

# **Robust validation of network designs under uncertain demands and failures**

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# Validating network design

- Network design today is ad-hoc, and validating design is usually an afterthought
  - Contrast: Tools for chip and software industry a \$10B business [McKeown, 2012]
- Much progress on verification of network data plane (e.g., reachability, security policy)
  - HSA, Veriflow, Batfish, NoD, etc.
- Our goal: Validating quantitative network properties
  - Formal approach to guarantee network performance (e.g., bandwidth, link utilization)
    - Under diverse failure/traffic scenarios
    - Use the formal approach to inform network design

# Why is network validation hard? (1)

- **Scenarios of interest are too many**
  - Exponentially many failure scenarios [Wang et al., Sigcomm '10, Liu et al., Sigcomm '14]
    - E.g., All possible simultaneous  $f$  link failures
    - All possible traffic demands – non-enumerable

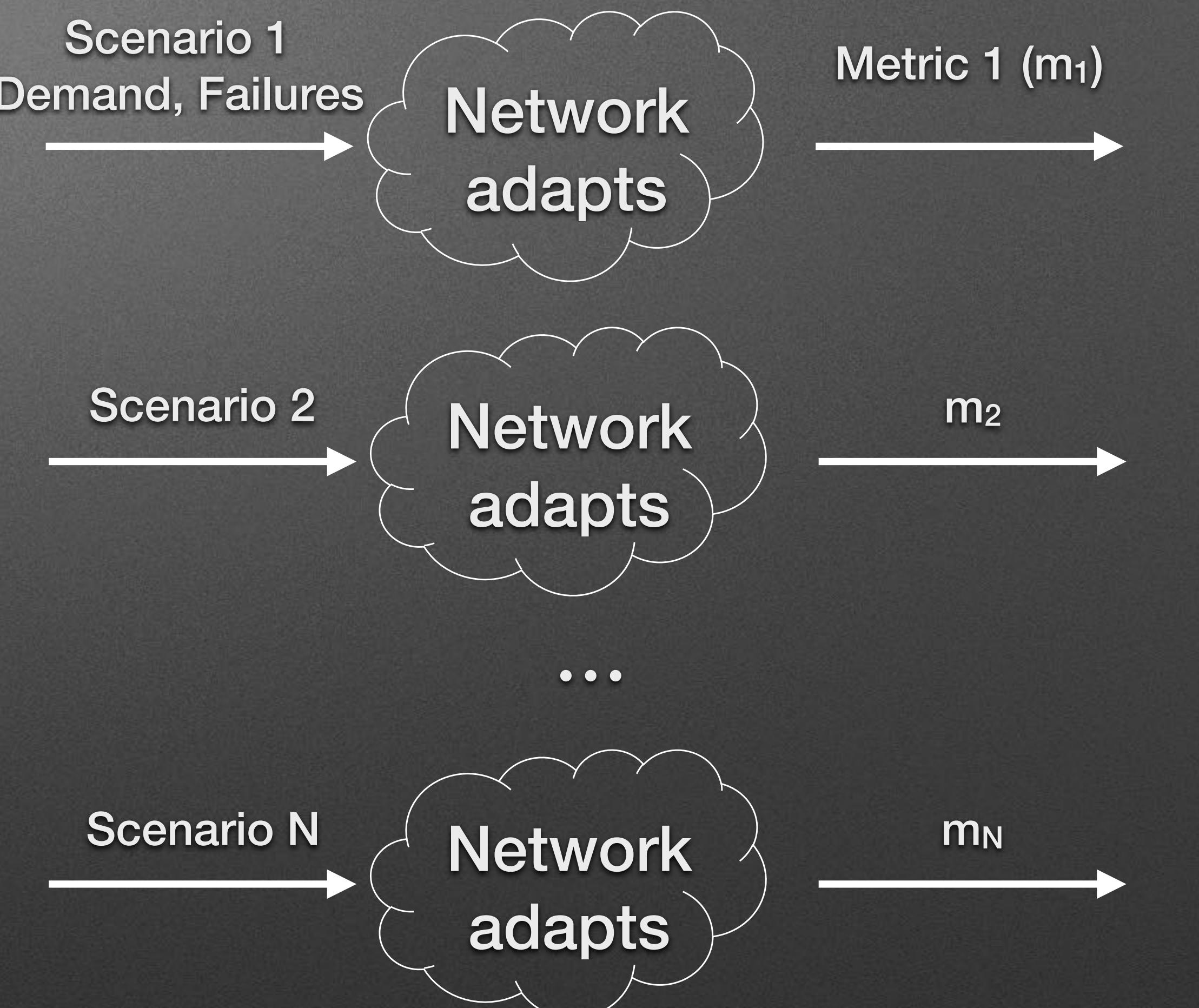
# Why is network validation hard? (2)

