

NETWORK AND WEB APPLICATION SECURITY LABORATORY EXPERIMENTS WITH SOLUTIONS

LIST OF EXPERIMENTS:

1. Apply DES algorithm for practical applications.
2. Apply AES algorithm for practical applications
3. Implement RSA Algorithm using HTML and JavaScript
4. Implement the Diffie-Hellman Key Exchange algorithm for a given problem.
5. Calculate the message digest of a text using the SHA-1 algorithm
6. Implement the SIGNATURE SCHEME - Digital Signature Standard.
7. Demonstrate intrusion detection system (ids) using any tool eg. Snort or any other s/w.
8. Automated Attack and Penetration Tools Exploring N-Stalker, a Vulnerability Assessment Tool
9. Defeating Malware - Building Trojans, Rootkit Hunter

DESCRIPTION OF MAJOR SOFTWARE USED

JAVA

Java is a high-level programming language originally developed by Sun Microsystems. Java runs on a variety of platforms, such as Windows, Mac OS, and the various versions of UNIX. Java programming were "Simple, Robust, Portable, Platform-independent, Secured, High Performance, Multithreaded, Architecture Neutral, Object-Oriented, Interpreted, and Dynamic". It was originally designed for developing programs for set-top boxes and handheld devices, but later became a popular choice for creating web applications.

Installation requires you to download an executable file available at the manual Java download page, which includes all the files needed for the complete installation at the user's discretion. There is no need to remain connected to the Internet during the installation. The file can also be copied to and installed on another computer that is not connected to the Internet. Administrative permission is required in order to install Java on Microsoft Windows

The Java syntax is similar to C++, but is strictly an object-oriented programming language. For example, most Java programs contain classes, which are used to define objects, and methods, which are assigned to individual classes. Java is also known for being stricter than C++, meaning variables and functions must be explicitly defined. This means Java source code may produce errors or "exceptions" more easily than other languages, but it also limits other types of errors that may be caused by undefined variables or unassigned types.

Unlike Windows executables (.EXE files) or Macintosh applications (.APP files), Java programs are not run directly by the operating system. Instead, Java programs are interpreted by the Java Virtual Machine, or JVM, which runs on multiple platforms. This means all Java programs are multiplatform and can run on different platforms, including Macintosh, Windows, and Unix computers. However, the JVM must be installed for Java applications or applets to run at all. Fortunately, the JVM is included as part of the Java Runtime Environment (JRE), which is available as a free download.

SNORT

Snort really isn't very hard to use, but there are a lot of command line options to play with, and it's not always obvious which ones go together well. This file aims to make using Snort easier for new users. Before we proceed, there are a few basic concepts you should understand about Snort. Snort can be configured to run in three modes:

- Sniffer mode, which simply reads the packets off of the network and displays them for you in a continuous stream on the console (screen).
- Packet Logger mode, which logs the packets to disk.
- Network Intrusion Detection System (NIDS) mode, which performs detection and analysis on network traffic. This is the most complex and configurable mode.

Snort is based on libpcap (for library packet capture), a **tool** that is widely **used** in TCP/IP traffic sniffers and analyzers. Through protocol analysis and content searching and matching, **Snort** detects attack methods, including denial of service, buffer overflow, CGI attacks, stealth port scans, and SMB probes.

N-STALKER

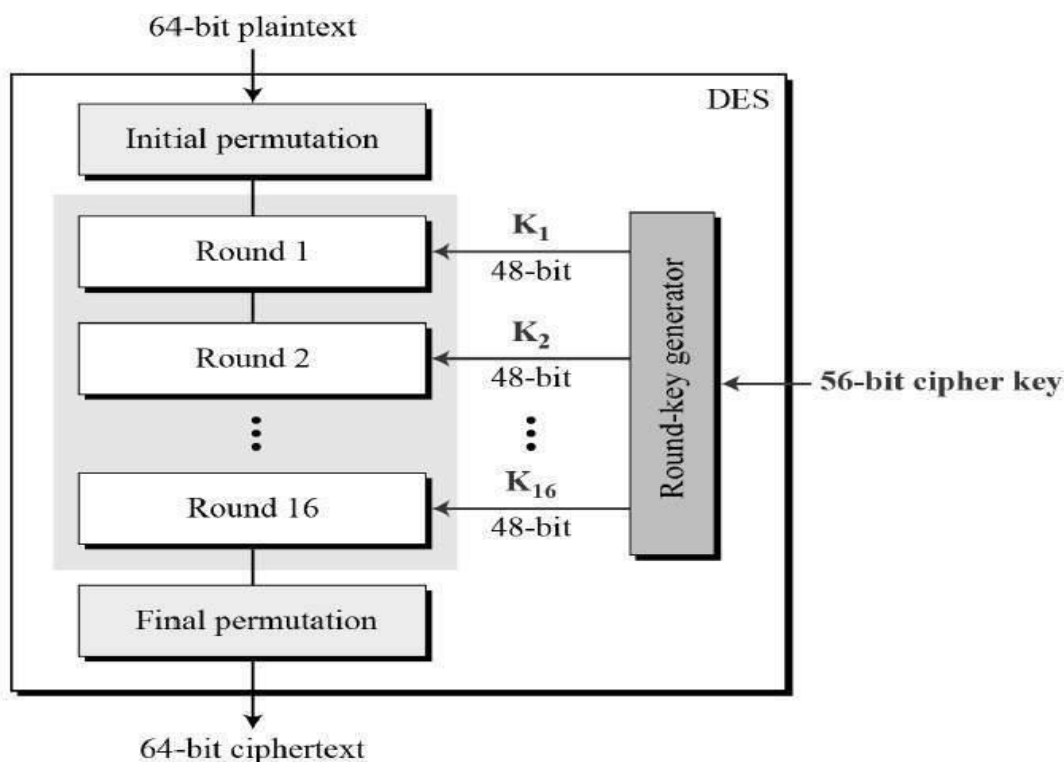
N-STALKER is a world leader in Web Application Security solutions since 2000. It has started providing the first commercial and most complete HTTP Security Scanner, holding the largest signatures database available in the market – more than 39,000 attack signatures. Our products are delivered to hundreds of customers distributed in more than 30 different countries around the world. Back in 2000, N-Stalker's challenge was to provide complete solutions for your Web server infrastructure, which ended up with the release of N-Stealth HTTP Security Scanner. Nowadays, N- Stalker is seeking to provide the most complete solution for your enterprise web applications, the N-Stalker Web Application Security Scanner Suite.

AIM:

To develop a program to implement Data Encryption Standard for encryption and decryption.

PRELAB DISCUSSION:

- The Data Encryption Standard (DES) is a symmetric-key block cipher published by the National Institute of Standards and Technology (NIST).
- DES is an implementation of a Feistel Cipher. It uses 16 round Feistel structure. The block size is 64-bit.
- Though, key length is 64-bit, DES has an effective key length of 56 bits, since 8 of the 64 bits of the key are not used by the encryption algorithm (function as check bits only).
- General Structure of DES is depicted in the following illustration

**ALGORITHM:**

1. Process the key.
 - i. Get a 64-bit key from the user.
 - ii. Calculate the key schedule.
 1. Perform the following permutation on the 64-bit key. The parity bits are discarded, reducing the key to 56 bits. Bit 1 of the permuted block is bit 57 of the original key, bit 2 is bit 49, and so on with bit 56 being bit 4 of the original key.
 2. Split the permuted key into two halves. The first 28 bits are called C[0] and the last 28 bits are called D[0].

3. Calculate the 16 subkeys. Start with $i = 1$.

1. Perform one or two circular left shifts on both $C[i-1]$ and $D[i-1]$ to get $C[i]$ and $D[i]$, respectively. The number of shifts per iteration are given in the table below.

Iteration # 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16

Left Shifts 1 1 2 2 2 2 2 2 1 2 2 2 2 2 2 1

2. Permute the concatenation $C[i]D[i]$ as indicated below. This will yield $K[i]$, which is 48 bits long.
3. Loop back to 1.ii.c.1 until $K[16]$ has been calculated.

2 Process a 64-bit data block.

- i. Get a 64-bit data block. If the block is shorter than 64 bits, it should be padded as appropriate for the application.
- ii. Perform the initial permutation on the data block.
- iii. Split the block into two halves. The first 32 bits are called $L[0]$, and the last 32 bits are called $R[0]$.
- iv. Apply the 16 subkeys to the data block. Start with $i = 1$.
 - a. Expand the 32-bit $R[i-1]$ into 48 bits according to the bit-selection function Expansion (E)
 - b. Exclusive-or $E(R[i-1])$ with $K[i]$.
 - c. Break $E(R[i-1]) \text{ xor } K[i]$ into eight 6-bit blocks. Bits 1-6 are $B[1]$, bits 7-12 are $B[2]$, and so on with bits 43-48 being $B[8]$.
 - d. Substitute the values found in the S-boxes for all $B[j]$. Start with $j = 1$. All values in the S-boxes should be considered 4 bits wide.
 - i. Take the 1st and 6th bits of $B[j]$ together as a 2-bit value (call it m) indicating the row in $S[j]$ to look in for the substitution.
 - ii. Take the 2nd through 5th bits of $B[j]$ together as a 4-bit value (call it n) indicating the column in $S[j]$ to find the substitution.
 - iii. Replace $B[j]$ with $S[j][m][n]$.
 - iv. Loop back to 2.iv.d.i until all 8 blocks have been replaced.
 - e. Permute the concatenation of $B[1]$ through $B[8]$
 - f. Exclusive-or the resulting value with $L[i-1]$. Thus, all together, your $R[i] = L[i-1] \text{ xor } P(S[1](B[1])...S[8](B[8]))$, where $B[j]$ is a 6-bit block of $E(R[i-1]) \text{ xor } K[i]$. (The function for $R[i]$ is written as, $R[i] = L[i-1] \text{ xor } f(R[i-1], K[i])$.)
 - g. $L[i] = R[i-1]$.
 - h. Loop back to 2.iv.a until $K[16]$ has been applied.
- v. Perform the final permutation on the block $R[16]L[16]$.

3. Decryption : Use the keys $K[i]$ in reverse order. That is, instead of applying $K[1]$ for the first iteration, apply $K[16]$, and then $K[15]$ for the second, on down to $K[1]$

PROGRAM:

DES :-

```
import javax.swing.*;
import java.security.SecureRandom;
```

```

import javax.crypto.Cipher;
import javax.crypto.KeyGenerator;
import javax.crypto.SecretKey;
import javax.crypto.spec.SecretKeySpec;
import java.util.Random ;
class DES {
byte[] skey = new byte[1000];
String skeyString;
static byte[] raw;
String inputMessage,encryptedData,decryptedMessage;
public DES() {
try {
generateSymmetricKey();
inputMessage=JOptionPane.showInputDialog(null,"Enter message to encrypt");
byte[] ibyte = inputMessage.getBytes();
byte[] ebyte=encrypt(raw, ibyte);
String encryptedData = new String(ebyte);
System.out.println("Encrypted message "+encryptedData);
JOptionPane.showMessageDialog(null,"Encrypted Data "+"\\n"+encryptedData);
byte[] dbyte= decrypt(raw,ebyte);
String decryptedMessage = new String(dbyte);
System.out.println("Decrypted message "+decryptedMessage);
JOptionPane.showMessageDialog(null,"Decrypted Data "+"\\n"+decryptedMessage);
}
catch(Exception e) {
System.out.println(e);
}
}
void generateSymmetricKey() {
try {
Random r = new Random();
intnum = r.nextInt(10000);
String knum = String.valueOf(num);
byte[] knumb = knum.getBytes();
skey=getRawKey(knumb);
skeyString = new String(skey);
System.out.println("DES Symmetric key = "+skeyString);
}
catch(Exception e) {
System.out.println(e);
}
}
private static byte[] getRawKey(byte[] seed) throws Exception {
KeyGeneratorkgen = KeyGenerator.getInstance("DES");
SecureRandomsr= SecureRandom.getInstance("SHA1PRNG");
sr.setSeed(seed);

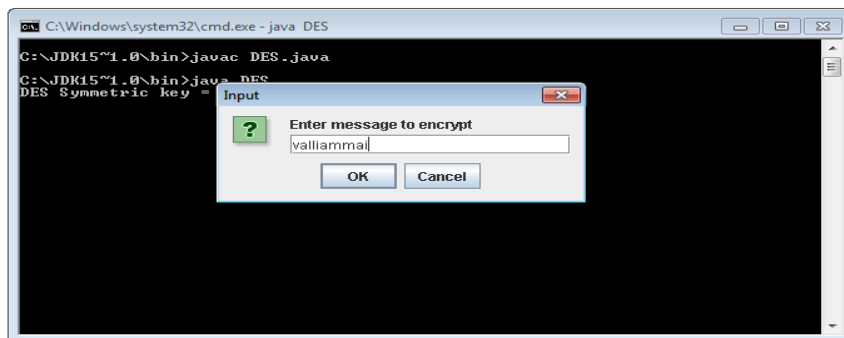
```

