

# SURP - Graph Theory

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## 1 Introduction

A 1-factor in a directed graph is a spanning 1-regular subgraph. Equivalently, it is a collection of disjoint directed cycles such that every vertex in the graph is in exactly one cycle. Call a 1-factor special if for any pair of vertices  $u, v$  such that both  $(u, v)$  and  $(v, u)$  are edges, either both edges are in the 1-factor, or none of them is.

A directed graph is symmetric if  $(b, a)$  is an edge whenever  $(a, b)$  is an edge. Also we define the  $(a, b)$ -star graph as a directed graph on  $a + b + 1$  vertices with the center vertex  $u$ , such that there are  $a$  vertices  $x_1, x_2, \dots, x_a$  with edges  $(u, x_i)$  for all  $i \in [a]$ , and  $b$  vertices  $y_1, y_2, \dots, y_b$  with edges  $(y_i, u)$  for all  $i \in [b]$ . A  $(a, b)$ -star decomposition of a directed graph is a partition of the edges of the graph into  $(a, b)$ -stars, and such a graph is called  $(a, b)$ -star decomposable.

In this article, I study the decomposition of directed regular graphs and directed bipartite graphs into star graphs, and make some observations regarding the relation between a graph being star decomposable and being factorizable into 1-factors or special 1-factors.

## 2 Star Decomposition of Regular Digraphs

**LEMMA 1:** For  $a > 1$ , if an  $(a + 1)$ -regular directed graph  $\mathcal{G}$  is  $(a, 1)$ -star decomposable, then every vertex in  $\mathcal{G}$  is the center of exactly one  $(a, 1)$ -star graph

*Proof.* First suppose some vertex is the center of at least two such star graphs. Then it will have at least  $2a$  outgoing edges, contrary to our assumption that  $\mathcal{G}$  is  $a + 1$ -regular. Now if  $\mathcal{G}$  has  $n$  vertices, then it must have a total of  $(a + 1)n$  directed edges. So it must have gotten divided into  $n$  star graphs, which means that there are  $n$  centers of these star graphs. This is only possible if each vertex is the center of exactly one star graph.  $\square$

**LEMMA 2:** Suppose we have a 3-regular symmetric directed graph  $\mathcal{G}$ . Then it has a  $(2, 1)$ -star decomposition if and only if the underlying undirected graph has a perfect matching.

*Proof. IF*

Suppose the graph has a perfect matching  $\mathcal{M}$ . Take an edge in  $\mathcal{M}$  joining vertices  $u$  and  $v$ . Then, as  $\mathcal{G}$  is symmetric, so it contains both edges  $(u, v)$  and  $(v, u)$ . Suppose  $x_1$  and  $y_1$  are the other two neighbors of  $u$ , and  $x_2$  and  $y_2$  are the other two neighbors of  $v$ . Then the graphs  $\{(u, x_1), (u, y_1), (v, u)\}$  and  $\{(v, x_2), (v, y_2), (u, v)\}$  are star graphs. We can do the same division for each edge of  $\mathcal{M}$ , and note that none of the edges are common to any two star graphs. Thus, we get a  $(2, 1)$ -star decomposition of  $\mathcal{G}$ .

**ONLY IF**

Now suppose  $\mathcal{G}$  is  $(2, 1)$ -star decomposable. Now take  $v$  as some vertex in  $\mathcal{G}$ . Then by Lemma 1, it is the center of a unique  $(2, 1)$ -star graph  $\{(v, x), (v, y), (u, v)\}$ . Then the star graph of which  $u$  is a center cannot have an edge towards  $v$ . So this star graph must have the edge  $(v, u)$ . And so  $u$  and  $v$  form a unique pair. We can then take such pairs and get the desired perfect matching.  $\square$

Note that if we have a symmetric directed graph, then it has a special 1-factor if and only if the underlying undirected graph has a perfect matching (since the only directed cycles allowed in a special 1-factor in a symmetric graph will be a directed cycle of length 2, which can be thought of as a simple edge in the undirected graph). This observation can be extended to a more general situation, as given in the section below.

### 3 Special 1-factors and Star Decomposition

**LEMMA 3:** For  $a \in \mathbb{N}$ , consider an  $(a + 1)$ -regular directed graph  $\mathcal{G}$ . Then-

1. If  $\mathcal{G}$  has a special 1-factor, then it also has an  $(a, 1)$ -star decomposition.
2. If  $a > 1$ , then  $\mathcal{G}$  has a 1-factor if it is  $(a, 1)$ -star decomposable.

*Proof.* 1. Consider a vertex  $v$  in  $\mathcal{G}$  such that it has outgoing edges towards  $x_1, x_2, \dots, x_{a+1}$  and incoming edges from  $y_1, y_2, \dots, y_{a+1}$ . Suppose in the special 1-factor, this vertex is a part of a directed cycle which passes through  $y_1 \rightarrow v \rightarrow x_1$ . Then we take one of the  $(a, 1)$ -stars as  $\{(y_2, v), (y_3, v), \dots, (y_{a+1}, v), (v, x_1)\}$ , and do this for all vertices. To prove that this works, note that none of the other edges adjacent to  $v$  can be a part of a star graph centered at  $v$ . Also this sort of division without running into an edge twice will always be possible since the vertex  $x_1$  cannot be the same as any of the vertices  $y_i$  for  $2 \leq i \leq a + 1$  (So each  $y_i$  has the star graph). Thus  $\mathcal{G}$  is  $(a, 1)$ -star decomposable.