- **Adaptation makes the problem intractable**
- Networks increasingly agile and flexible in adaptation
  - E.g., SDNs and NFVs
- Tools exist to bound worst case performance
  - E.g., robust optimization, and oblivious routing  
[Applegate et al., Sigcomm '03]
  - Assume networks do not adapt, or consider limited forms of adaptation to make problem tractable



# Our work

- General framework for network validation
  - Find the **worst** performance of the network across **all** scenarios assuming network can adapt in **best** fashion for each scenario
  - Handles intractable problems drawing on cutting-edge optimization technique
  - Applies to **network synthesis**



# Example: Failure validation

Uncertainty Set	Adaptations	Performance metric
• All $f$ or fewer link failures	• Flexible rerouting (multi-commodity flow)	• Utilization of most congested link

## Problem:

- Given up to  $f$  links may simultaneously fail, what is the worst-case utilization of any link across all failure scenarios?

# Formal formulation of a network validation problem

$$\max_{x \in X} \min_{y \in Y(x)} F(x, y)$$

Uncertainty Set      Adaptations      Performance metric  
Less is better

```
graph TD; X[x ∈ X] --> Min[min<br/>y ∈ Y(x)]; Min --> F[F(x, y)]; F -- "Less is better" --> Max[Max<br/>x ∈ X];
```

## Example: Validation under failures

X	Set of failures
Y(x)	Feasible routing of demands under given failure
F(x, y)	Utilization of most congested link

Inner problem: For a fixed scenario - Easy to compute online (LP)  
E.g., multi-commodity flow

Outer problem: Potentially hard since large number of scenarios

# Wide applicability of framework

## Uncertainty Set

- All  $f$  or fewer link failures
- Shared risk link group
- Weighted averages of historical demands

## Adaptations

- Flexible rerouting (multi-commodity flow)
- Rerouting constrained to pre-selected tunnels
- Constrain with middlebox traversal requirements

## Performance metric

- Utilization of most congested link
- Bandwidth of business critical applications

# Reformulating the problem

$$\max_{x \in X} \quad \min_{y \in Y(x)} \quad F(x, y)$$



LP dualization

$$\max_{\lambda, v, x} \quad F'(\lambda, v, x)$$

# Failure validation: Formulation

$$\max_{v, \lambda, x} \sum_{t, i \neq t} d_{it} (v_{it} - v_{tt})$$

$$s.t. \quad v_{it} - v_{jt} \leq \lambda_{ij} \quad \forall t, \langle i, j \rangle \in E$$

$$\sum_{\langle i, j \rangle \in E} \lambda_{ij} c_{ij} (1 - x_{ij}^f) = 1$$

$$x^f \in X$$

$$x_{ij}^f \in \{0, 1\}; \quad \lambda_{ij} \geq 0; \quad \langle i, j \rangle \in E$$

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$$x^f \in X$$
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Depends on failure model of interests  
• E.g. simultaneous  $f$  link failures

$$\sum_{i, j} x_{ij}^f = f$$

# Failure validation: Formulation

$$\max_{v, \lambda, x} \sum_{t, i \neq t} d_{it} (v_{it} - v_{tt})$$

$$s.t. \quad v_{it} - v_{jt} \leq \lambda_{ij} \quad \forall t, \langle i, j \rangle \in E$$

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$$x^f \in X$$

$$x_{ij}^f \in \{0, 1\}; \quad \lambda_{ij} \geq 0; \quad \langle i, j \rangle \in E$$

Can be converted to mixed-inter linear program.  
In general, validation problems could be non-linear.