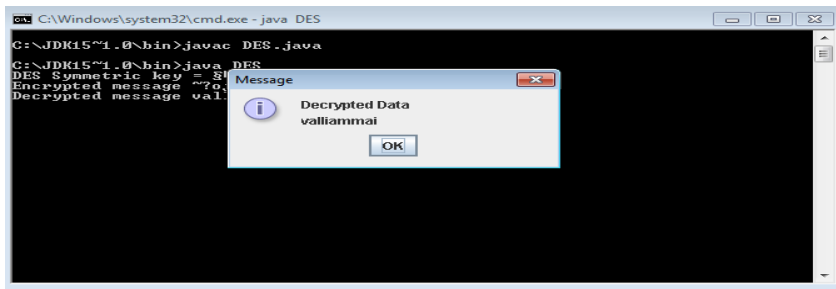
```

kgen.init(56, sr);
SecretKeyskey = kgen.generateKey();
raw = skey.getEncoded();
return raw;
}
private static byte[] encrypt(byte[] raw, byte[] clear) throws Exception {
SecretKeySpeckeySpec = new SecretKeySpec(raw, "DES");
Cipher cipher = Cipher.getInstance("DES");
cipher.init(Cipher.ENCRYPT_MODE, skeySpec);
byte[] encrypted = cipher.doFinal(clear);
return encrypted;
}
private static byte[] decrypt(byte[] raw, byte[] encrypted) throws Exception {
SecretKeySpeckeySpec = new SecretKeySpec(raw, "DES");
Cipher cipher = Cipher.getInstance("DES");
cipher.init(Cipher.DECRYPT_MODE, skeySpec);
byte[] decrypted = cipher.doFinal(encrypted);
return decrypted;
}
public static void main(String args[]) {
DES des = new DES();
}
}

```

OUTPUT:





```
C:\Windows\system32\cmd.exe - java DES
C:\JDK15\1.0\bin>javac DES.java
C:\JDK15\1.0\bin>java DES
DES Symmetric key = 81
Encrypted message ~79
Decrypted message val3
Decrypted Data
valliannmai
OK
```

VIVA QUESTIONS (PRELAB and POSTLAB):

1. DES follows which basic stream cipher?
2. The DES Algorithm Cipher System consists of how many rounds (iterations) each with a round key?
3. What is the key length of the DES algorithm?
4. In the DES algorithm, although the key size is 64 bits only 48bits are used for the encryption procedure, the rest are parity bits. Is it true or false?
5. In the DES algorithm, what is the size of the round key and the Round Input?
6. In the DES algorithm how the Round Input is expanded to 48 ?
7. What is size of the Initial Permutation table/matrix
8. How many unique substitution boxes are in DES after the 48 bit XOR operation?
9. In the DES algorithm the 64 bit key input is shortened to 56 bits by ignoring every 4th bit. Is it true or false?

RESULT:

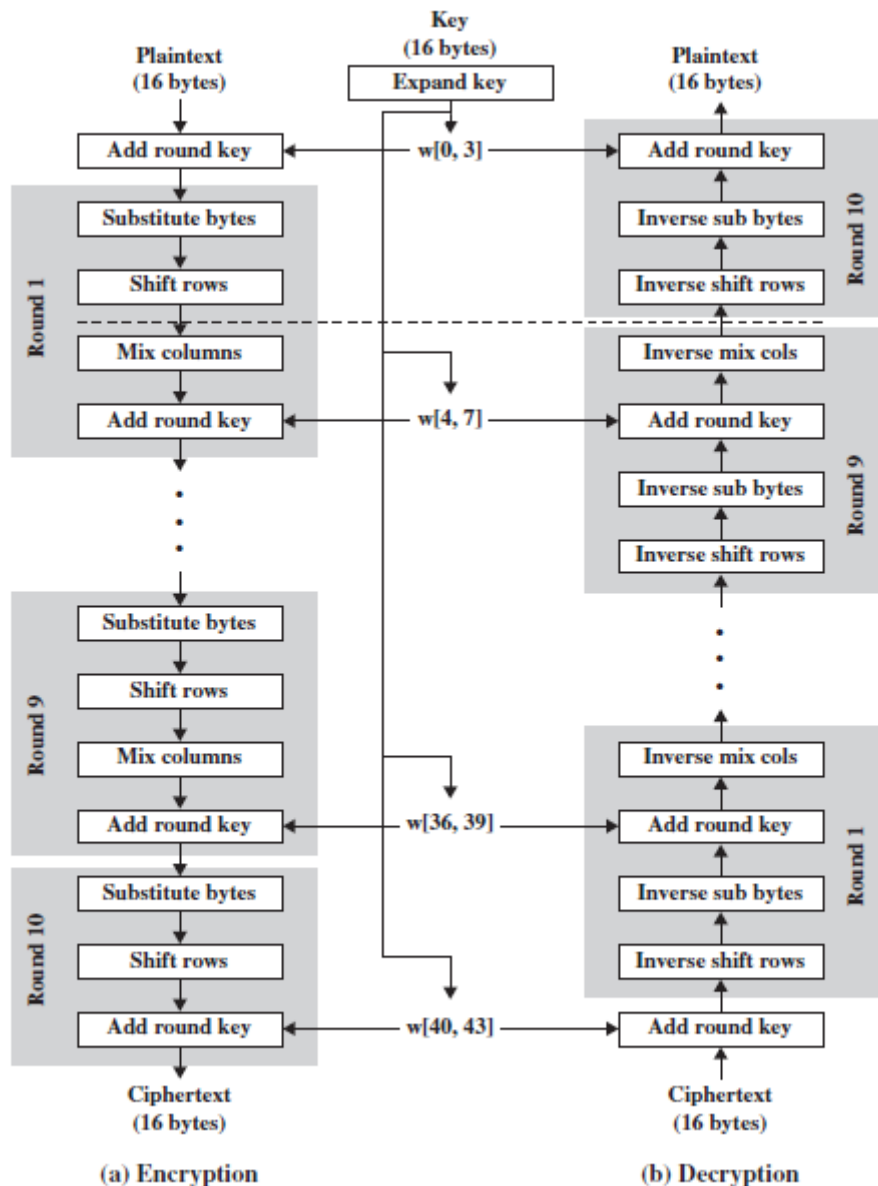
Thus the program to implement DES encryption technique was developed and executed successfully

AIM:

To develop a program to implement Advanced Encryption Standard for encryption and decryption.

PRELAB DISCUSSION:

The cipher takes a plaintext block size of 128 bits, or 16 bytes. The key length can be 16, 24, or 32 bytes (128, 192, or 256 bits). The algorithm is referred to as AES-128, AES-192, or AES-256, depending on the key length.



The input to the encryption and decryption algorithms is a single 128-bit block. In FIPS PUB 197, this block is depicted as a 4×4 square matrix of bytes. This block is copied into the **State** array, which is modified at each stage of encryption or decryption. After the final stage, **State** is copied to an output matrix. Similarly, the key is depicted as a square matrix of bytes. This key is then expanded

into an array of key schedule words. Each word is four bytes, and the total key schedule is 44 words for the 128-bit key. Note that the ordering of bytes within a matrix is by column. So, for example, the first four bytes of a 128-bit plaintext input to the encryption cipher occupy the first column of the **in** matrix, the second four bytes occupy the second column, and so on. Similarly, the first four bytes of the expanded key, which form a word, occupy the first column of the **w** matrix.

The cipher consists of N rounds, where the number of rounds depends on the key length: 10 rounds for a 16-byte key, 12 rounds for a 24-byte key, and 14 rounds for a 32-byte key. The first $N - 1$ rounds consist of four distinct transformation functions: SubBytes, ShiftRows, MixColumns, and AddRoundKey. The final round contains only three transformations, and there is a initial single transformation (AddRoundKey) before the first round, which can be considered Round 0. Each transformation takes one or more $4 * 4$ matrices

PROGRAM:

```
package com.includehelp.stringsample;
```

```
import java.util.Base64;
```

```
import java.util.Scanner;
```

```
import javax.crypto.Cipher;
```

```
import javax.crypto.spec.IvParameterSpec;
```

```
import javax.crypto.spec.SecretKeySpec;
```

```
/**
```

```
 * Program to Encrypt/Decrypt String Using AES 128 bit Encryption Algorithm
```

```
 */
```

```
public class EncryptDecryptString
```

```
{
```

```
    private static final String encryptionKey      = "ABCDEFGHJKLMNOP";
```

```
    private static final String characterEncoding   = "UTF-8";
```

```
    private static final String cipherTransformation = "AES/CBC/PKCS5PADDING";
```

```
    private static final String aesEncryptionAlgorithem = "AES";
```

```
/**
```

```
 * Method for Encrypt Plain String Data
```

```
 * @param plainText
```

```
 * @return encryptedText
```

```
 */
```

```
    public static String encrypt(String plainText)
```

```
    {
```

```
        String encryptedText = "";
```

```
        try
```

```
        {
```

```
            Cipher cipher = Cipher.getInstance(cipherTransformation);
```

```
            byte[] key      = encryptionKey.getBytes(characterEncoding);
```

```
            SecretKeySpec secretKey = new SecretKeySpec(key, aesEncryptionAlgorithem);
```

```
            IvParameterSpec ivparameterspec = new IvParameterSpec(key);
```

```
            cipher.init(Cipher.ENCRYPT_MODE, secretKey, ivparameterspec);
```

```
            byte[] cipherText = cipher.doFinal(plainText.getBytes("UTF8"));
```

```

        Base64.Encoder encoder = Base64.getEncoder();
        encryptedText = encoder.encodeToString(cipherText);
    } catch (Exception E)
    {
        System.err.println("Encrypt Exception : "+E.getMessage());
    }
    return encryptedText;
}

/**

* Method For Get encryptedText and Decrypted provided String
* @param encryptedText
* @return decryptedText
*/
public static String decrypt(String encryptedText)
{
    String decryptedText = "";
    try
    {
        Cipher cipher = Cipher.getInstance(cipherTransformation);
        byte[] key = encryptionKey.getBytes(characterEncoding);
        SecretKeySpec secretKey = new SecretKeySpec(key, aesEncryptionAlgorithm);
        IvParameterSpec ivparameterspec = new IvParameterSpec(key);
        cipher.init(Cipher.DECRYPT_MODE, secretKey, ivparameterspec);
        Base64.Decoder decoder = Base64.getDecoder();
        byte[] cipherText = decoder.decode(encryptedText.getBytes("UTF8"));
        decryptedText = new String(cipher.doFinal(cipherText), "UTF-8");
    } catch (Exception E)
    {
        System.err.println("decrypt Exception : "+E.getMessage());
    }
    return decryptedText;
}

public static void main(String[] args)
{
    Scanner sc = new Scanner(System.in);
    System.out.println("Enter String : ");
    String plainString = sc.nextLine();

    String encryptStr = encrypt(plainString);
    String decryptStr = decrypt(encryptStr);

    System.out.println("Plain String : "+plainString);
    System.out.println("Encrypt String : "+encryptStr);
    System.out.println("Decrypt String : "+decryptStr);
} }

```

OUTPUT:

Enter String : Hello World

Plain String : Hello World

Encrypt String : IMfL/ifkuvkZwG/v2bn6Bw==

Decrypt String : Hello World

VIVA QUESTIONS (PRELAB and POSTLAB):

This set of Cryptography Multiple Choice Questions & Answers (MCQs) focuses on “The Data Encryption Standard (DES) and It’s Strength”.

1. DES follows

a) Hash Algorithm b) Caesars Cipher c) Feistel Cipher Structure d) SP Networks

2. The DES Algorithm Cipher System consists of _____ rounds (iterations) each with a round key

a) 12 b) 18 c) 9 d) 16

3. The DES algorithm has a key length of

a) 128 Bits b) 32 Bits c) 64 Bits d) 16 Bits

4. In the DES algorithm, although the key size is 64 bits only 48bits are used for the encryption procedure, the rest are parity bits.

a) True b) False

5. In the DES algorithm the round key is _____ bit and the Round Input is _____ bits.

a) 48, 32 b) 64,32 c) 56, 24 d) 32, 32

6. In the DES algorithm the Round Input is 32 bits, which is expanded to 48 bits via _____

a) Scaling of the existing bits b) Duplication of the existing bits
c) Addition of zeros d) Addition of ones

7. The Initial Permutation table/matrix is of size

a) 16×8 b) 12×8 c) 8×8 d) 4×8

8. The number of unique substitution boxes in DES after the 48 bit XOR operation are

a) 8 b) 4 c) 6 d) 12

9. In the DES algorithm the 64 bit key input is shortened to 56 bits by ignoring every 4th bit.

a) True b) False

RESULT:

Thus the program to implement AES encryption technique was developed and executed successfully.

AIM:

Develop a program to implement RSA algorithm for encryption and decryption. This cryptosystem is one the initial system. It remains most employed cryptosystem even today. The system was invented by three scholars **Ron Rivest, Adi Shamir, and Len Adleman** and hence, it is termed as RSA cryptosystem. The two aspects of the RSA cryptosystem, firstly generation of key pair and secondly encryption-decryption algorithms

PRELAB DISCUSSION:**Generation of RSA Key Pair**

- Each person or a party who desires to participate in communication using encryption needs to generate a pair of keys, namely public key and private key.
- The process followed in the generation of keys is described below –
- Generate the RSA modulus (n)
 - Select two large primes, p and q.
 - Calculate $n=p*q$. For strong unbreakable encryption, let n be a large number, typically a minimum of 512 bits.
- Find Derived Number (e)
 - Number e must be greater than 1 and less than $(p - 1)(q - 1)$.
 - There must be no common factor for e and $(p - 1)(q - 1)$ except for 1. In other words two numbers e and $(p - 1)(q - 1)$ are coprime.
- Form the public key
 - The pair of numbers (n, e) form the RSA public key and is made public.
 - Interestingly, though n is part of the public key, difficulty in factorizing a large prime number ensures that attacker cannot find in finite time the two primes (p & q) used to obtain n. This is strength of RSA.
- Generate the private key
 - Private Key d is calculated from p, q, and e. For given n and e, there is unique number d.
 - Number d is the inverse of e modulo $(p - 1)(q - 1)$. This means that d is the number less than $(p - 1)(q - 1)$ such that when multiplied by e, it is equal to 1 modulo $(p - 1)(q - 1)$.
- This relationship is written mathematically as follows $ed = 1 \text{ mod } (p - 1)(q - 1)$
- The Extended Euclidean Algorithm takes p, q, and e as input and gives d as output.

