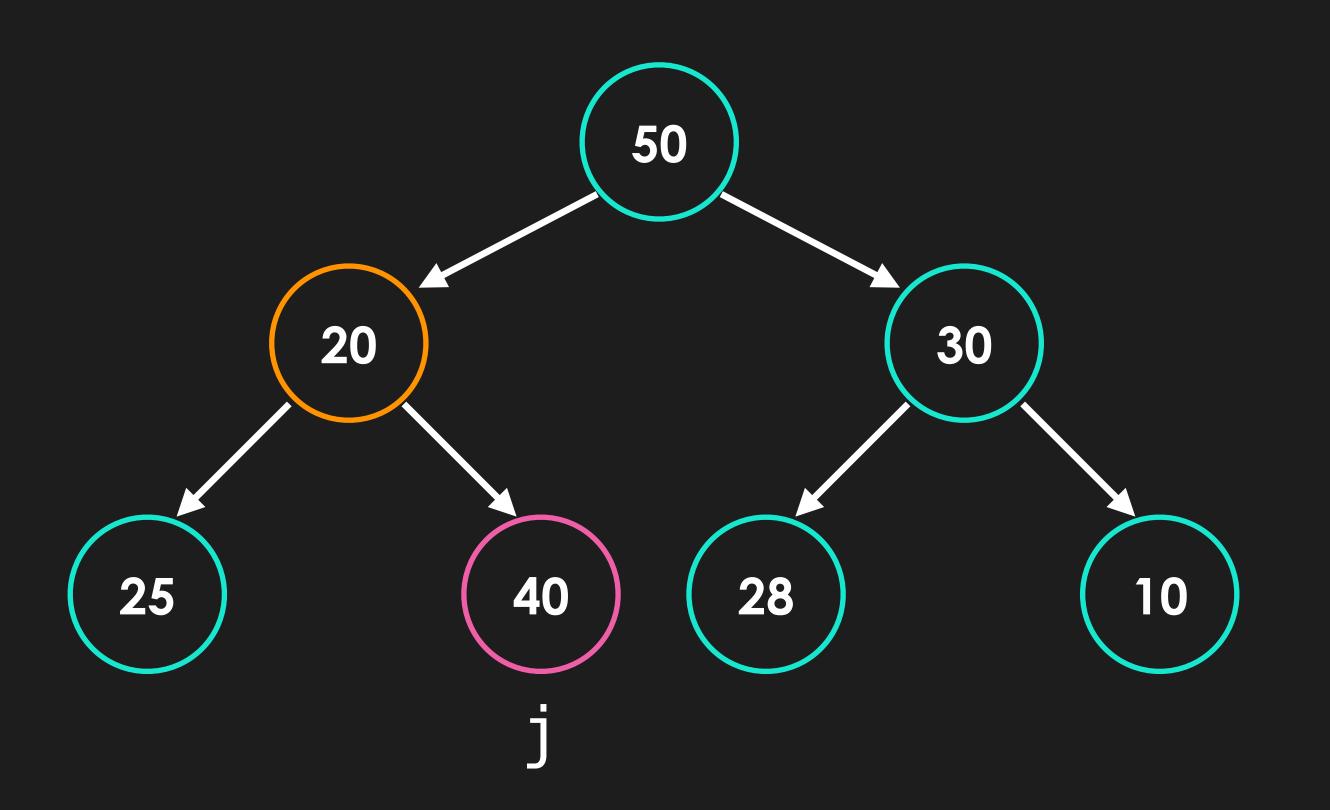




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    if self.greater(index, j):
        break
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```