# Solution approach

- Focus on upper bounds (**relaxation**)
  - Intractable problems – hard to solve to optimality
  - Upper bounds sufficient for validation use
- Goal: Develop a general approach
  - Applicable to diverse validation problems (e.g., validating failures, demands...)
  - Yet, amenable to problem-specific structure
- Use cutting-edge techniques from non-linear optimization

# Tractable relaxations: RLT

- RLT relaxations: general approach to relax non-convex problems into tractable LPs
  - Family of relaxations
  - Higher levels of hierarchy
    - Converge to optimal value of the non-convex problem
    - Incur higher complexity
- For scalability, focus on the first level

# RLT relaxation: example

$$\begin{array}{ll} \min_{x,y} & \boxed{xy} - x + y \rightarrow z - x + y \\ & \boxed{x - 2 \geq 0; \quad y - 3 \geq 0} \rightarrow \boxed{xy} - 2y - 3x + 6 \geq 0 \\ & 3 - x \geq 0; \quad 4 - y \geq 0 \end{array}$$

↓  
 $z$

Relaxation steps:

1. Multiply constraints with each other
2. Replace products of variables  $xy, x^2, y^2$  by new variables

# Our results on effectiveness of RLT

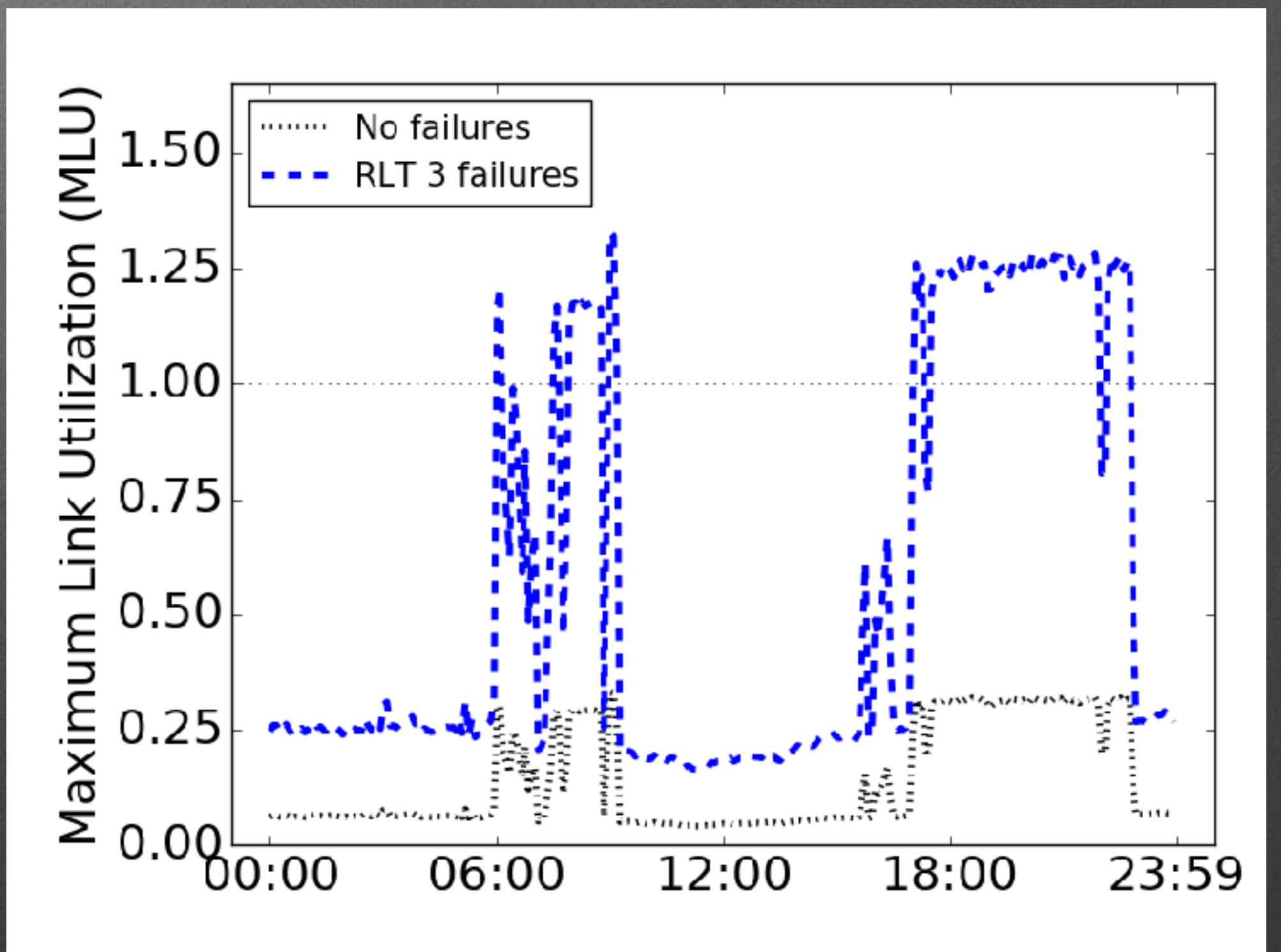
- Compare RLT with two theoretical benchmarks
  - Both bound worst case performance across failures/demands, but with limited network adaptation
  - Oblivious routing [Applegate, et al., Sigcomm '03; Wang, et al., Sigcomm '06, etc.]
  - Affine adaptation: a generalization of oblivious routing, studied in robust optimization
- Our results show
  - First-level RLT dominate oblivious/affine adaptations
  - Better results possible by exploiting problem-specific structure combined with RLT

