## Solving Recurrence Relations

We want to find an explicit formula to a given recurrence relation rather than having to calculate each term of the sequence one by one.

Defn: A linear homogeneous recurrence relation of degree k with constant coefficients is a recurrence relation of the form

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \ldots + c_k a_{n-k}, \quad (\ddagger)$$

where  $c_1, c_2, \ldots, c_k \in \mathbb{R}$  and  $c_k \neq 0$ .

"Linear" means that  $a_{n-1}, a_{n-2}, \ldots, a_{n-k}$  appear in separate terms and to the first power.

"Homogeneous" means that there is no constant term.

"Degree k" denotes that the expression for  $a_n$  depends on the previous k terms  $a_{n-1}, a_{n-2}, \ldots, a_{n-k}$  (some of these terms may not appear, but  $a_{n-k}$  has to).

Such a recurrence relation is completely (and uniquely) determined by the values of the k initial conditions  $a_1 = t_1, a_2 = t_2, \ldots, a_k = t_k$ .

## Examples:

npies:
$$a_{n} = 7a_{n-1} + 8a_{n-2} - 3a_{n-4} + 2a_{n-5}$$
recurrence relation of degree 5

$$a_{n} = c_{n}a_{n-1} + c_{n}a_{n-2} + \dots + c_{n}a_{n-k}$$

$$c_{n} = c_{n}a_{n-1} + c_{n}a_{n-k} + c_{n}a_{n-k} + c_{n}a_{n-k}$$

$$c_{n} = c_{n}a_{n-k} + c_{n}a_{n-k} + c_{n}a_{n-k} + c_{n}a_{n-k} + c_{n}a_{n-k}$$

$$c_{n} \neq 0$$
Fact:  $a_{n} = r^{n}$  is a solution to the recurrence relation (‡) if and

only if

 $r^n = c_1 r^{n-1} + c_2 r^{n-2} + \ldots + c_k r^{n-k}.$ 

We divide both sides by  $r^{n-k}$  and then subtract the right hand side from the left, and we get

This equation is called the characteristic equation of the recurrence relation (‡).

So  $a_n = r^n$  is a solution to (‡) if and only if r is a solution of the characteristic equation.

The solutions of the characteristic equation are called the characteristic roots of the recurrence relation.

We will use characteristic roots to explicitly solve (some) given recurrence relations.

Solving linear homogeneous recurrence relations of degree 2 with constant coefficients:

We state a couple of very useful theorems without proof. (For the proofs you may consult most discrete mathematics textbooks – Rosen is a classic. These notes also cover the topic properly.)

Suppose a sequence  $\{a_n\} = a_1, a_2, a_3, \dots$  satisfies a recurrence relation  $a_k = Aa_{k-1} + Ba_{k-2}$  for all integers  $k \geq 3$ , where A and  $B \neq 0$ are some constants.

If the characteristic equation  $r^2 - Ar - B = 0$  has two distinct real roots  $\alpha$  and  $\beta$ , then  $\{a_n\}$  is given by the explicit formula

$$a_n = K_1 \alpha^n + K_2 \beta^n,$$

where  $K_1$  and  $K_2$  are chosen so that

$$a_1 = K_1 \alpha + K_2 \beta$$

$$a_2 = K_1 \alpha^2 + K_2 \beta^2.$$

then ...  $a_0 = K_1 \alpha^0 + K_2 \beta^0$   $= K_1 + K_2$   $a_1 = K_1 \alpha^1 + K_2 \beta^1$ 

**Example:** Solve the recurrence relation  $a_n = a_{n-1} + 2a_{n-2}$   $(n \ge 3)$  with the initial conditions  $a_1 = 0$  and  $a_n = 0$ 

with the limital conditions 
$$a_1 = 0$$
 and  $a_2 = 0$ .

$$a_1 = a_{1-1} + 2a_{1-2}$$

$$a_2 = c_{1-1} + 2c_{1-2}$$

$$a_3 = c_{1-1} + 2c_{1-2}$$

$$a_4 = c_{1-1} + 2c_{1-2}$$

$$a_5 = c_{1-1} + 2c_{1-2}$$

$$a_7 = c_{1-1} + 2c_{1-1}$$

$$a_7 = c_{1-1} + 2c_$$

**Example**: Solve the Fibonacci sequence  $F_n = F_{n-1} + F_{n-2}$   $(n \ge 3)$ ,  $F_1 = 1$  and  $F_2 = 2$ .

exercise

Theorem: (single root)

, could start at a too.