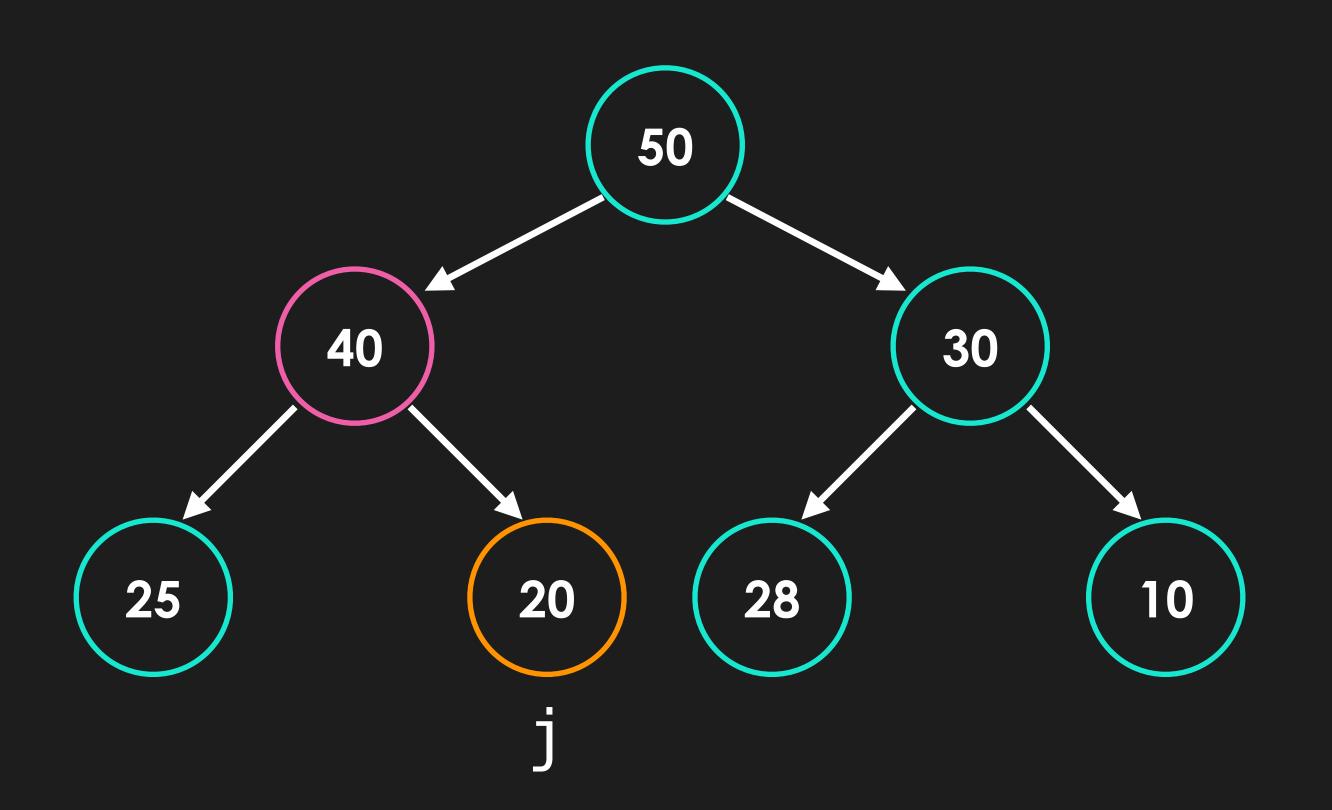




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    self.greater(right(index), j)):
        j = right(index)
    if self.greater(index, j):
        break
    else:
        swap(self.heap, j, index)
        index = j
```



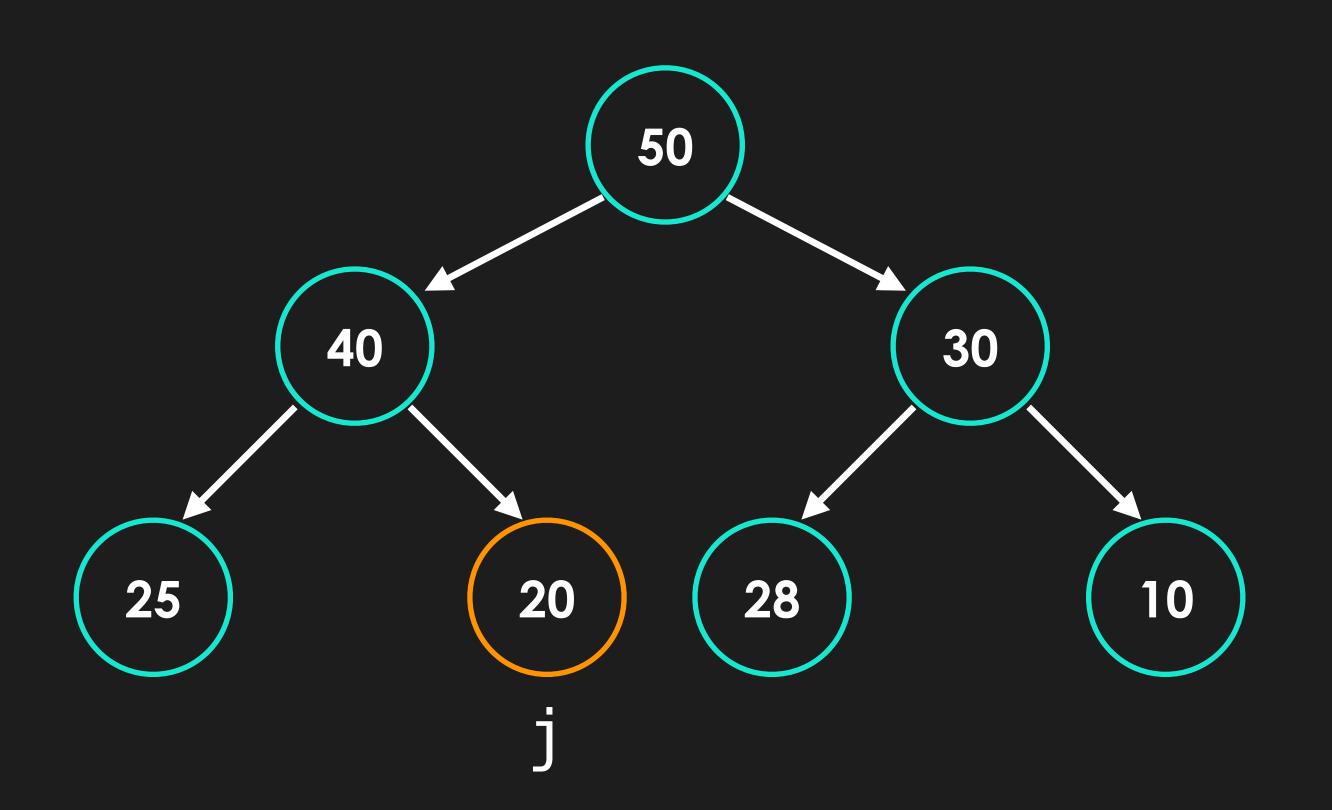




```
while (2 * index <= self.size):</pre>
    j = left(index)
    if (j < self.size and</pre>
    self.greater(right(index), j)):
        j = right(index)
    if self.greater(index, j):
        break
    else:
        swap(self.heap, j, index)
        index = j
```



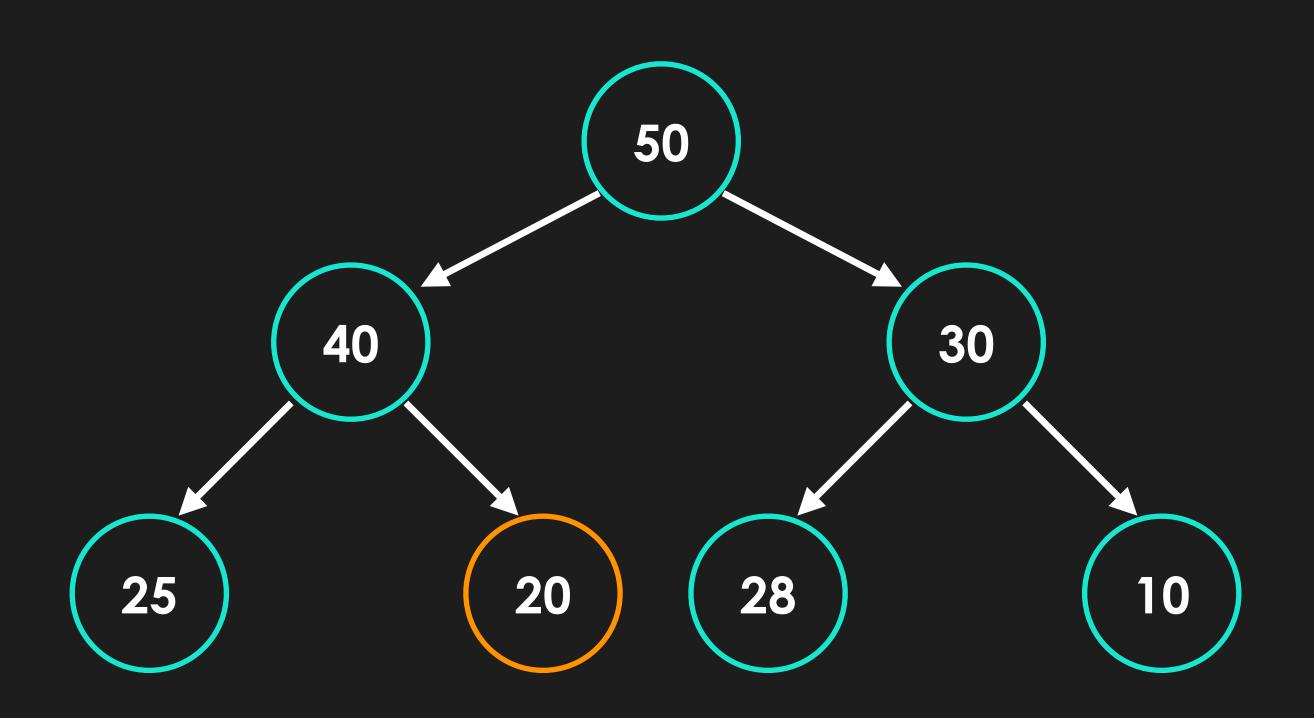




```
while (2 * index <= self.size):</pre>
    j = left(index)
    if (j < self.size and</pre>
    self.greater(right(index), j)):
        j = right(index)
    if self.greater(index, j):
        break
    else:
        swap(self.heap, j, index)
        index = j
```







```
while (2 * index <= self.size):
    j = left(index)

if (j < self.size and
    self.greater(right(index), j)):
        j = right(index)

if self.greater(index, j):
        break

else:
        swap(self.heap, j, index)
        index = j</pre>
```

heap	50	40	30	20	25	28	10		
							7	9	10



sink

```
def sink(self, index):
    while (2 * index <= self.size):</pre>
    # initialize j to be left child
        j = left(index)
        # compare left and right
        if (j < self.size and self.greater(right(index), j)):</pre>
            j = right(index)
        if self.greater(index, j):
            break
        else:
            swap(self.heap, j, index)
            index = j
    return
```



getMax

```
def getMax(self):
    if self.size == 0:
        return None
    # SAVE MAX ITEM FOR REMOVAL
    maxItem = self.heap[1]
    # SWAP MAX ITEM
    swap(self.heap, self.positions, 1, self.size)
    # REMOVE ITEM
    self.heap[self.size] = None
    self.size -= 1
    # SINK FIRST ITEM
    self.sink(1)
    return maxItem
```



Putting it all together

```
pq = MaxHeap(10)
pq.insert(30, "Google")
pq.insert(20, "Amazon")
pq.insert(60, "Apple")
pq.insert(40, "Microsoft")
pq.insert(25, "Netflix")

while pq.size != 0:
    maxItem = pq.getMax()
    print(maxItem.key, maxItem.value)
```

```
pq = MaxHeap(10)
pq.insert(30, "Google")
pq.insert(20, "Amazon")
pq.insert(60, "Apple")
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pq.insert(25, "Netflix")

while pq.size != 0:
    maxItem = pq.getMax()
    print(maxItem.key, maxItem.value)
```