# Evaluation

- **Real topologies**
  - Abilene, GEANT, and ANS (from The Internet Topology Zoo)
- **Real and synthetic traffic matrices**
  - Real trace: 6-month end-to-end demand on Abilene
  - Synthetic: Gravity model

# Results: Effectiveness of RLT

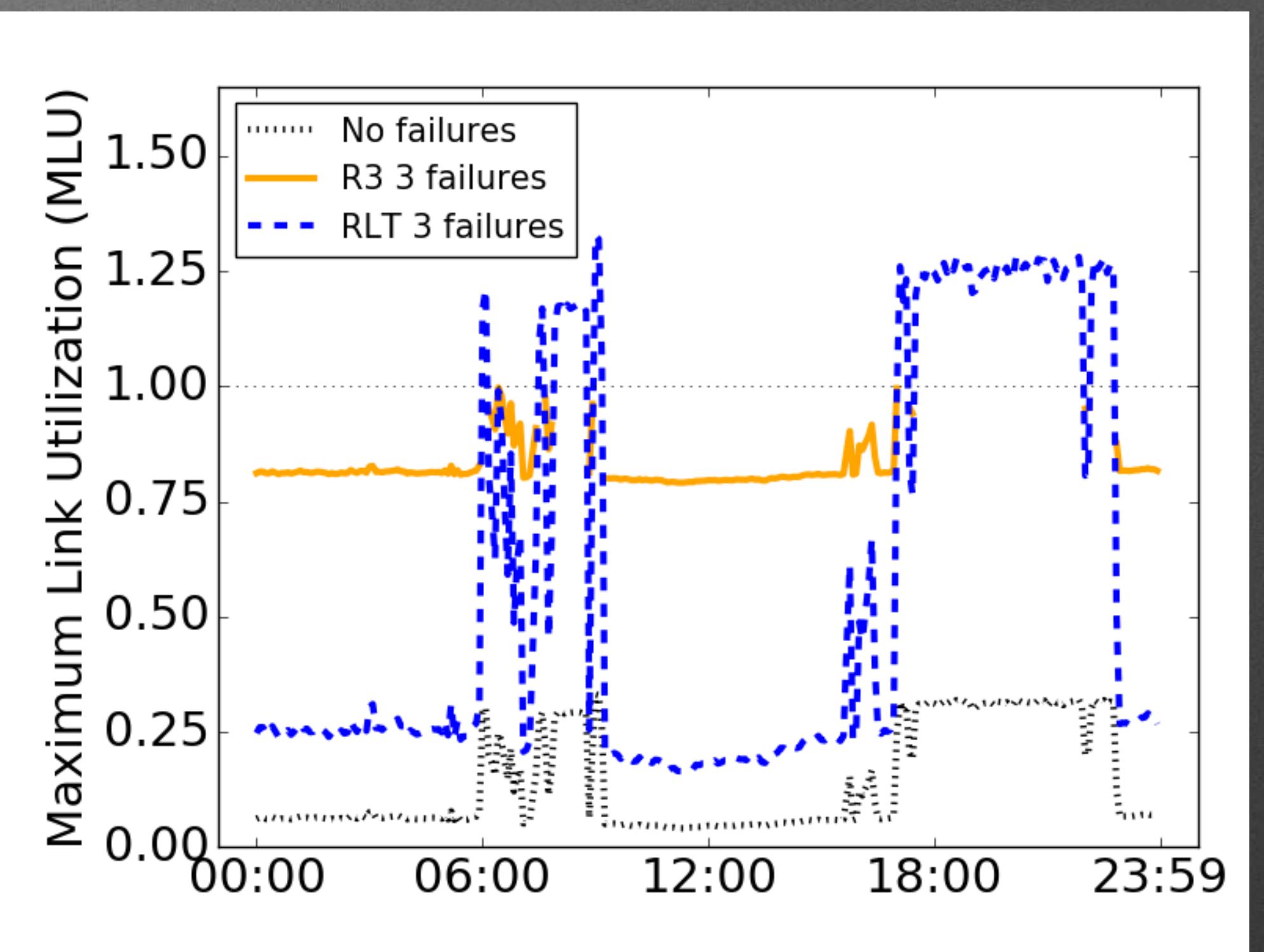
- Compare maximum link utilization (MLU)
  - The optimal IP scheme vs. our RLT relaxation LP
  - **RLT matches optimal in all our experiments**