Suppose a sequence  $\{a_n\} = a_1, a_2, a_3, \dots$  satisfies a recurrence relation  $a_k = Aa_{k-1} + Ba_{k-2}$  for all integers  $k \geq 3$ , where A and  $B \neq 0$  are some constants.

If the characteristic equation  $r^2 - Ar - B = 0$  has a single (real) root  $\alpha$ , then  $\{a_n\}$  is given by the explicit formula

$$a_n = K_1 \alpha^n + K_2 n \alpha^n,$$

where  $K_1$  and  $K_2$  are chosen so that

$$a_1 = K_1 \alpha + K_2 \alpha$$

$$a_2 = K_1 \alpha^2 + 2K_2 \alpha^2.$$

**Example**: Solve the recurrence relation  $a_n = 6a_{n-1} - 9a_{n-2}$   $(n \ge 3)$  with the initial conditions  $a_1 = 9$  and  $a_2 = 36$ .

$$a_{n} = 6a_{n-1} - 9a_{n-2}$$

$$\frac{\Gamma^{n}}{\Gamma^{n-2}} = \frac{6\Gamma^{n-1} - 9\Gamma^{n-2}}{6\Gamma^{n-2}}$$

$$\frac{\Gamma^{n-2}}{\Gamma^{n-2}} = \frac{6\Gamma^{n-2}}{6\Gamma^{n-2}}$$

$$r^2 - 6r + 9 = 0$$
 $(r-3)^2 = 0$ 

$$a_n = K_1 \alpha^n + K_2 n \alpha^n$$

$$q = a_1 = K_1 \cdot 3^2 + K_2 \cdot 3^2 + K_3 \cdot 3^2 + K_3 \cdot 3^2 + 2 \cdot K_2 \cdot 3^2 + 2 \cdot K_3 \cdot 3^2 + 2$$

$$I: 9 = 3K_1 + 3K_2$$
  
 $-3I + II: 9 = 9K_2 = 7K_2 = 1$ 

108

$$7 a_n = 2 \cdot 3^n + 1 \cdot n \cdot 3^n$$
 $a_n = (2+n) \cdot 3^n$ 

formula, checke formula, checke your formula for Generalization to degree > k: The last two theorems can actually be generalized to higher degree linear homogeneous recurrence relations.

**Theorem**: (degree k with k distinct roots)

Suppose a sequence  $\{a_n\}$  satisfies a recurrence relation of degree k with the characteristic equation  $r^k - c_1 r^{k-1} - c_2 r^{k-2} - \ldots - c_k = 0$   $(c_1, \ldots, c_k \text{ are constants})$  with k distinct real roots  $\alpha_1, \ldots, \alpha_k$ , then  $\{a_n\}$  is given by the explicit formula

$$a_n = K_1 \alpha_1^n + K_2 \alpha_2^n + \ldots + K_k \alpha_k^n,$$

where the constants  $K_1, K_2, \ldots, K_k$  are determined by setting up a system of k linear equations using the values of the initial conditions  $a_1, a_2, \ldots, a_k$  (just like we did for degree 2 recurrences).

**Theorem**: (degree k, not all roots distinct)

Suppose a sequence  $\{a_n\}$  satisfies a recurrence relation of degree k with the characteristic equation  $r^k - c_1 r^{k-1} - c_2 r^{k-2} - \ldots - c_k = 0$   $(c_1, \ldots, c_k \text{ are constants})$  with j distinct real roots  $\alpha_1, \ldots, \alpha_j$  of positive multiplicities  $m_1, \ldots, m_j$   $(m_1 + \ldots + m_j = k)$ , then  $\{a_n\}$  is given by the explicit formula

$$a_n = (K_{1,1} + K_{1,2}n \dots + K_{1,m_1}n^{m_1-1})\alpha_1^n$$

$$+ (K_{2,1} + K_{2,2}n + \dots + K_{2,m_2}n^{m_2-1})\alpha_2^n$$

$$+ \dots + (K_{j,1} + K_{j,2}n + \dots + K_{j,m_j}n^{m_j-1})\alpha_j^n,$$

where the constants  $K_{i,j}$  are determined by setting up a system of k linear equations using the values of the initial conditions  $a_1, a_2, \ldots, a_k$  (just like we did for degree 2 recurrences).

## Derangements

**Problem:** There are n guests at a party. The guest list has the guests numbered 1, 2, ..., n. Each guest has a hat, which they leave at the coatroom when they enter. After the party, each guest takes a hat at random [well, don't ask me why, but you may find an interesting discussion here on the question of whether word problems in mathematics should actually make sense]. How likely is it that no guest gets their own hat?

Let us number the hats accordingly with 1, 2, ..., n such that the hat belonging to guest i is labelled i as well.