ALGORITHM:**1. Key Generation**

- i. Choose two distinct prime numbers p and q.
- ii. Find n such that $n = pq$, n will be used as the modulus for both the public and private keys.
- iii. Find the totient of n, $\phi(n)$ $\phi(n)=(p-1)(q-1)$
- iv. Choose an e such that $1 < e < \phi(n)$, and such that e and $\phi(n)$ share no divisors other than 1 (e and $\phi(n)$ are relatively prime). e is kept as the public

key exponent

- v. Determine d (using modular arithmetic) which satisfies the congruence relation

$$de \equiv 1 \pmod{\phi(n)}.$$

The public key has modulus n and the public (or encryption) exponent

- e. The private key has modulus n and the private (or decryption) exponent d , which is kept secret.

2. Encryption

$$c \equiv m^e \pmod{n}.$$

3. Decryption:

$$m \equiv c^d \pmod{n}.$$

4. Stop.

PROGRAM:

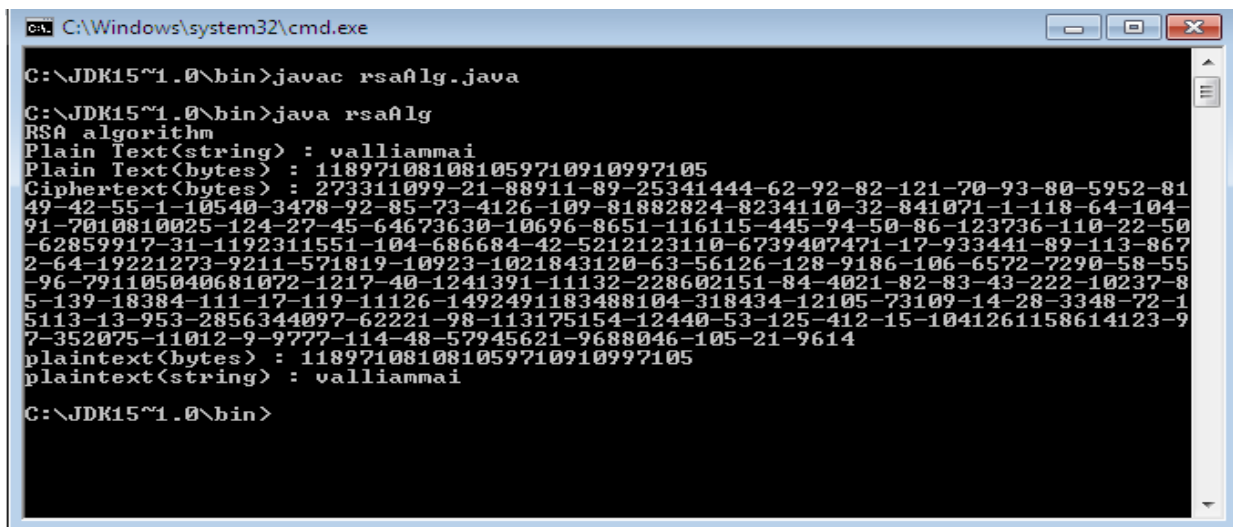
```
import java.math.BigInteger;
import java.util.Random;
import java.io.*;
class rsaAlg
{
private BigInteger p, q, n, phi, e, d; /* public key components */
private int bitLen = 1024;
private int blkSz = 256; /* block size in bytes */
private Random rand;
/* convert bytes to string */
private static String bytesToString(byte[] encrypted)
{
String str = "";
for (byte b : encrypted)
{
str += Byte.toString(b);
}
return str;
}
/* encrypt message */
public byte[] encrypt(byte[] msg)
{
return (new BigInteger(msg)).modPow(e, n).toByteArray();
}
/* decrypt message */
public byte[] decrypt(byte[] msg)
{
return (new BigInteger(msg)).modPow(d, n).toByteArray();
}
/* calculate public key components p, q, n, phi, e, d */
public rsaAlg()
{
```

```

rand = new Random();
p = BigInteger.probablePrime(bitLen, rand);
q = BigInteger.probablePrime(bitLen, rand);
n = p.multiply(q);
phi = p.subtract(BigInteger.ONE).multiply(q.subtract(BigInteger.ONE));
e = BigInteger.probablePrime(bitLen/2, rand);
while (phi.gcd(e).compareTo(BigInteger.ONE) > 0 &&
e.compareTo(phi) < 0)
{
e.add(BigInteger.ONE);
}
d = e.modInverse(phi);
}
public rsaAlg (BigInteger e, BigInteger d, BigInteger n)
{
this.e = e;
this.d = d;
this.n = n;
}
public static void main (String[] args) throws java.lang.Exception
{
rsaAlg rsaObj = new rsaAlg();
String msg = "Hello world! Security Laboratory";
System.out.println("simulation of RSA algorithm");
System.out.println("message(string) : " + msg);
System.out.println("message(bytes) : " +
bytesToString(msg.getBytes()));
/* encrypt test message */
byte[] ciphertext = rsaObj.encrypt(msg.getBytes());
System.out.println("ciphertext(bytes) : " + bytesToString(ciphertext));
/* decrypt ciphertext */
byte[] plaintext = rsaObj.decrypt(ciphertext);
System.out.println("plaintext(bytes) : " + bytesToString(plaintext));
System.out.println("plaintext(string) : " + new String(plaintext)); } }

```

Output:



```
C:\Windows\system32\cmd.exe

C:\JDK15~1.0\bin>javac rsaAlg.java
C:\JDK15~1.0\bin>java rsaAlg
RSA algorithm
Plain Text(string) : valliammai
Plain Text(bytes) : 118971081081059710910997105
Ciphertext(bytes) : 273311099-21-88911-89-25341444-62-92-82-121-70-93-80-5952-81
49-42-55-1-10540-3478-92-85-73-4126-109-81882824-8234110-32-841071-1-118-64-104-
91-7010810025-124-27-45-64673630-10696-8651-116115-445-94-50-86-123736-110-22-50
-62859917-31-1192311551-104-686684-42-5212123110-6739407471-17-933441-89-113-867
2-64-19221273-9211-571819-10923-1021843120-63-56126-128-9186-106-6572-7290-58-55
-96-791105040681072-1217-40-1241391-11132-228602151-84-4021-82-83-43-222-10237-8
5-139-18384-111-17-119-11126-1492491183488104-318434-12105-73109-14-28-3348-72-1
5113-13-953-2856344097-62221-98-113175154-12440-53-125-412-15-1041261158614123-9
7-352075-11012-9-9777-114-48-57945621-9688046-105-21-9614
plaintext(bytes) : 118971081081059710910997105
plaintext(string) : valliammai
C:\JDK15~1.0\bin>
```

VIVA QUESTIONS (PRELAB and POSTLAB):

1. What Is RSA?
2. Are Strong Primes Necessary In RSA?
3. Can Users Of RSA Run Out Of Distinct Primes?
4. What Are The Alternatives To RSA?
5. How fast is RSA?
6. What would it take to break RSA?
7. How large a modulus key should be used in RSA?
8. How large should the primes be?
9. How do you know if a number is prime?
10. How is RSA used for authentication in practice?
11. What do you mean by Secret Key Cryptography and Public Key Cryptography? How they are different from one another?
12. What exactly do you know about RSA?

RESULT:

Thus the program for implementation of RSA algorithm was executed and verified successfully.

AIM:

Develop a program to implement Diffie Hellman Key Exchange Algorithm for encryption and Decryption.

PRELAB DISCUSSION:

Diffie–Hellman key exchange (D–H) is a specific method of securely exchanging cryptographic keys over a public channel and was one of the first public-key protocols. The Diffie–Hellman key exchange method allows two parties that have no prior knowledge of each other to jointly establish a shared secret key over an insecure channel. This key can then be used to encrypt subsequent communications using a symmetric key cipher. This algorithm uses arithmetic modulus as the basis of its calculation. Suppose Alice and Bob follow this key exchange procedure with Eve acting as a man in middle interceptor (or the bad guy).

Here are the calculation steps followed in this algorithm that make sure that Eve never gets to know the final keys through which actual encryption of data takes place. First, both Alice and Bob agree upon a prime number and another number that has no factor in common. Let's call the prime number as **p** and the other number as **g**. Note that **g** is also known as the generator and **p** is known as prime modulus. Now, since Eve is sitting in between and listening to this communication so Eve also gets to know **p** and **g**. Now, the modulus arithmetic says that $r = (g \text{ to the power } x) \bmod p$. So **r** will always produce an integer between 0 and **p**. The first trick here is that given **x** (with **g** and **p** known), it's very easy to find **r**. But given **r** (with **g** and **p** known) it's difficult to deduce **x**. One may argue that this is not that difficult to crack but what if the value of **p** is a very huge prime number? Well, if this is the case then deducing **x** (if **r** is given) becomes almost next to impossible as it would take thousands of years to crack this even with supercomputers. This is also called the discrete logarithmic problem. Coming back to the communication, all the three Bob, Alice and Eve now know **g** and **p**. Now, Alice selects a random private number **xa** and calculates $(g \text{ to the power } xa) \bmod p = ra$.

This

resultant **ra** is sent on the communication channel to Bob. Intercepting in between, Eve also comes to know **ra**. Similarly Bob selects his own random private number **xb**, calculates $(g \text{ to the power } xb) \bmod p = rb$ and sends this **rb** to Alice through the same communication channel. Obviously Eve also comes to know about **rb**. So Eve now has information about **g**, **p**, **ra** and **rb**. Now comes the heart of this algorithm. Alice calculates $(rb \text{ to the power } xa) \bmod p = \text{Final key}$ which is equivalent to $(g \text{ to the power } (xa * xb)) \bmod p$. Similarly Bob calculates $(ra \text{ to the power } xb) \bmod p = \text{Final key}$ which is again equivalent to $(g \text{ to the power } (xb * xa)) \bmod p$. So both Alice and Bob were able to calculate a common **Final key** without sharing each other's private random number and Eve sitting in between will not be able to determine the **Final key** as the private numbers were never transferred.

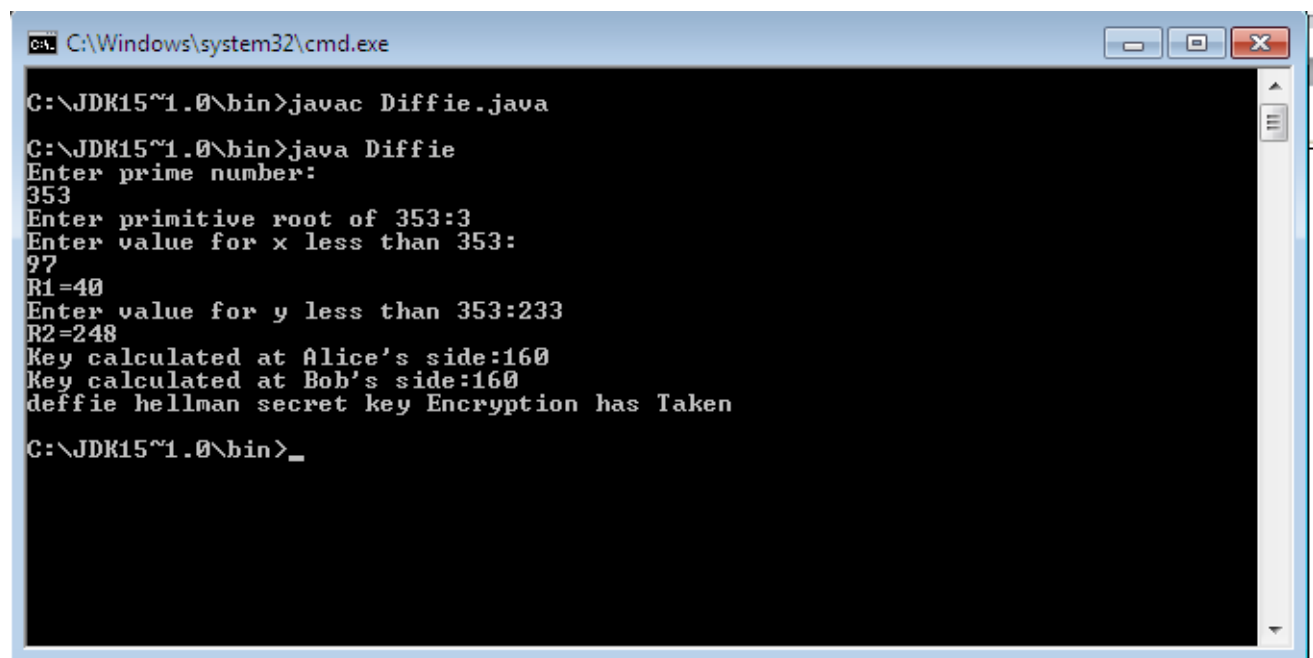
ALGORITHM

1. Global Public Elements:
Let q be a prime number and α where $\alpha < q$ and α is a primitive root of q .
2. User A Key Generation:
 - i. Select private X_A where $X_A < q$
 - ii. Calculate public Y_A where $Y_A = \alpha^{X_A} \bmod q$
3. User B Key Generation:
 - i. Select private X_B where $X_B < q$
 - ii. Calculate public Y_B where $Y_B = \alpha^{X_B} \bmod q$
4. Calculation of Secret Key by User
A: $K = (Y_B)^{X_A} \bmod q$
5. Calculation of Secret Key by User
B: $K = (Y_A)^{X_B} \bmod q$