```
Apple 60
Microsoft 40
Google 30
Netflix 25
Amazon 20
```

Note: A heap doesn't **necessarily have to contain key-value pairs**. It can contain, for instance, just a **simple array of numbers**. The logic of a heap to extract min / max is essential to many logical processes!

Lab Session 1

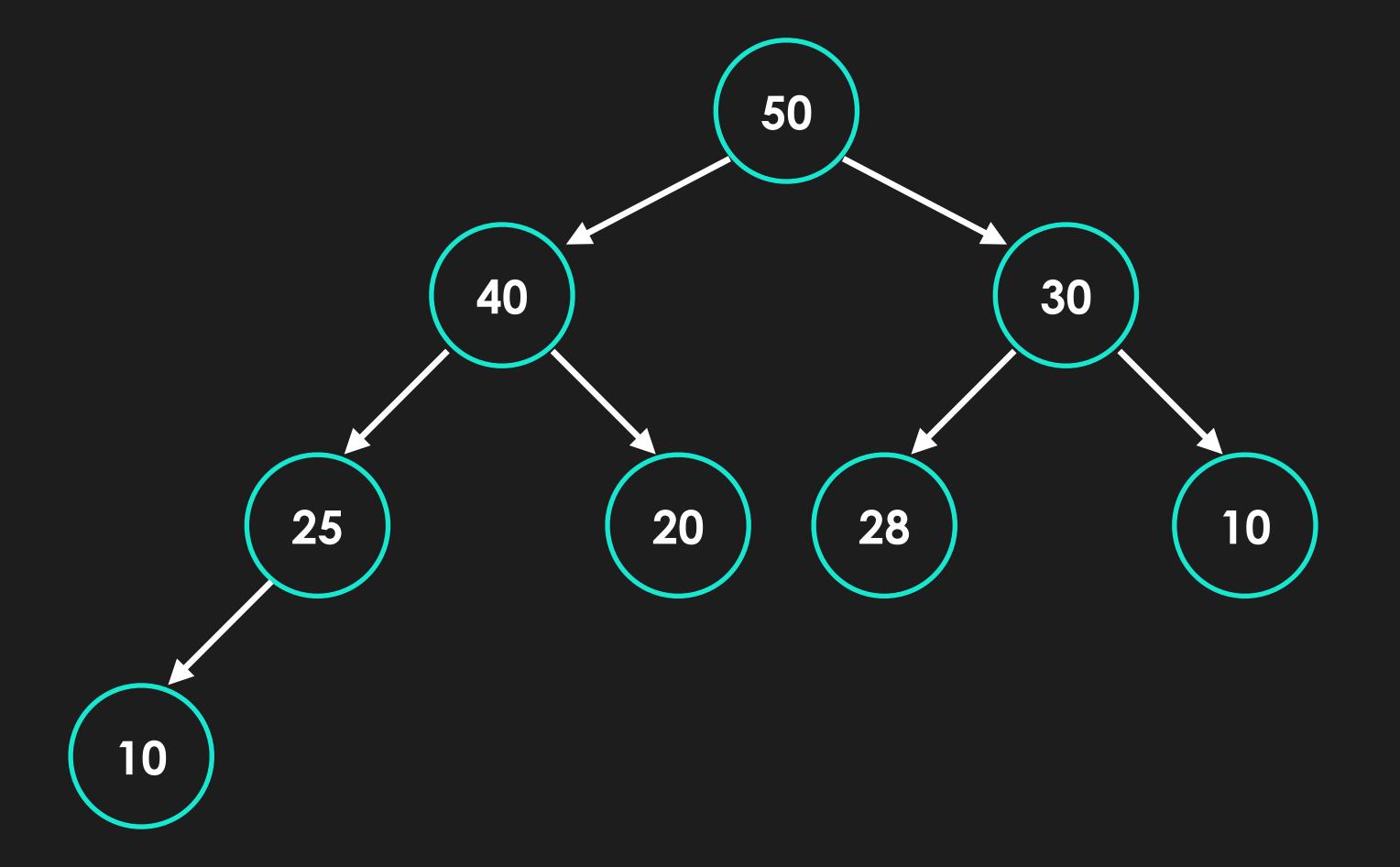
- Implement min_heap1.py
- To test, run `python utils/mh1_test.py`



Analysis of Heap Operations

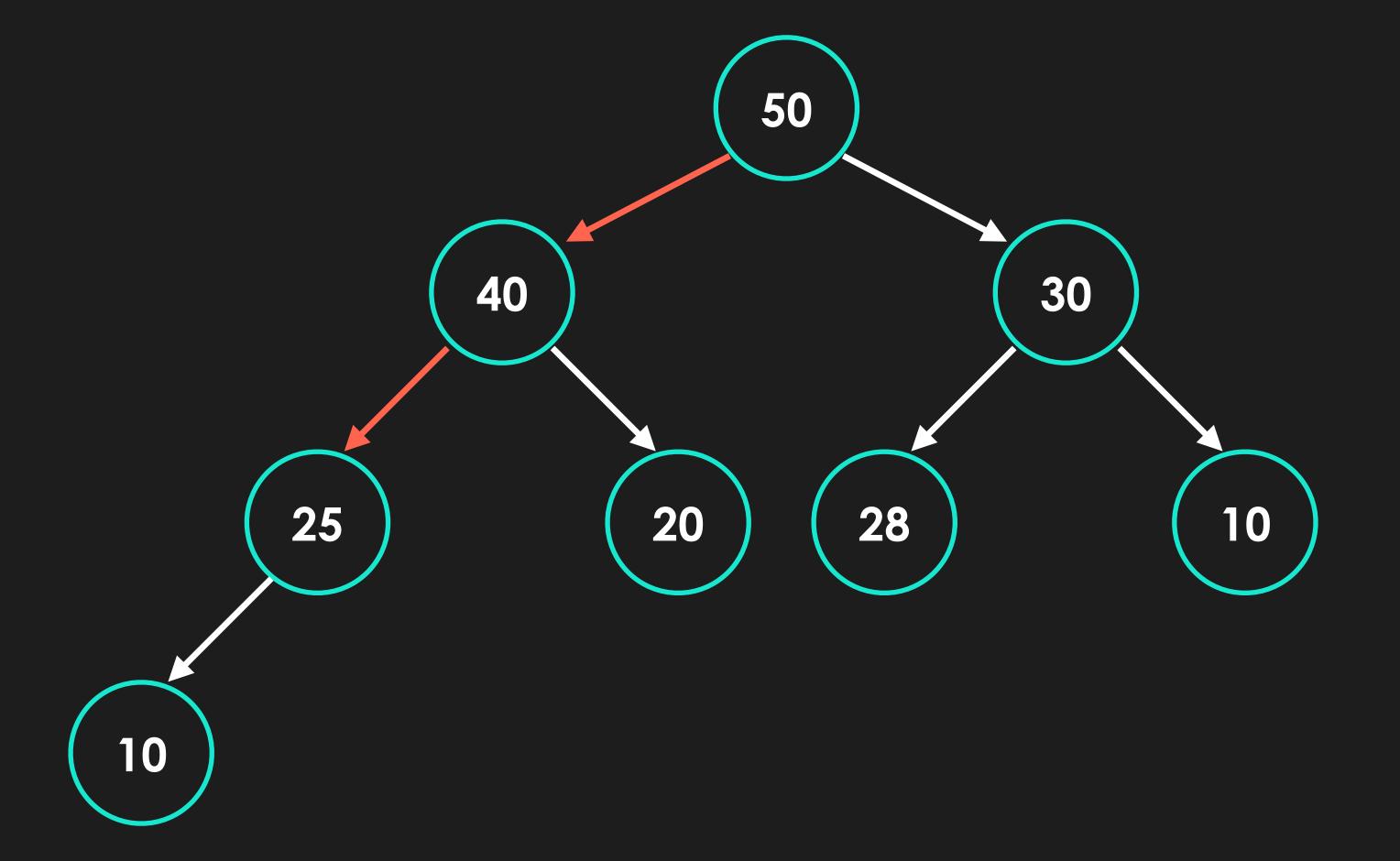
Tree Height

Maximum Tree Height





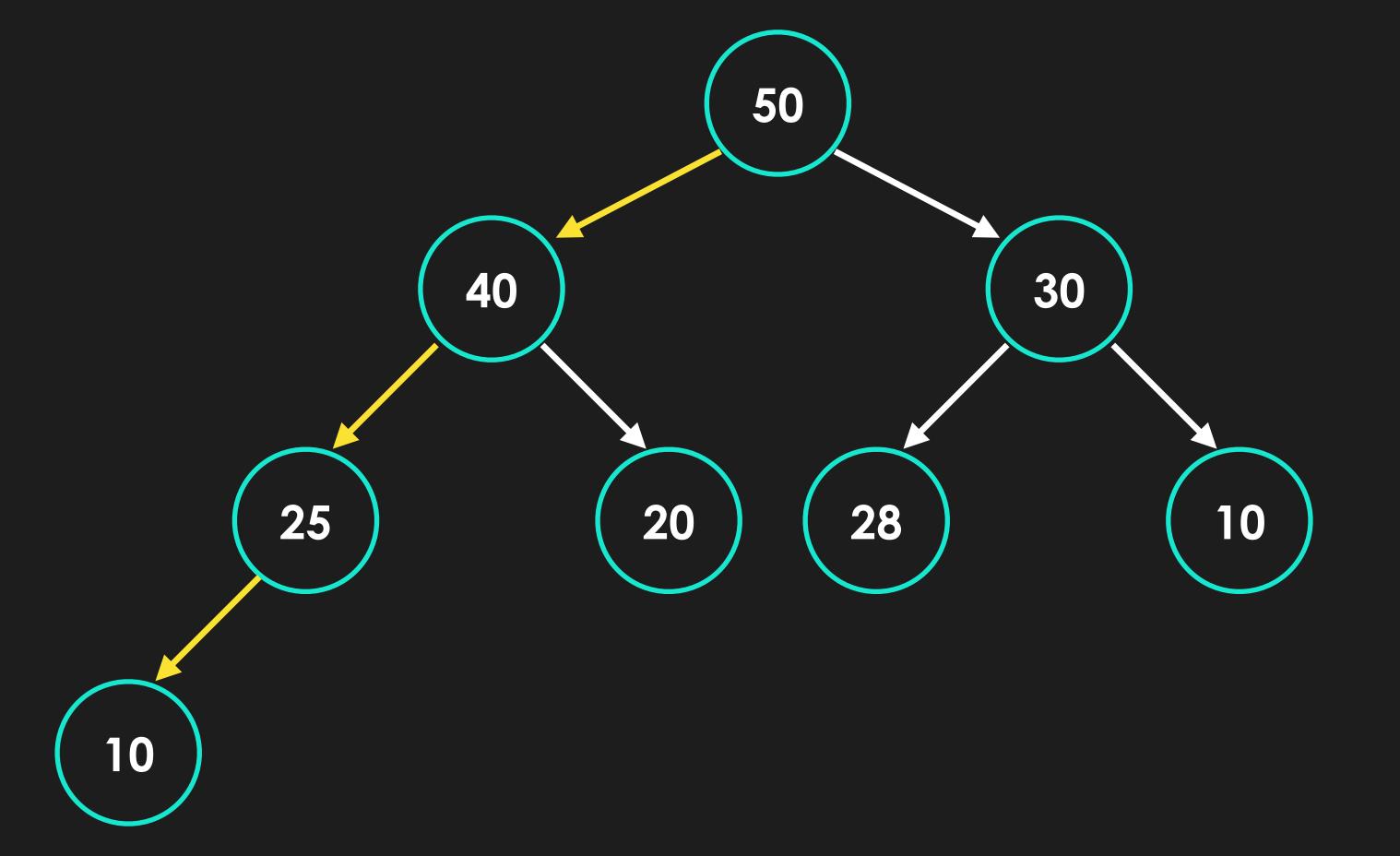
Maximum Tree Height 4 nodes: 2



Maximum Tree Height

4 nodes: 2

8 nodes: 3

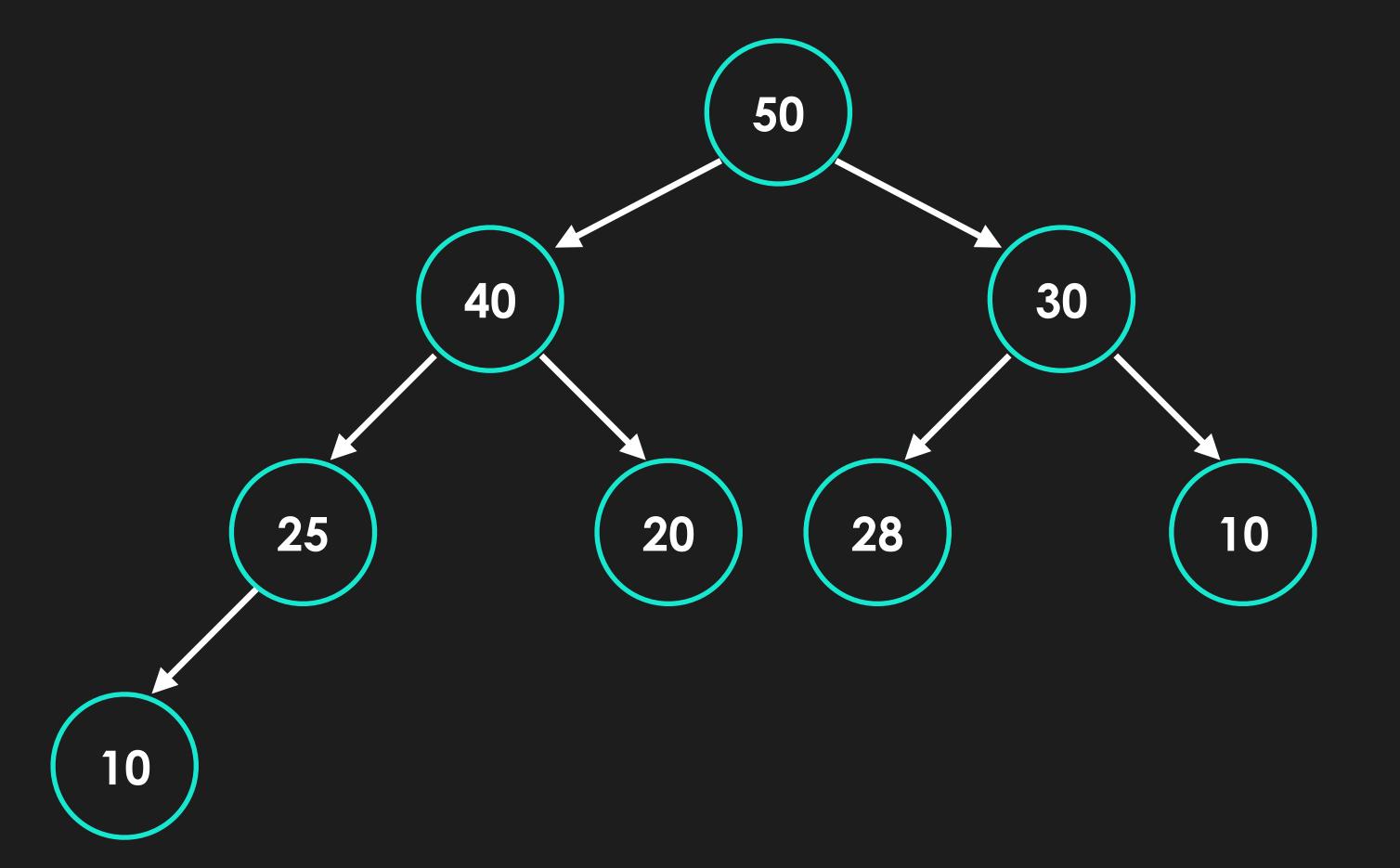


Maximum Tree Height

4 nodes: 2

8 nodes: 3

N nodes: logN (base 2)



Time complexity for sink & swim

- Both sink and swim have worst case time complexity of O (logN)
- This is because at most logN nodes are visited!

Time complexity for heap insertion

Assuming no **resize**, adding an extra item will take **constant time**, and calling **swim** on the new item will take **logN**. Therefore, heap insertion has **logN** worst case time complexity!

Time complexity for heap extraction

For extraction, swapping the first and last element takes constant time.
 Performing sink will take logN time. Therefore, heap extraction has worst case time complexity of logN

Comparison of Priority Queue Data Structures

	Worst C	ase	Average Case		
	Binary Heap	List	Binary Heap	List	
Insert	logN	1	logN	1	
Extract	logN	N	logN	Ν	

Note: If your application will have a huge number of inserts but very few removals of highest priority, lists / arrays would be the way to go because you can insert in **constant time**

MinHeap

Imagine if now instead of tracking returns, we want to track risk, and therefore keep track of investments with the lowest risk

Google	30%
Amazon	20%
Apple	60%
Microsoft	40%
Netflix	25%

We can use a min heap, and to construct one, all we have to make is one small change:

```
def greater(self, index1, index2):
   if self.heap[index1].value > self.heap[index2].value:
      return True
   else:
      return False
```



We can use a min heap, and to construct one, all we have to make is one small change:

Our comparisons now treat lower values with higher priority!



```
pq = MaxHeap(10)
pq.insert(30, "Google")
pq.insert(20, "Amazon")
pq.insert(60, "Apple")
pq.insert(40, "Microsoft")
pq.insert(25, "Netflix")

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pq.insert(25, "Netflix")

while pq.size != 0:
    maxItem = pq.getMax()
    print(maxItem.key, maxItem.value)
```

```
Amazon 20
Netflix 25
Google 30
Microsoft 40
Apple 60
```

Applications of Priority Queues

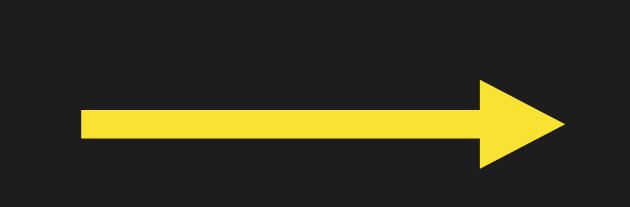
Applications of Priority Queues

- 1. Rank (Get kth largest element in a list)
- 2. Greedy Algorithms (E.g. Dijkstra's Shortest Path)
- 3. Huffman Compression
- 4. Stock Exchange Matching (match buyers to sellers, highest bid)

Priority Queues with changing values

Priority Queues with changing values

Google	30%
Amazon	20%
Apple	60%
Microsoft	40%
Netflix	25%



Google	30%
Amazon	20%
Apple	60%
Microsoft	10%
Netflix	90%

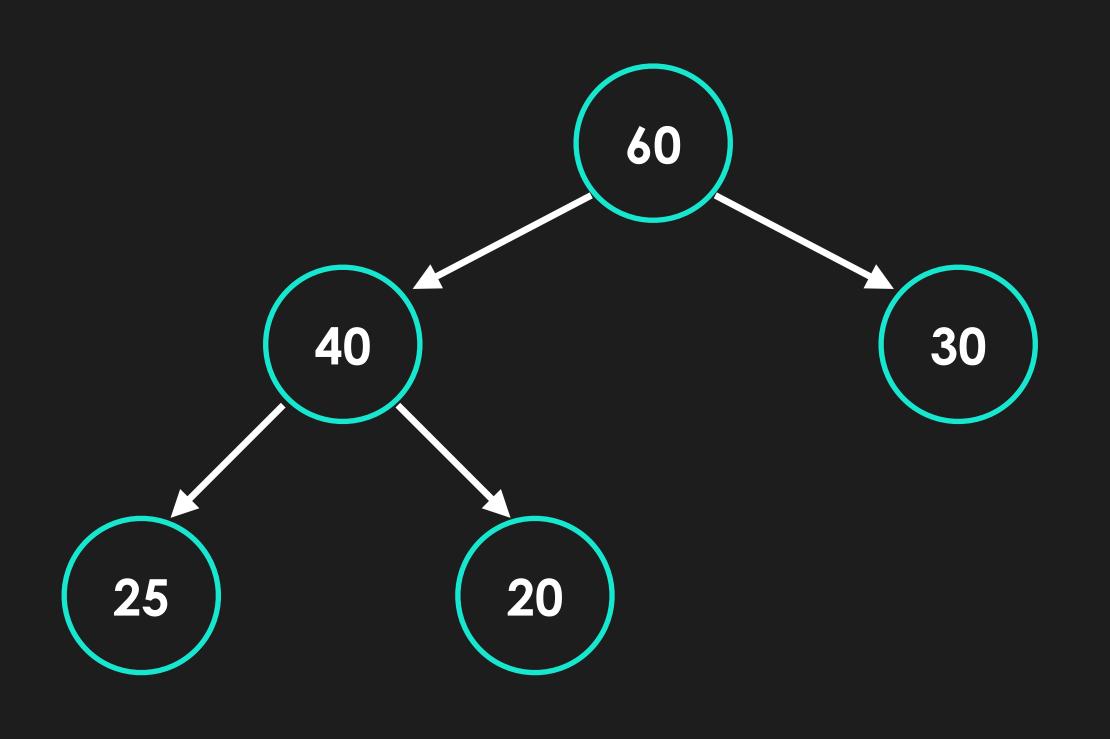
Priority Queues with changing values

Google	30%	Google	30%