2. Now suppose  $a > 1$  and  $\mathcal{G}$  has a  $(a, 1)$ -star decomposition. By Lemma 1, each vertex  $v$  in  $\mathcal{G}$  is the center of exactly one  $(a, 1)$ -star graph. Let this star graph has an edge  $(v, x)$ , and let  $(y, v)$  be the only incoming edge to  $v$  which is not a part of this star graph. Then we consider the  $(1, 1)$ -star  $y \rightarrow v \rightarrow x$ , and take the union of all such  $(1, 1)$ -stars to get the desired 1-factor. Here we use the fact that  $(y, v)$  will be a part of the  $(1, 1)$ -star centered at  $y$ , and so we'll get directed cycles by this union operation. □

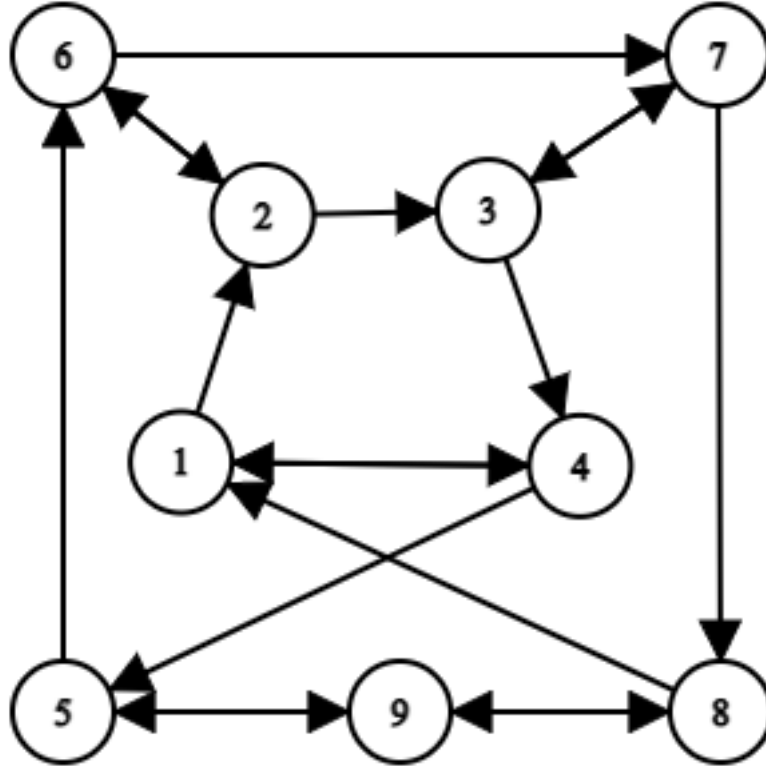


Figure 1: 1-factorizable and  $(1, 1)$ -star docomposable 2-regular digraph with no special 1-factor

## 4 1-factors in Bipartite Digraphs

**LEMMA 4:** A bipartite directed graph  $\mathcal{G}$  with bipartite sets  $X$  and  $Y$  has a 1-factor if and only if the subgraph consisting of edges from  $X$  to  $Y$  and the subgraph consisting of edges from  $Y$  to  $X$  both have a perfect matching.

*Proof.* **IF**

First assume the two subgraphs have perfect matchings. Then for every vertex  $u$  in  $X$ , there is a vertex  $v$  in  $Y$  and  $w$  in  $X$  ( $w$  might be the same as  $x$ ) such that  $(u, v)$  is a perfect matched pair in the first subgraph, and  $(v, w)$  a perfect matched pair in the second one. We can then divide the whole graph into such triplets  $(u, v, w)$  for any  $u$  in  $X$ . Joining these edges gives us the desired 1-factor.

**ONLY IF**

Suppose  $\mathcal{G}$  has a 1-factor. Then for each vertex  $u$  in  $X$ , there is a unique vertex  $v$  in  $Y$  such that  $(u, v)$  is a directed edge in the 1-factor. Dividing  $\mathcal{G}$  into these pairs  $\{u, v\}$  gives the desired perfect matching for the subgraph consisting of edges from  $X$  to  $Y$ . Similar argument works for the other direction. Note that, as a corollary, this also gives that  $|X| = |Y|$ .  $\square$

**COROLLARY:** Any  $d$ -regular bipartite directed graph, for  $d \geq 1$ , has a 1-factor.

*Proof.* Due to the above lemma, we just need to show the existence of a perfect matching in the two subgraphs which only have edges from one bipartition to another and vice-versa. Take one such subgraph, say from bipartition  $X$  to  $Y$ . Then for any subset  $S$  of  $X$ , the number of edges adjacent to  $S$  and present in this subgraph is  $d \cdot |S|$ . If  $N(S)$  is the neighborhood of  $S$  in this subgraph, then each vertex in  $N(S)$  has exactly  $d$  incoming edges, so we get

$$d \cdot |S| \leq d \cdot |N(S)| \Rightarrow |S| \leq |N(S)|$$

Then Hall's Marriage Theorem gives us that this subgraph has a perfect matching. By a similar argument, the other subgraph also has a perfect matching, and so a  $d$  regular bipartite graph has a 1-factor.  $\square$

The above lemma also gives us a polynomial time algorithm to find a 1-factor in a bipartite directed graph. We can simply find perfect matchings in the subgraph consisting of edges from  $X$  to  $Y$  and  $Y$  to  $X$ , and then for each pair  $\{u, v\}$  in the first matching, consider the corresponding pair  $\{v, w\}$  in the second matching. Then we take the path  $u \rightarrow v \rightarrow w$ , and continue the same operation for  $w$ . Finally the whole graph will be divided into various directed cycles, which gives us the desired 1-factor.