PROGRAM:

```
import java.io.*;
import java.math.BigInteger;
class Diffie
{
    public static void main(String[] args) throws IOException
    {
        BufferedReader br = new BufferedReader(new InputStreamReader(System.in));
        System.out.println("Enter prime number:");
        BigInteger p = new BigInteger(br.readLine());
        System.out.print("Enter primitive root of "+p+":");
        BigInteger g = new BigInteger(br.readLine());
        System.out.println("Enter value for x less than "+p+":");
        BigInteger x = new BigInteger(br.readLine());
        BigInteger R1 = g.modPow(x, p);
        System.out.println("R1="+R1);
        System.out.print("Enter value for y less than "+p+":");
        BigInteger y = new BigInteger(br.readLine());
        BigInteger R2 = g.modPow(y, p);
        System.out.println("R2="+R2);
        BigInteger k1 = R2.modPow(x, p);
        System.out.println("Key calculated at Sender's side: "+k1);
        BigInteger k2 = R1.modPow(y, p);
        System.out.println("Key calculated at Receiver's side: "+k2);
        System.out.println("diffie hellman secret key Encryption has Taken");
    }
}
```

OUTPUT:



```
C:\Windows\system32\cmd.exe

C:\JDK15~1.0\bin>javac Diffie.java
C:\JDK15~1.0\bin>java Diffie
Enter prime number:
353
Enter primitive root of 353:3
Enter value for x less than 353:
97
R1=40
Enter value for y less than 353:233
R2=248
Key calculated at Alice's side:160
Key calculated at Bob's side:160
deffie hellman secret key Encryption has Taken

C:\JDK15~1.0\bin>_
```

VIVA QUESTIONS (PRELAB and POSTLAB):

1. What's the difference between Diffie-Hellman and RSA?
2. Does Diffie Hellman guarantee secrecy?
3. Why is RSA preferred over Diffie-Hellman if they are both used to establish shared key?
4. Are there any one way operations that could be used for Diffie-Hellman post quantum?
5. Why is Diffie-Hellman required when RSA is already used for key exchange in TLS?
6. What is Authenticated Diffie-Hellman Key Agreement?
7. How secure is ECDH if the public keys are never shared?
8. Which is better when the secret is leaked, RSA or Diffie-Hellman?
9. What role does RSA play in DH-RSA cipher suite?
10. Why is Diffie-Hellman used alongside public keys?
11. Is Diffie-Hellman key exchange based on one-way function or trapdoor function?

RESULT:

Thus the program to implement Diffie-Hellman Key Exchange algorithm was developed and executed successfully

AIM:

Develop a program to implement Secure Hash Algorithm (SHA-1)

PRELAB DISCUSSION:

In cryptography, SHA-1 (Secure Hash Algorithm 1) is a cryptographic hash function which takes an input and produces a 160-bit (20-byte) hash value known as a message digest - typically rendered as a hexadecimal number, 40 digits long.

Secure Hashing Algorithms, also known as SHA, are a family of cryptographic functions designed to keep data secured. It works by transforming the data using a hash function: an algorithm that consists of bitwise operations, modular additions, and compression functions. The hash function then produces a fixed size string that looks nothing like the original. These algorithms are designed to be one-way functions, meaning that once they're transformed into their respective hash values, it's virtually impossible to transform them back into the original data. A few algorithms of interest are SHA-1, SHA-2, and SHA-5, each of which was successively designed with increasingly stronger encryption in response to hacker attacks. SHA-0, for instance, is now obsolete due to the widely exposed vulnerabilities.

A common application of SHA is to encrypting passwords, as the server side only needs to keep track of specific user's hash value, rather than the actual password. This is helpful in case an attacker hacks the database, as they will only find the hashed functions and not the actual passwords, so if they were to input the hashed value as a password, the hash function will convert it into another string and subsequently deny access. Additionally, SHA exhibit the avalanche effect, where the modification of very few letters being encrypted cause a big change in output; or conversely, drastically different strings produce similar hash values. This effect causes hash values to not give any information regarding the input string, such as its original length. In addition, SHAs are also used to detect the tampering of data by attackers, where if a text file is slightly changed and barely noticeable, the modified file's hash value will be different than the original file's hash value, and the tampering will be rather noticeable.

ALGORITHM:

1. Append Padding Bits: Message is "padded" with a 1 and as many 0's as necessary to bring the message length to 64 bits less than an even multiple of 512.
2. Append Length: 64 bits are appended to the end of the padded message. These bits hold the binary format of 64 bits indicating the length of the original message.
3. Prepare Processing Functions: SHA1 requires 80 processing functions defined as:

$$f(t;B,C,D) = (B \text{ AND } C) \text{ OR } ((\text{NOT } B) \text{ AND } D) \quad (0 \leq t \leq 19)$$

$$f(t;B,C,D) = B \text{ XOR } C \text{ XOR } D \quad (20 \leq t \leq 39)$$

$$f(t;B,C,D) = (B \text{ AND } C) \text{ OR } (B \text{ AND } D) \text{ OR } (C \text{ AND } D) \quad (40 \leq t \leq 59)$$

$$f(t;B,C,D) = B \text{ XOR } C \text{ XOR } D \quad (60 \leq t \leq 79)$$

4. Prepare Processing Constants: SHA1 requires 80 processing constant words defined as:

$$K(t) = 0x5A827999 \quad (0 \leq t \leq 19)$$

$$K(t) = 0x6ED9EBA1 \quad (20 \leq t \leq 39)$$

$$K(t) = 0x8F1BBCDC \quad (40 \leq t \leq 59)$$

$$K(t) = 0xCA62C1D6 \quad (60 \leq t \leq 79)$$
5. Initialize Buffers: SHA1 requires 160 bits or 5 buffers of words (32 bits):

$$H0 = 0x67452301 \quad H1 = 0xEFCDAB89$$

$$H2 = 0x98BADCFE \quad H3 = 0x10325476$$

$$H4 = 0xC3D2E1F0$$
6. Processing Message in 512-bit blocks (L blocks in total message)
 - i. This is the main task of SHA1 algorithm which loops through the padded and appended message in 512-bit blocks.
 - ii. Input and predefined functions: $M[1, 2, \dots, L]$: Blocks of the padded and appended message $f(0;B,C,D)$, $f(1;B,C,D)$, ..., $f(79;B,C,D)$: 80 Processing Functions

$$K(0), K(1), \dots, K(79)$$
: 80 Processing Constant Words

$$H0, H1, H2, H3, H4, H5$$
: 5 Word buffers with initial values
7. For loop on $k = 1$ to L
 1. $(W(0), W(1), \dots, W(15)) = M[k]$ /* Divide $M[k]$ into 16 words */
8. For $t = 16$ to 79 do:

$$W(t) = (W(t-3) \text{ XOR } W(t-8) \text{ XOR } W(t-14) \text{ XOR } W(t-16)) \lll 1$$

$$A = H0, B = H1, C = H2, D = H3, E = H4$$
 For $t = 0$ to 79 do:

$$\text{TEMP} = A \lll 5 + f(t;B,C,D) + E + W(t) + K(t)$$

$$E = D, D = C, C = B \lll 30, B = A, A = \text{TEMP}$$
 End of for loop

$$H0 = H0 + A, H1 = H1 + B, H2 = H2 + C, H3 = H3 + D, H4 = H4 + E$$
 End of for loop

PROGRAM

```
import java.security.*;
public class SHA1 {
public static void main(String[] a) {
try {
MessageDigest md = MessageDigest.getInstance("SHA1");
String input = "srm";
md.update(input.getBytes());
byte[] output = md.digest();
System.out.println();
System.out.println("SHA1(\""+input+"\") = " + bytesToHex(output));
input = "vec";
md.update(input.getBytes());
output = md.digest();
```

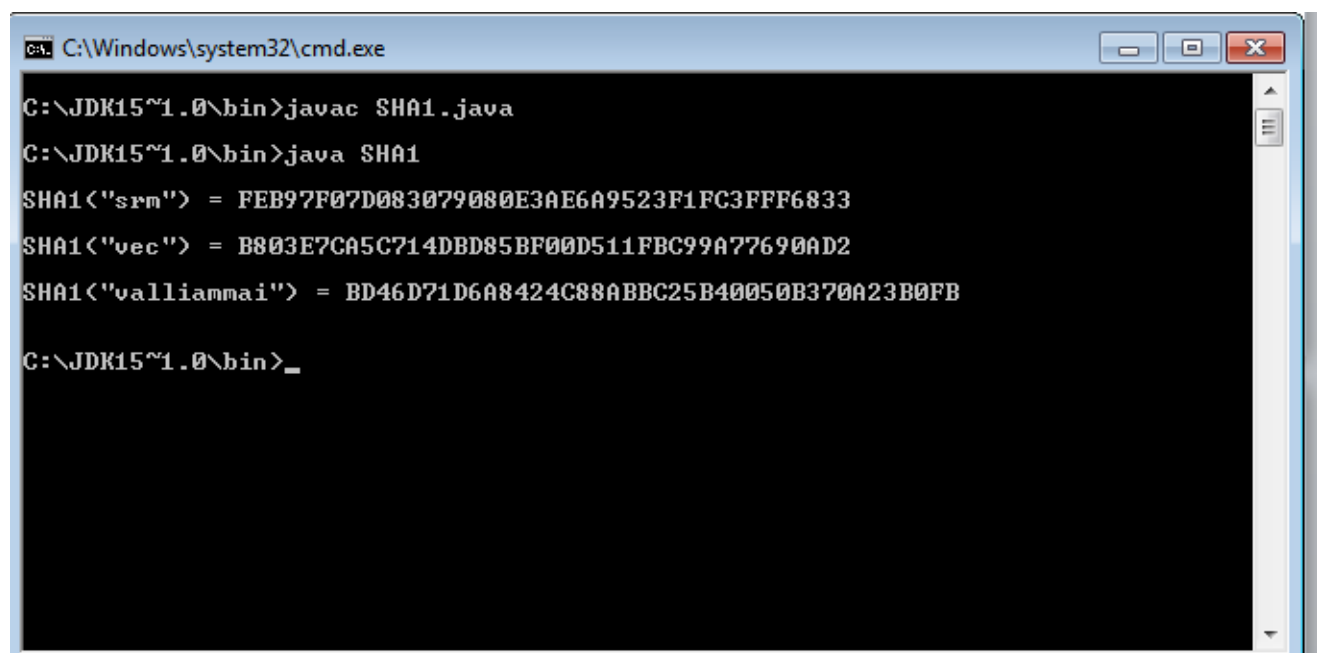
```

System.out.println();
System.out.println("SHA1(\""+input+"\") = " +bytesToHex(output));
input = "valliammai";
md.update(input.getBytes());
output = md.digest();
System.out.println();
System.out.println("SHA1(\""+input+"\") = " +bytesToHex(output));
System.out.println(""); }
catch (Exception e) {
System.out.println("Exception: " +e);
}
}

public static String bytesToHex(byte[] b) {
    char hexDigit[] = {'0', '1', '2', '3', '4', '5', '6', '7', '8', '9', 'A', 'B', 'C', 'D', 'E', 'F'};
    StringBuffer buf = new StringBuffer();
    for (int j=0; j<b.length; j++) {
        buf.append(hexDigit[(b[j] >> 4) & 0x0f]);
        buf.append(hexDigit[b[j] & 0x0f]); }
    return buf.toString(); }
}

```

OUTPUT:



The screenshot shows a Windows command prompt window titled "C:\Windows\system32\cmd.exe". The user has navigated to the directory "C:\JDK15~1.0\bin" and executed the following commands:

```

C:\JDK15~1.0\bin>javac SHA1.java
C:\JDK15~1.0\bin>java SHA1

```

The output of the program is displayed as follows:

```

SHA1<'srm'> = FEB97F07D083079080E3AE6A9523F1FC3FFF6833
SHA1<'vec'> = B803E7CA5C714DBD85BF00D511FBC99A77690AD2
SHA1<'valliammai'> = BD46D71D6A8424C88ABBC25B40050B370A23B0FB

```

The prompt then shows the user entering an underscore character: "C:\JDK15~1.0\bin>_".

VIVA QUESTIONS (PRELAB and POSTLAB):

1. SHA-1 produces a hash value of how many bits?
2. What is the number of round computation steps in the SHA-256 algorithm?
3. In SHA-512, the message is divided into blocks size of how many bits for the hash computation.
4. What is the maximum length of the message (in bits) that can be taken by SHA-512?
5. What is the length of the message in SHA-512 after it is padded?
6. Describe the big-endian format?
7. In SHA-512, the registers 'a' to 'h' are obtained by taking the first 64 bits of the fractional parts of the cube roots of the first 8 prime numbers. Is it true or false?
8. What is the size of W (in bits) in the SHA-512 processing of a single 1024- bit block?
9. In the SHA-512 processing of a single 1024- bit block, how the round constants are obtained?
10. What is the maximum length of the message (in bits) that can be taken by SHA-512?
11. What does the figure represent?
12. Among the registers 'a' to 'h' how many involve permutation in each round?

RESULT:

Thus the program to implement Secure Hash Algorithm was developed and executed successfully.

AIM:

To write a program to implement the digital signature scheme in java

PRELAB DISCUSSION:

Digital signatures are based on public key cryptography, also known as asymmetric cryptography. Using a public key algorithm such as RSA, one can generate two keys that are mathematically linked: one private and one public. To create a digital signature, signing software (such as an email program) creates a one-way hash of the electronic data to be signed. The private key is then used to encrypt the hash. The encrypted hash -- along with other information, such as the hashing algorithm -- is the digital signature. The reason for encrypting the hash instead of the entire message or document is that a hash function can convert an arbitrary input into a fixed length value, which is usually much shorter. This saves time since hashing is much faster than signing. The value of the hash is unique to the hashed data. Any change in the data, even changing or deleting a single character, results in a different value. This attribute enables others to validate the integrity of the data by using the signer's public key to decrypt the hash. If the decrypted hash matches a second computed hash of the same data, it proves that the data hasn't changed since it was signed. If the two hashes don't match, the data has either been tampered with in some way (integrity) or the signature was created with a private key that doesn't correspond to the public key presented by the signer (authentication).

