



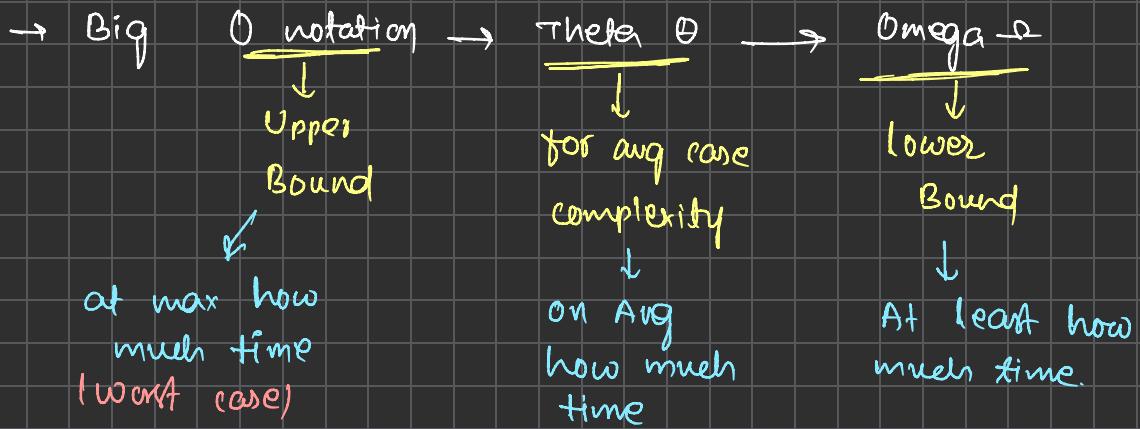
## \* Time & Space complexity.

### \* Time complexity:

→ It is the amount of time taken by an algorithm to run.

→ as a function of length of the input.

why? ← for making better program  
    └ comparison of algo



→ Constant time  $\rightarrow O(1)$

→ Linear time  $\rightarrow O(n)$

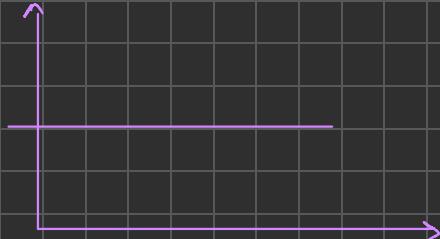
→ Logarithmic time  $\rightarrow O(\log n)$

→ Quadratic time  $\rightarrow O(n^2)$

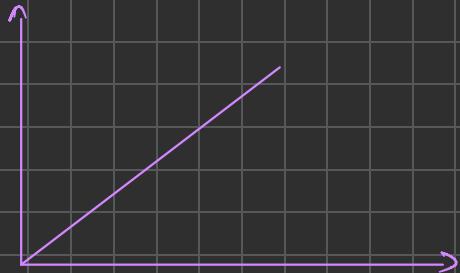
→ Cubic time  $\rightarrow O(n^3)$

## Graphs:

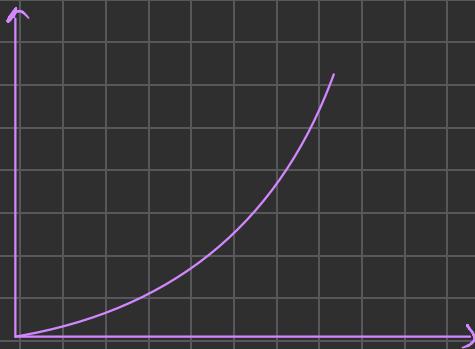
→  $O(1)$



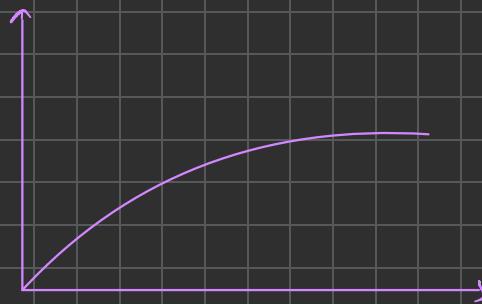
→  $O(n)$



→  $O(n^2)$



→  $O(\log n)$



$O(n!)$

$O(2^n)$

$O(n^3)$

$O(n^2)$

$O(n \log n)$

$O(n)$

$O(\log n)$

$O(1)$

↑  
Complexity

Questions. Ans.

$$f(n) \rightarrow 2n^2 + 3n \rightarrow O(n^2)$$

$$f(n) \rightarrow 4n^4 + 3n^3 \rightarrow O(n^4)$$

$$f(n) \rightarrow n^2 + \log n \rightarrow O(n^2)$$

$$f(n) \rightarrow 12001 \rightarrow O(1)$$

$$f(n) \rightarrow 3n^3 + 2n^2 + 5 \rightarrow O(n^3)$$

$$f(n) \rightarrow \frac{n^3}{300} \rightarrow O(n^3)$$

$$f(n) \rightarrow 5n^2 + \log n \rightarrow O(n^2)$$

$$f(n) \rightarrow n/4 \rightarrow O(n)$$

Q.  $\text{for } i = 0 \rightarrow n \text{ do}$   
|  
|  
|  
y  
 $\left. \begin{array}{c} \\ \\ \\ \end{array} \right\} O(n) \rightarrow \text{print Array}$

Q.  $\text{for } i = 0 \rightarrow n/2 \text{ do}$   
|  
|  
|  
y  
 $\left. \begin{array}{c} \\ \\ \\ \end{array} \right\} O(n/2) = O(n) \rightarrow \text{Swap.}$

Q.  $\text{for } i = 0 \rightarrow n \text{ do}$   
|  
|  
y  
 $\left. \begin{array}{c} \\ \\ \end{array} \right\} O(n) \rightarrow \text{Linear Search.}$

## Q. time complexity?

inf     $a=0, b=0;$

for (int  $i=0; i < n; i++$ ) {  
     $a = a + \text{rand}();$  }     $\left\{ O(n) \right\}$

y

for (int  $j=0; j < m; j++$ ) {  
     $b = b + \text{rand}();$  }     $\left\{ O(m) \right\}$

y

$O(n) + O(m)$   
 $= O(n+m)$

Q.  $\text{for } (0 \rightarrow n) \{ \rightarrow O(n)$

$\text{for } (0 \rightarrow n) \{$   
        :  
     $\} O(n) \} O(n) \times O(n) = O(n^2)$

y

Q  $\text{for } (0 \rightarrow n) \{ \rightarrow O(n)$

$\text{for } (n \rightarrow i) \{ \rightarrow \text{worst case } n \rightarrow 0$   
        :  
     $\} O(n) \} O(n) \times O(n)$   
 $\Rightarrow O(n^2)$

→ We will see recursive time complexities in future video

→ Stuck is TLE

$10^8$  operation rule  $\rightarrow$  most of the modern machine can perform  $10^8$  operation/seconds.

### Time Complexity.

$\leq [10 \dots 11]$	$O(n!)$ , $O(n^6)$
$\leq [15 \dots 16]$	$O(2^n \times n^2)$
$< 100$	$O(n^4)$
$< 400$	$O(n^3)$
$< 2000$	$O(n^2 + \log n)$
$< 10^4$	$O(n^2)$
$< 10^6$	$O(n \log n)$
$< 10^8$	$O(n!)$ , $O(\log n)$

### \* Space Complexity:

$\rightarrow$  if there is fixed size then the Space complexity will be constant.

```
→ for (0-n) {  
    vector<int> v(n);  
    for (0-n) {  
        y  
    }  
}
```

here in every iteration the vector will be created of size  $n$ . so space complexity =  $O(n)$