Amazon	20%	Amazon	20%
Apple	60%	Apple	60%
Microsoft	40%	Microsoft	10%
Netflix	25%	Netflix	90%

If our keys' values change, how can we support this in a Heap structure?



We can create a dictionaries (positions), to store the index of each key in the heap

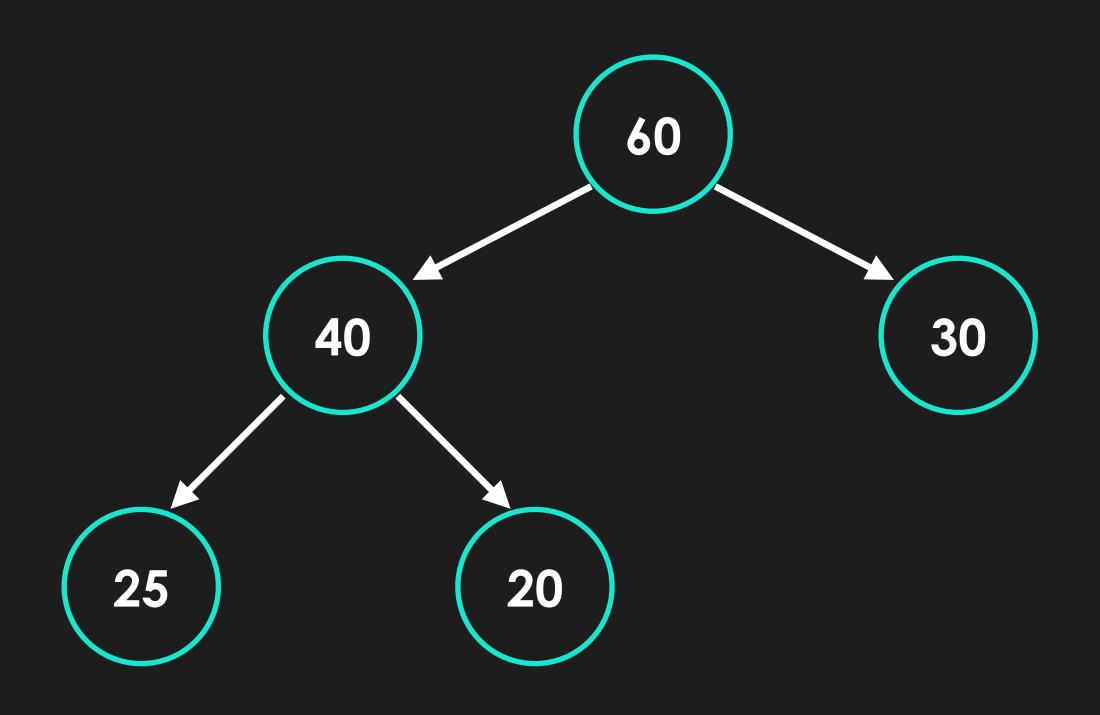


heap		60	40	30	20	25
	0	1	2	3	4	5

key	value
"Apple"	1
"Microsoft"	2
"Google"	3
"Netflix"	5
"Amazon"	4



STEP 1: Check if key exists



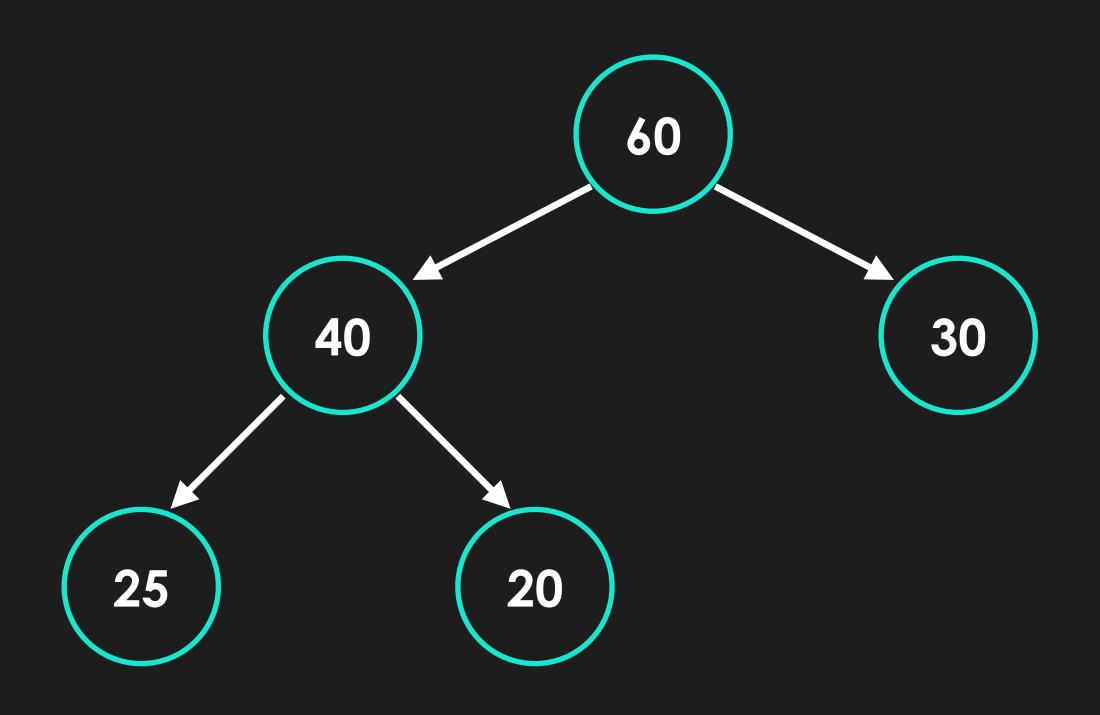
heap		60	40	30	25	20
	0	1	2	3	4	5

pq.insert("Microsoft", 10)

key	value
"Apple"	1
"Microsoft"	2
"Google"	3
"Netflix"	5
"Amazon"	4



STEP 1: Check if key exists



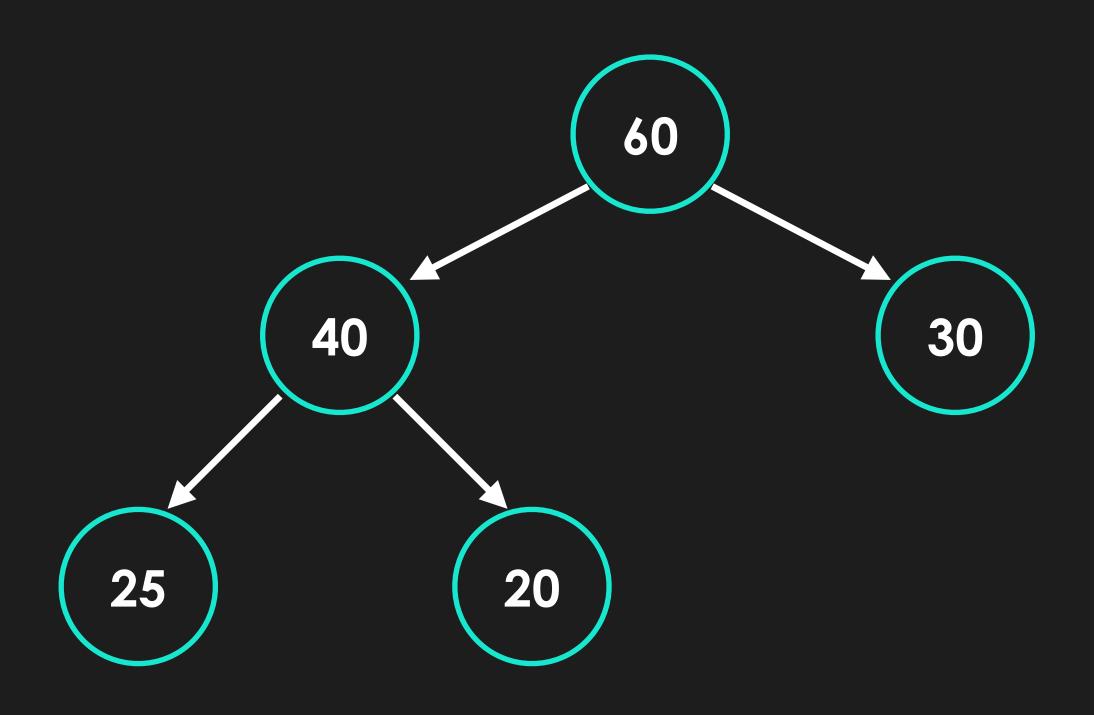
heap		60	40	30	25	20
	0	1	2	3	4	5

pq.insert("Microsoft", 10)

key	value
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"Microsoft"	2
"Google"	3
"Netflix"	4
"Amazon"	5



STEP 2: Get index based on positions dict



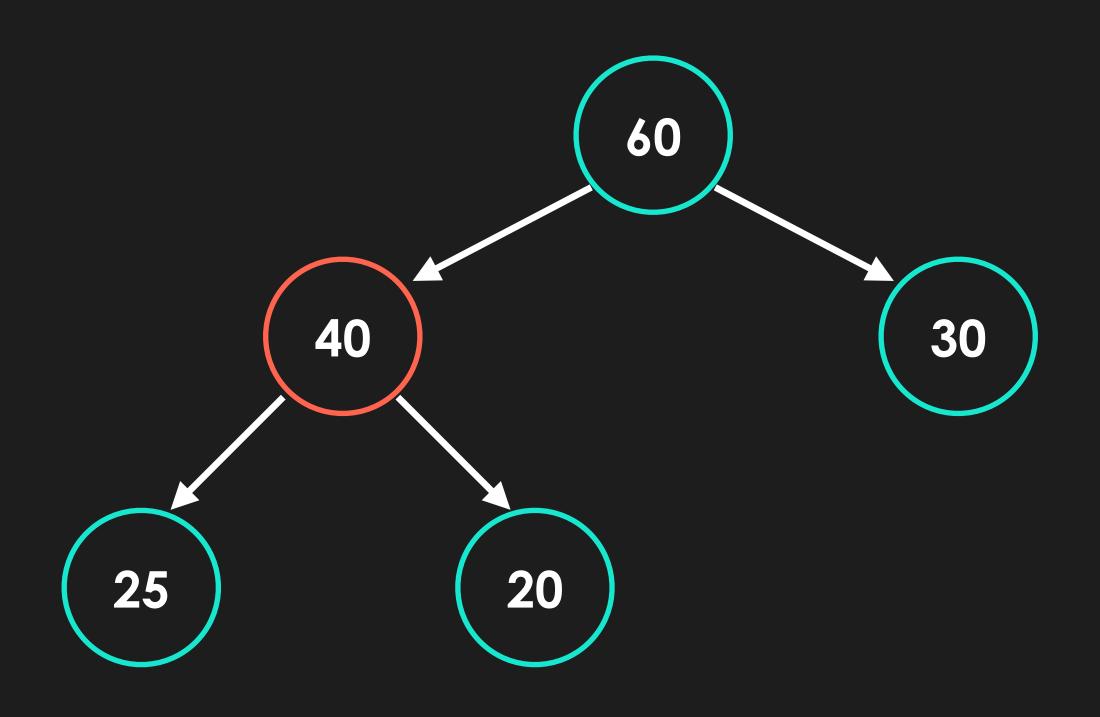
heap		60	40	30	25	20
	0	1	2	3	4	5

pq.insert("Microsoft", 10)

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"Netflix"	4
"Amazon"	5



STEP 2: Get index based on positions dict



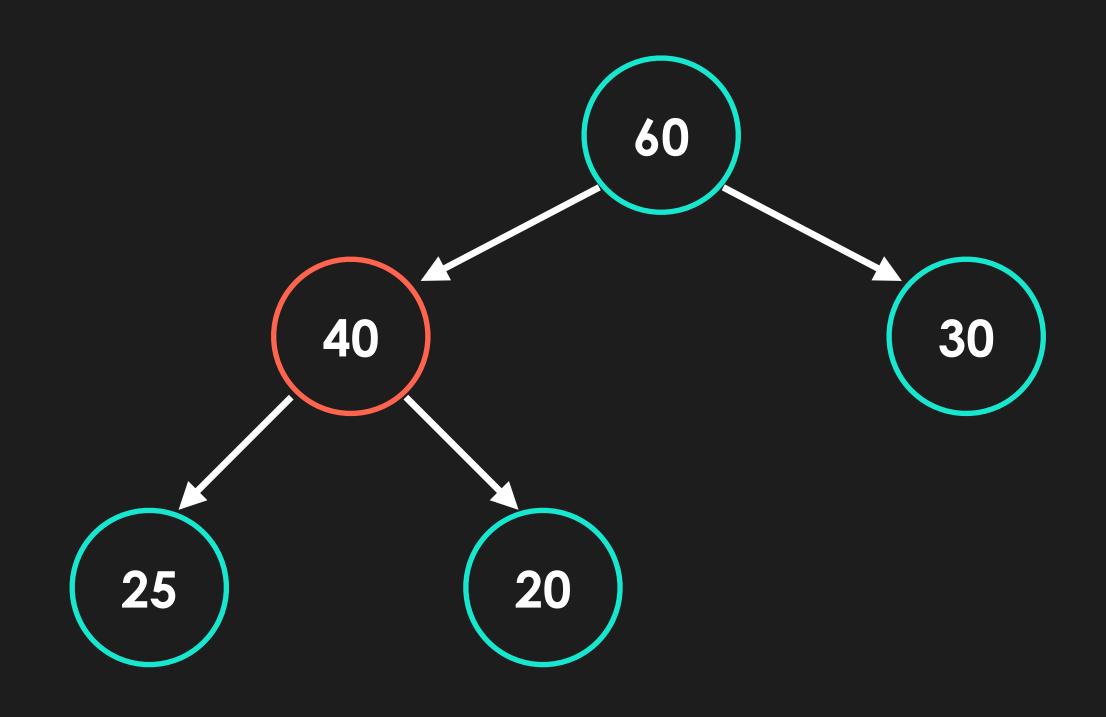
heap		60	40	30	25	20
	0	1	2	3	4	5

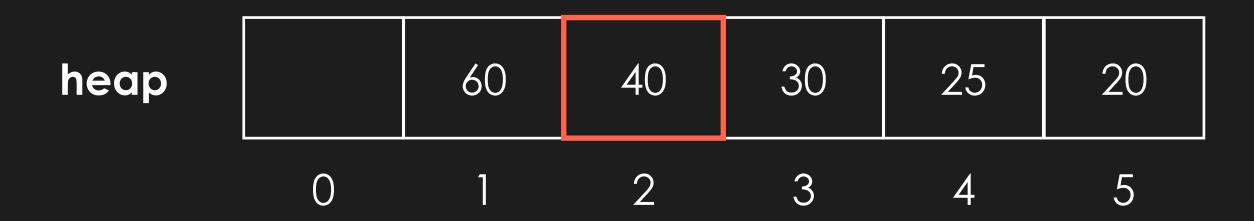
pq.insert("Microsoft", 10)

key	value