A digital signature can be used with any kind of message -- whether it is encrypted or not -- simply so the receiver can be sure of the sender's identity and that the message arrived intact. Digital signatures make it difficult for the signer to deny having signed something (non-repudiation) -- assuming their private key has not been compromised -- as the digital signature is unique to both the document and the signer, and it binds them together. A digital certificate, an electronic document that contains the digital signature of the certificate-issuing authority, binds together a public key with an identity and can be used to verify a public key belongs to a particular person or entity. Most modern email programs support the use of digital signatures and digital certificates, making it easy to sign any outgoing emails and validate digitally signed incoming messages. Digital signatures are also used extensively to provide proof of authenticity, data integrity and non- repudiation of communications and transactions conducted over the Internet.

ALGORITHM:

1. Choose a prime number q , which is called the prime divisor.
2. Choose another prime number p , such that $p-1 \bmod q = 0$. p is called the prime modulus.
3. Choose an integer g , such that $1 < g < p$, $g^{q-1} \bmod p = 1$ and $g = h^{((p-1)/q)} \bmod p$. q is also called g 's multiplicative order modulo p .
4. Choose an integer, such that $0 < x < q$.
5. Compute y as $g^x \bmod p$.
6. Package the public key as $\{p, q, g, y\}$, $\{p, q, g, x\}$.

7. Generate the message digest h , using a hash algorithm like SHA1.
8. Generate a random number k , such that $0 < k < q$.
9. Compute r as $(g^{**k} \bmod p) \bmod q$. If $r = 0$, select a different k .
10. Compute i , such that $k * i \bmod q = 1$. i is called the modular multiplicative inverse of k modulo q .
11. Compute $s = i * (h + r * x) \bmod q$. If $s = 0$, select a different k .
12. Package the digital signature as $\{r, s\}$.
13. Generate the message digest h , using the same hash algorithm.
14. Compute w , such that $s * w \bmod q = 1$. w is called the modular multiplicative inverse of s modulo q .
15. Compute $u1 = h * w \bmod q$. Compute $u2 = r * w \bmod q$.
16. Compute $v = (((g^{**u1}) * (y^{**u2})) \bmod p) \bmod q$.
17. If $v == r$, the digital signature is valid.

PROGRAM

```
import java.util.*;
import java.math.BigInteger;
class dsaAlg
{
    final static BigInteger one = new BigInteger("1");
    final static BigInteger zero = new BigInteger("0");

    /* incrementally tries for next prime */
    public static BigInteger getNextPrime(String ans)
    {
        BigInteger test = new BigInteger(ans);
        while (!test.isProbablePrime(99))
        {
            test = test.add(one);
        }
        return test;
    }

    /* finds largest prime factor of n */
    public static BigInteger findQ(BigInteger n)
    {
        BigInteger start = new BigInteger("2");
        while (!n.isProbablePrime(99))
        {
            while (!(n.mod(start)).equals(zero))
            {
                start = start.add(one);
            }
        }
    }
}
```

```

        n = n.divide(start);
    }
    return n;
}

/* finds a generator mod p */
public static BigInteger getGen(BigInteger p, BigInteger q, Random r)
{
    BigInteger h = new BigInteger(p.bitLength(), r);
    h = h.mod(p);
    return h.modPow((p.subtract(one)).divide(q), p);
}

public static void main (String[] args) throws java.lang.Exception
{
    Random randObj = new Random();

    /* establish the global public key components */
    BigInteger p = getNextPrime("10600"); /* approximate prime */
    BigInteger q = findQ(p.subtract(one));
    BigInteger g = getGen(p,q,randObj);

    /* public key components */
    System.out.println("Digital Signature Algorithm");
    System.out.println("global public key components are:");
    System.out.println("p is: " + p);
    System.out.println("q is: " + q);
    System.out.println("g is: " + g);

    /* find the private key */
    BigInteger x = new BigInteger(q.bitLength(), randObj);
    x = x.mod(q);

    /* corresponding public key */
    BigInteger y = g.modPow(x,p);

    /* random value message */
    BigInteger k = new BigInteger(q.bitLength(), randObj);
    k = k.mod(q);

    /* randomly generated hash value and digital signature */
    BigInteger r = (g.modPow(k,p)).mod(q);
    BigInteger hashVal = new BigInteger(p.bitLength(), randObj);
    BigInteger kInv = k.modInverse(q);
    BigInteger s = kInv.multiply(hashVal.add(x.multiply(r)));
    s = s.mod(q);
}

```

```

/* secret information */
System.out.println("secret information are:");
System.out.println("x (private) is: " + x);
System.out.println("k (secret) is: " + k);
System.out.println("y (public) is: " + y);
System.out.println("h (rndhash) is: " + hashVal);
System.out.println("Generating digital signature:");
System.out.println("r is : " + r);
System.out.println("s is : " + s);

```

```

/*verify the digital signature */
BigInteger w = s.modInverse(q);
BigInteger u1 = (hashVal.multiply(w)).mod(q);
BigInteger u2 = (r.multiply(w)).mod(q);
BigInteger v = (g.modPow(u1,p)).multiply(y.modPow(u2,p));
v = (v.mod(p)).mod(q);
System.out.println("verifying digital signature (checkpoints):");
System.out.println("w is : " + w);
System.out.println("u1 is: " + u1);
System.out.println("u2 is: " + u2);
System.out.println("v is : " + v);
if (v.equals(r))
{
    System.out.println("success: digital signature is verified! " + r);
}
else
{
    System.out.println("error: incorrect digital signature");
}

```

```

}

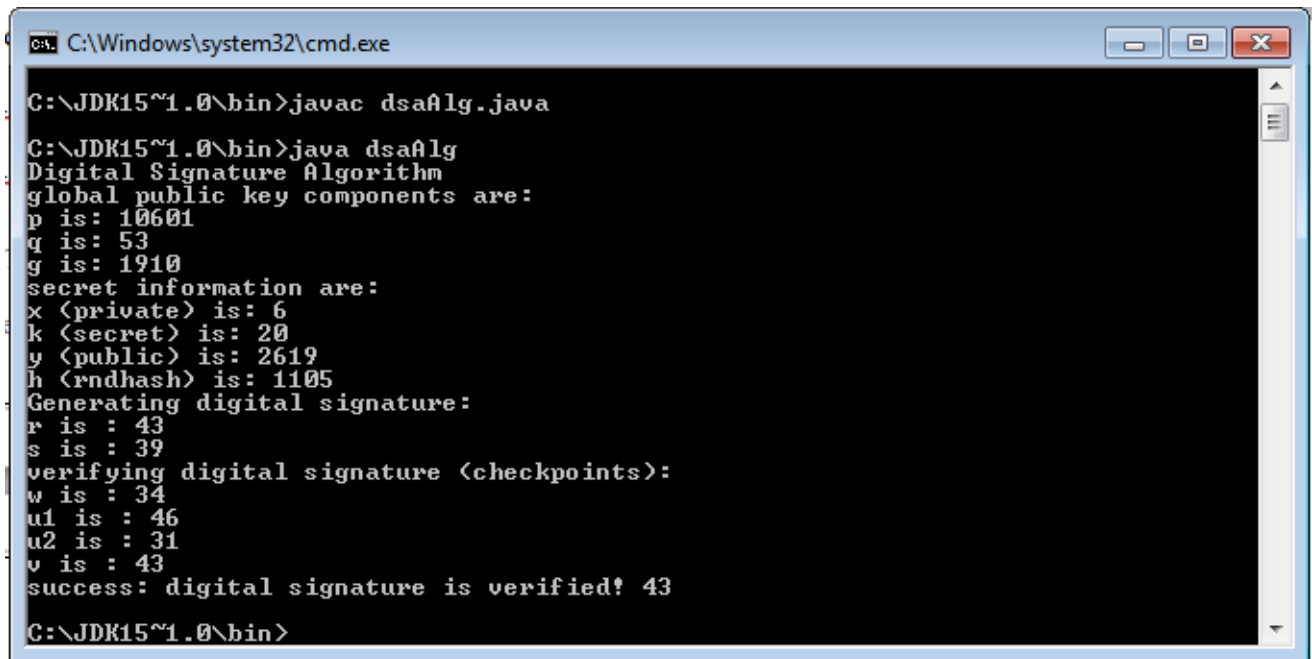
```

```

}

```

OUTPUT



```
C:\Windows\system32\cmd.exe

C:\JDK15~1.0\bin>javac dsaAlg.java

C:\JDK15~1.0\bin>java dsaAlg
Digital Signature Algorithm
global public key components are:
p is: 10601
q is: 53
g is: 1910
secret information are:
x (private) is: 6
k (secret) is: 20
y (public) is: 2619
h (rndhash) is: 1105
Generating digital signature:
r is : 43
s is : 39
verifying digital signature (checkpoints):
w is : 34
u1 is : 46
u2 is : 31
v is : 43
success: digital signature is verified! 43

C:\JDK15~1.0\bin>
```

VIVA QUESTIONS (PRELAB and POSTLAB):

1. What is a digital signature?
2. What does a digital signature look like?
3. What is an electronic document?
4. Does that mean that the authenticity of any electronic document can be verified by a digital signature?
5. What is it like to actually sign an electronic document?
6. Can you actually see the signer's handwritten signature?
7. How do I get a digital signature certificate?
8. What is a certificate? What does it mean to "publish" a certificate?
9. How am I identified as the signer?
10. If my private key is stored on my computer, can't someone sign the documents without my permission by getting access to the computer?
11. Can a digital signature be forged?
12. What are the responsibilities and the liability of a digital signature certificate subscriber?
13. What are the practical uses of a digital signature?

RESULT:

Thus the program to implement Digital Signature was developed and executed successfully

AIM:

To demonstrate intrusion detection system (ids) using the tool snort

PRELAB DISCUSSION and PROCEDURE:**1. Configure and Use Snort IDS on Windows**

Steps to configure Snort on Windows machine and how to use it for detection of attacks.

2. Steps:

1. Download Snort from "<http://www.snort.org/>" website.
2. Also download Rules from the same website. You need to sign up to get rules for registered users.
3. Click on the Snort_(version-number)_Installer.exe file to install it. By-default it will install snort in the "C:\Snort" directory.
4. Extract downloaded Rules file: snortrules-snapshot-(number).tar.gz
5. Copy all files from the "rules" directory of the extracted folder and paste them into "C:\Snort\rules" directory.
6. Copy "snort.conf" file from the "etc" directory of the extracted folder and paste it into "C:\Snort\etc" directory. Overwrite existing file if there is any.
7. Open command prompt (cmd.exe) and navigate to directory "C:\Snort\bin" directory.
8. To execute snort in sniffer mode use following command:

```
snort -dev -i 2
```

-i indicate interface number.

-dev is used to run snort to capture packets.

To check interface list use following command: snort -W

9. To execute snort in IDS mode, we need to configure a file "snort.conf" according to our network environment.

10. Set up network address we want to protect in snort.conf file. To do that look for "HOME_NET" and add your IP address.

```
var HOME_NET 10.1.1.17/8
```

11. You can also set addresses or DNS_SERVERS, if you have any. otherwise go to the next step.

12. Change RULE_PATH variable with the path of rules directory.

```
var RULE_PATH c:\snort\rules
```

13. Change the path of all libraries with the name and path on your system. or change path of snort_dynamicpreprocessorvariable.

or file C:\Snort\lib\snort_dynamicccpreprocessor\sf_dcerpc.dll

You need to do this to all library files in the "C:\Snort\lib" directory. The old path might be something like: "/usr/local/lib/...". you need to replace that path with you system path.

14. Change path of the "dynamicengine" variable value in the "snort.conf" file with the path of your system. Such as:

dynamicengine C:\Snort\lib\snort_dynamicengine\sf_engine.dll

15 Add complete path for "include classification.config" and "include reference.config" files.

include c:\snort\etc\classification.config

include c:\snort\etc\reference.config

16. Remove the comment on the line to allow **ICMP** rules, if it is already commented.

include \$RULE_PATH/icmp.rules

17. Similarly, remove the comment of ICMP-info rules comment, if it is already commented.

include \$RULE_PATH/icmp-info.rules

18 To add log file to store alerts generated by snort, search for "output log" test and add following line:

output alert_fast: snort-alerts.ids

19. Comment whitelist \$WHITE_LIST_PATH/white_list.rules and blacklist

\$BLACK_LIST_PATH/black_list.rules lines. Also ensure that you add change the line above

\$WHITE_LIST_PATH

Change nested_ip inner , \ to nested_ip inner #, \

20. Comment following lines:

#preprocessor normalize_ip4

#preprocessor normalize_tcp: ips ecn stream

#preprocessor normalize_icmp4

#preprocessor normalize_ip6

#preprocessor normalize_icmp6

21. Save the "snort.conf" file and close it.

22. Go to the "C:\Snort\log" directory and create a file: snort-alerts.ids

23. To start snort in IDS mode, run following command:

snort -c c:\snort\etc\snort.conf -l c:\snort\log -i 2

Above command will generate log file that will not be readable without using a tool. To read it use following command:

C:\Snort\Bin> snort -r ../log/log-filename

To generate Log files in ASCII mode use following command while running snort in IDS mode:

snort -A console -i2 -c c:\Snort\etc\snort.conf -l c:\Snort\log -K ascii

24. Scan the computer running snort from another computer using PING or launch attack. Then check snort-alerts.ids file the log folder.