Abilene Network – 3 link failures

# Results: Effectiveness of RLT

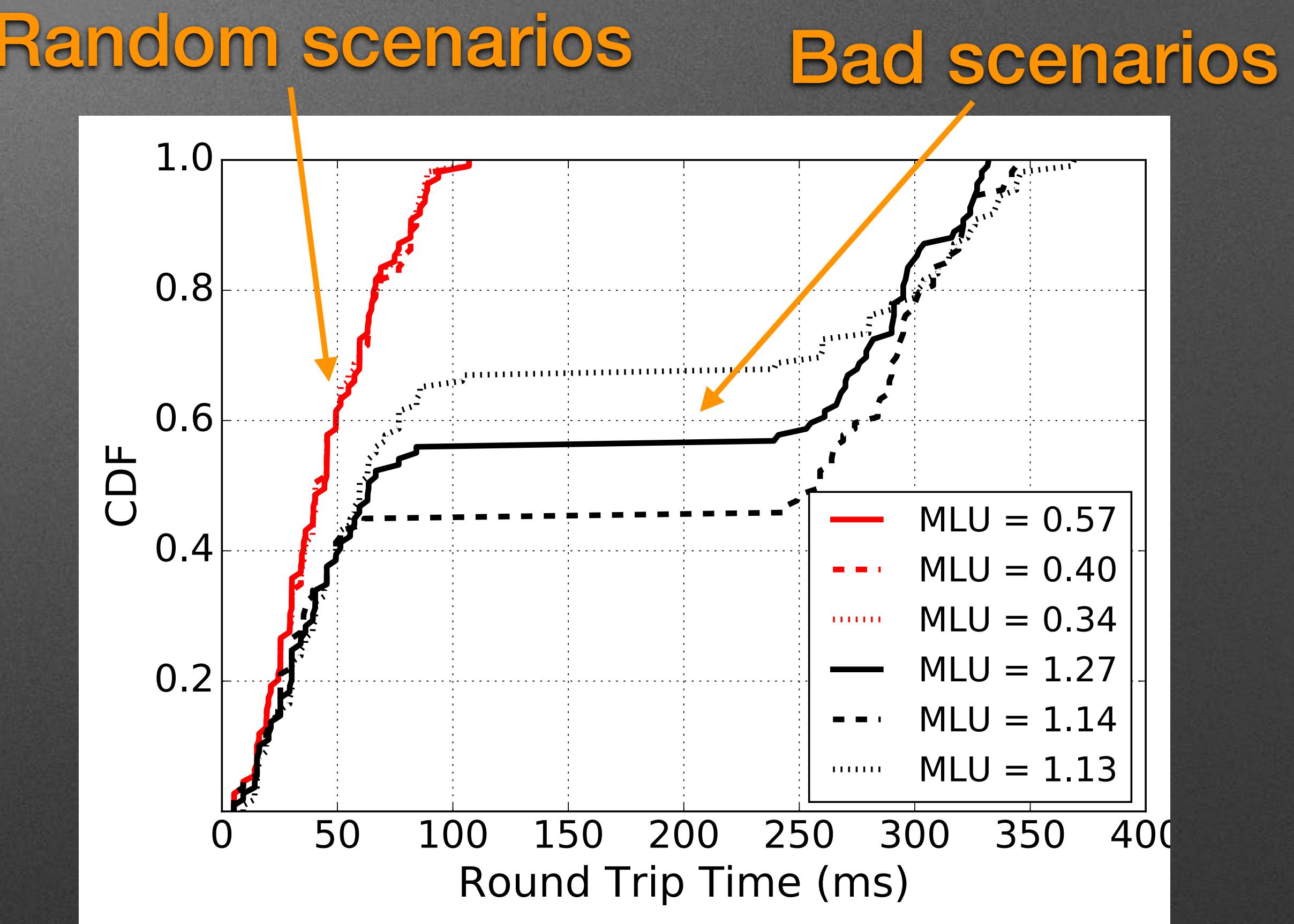
- Compare with R3 [Wang et al., Sigcomm '10]
  - Determines if  $\text{MLU} < 1$  under  $f$  failures
  - Gives a valid bound only when  $\text{MLU} < 1$
  - Based on oblivious approach
- Our result
  - First-level RLT dominates R3 whenever R3 provides a valid bound
- Other advantages of our approach
  - Useful to detect bad failure scenarios, and the amount of exceeded link capacity
  - Generalizes to other validation problems



Abilene Network – 3 link failures

# Using framework to detect bad failures

- Framework allows finding failures that impact the network the most
  - Random search not efficient
  - Only 0.05% of 3-failure scenarios are bad ( $MLU > 1$ )
- Emulate to understand latency behavior



Emulated Abilene traffic matrix  
with Mininet, and ONOS controller

# Results: running time

- RLT relaxation LP vs. optimal IP (IP run for 2 hours)
- On scaled GEANT network (32 nodes, 1000 edges), 3 link failures:
  - RLT finished in **608** seconds, whereas IP finished in **3890** seconds
  - Only 60% of the IP instances completed in 2 hours
- Our RLT relaxation LP doesn't degrade with larger number of failures

# Example: Tunnel selection validation

## Uncertainty Set

- All  $f$  or fewer link failures
- Shared risk link group
- Weighted averages of historical demands

## Adaptations

- Flexibly rerouting (Multi-commodity flow)
- Rerouting constrained to pre-selected tunnels
- Constrain with middlebox traversal requirements

## Performance metric

- Utilization of most congested link
- Bandwidth of business critical applications

## Problem:

- For a given choice of tunnels, are utilizations of all links across all traffic demands of interest within acceptable limits?

