

Example: How many edges are there in a graph with 5 vertices and degree sequence (2,3,3,4,4)?

$$\text{degree sum} = 2 + 3 + 3 + 4 + 4 = 16$$

By Handshake Theorem, $16 = 2|E|$

$$\Rightarrow |E| = \frac{16}{2} = 8$$

There are 8 edges

Example: A graph with 8 edges has twice as many vertices of degree 3 as there are vertices of degree 2 (and no other vertices). How many vertices are of degree 2?

$|V_3|$ 8 edges
of vertices of degree 3: $2x$
of " " degree 2: x
 $|V_2|$

Find # of vertices of degree 2.

By Handshake theorem, the degree sum is $8 \cdot 2 = 16$.

$$16 = 3(2x) + 2(x)$$

$$16 = 6x + 2x = 8x \Rightarrow x = 2$$

\Rightarrow So there are 2 vertices of degree 2 and 4 vertices of degree 3.

Corollary: Every graph has an even number of vertices of odd degree.
(odd vertices: vertices of odd degree)
(even vertices: vertices of even degree)

Proof:

By the Handshake Theorem, $\sum_{u \in V} \deg(u) = 2|E|$

$$\text{Then } \sum_{\substack{u \in V \\ u \text{ odd vertex}}} \deg(u) + \sum_{\substack{u \in V \\ u \text{ even vertex}}} \deg(u) = 2|E|$$

$\underbrace{\sum_{\substack{u \in V \\ u \text{ odd vertex}}} \deg(u)}_{\substack{\text{each vertex} \\ \text{contributes} \\ \text{an odd number} \\ \text{to the sum}}} + \underbrace{\sum_{\substack{u \in V \\ u \text{ even vertex}}} \deg(u)}_{\substack{\text{even} \\ \text{number}}} = 2|E|$
138 (1)

two times something
this is an
even
number
(2)

and the sum has to be even because (1) and (2) are even.

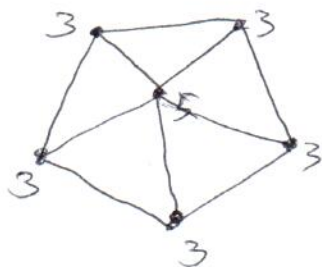
So there has to be an even number of vertices here (even number of odd vertices)

Exercise: Does there exist a simple graph with the following degree sequence?

If so, draw a picture of such a graph. If not, explain why.

(i): (1, 2, 2, 4, 5, 5)
 no such ^(simple) graph exists because ^{by} the corollary ~~requires that~~ there must be an even number of odd-degree vertices. This degree sequence has 3 vertices of odd degree.

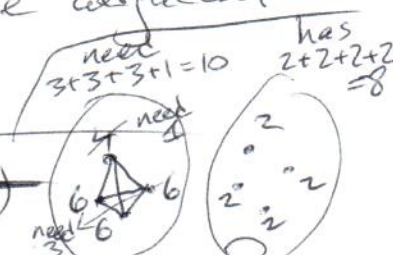
(ii): (3, 3, 3, 3, 3, 5)



Yes. There it is.

(iii): (0, 2, 2, 3, 4, 5) 6 vertices

There's a vertex v of degree 5, so v is adjacent with all other vertices, but there's a vertex of degree 0, which cannot be adjacent to v . So, this doesn't work.



~~(1, 1, 3, 3, 4, 6, 6, 6)~~ NO! ~~(1, 1, 3, 3, 4, 6, 6, 6)~~

similar idea group 4, 6, 6, 6 vertices

(2, 2, 2, 2, 4, 6, 6, 6) left too many edges to join to 2, 2, 2, 2

~~139~~ (2, 2, 2, 2, 2, 6, 6, 6) ?

NO! need 4+4+4=12 more degree



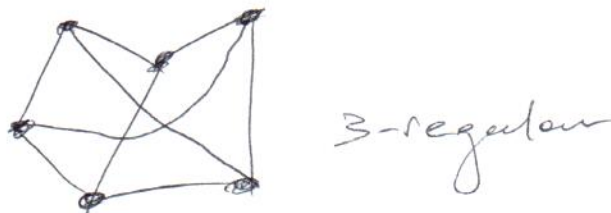
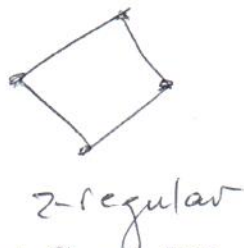
degree 6 vertices



can contribute 2+2+2+2+2=10 degree

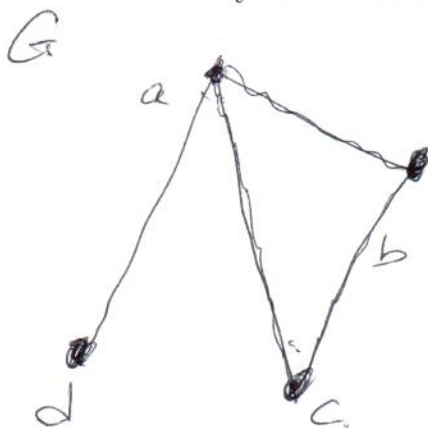
Some Special Graphs

A graph G is said to be k -**regular** if each vertex in G has degree k .



Let G and H be simple graphs. We say that H is a subgraph of G if $V(H) \subseteq V(G)$ and $E(H) \subseteq E(G)$.

Example: Let G be a graph with vertex set $V = \{a, b, c, d\}$ and edge set $E = \{ab, ac, ad, bc\}$. Find all subgraphs of the graph G with exactly 3 vertices and 2 edges:



subgraph:

vertices: a, b, c ①

a, b, d ②

a, c, d ③

b, c, d ④

① a, b, c

edges: ab, ac G_1

ab, bc G_2

ac, bc G_3

② a, b, d : vertices

edges: ab, ad G_4

③ a, c, d : vertices

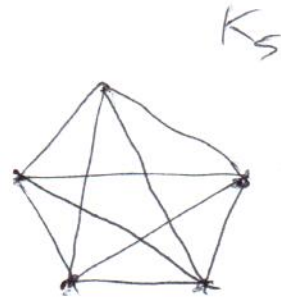
edges: ac, ad G_5

④ b, c, d : vertices

no subgraph with 2 edges

A complete graph K_n :

n vertices
each vertex adjacent
to all others



K_n has edges? $\binom{n}{2} = \frac{n \cdot (n-1)}{2}$
how many

each
vertex
has
degree $n-1$

A path P_n of length n :

A cycle C_n of length n :