VIVA QUESTIONS (PRE LAB and POSTLAB)

1. When discussing IDS/IPS, what is a signature?
2. "Semantics-aware" signatures automatically generated by Nemean are based on traffic at which two layers?
3. Which of the following is used to provide a baseline measure for comparison of IDSes?
4. What Is Ips And Ids?
5. What Are The Functions Of Intrusion Detection?
6. What Is Ids In Networking?
7. Explain Host Based (hids)?
8. What Is An Intrusion Detection System?
9. What is the advantage of anomaly detection?
10. Define a false positive.
11. One of the most obvious places to put an IDS sensor is near the firewall. Where exactly in relation to the firewall is the most productive placement?
12. What is the purpose of a shadow honeypot?
13. At which two traffic layers do most commercial IDSes generate signatures?
14. An IDS follows a two-step process consisting of a passive component and an active component.
15. Which of the following is part of the active component?

RESULT:

Thus the intrusion detection system (ids) using the tool snort program was demonstrated and verified successfully.

AIM:

To explore automated and penetration tools on network (KF Sensor)

PRELAB DISCUSSION:**HONEYPOTS**

When it comes to computer security, honeypots are all the rage. Honeypots can detect unauthorized activities that might never be picked up by a traditional intrusion detection system. Furthermore, since almost all access to a honeypot is unauthorized, nearly everything in a honeypot's logs is worth paying attention to. Honeypots can act as a decoy to keep hackers away from your production servers. At the same time though, a honeypot can be a little tricky to deploy. In this article, I will walk you through the process of deploying a honeypot.

INTRODUCTION

There are many different types of honeypot systems. Honeypots can be hardware appliances or they can be software based. Software based firewalls can reside on top of a variety of operating systems. For the most part though, honeypots fall into two basic categories; real and virtual.

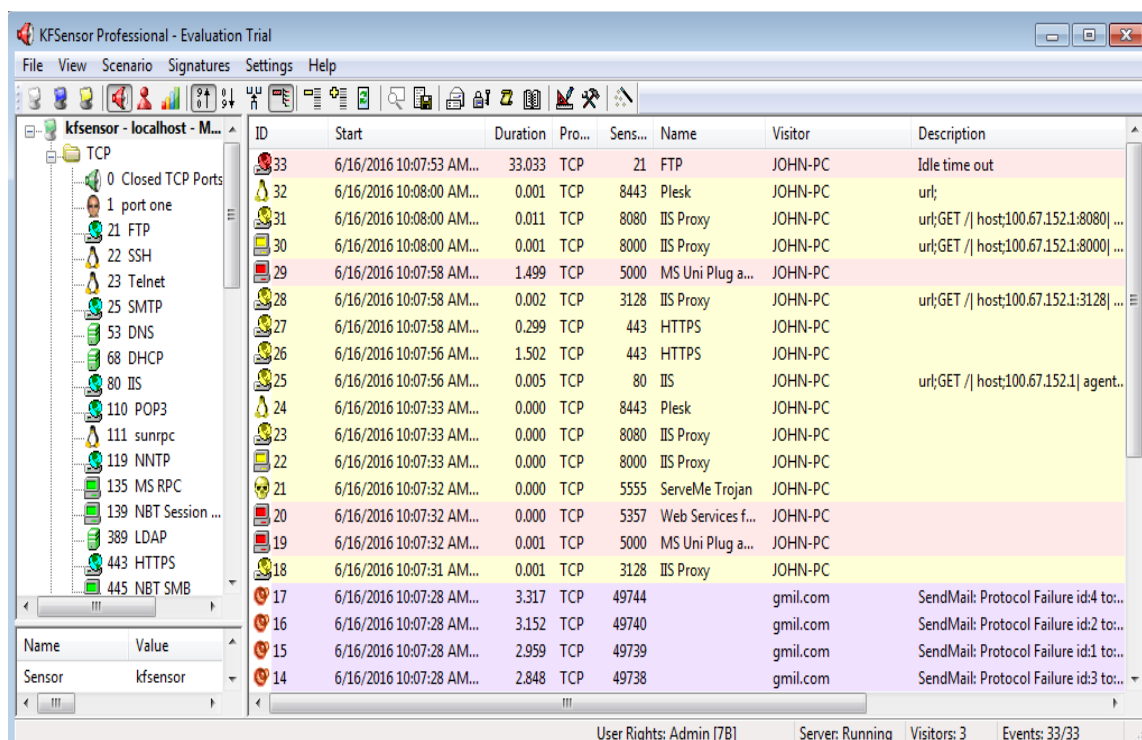
A virtual honeypot is essentially an emulated server. There are both hardware and software implementations of virtual honeypots. For example, if a network administrator was concerned that someone might try to exploit an FTP server, the administrator might deploy a honeypot appliance that emulates an FTP server.

Downloading and installing KF Sensor

- The KF Sensor download consists of a 1.7 MB self-extracting executable file.
- Download the file and copy it into an empty folder on your computer.
- When you double click on the file, it will launch a very basic Setup program.
- The only thing special that you need to know about the Setup process is that it will require a reboot

Using KFSensor

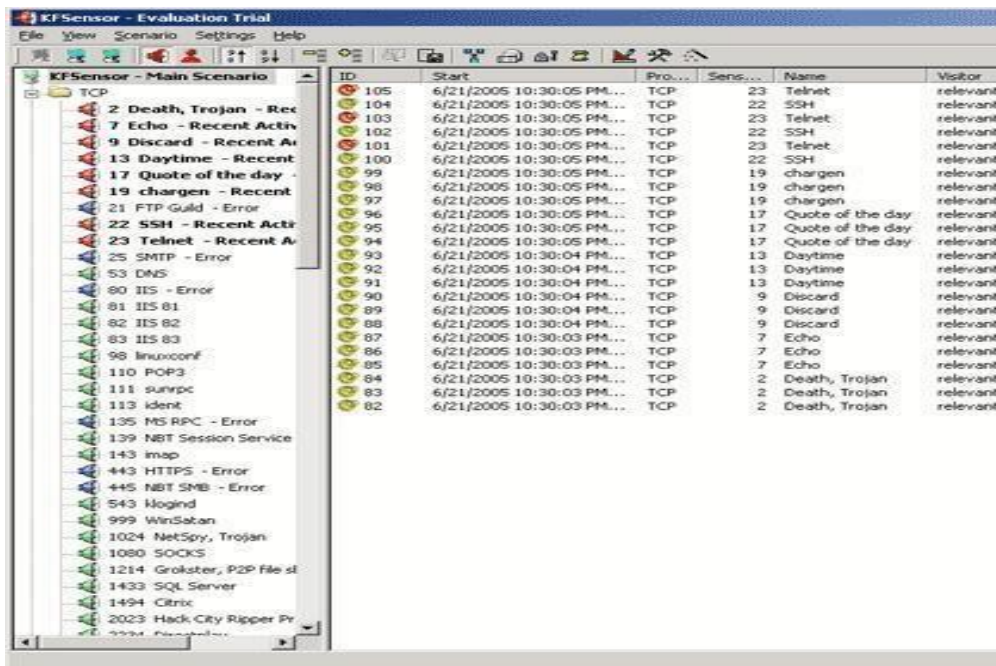
Step1: You will see the main KFSensor screen shown



- As you can see, the column on the left contains a list of port numbers and what the port is typically used for.
- If the icon to the left of a port listing is green, it means that KFSensor is actively monitoring that port for attacks.
- If the icon is blue, it means that there has been an error and KFSensor is not watching for exploits aimed at that particular port.

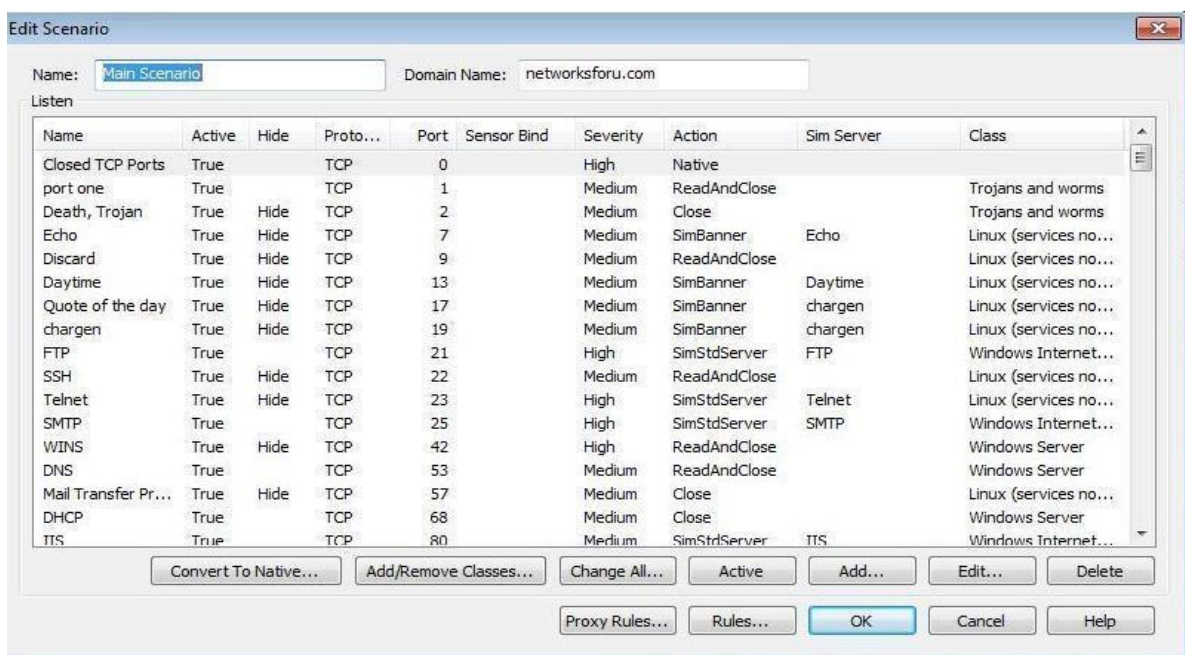
Testing the software

- Once you've got the software up and running, one of the best things that you can do is to test the software by launching a port scan against the machine that's running KFSensor.
 - For the port scan, we using the HostScan.
- It simply scans a block of IP addresses, looking for open ports. Figure B shows how the KFSensor reacts to a partial port scan.
- If you look at Figure B, you will notice that the icons next to ports that were scanned turn red to indicate recent activity.



Modifying the Honeypot's behavior

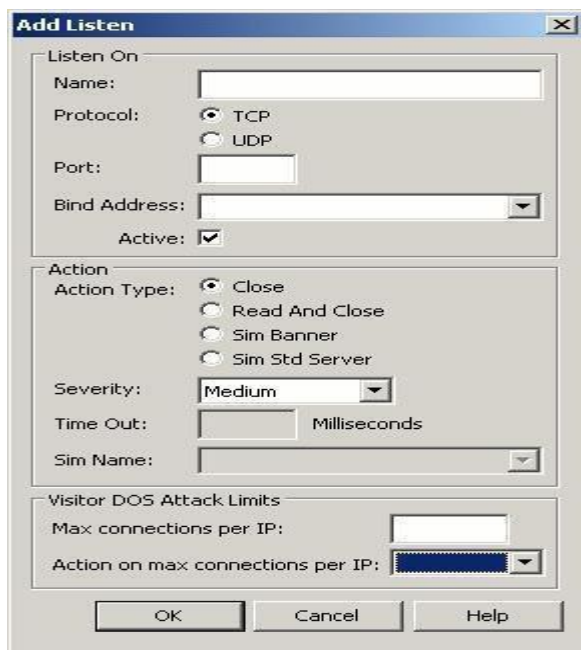
- To create or modify rules, select the Edit Active Scenario command from the scenario menu.
- When you do, you will see a dialog box which contains a summary of all of the existing rules.
- You can either select a rule and click the Edit button to edit a rule, or you can click the Add button to create a new rule.
- Both procedures work similarly.



Click the Add button and you will see the Add Listen dialog box, shown in Figure D.

- The first thing that this dialog box asks for is a name. This is just a name for the rule.
- Pick something descriptive though, because the name that you enter is what will show up in the logs whenever the rule is triggered.

