

# Exercises for Introduction to Quantum Computing

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In [ ]: import numpy as np
import qiskit
import qiskit.circuit.library as circuit
import qiskit.quantum_info as qi
from qiskit import QuantumCircuit, QuantumRegister
from qiskit import IBMQ, Aer
from qiskit.providers.ibmq import least_busy
from qiskit import QuantumCircuit, transpile, execute
import math
from qiskit.visualization import plot_histogram
from qiskit.circuit.library import QFT
from qiskit.providers.aer import Aer
```

## 1. Multiplication by Quantum Fourier Transform

In this exercise sheet we would like to implement a Quantum multiplier i.e. a Quantum circuit that could multiply two numbers stored in two quantum registers upto modulo  $2^n$  where  $n$  is the number of Qubits in either of the Quantum registers.

a) The position operator acting on the first quantum register is defined as,

$$\hat{x} = \sum_x x |x\rangle \langle x|$$

and the corresponding momentum operator acting on the first is defined to be,

$$\hat{p}_x = F_x^\dagger \hat{x} F_x$$

where  $F_x \equiv \text{qFT}_x$ . We have two quantum registers,

$$|x\rangle |y\rangle = |x, y\rangle$$

We are to prove that,

$$e^{i\hat{x}\hat{p}_y} = F_y^\dagger e^{i\hat{x}\hat{y}} F_y$$

We shall first expand out the momentum operator in the L.H.S,

$$e^{i\hat{x}\hat{p}_y} = e^{i\hat{x}F_y^\dagger \hat{y} F_y}$$

Since  $\hat{x}$  does not act on  $y$ , we can rewrite this to be,

$$e^{i\hat{x}\hat{p}_y} = e^{F_y^\dagger (i\hat{x}\hat{y}) F_y}$$

We know that,

$$e^{P.A.P^{-1}} = P e^A P^{-1}$$

where  $A$  is some square matrix and  $P$  is an invertible matrix. We see that this identity is exactly what we need, substituting  $A = i\hat{x}\hat{y}$  and  $P = F_y^\dagger$ , we have

$$e^{i\hat{x}\hat{p}_y} = F_y^\dagger e^{i\hat{x}\hat{y}} F_y$$

Hence, proved.

b) We are to show that,

$$e^{ixy} = \prod_{k,l=1}^n e^{ix_k y_l 2^{n-k-l}}$$

Let's begin by writing down  $x$  and  $y$  in their binary representations

$$x = \sum_n^{k=1} x_k 2^{n-k}$$

$$y = \sum_n^{l=1} y_l 2^{n-l}$$

thus, we have,

$$e^{ixy} = e^{i \sum_{k,l=1}^n x_k 2^{n-k} \cdot y_l 2^{n-l}}$$

$$e^{ixy} = e^{i \sum_{l,k}^n x_k y_l 2^{n-k-l}}$$

$$e^{ixy} = \prod_{k,l=1}^n e^{ix_k y_l 2^{n-k-l}}$$

Hence, proven.

c) Using the previous results, we are to show that,

$$e^{\frac{2\pi i \hat{x} \hat{p}_y}{2^n}} |x, y\rangle = |x, x + y \mod 2^n\rangle$$

using the result from (a) let us rewrite the L.H.S,

$$e^{\frac{2\pi i \hat{x} \hat{p}_y}{2^n}} |x, y\rangle = F_y^\dagger e^{\frac{2\pi i \hat{x} y}{2^n}} F_y |x, y\rangle$$

acting with  $F_y$  on  $|x, y\rangle$ , we get,

$$e^{\frac{2\pi i \hat{x} \hat{p}_y}{2^n}} |x, y\rangle = \frac{F_y^\dagger}{2^{\frac{n}{2}}} \sum_{k=1}^n e^{\frac{2\pi i \hat{x} k}{2^n}} e^{\frac{2\pi i k y}{2^n}} |x, k\rangle$$

from (b) we know that this can then be written as,

$$e^{\frac{2\pi i \hat{x} \hat{p}_y}{2^n}} |x, y\rangle = \frac{F_y^\dagger}{2^{\frac{n}{2}}} \sum_{k=1}^n e^{\frac{2\pi i \hat{x} k}{2^n}} e^{\frac{2\pi i k(x+y)}{2^n}} |x, k\rangle$$

now we apply the inverse Fourier transform to get,

$$e^{\frac{2\pi i \hat{x} \hat{p}_y}{2^n}} |x, y\rangle = \frac{1}{2^n} \sum_{k, m=0}^{2^n-1} e^{\frac{2\pi i k[(x+y)-m]}{2^n}} |m\rangle$$

we use the identity,

$$\delta_{b,c} = \frac{1}{N} \sum_{j=1}^N e^{\frac{2\pi i j \cdot (b-c)}{N}}$$

Thus we have,

$$e^{\frac{2\pi i \hat{x} \hat{p}_y}{2^n}} |x, y\rangle = |x, x + y \mod 2^n\rangle$$

d) Now, we shall consider three quantum registers  $|x\rangle$ ,  $|y\rangle$ , and  $|z\rangle$  with  $n$  qubits each. We shall attempt to construct a multiply add operation of the form,

$$|x, y, z\rangle \rightarrow |x, y, z + xy \bmod 2^n\rangle$$

First, we need to construct a circuit that multiplies  $x$  and  $y \pmod{2^n}$ . We know that multiplying  $x$  by  $y$  simply implies that, we are adding  $x$  to itself  $y$  times. Thus, a multiplier circuit is simply a quantum adder circuit between the state  $|0\rangle$  and  $|x\rangle$  performed iteratively  $y$  times (this is implemented either by adding a controlled not between the  $|y\rangle$  state and the adder for every iteration and decrementing after each iteration or via controlled weighted addition). The addition of  $z$  to the product can then be implemented via an adder circuit, let us prove this now, we saw from (c), that the addition operation can be represented via,

$$e^{\frac{2\pi i \hat{x} \hat{p}_y}{2^n}} |x, y\rangle = |x, x + y \bmod 2^n\rangle$$

to multiply  $x$  and  $y$ , we simply do,

$$\left(e^{\frac{2\pi i \hat{x} \hat{p}_2}{2^n}}\right)^y |x, 0, z\rangle = |x, xy, z\rangle$$

Where the 0 here actual stands for  $|0\rangle^{\otimes n}$  i.e. we have initialized the second register i.e. the accumulator to be in the  $|0\rangle$  state. The second register now contains the product  $xy \bmod 2^n$ . The addition with  $z$  can then be written in the form,

$$e^{\frac{2\pi i \hat{y} \hat{p}_z}{2^n}} \left(e^{\frac{2\pi i \hat{x} \hat{p}_2}{2^n}}\right)^y |x, 0, z\rangle = |x, xy, z + xy \bmod 2^n\rangle$$

e) We shall now implement this multiply-add operation in Qiskit.