"Apple"	1
"Microsoft"	2
"Google"	3
"Netflix"	4
"Amazon"	5



STEP 3: Assign new value to item



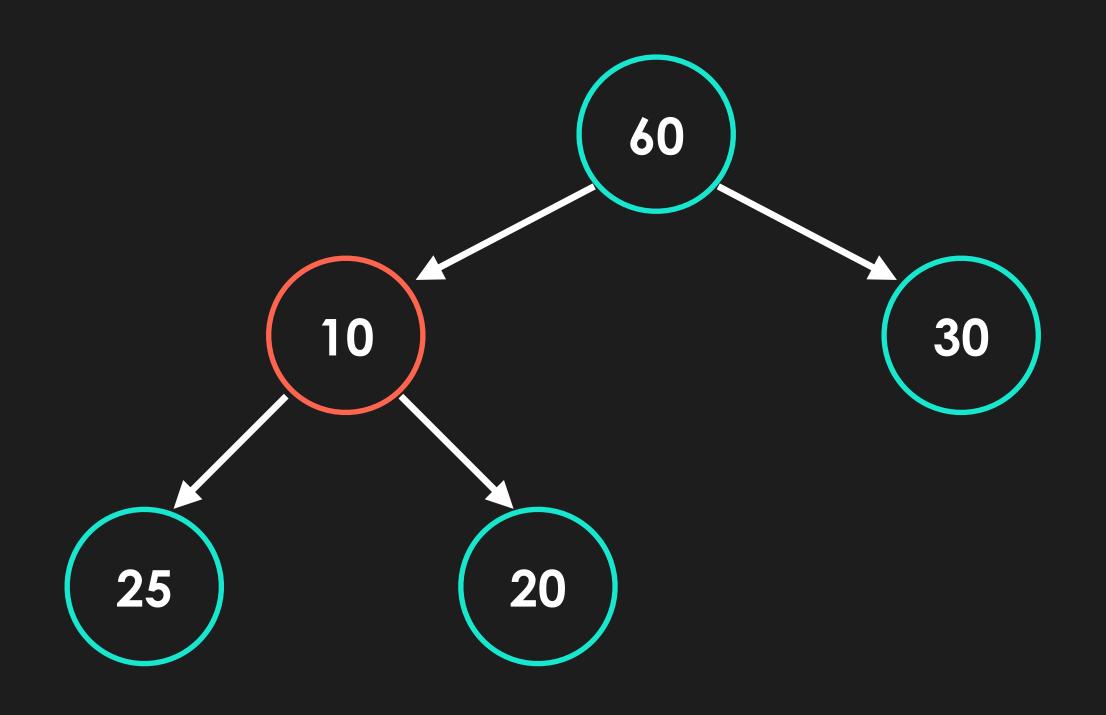


pq.insert("Microsoft", 10)

key	value
"Apple"	1
"Microsoft"	2
"Google"	3
"Netflix"	4
"Amazon"	5



STEP 3: Assign new value to item



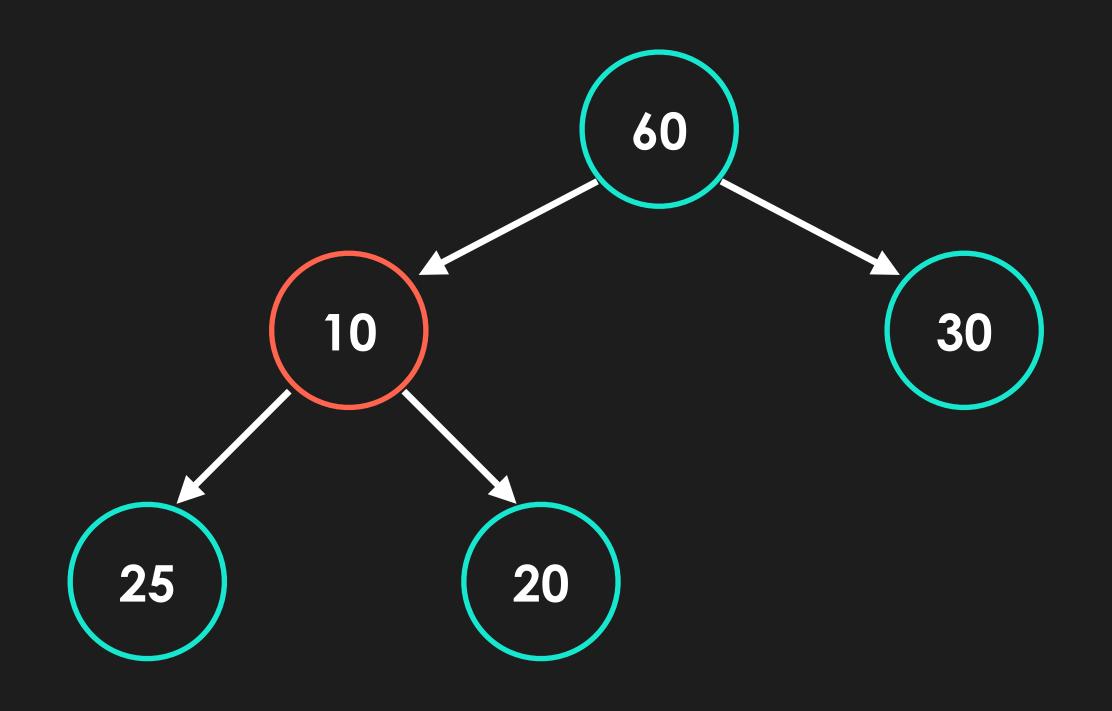
heap		60	10	30	25	20
	0	1	2	3	4	5

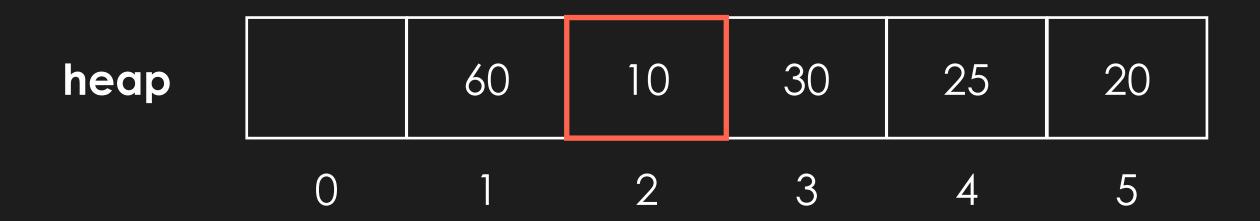
pq.insert("Microsoft", 10)

key	value
"Apple"	1
"Microsoft"	2
"Google"	3
"Netflix"	4
"Amazon"	5



STEP 4: If parent value is lower priority, perform swim, else, perform sink



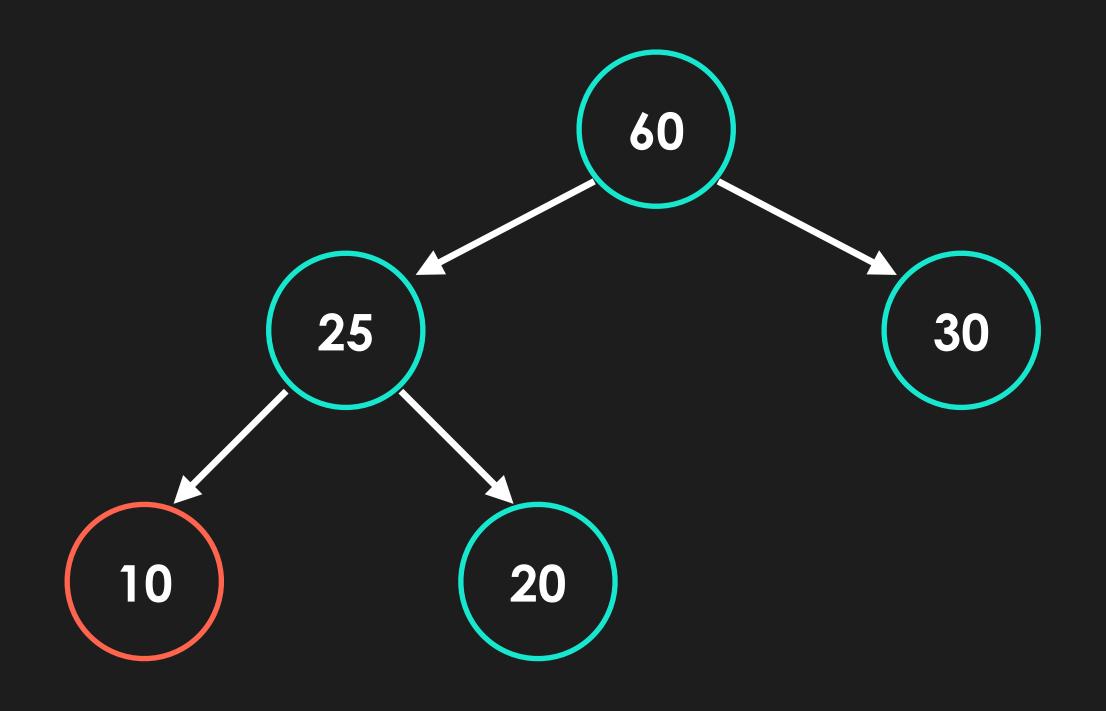


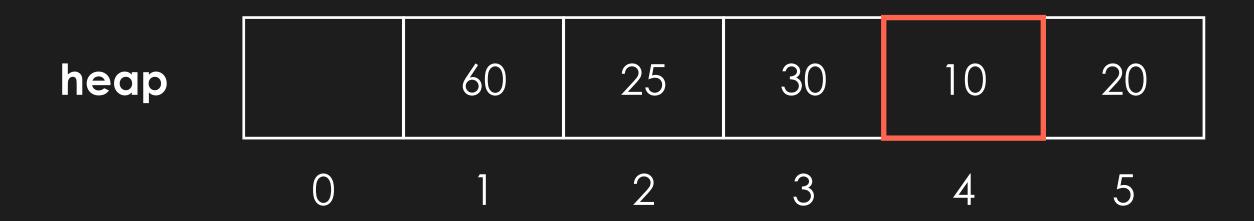
pq.insert("Microsoft", 10)

key	value
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"Netflix"	4
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STEP 4: If parent value is lower priority, perform swim, else, perform sink



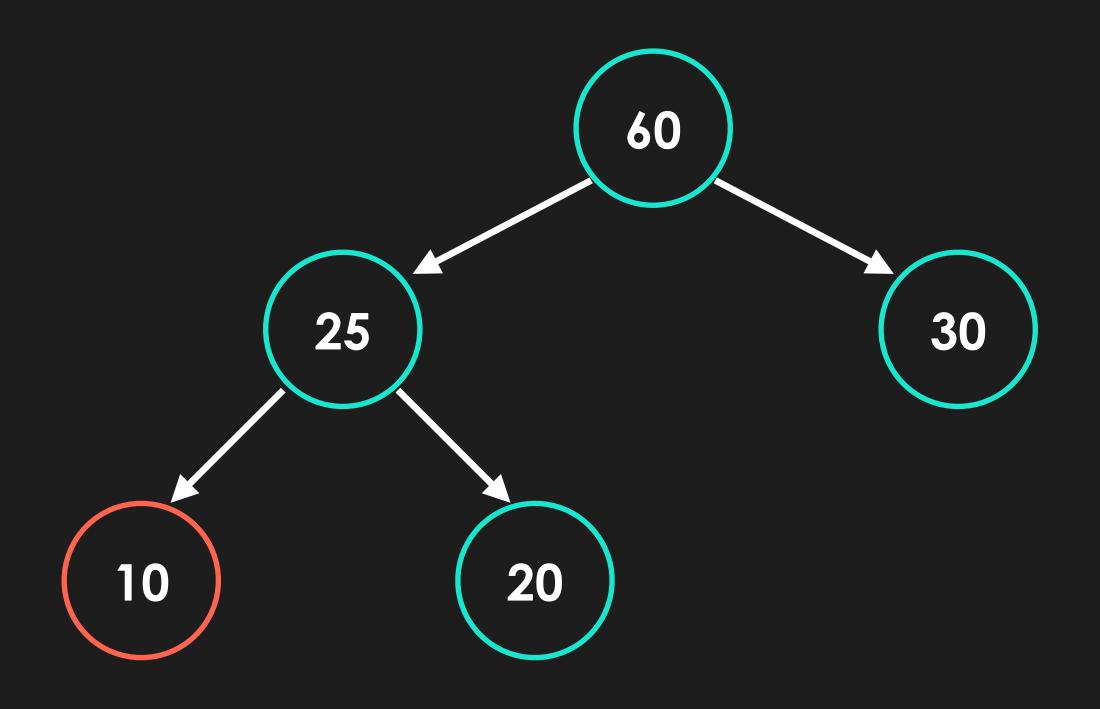


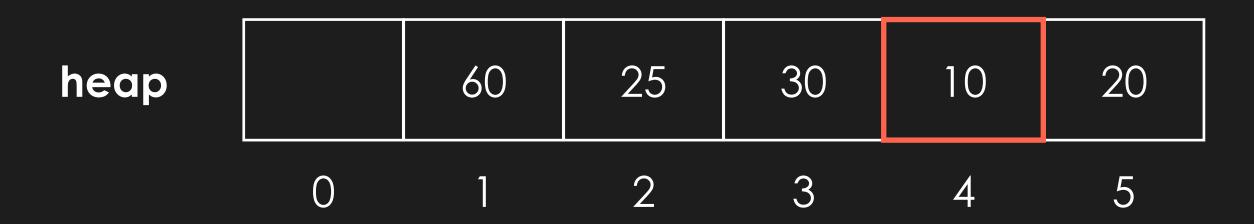
pq.insert("Microsoft", 10)

key	value
"Apple"	1
"Amazon"	2
"Google"	3
"Microsoft"	4
"Amazon"	5



Note: When swapping, you must now also swap in positions dictionary!





pq.insert("Microsoft", 10)

key	value
"Apple"	1
"Amazon"	2
"Google"	3
"Microsoft"	4
"Amazon"	5



Implementation



constructor

```
def __init__ (self, maxsize):
    self.maxsize = maxsize
    self.size = 0
    self.heap = [None] * (maxsize + 1)
    self.positions = {}
```



swap

```
def swap (array, positions, index1, index2):
    key1 = array[index1].key
    key2 = array[index2].key
    positions[key1], positions[key2] = positions[key2], positions[key1]
    array[index1], array[index2] = array[index2], array[index1]
```



insert

```
def insert(self, newKey, newValue):
    if self.size >= self.maxsize:
        print('size limit reached')
        return
    # HANDLE EXISTING KEY, NEW VALUE
    if newKey in self.positions:
        currentIndex = self.positions[newKey]
        self.heap[currentIndex].value = newValue
        if self.greater(currentIndex, parent(currentIndex)):
            self.swim(currentIndex)
        else:
            self.sink(currentIndex)
        return
    self.size += 1
    self.heap[self.size] = HeapItem(newKey, newValue)
    self.positions[newKey] = self.size
    self.swim(self.size)
```

When would we need changing values?

When would we need changing values?

In Dijkstra's algorithm, if we have a heap that handles changing values, we can achieve a faster time complexity for the algorithm!

Lab Session 1

- Implement min_heap2.py
- To test, run `python utils/mh2_test.py`