Click on Add Button

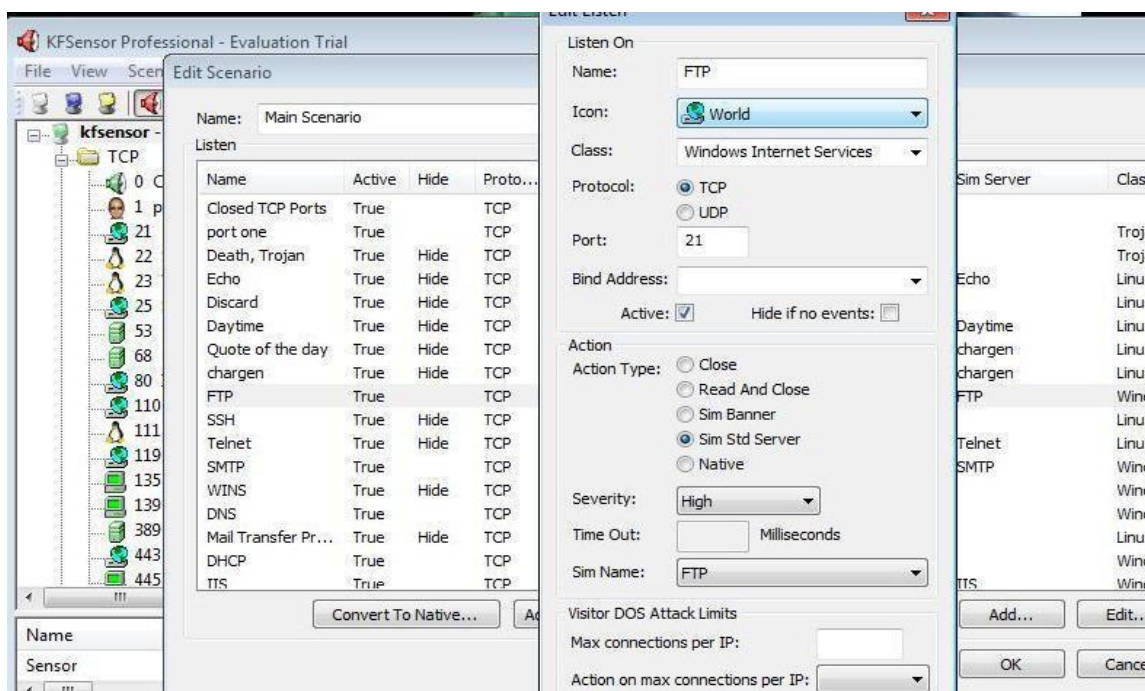


The 'Add Listen' dialog box is used to configure a new listener rule. It contains the following fields and options:

- Listen On:**
 - Name: Text input field.
 - Protocol: Radio buttons for TCP (selected) and UDP.
 - Port: Text input field.
 - Bind Address: Dropdown menu.
 - Active: Checked checkbox.
- Action:**
 - Action Type: Radio buttons for Close (selected), Read And Close, Sim Banner, and Sim Std Server.
 - Severity: Dropdown menu set to Medium.
 - Time Out: Text input field followed by 'Milliseconds'.
 - Sim Name: Dropdown menu.
- Visitor DOS Attack Limits:**
 - Max connections per IP: Text input field.
 - Action on max connections per IP: Dropdown menu.

Buttons at the bottom: OK, Cancel, Help.

Click on Edit Button



The screenshot shows the KFSensor Professional interface. On the left is a tree view of sensors. The main area is divided into two panes:

- Edit Scenario:** A table listing various sensors and their status.
- Edit Listen:** A detailed configuration window for a specific listener rule.

Name	Active	Hide	Proto...
Closed TCP Ports	True		TCP
port one	True		TCP
21	True		TCP
22	True	Hide	TCP
23	True	Hide	TCP
25	True	Hide	TCP
53	True	Hide	TCP
68	True	Hide	TCP
80	True	Hide	TCP
110	True	Hide	TCP
111	True	Hide	TCP
119	True	Hide	TCP
135	True	Hide	TCP
139	True	Hide	TCP
389	True	Hide	TCP
443	True	Hide	TCP
445	True	Hide	TCP

The 'Edit Listen' window for the 'FTP' rule shows the following configuration:

- Listen On:** Name: FTP, Icon: World, Class: Windows Internet Services, Protocol: TCP, Port: 21, Bind Address: (blank), Active: checked, Hide if no events: unchecked.
- Action:** Action Type: Sim Std Server (selected), Severity: High, Time Out: (blank) Milliseconds, Sim Name: FTP.
- Visitor DOS Attack Limits:** Max connections per IP: (blank), Action on max connections per IP: (blank).

Buttons at the bottom of the 'Edit Listen' window: Add..., Edit..., OK, Cancel.

- The next few fields are protocol, port, and Bind Address. These fields allow you to choose what the rule is listening for. For example, you could configure the rule to listen to TCP port 1023 on IP address 192.168.1.100. The bind address portion of the rule is optional though. If you leave the bind address blank, the rule will listen across all of the machine's NICs.
- Now that you have defined the listener, it's time to configure the action that the rule takes when traffic is detected on the specified port. Your options are close, read and close, Sim Banner, and SimStd Server.
- The close option tells the rule to just terminate the connection. Read and close logs the information and then terminates the connection. The SimStd Server and Sim Banner options

pertain to server emulation. The Sim Banner option allows you to perform a very simple server emulation, such as what you might use to emulate an FTP server.

- The Sim STD Server option allows you to emulate a more complex server, such as an IIS server.
- If you choose to use one of the sim options, you will have to fill in the simulator's name just below the Time Out field.
- The other part of the Action section that's worth mentioning is the severity section. KFSensor treated some events as severe and other events as a more moderate threat. The dialog box's Severity drop down list allows you to determine what level of severity should be associated with the event that you are logging.
- The final portion of the Add Listen dialog box is the Visitor DOS Attack Limits section. This section allows you to prevent denial of service attacks against KFSensor. You can determine the maximum number of connections to the machine per IP address (remember that this applies on a per rule basis).
- If your threshold is exceeded, you can choose to either ignore the excessive connections or you can lock out the offending IP address.
- Now that you have configured the new rule, select the Active Button to Enable/Disable. The new rule should now be in effect.

VIVA QUESTIONS (PRE LAB and POSTLAB):

1. How to build and use a Honeypot
2. List the types of interactions
3. Define honey tokens
4. Give the advantages of honey pot
5. Mention the protocols used by IPSec to provide security.
6. How are the passwords stored in password file in UNIX operating system?
7. Which system is used to protect credit card transactions on the internet?
8. What is meant by Block Cipher?
9. Define Plain Text.

RESULT:

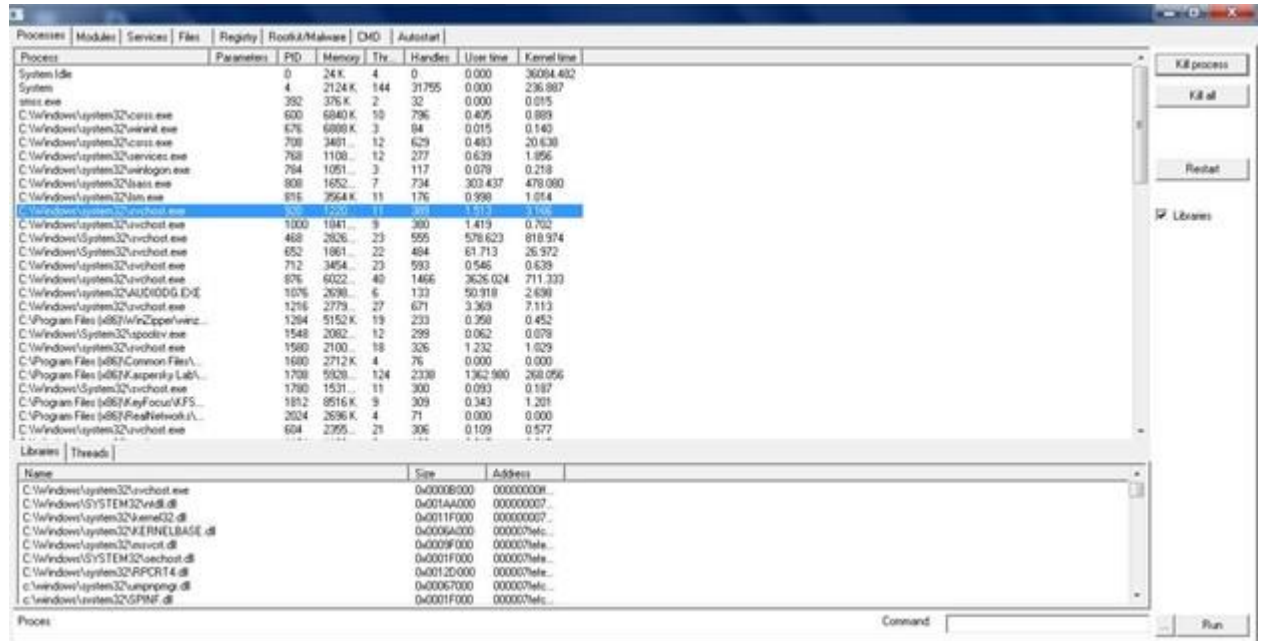
Thus the program to explore automated attack and penetration tools has been completed successfully

AIM

Root kit is a stealth type of malicious software designed to hide the existence of certain process from normal methods of detection and enables continued privileged access to a computer.

PRELAB DISCUSSION and PROCEDURE :

- Download Rootkit Tool from GMER website. www.gmer.net
- This displays the Processes, Modules, Services, Files, Registry, RootKit/Malwares, Autostart, CMD of local host.
- Select Processes menu and kill any unwanted process if any. Modules menu displays the various system files like .sys, .dll
- Services menu displays the complete services running with Autostart, Enable, Disable, System, Boot.
- Files menu displays full files on Hard-Disk volumes.
- Registry displays **Hkey_Current_user** and **Hkey_Local_Machine**. Rootkits/Malwares scans the local drives selected.
- **Autostart** displays the registry base Autostart applications.
- CMD allows the user to interact with command line utilities or Registry.



Processes	Modules	Services	Files	Registry	Rootkit/Malware	CMD	Autostart
Name	File	Address	Size				
taskmgr.exe	SystemRoot\system32\taskmgr.exe	0000000000000000	6164960				
hal.dll	SystemRoot\system32\hal.dll	0000000000000000	259008				
kernel.dll	SystemRoot\system32\kernel.dll	0000000000000000	40960				
update_GenuineInte...	SystemRoot\system32\update_GenuineInte...	0000000000000000	327684				
PSHED.dll	SystemRoot\system32\PSHED.dll	0000000000000000	81920				
CLFS.SYS	SystemRoot\system32\CLFS.SYS	0000000000000000	385024				
CI.dll	SystemRoot\system32\CI.dll	0000000000000000	479232				
Wd01000.sys	SystemRoot\system32\drivers\Wd01000.sys	0000000000000000	794624				
WdFLDR.SYS	SystemRoot\system32\drivers\WdFLDR.SYS	0000000000000000	67536				
HT.sys	SystemRoot\system32\DRIVERS\HT.sys	0000000000000000	7741440				
ACPI.sys	SystemRoot\system32\drivers\ACPI.sys	0000000000000000	356252				
WMLIB.SYS	SystemRoot\system32\drivers\WMLIB.SYS	0000000000000000	36864				
msadrv.sys	SystemRoot\system32\drivers\msadrv.sys	0000000000000000	40960				
pci.sys	SystemRoot\system32\drivers\pci.sys	0000000000000000	208896				
vdhroot.sys	SystemRoot\system32\drivers\vdhroot.sys	0000000000000000	53248				
usb3hci.sys	SystemRoot\system32\DRIVERS\usb3hci.sys	0000000000000000	40960				
on_fm.sys	SystemRoot\system32\DRIVERS\on_fm.sys	0000000000000000	248896				
partmgr.sys	SystemRoot\system32\drivers\partmgr.sys	0000000000000000	86016				
votmg.sys	SystemRoot\system32\drivers\votmg.sys	0000000000000000	86016				
votmgx.sys	SystemRoot\system32\drivers\votmgx.sys	0000000000000000	376832				
mountmgr.sys	SystemRoot\system32\drivers\mountmgr.sys	0000000000000000	106496				
atapi.sys	SystemRoot\system32\drivers\atapi.sys	0000000000000000	36864				
ataport.SYS	SystemRoot\system32\drivers\ataport.SYS	0000000000000000	172102				
usahci.sys	SystemRoot\system32\drivers\usahci.sys	0000000000000000	45956				
PODDECI.SYS	SystemRoot\system32\drivers\PODDECI.SYS	0000000000000000	65536				
usb3hub.sys	SystemRoot\system32\DRIVERS\usb3hub.sys	0000000000000000	2520640				
usbport.sys	SystemRoot\system32\DRIVERS\usbport.sys	0000000000000000	409600				
andata.sys	SystemRoot\system32\drivers\andata.sys	0000000000000000	45956				
fbmg.sys	SystemRoot\system32\drivers\fbmg.sys	0000000000000000	211296				
libata.sys	SystemRoot\system32\drivers\libata.sys	0000000000000000	81920				
Nifs.sys	SystemRoot\system32\Drivers\Nifs.sys	0000000000000000	1740800				
uspio.sys	SystemRoot\system32\Drivers\uspio.sys	0000000000000000	385024				
kscedd.sys	SystemRoot\system32\Drivers\kscedd.sys	0000000000000000	110592				
crq.sys	SystemRoot\system32\Drivers\crq.sys	0000000000000000	466944				
pcw.sys	SystemRoot\system32\drivers\pcw.sys	0000000000000000	69632				
Flt_Parc.sys	SystemRoot\system32\Drivers\Flt_Parc.sys	0000000000000000	40960				
ndis.sys	SystemRoot\system32\drivers\ndis.sys	0000000000000000	995328				
NETIO.SYS	SystemRoot\system32\drivers\NETIO.SYS	0000000000000000	283216				
ksacdk.sys	SystemRoot\system32\Drivers\ksacdk.sys	0000000000000000	176128				
tcpip.sys	SystemRoot\system32\drivers\tcpip.sys	0000000000000000	2053096				
fwpmclnt.sys	SystemRoot\system32\drivers\fwpmclnt.sys	0000000000000000	299008				
vmtoolsd.sys	SystemRoot\system32\drivers\vmtoolsd.sys	0000000000000000	65536				