```
In [ ]: ccp = circuit.UGate(0.1,0,0).control(2, ctrl_state="11")
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In [ ]: n = 2 #number of qubits in one register

x = QuantumRegister(n)
y = QuantumRegister(n)
ancilla = QuantumRegister(n)
z = QuantumRegister(n)

qc = QuantumCircuit(x,y,ancilla,z)

param = np.pi

ccp = circuit.PhaseGate(param).control(2, ctrl_state="11")

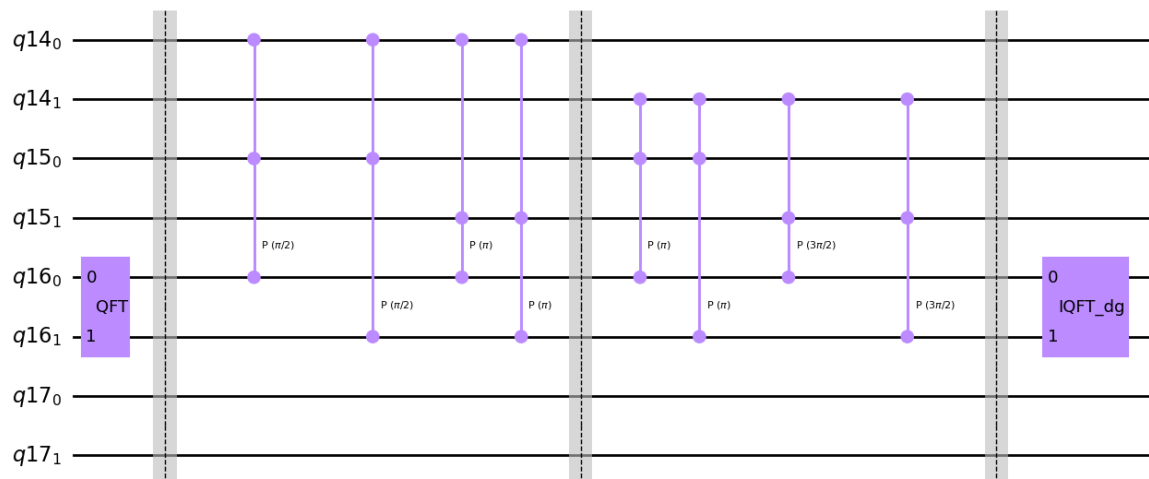
qc.compose(QFT(n, inverse=False, do_swaps=True), [3*n-i for i in range(n,0,-1)], in
qc.barrier()

for i in range(n):
    for j in range(n):
        for k in range(n):
            param = (2*np.pi*((3*n)-3-(int(i)+int(j)+int(k))))/(2**(n))
            ccp = circuit.PhaseGate(param).control(2, ctrl_state="11")
            qc.append(ccp, [x[i],y[j],ancilla[k]])
        qc.barrier()

qc.compose(QFT(n, inverse=True,do_swaps=True), [3*n-i for i in range(n,0,-1)], inpl
qc.draw('mpl')

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Out [ ]:



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In [ ]: state_2 = qi.Statevector.from_label('000101')
state_2 = state_2.evolve(qc)
(state_2.draw(output='latex'))

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Out [ ]:

$|110101\rangle$