A <u>derangement</u> of  $\{1, 2, ..., n\}$  is a permutation  $\pi$  of  $\{1, 2, ..., n\}$  such that  $\pi(i) \neq i$  for each i  $(1 \leq i \leq n)$ .

Let us denote by  $d_n$  the number of derangements of  $\{1, 2, \ldots, n\}$ .

So we need to find  $d_n/n!$ .

Express  $d_n$  as a recurrence relation.

$$d_1 = 0, d_2 = 1$$
 $d_3 = 2$ 
 $d_4 = ?$ 
 $2 \cdot 143$ 
 $2 \cdot 341$ 
 $2 \cdot 4 \cdot 13$ 
 $3 \cdot 142$ 
 $3 \cdot 412$ 
 $3 \cdot 421$ 
 $4 \cdot 123$ 
 $43 \cdot 12$ 
 $43 \cdot 12$ 

$$d_1 = 0$$
 $d_2 = 1$  (2,1)

 $d_3 = 2$  (2,3,1), (3,1,2)

2 N	
m m	
forbidden 3 hart 4 is forbidden forbidden	
× × × × × × × × × × × × × × × × × × ×	
w w	
2 E 4	44
- d +	4
fortisher fortisher haden taken	hat taken Case 2

Forbid exactly one hat for each guest

We could have stated the same problem requiring that each guest.

We could have stated the same problem requiring that each guest.

We could have stated the same problem requiring that each guest has exactly 1 forbidden choice for hats, say f(i), and that each hat is forbidden for exactly one guest. (In our case for guest i the forbidden hat is the hat labelled i, so f(i) = i.)

Now assume that guest 1 takes hat  $i \neq f(1)$ . There are n-1 choices for i.

Consider guest j whose forbidden hat is i.

Two scenarios are possible, depending on whether or not guest j takes hat f(i) in return:

(i) Guest j takes hat f(i). Now the problem reduces to n-2 guests and n-2 hats where each guest has a forbidden hat (among the remaining hats) and each remaining hat is forbidden for exactly one remaining guest. This can be done in  $d_{n-2}$  ways.

(ii) Guest j does not take hat f(i). Now the problem reduces to n-1 guests and n-1 hats where each guest has a forbidden hat (forbidden hat for guest j is now f(i)) and each remaining hat is forbidden for exactly one remaining guest. This can be done in  $d_{n-1}$  ways.

We have found

$$d_n = (n-1)(d_{n-1} + d_{n-2}).$$
111

$$d_5 = (5-1)(9+2) = 44$$
  
 $d_6 = (6-1)(44+9) = 265$ 

We can now confirm that  $d_4 = 9$ ,  $d_5 = 44$ , etc.

Note that the recurrence relation we have found for  $d_n$  is not linear, so cannot apply the technique we learned earlier on for finding an explicit formula for  $d_n$ .

We'll take another approach.  $d_n - nd_{n-1} = -d_{n-1} + nd_{n-2} - d_{n-2}$ 

$$d_{n-1}(n-1)d_{n-2}$$

$$= -(d_{n-2}-(n-2)d_{n-3})$$

$$d_n = (n-1)(d_{n-1}+d_{n-2})$$

$$d_n = -(d_{n-1}-(n-1)d_{n-2})$$

$$d_n = -(d_{n-1}-(n-1)d_{n-2})$$
Observe that the expression on the right is the negative of the

Observe that the expression on the right is the negative of the expression on the left with n replaced by n-1; this helps to iterate easily.  $d_{n-1} - (d_{n-1} - (n-1)d_{n-2})$ 

$$d_{n} - nd_{n-1} = -(d_{n-1} - (n-1)d_{n-2})$$

$$= (-1)^{2}(d_{n-2} - (n-2)d_{n-3})$$

$$= (-1)^{3}(d_{n-3} - (n-3)d_{n-4})$$

$$= (-1)^{3}(d_{n-3} - (n-3)d_{n-4})$$

$$= (-1)^3 (d_{n-3} - (n-3)d_{n-4})$$

$$\Rightarrow d_n - nd_{n-1} = (-1)^{n-2}(d_2 - 2d_1)$$
With  $d_1 = 0$  and  $d_2 = 1$ , we have

$$\Rightarrow d_n - nd_{n-1} = (-1)^{n-2} = (-1)^n. \qquad (-1)^{n-2} = (-1)^{n-2}. (-1)^n$$

This is better, but we still don't have an explicit formula. We make the following claim.

$$d_n=n!\sum_{m=0}^nrac{(-1)^m}{m!}$$
 ,  $\Lambda$   $\Xi$   $I$ 

112