GMER 2.2.19882 WINDOWS 6.1.7601 Service Pack 1 x64

AntiVirus: http://www.avast.com

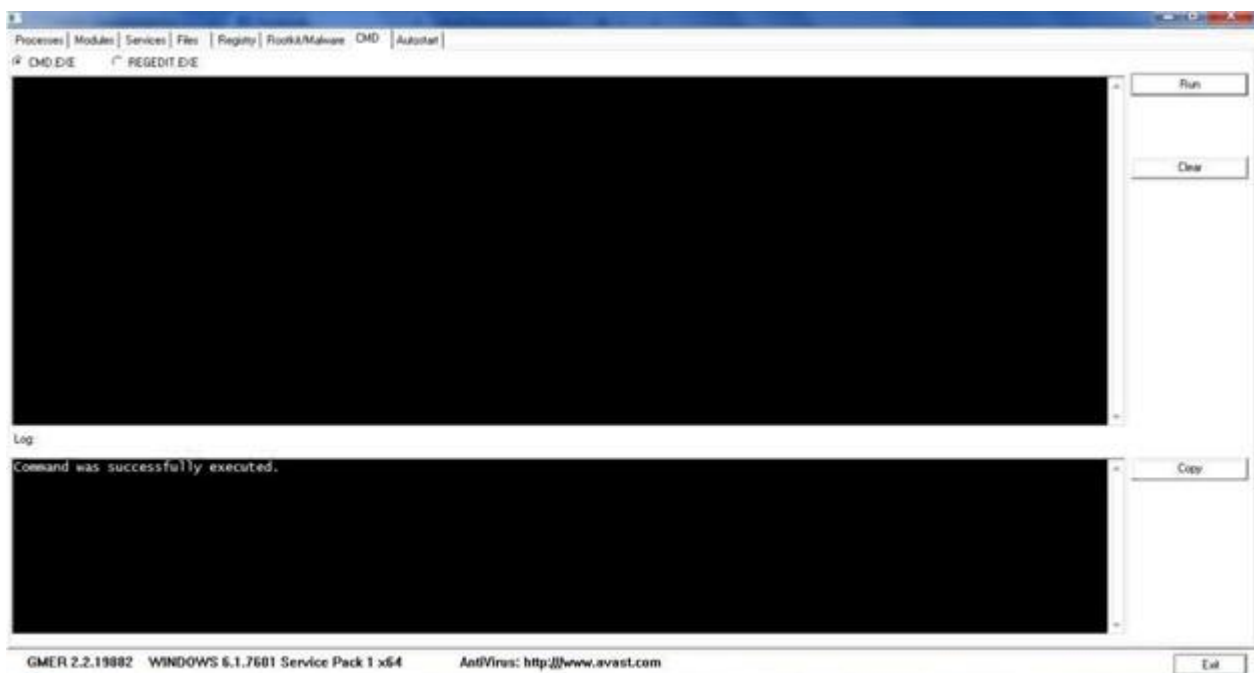
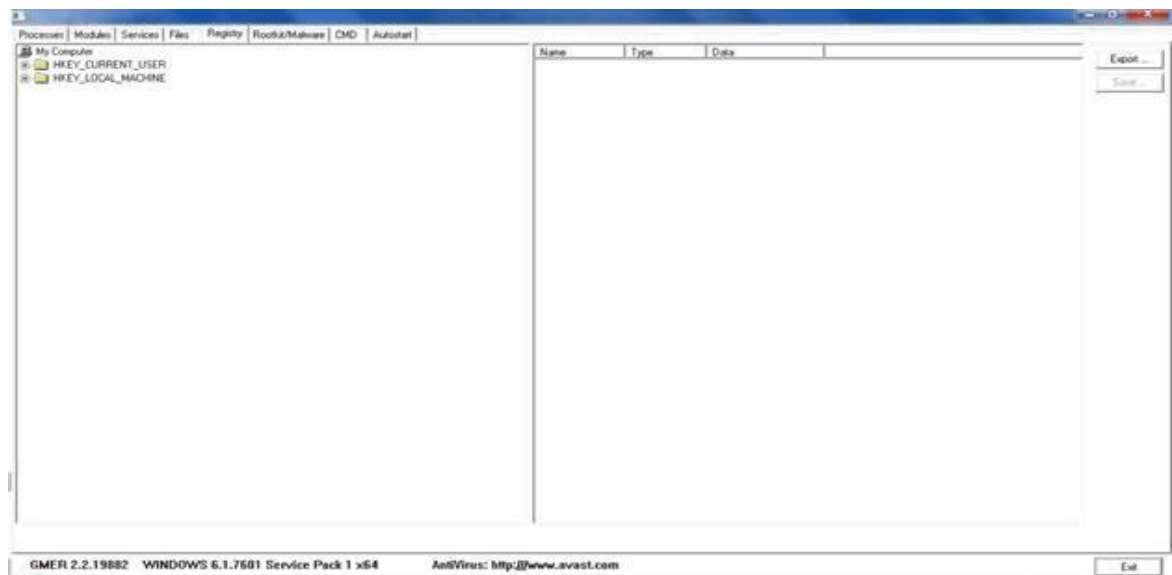
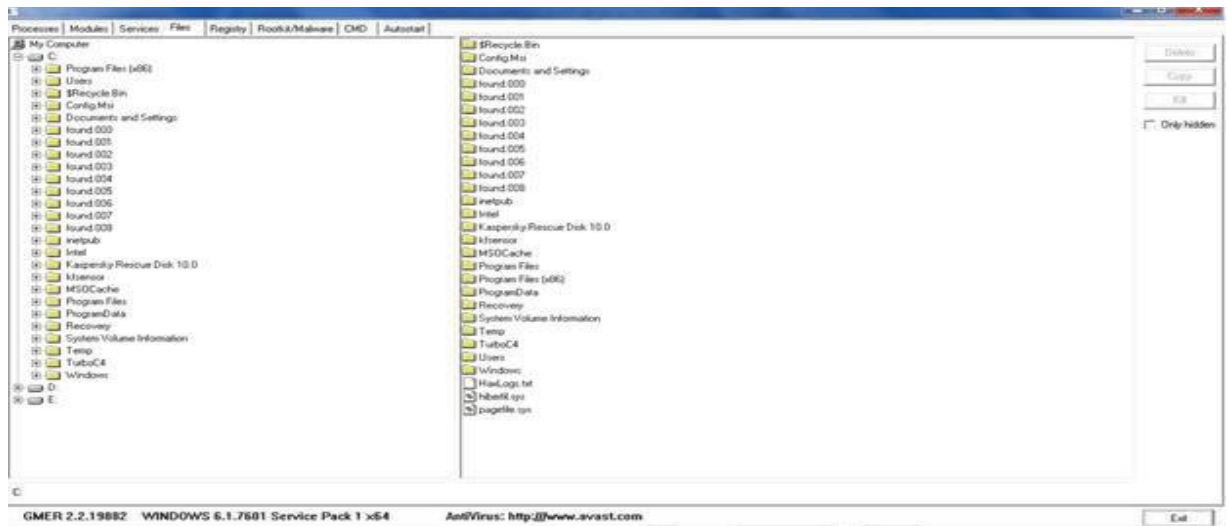
Exit

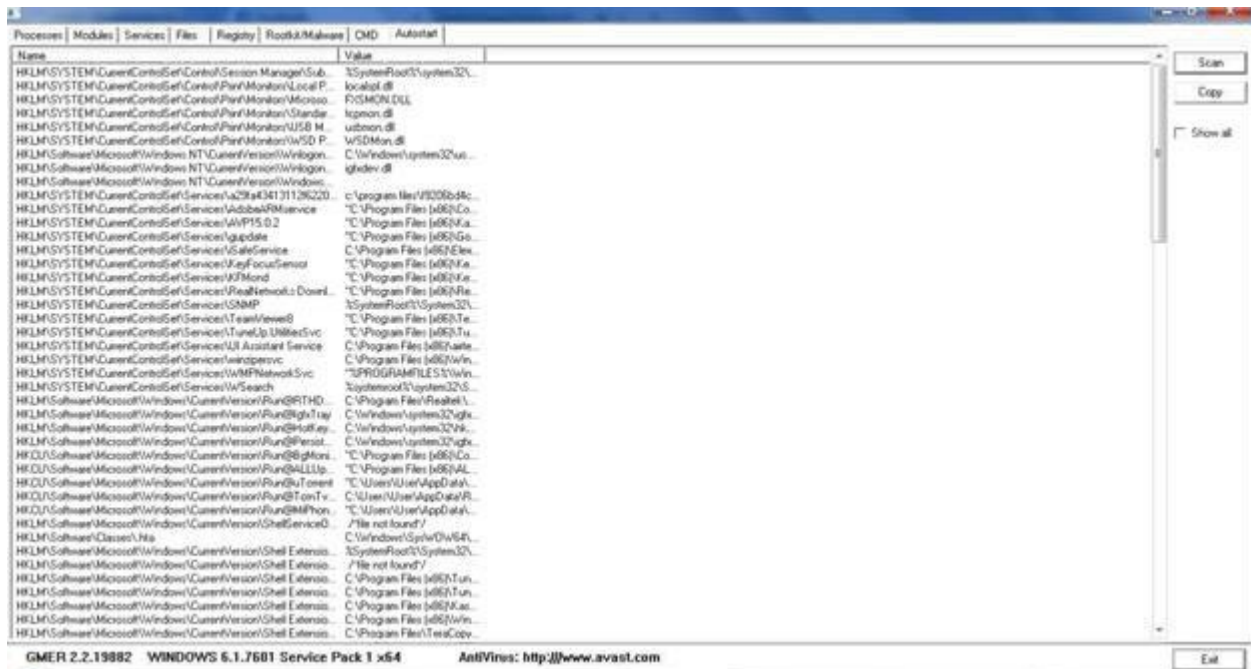
Processes	Modules	Services	Files	Registry	Rootkit/Malware	CMD	Autostart
Name	Start	File name	Description				
ac29434131126...	AUTO	c:\program files\ac294341311262006-500e6...	ac294341311262006-500e6...				
ACPI	BOOT	system32\drivers\ACPI.sys	Microsoft ACPI Driver				
AcpiFw	MANUAL	SystemRoot\system32\drivers\acpifw.sys	ACPI Power Meter Driver				
AdobeFlashPlayer	AUTO	"C:\Program Files\Adobe\Flash Player\...	Adobe Acrobat Reader keeps you Adobe softw...				
adp3hci	MANUAL	C:\Windows\System32\Drivers\adp3hci.sys	This service keeps your Adobe Flash Player inst...				
adp320	MANUAL	SystemRoot\system32\drivers\adp320.sys					
adp320	MANUAL	SystemRoot\system32\drivers\adp320.sys					
AsiLockupSvc	MANUAL	SystemRoot\system32\drivers\asi.sys					
AFD	SYSTEM	SystemRoot\system32\drivers\afd.sys					
agp440	MANUAL	SystemRoot\system32\drivers\agp440.sys	Intel AGP Bus Filter				
ALG	MANUAL	SystemRoot\system32\alg.exe					
alide	MANUAL	SystemRoot\system32\drivers\alide.sys					
amd64	MANUAL	SystemRoot\system32\drivers\amd64.sys	AMD K8 Processor Driver				
amdppm	MANUAL	SystemRoot\system32\drivers\amdppm.sys	AMD Processor Driver				
andata	MANUAL	SystemRoot\system32\drivers\andata.sys					
andbx	MANUAL	SystemRoot\system32\drivers\andbx.sys					
andata	BOOT	system32\drivers\andata.sys					
AppHostSvc	AUTO	SystemRoot\system32\drivers\apphost.sys					
AppID	MANUAL	SystemRoot\system32\drivers\appid.sys					
AppIDSvc	MANUAL	SystemRoot\system32\drivers\appid.sys					
AppInfo	MANUAL	SystemRoot\system32\drivers\appid.sys					
AppMgr	MANUAL	SystemRoot\system32\drivers\appid.sys					
arc	MANUAL	SystemRoot\system32\drivers\arc.sys					
arc4	MANUAL	SystemRoot\system32\drivers\arc4.sys					
ASPNET_4.0.30...	DISABLED	aspnet_counters.dll	Provides support for out-of-process session state...				
aspnet_state	MANUAL	system32\DRIVERS\aspnetmon.sys					
atapi	BOOT	system32\drivers\atapi.sys	IDE Channel				
AudioEndpointMa...	AUTO	SystemRoot\system32\drivers\audio.sys					
AudioSrv	AUTO	SystemRoot\system32\drivers\audio.sys					
AVPIS.D	AUTO	"C:\Program Files\Avast\Avast GUI\aspm...	Provides computer protection against viruses, da...				
AvastV	MANUAL	SystemRoot\system32\drivers\avast.sys					
bcbdrv	MANUAL	SystemRoot\system32\drivers\bcbdrv.sys	Broadcom NetXtreme II VPD				
bcbdrv	MANUAL	system32\DRIVERS\bcbdrv.sys	Broadcom NetXtreme Gigabit Ethernet - NDIS 6.0				
BatC	MANUAL	C:\Windows\system32\drivers\batc.sys					
BDE SVC	MANUAL	SystemRoot\system32\drivers\bde.sys					
Beep	MANUAL	C:\Windows\system32\drivers\beep.sys	Beep				
BFE	AUTO	SystemRoot\system32\drivers\bfe.sys					

GMER 2.2.19882 WINDOWS 6.1.7601 Service Pack 1 x64

AntiVirus: http://www.avast.com

Exit





VIVA QUESTIONS (PRE LAB and POSTLAB):

1. List any two web security threats.
2. List any two design goals for a firewall.
3. Define Security Mechanism.
4. What is S/MIME?
5. Define IPSec.
6. What are the key features of SET?
7. What is the use of public key encryption scheme?
8. Identify the possible threats for RSA algorithm.
9. List out the general schemes for the distribution of public keys.
10. What are the areas where Kerberos Version 5 addresses the limitation of Version 4?

RESULT:

Thus the Rootkits tool was installed and its various options were verified successfully

Best wishes,
Dr. Mustafa Basthikodi