# Tunnel selection: Results

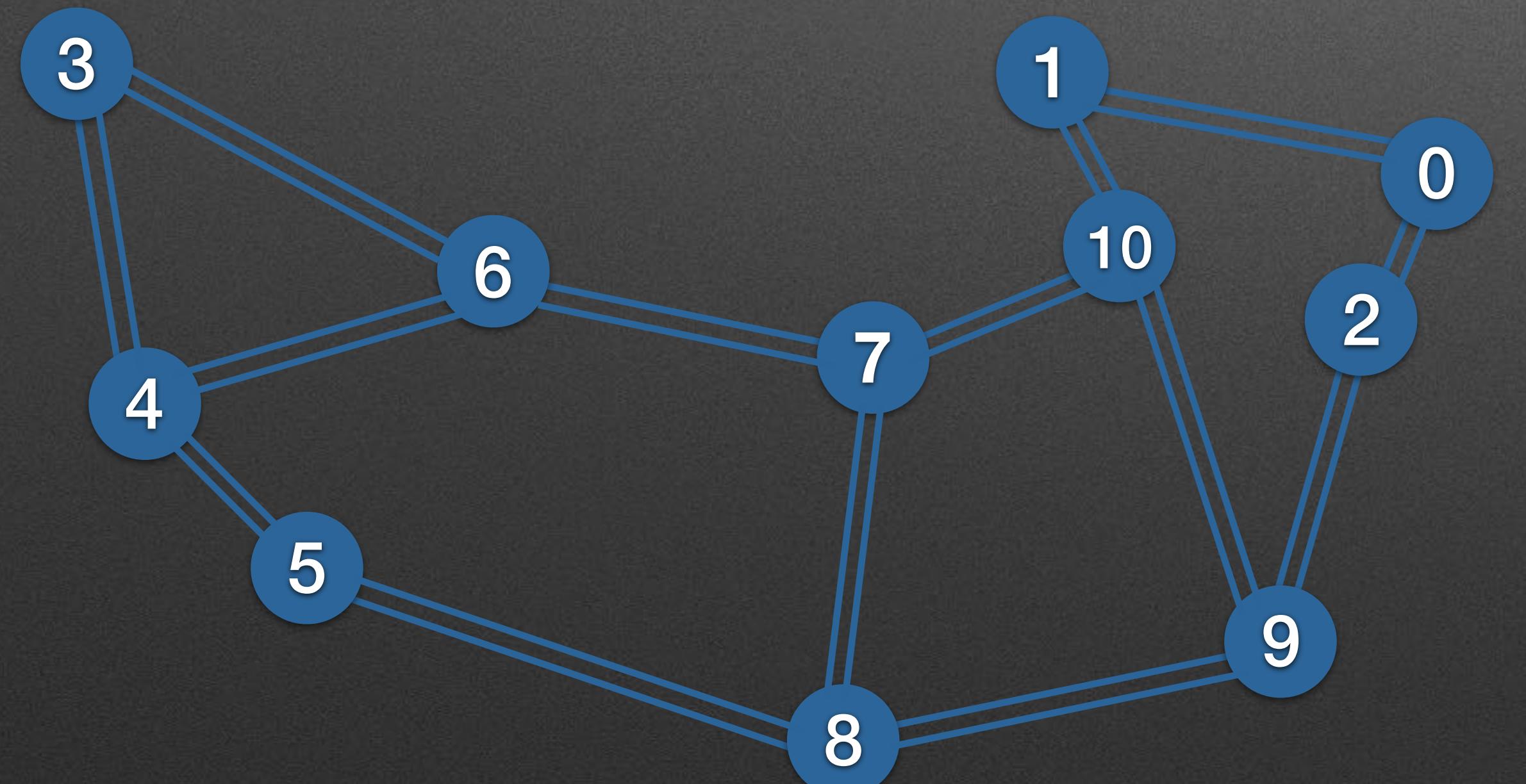
- **Predicted demand:** weighted averages of historical matrices
  - Validation problem is an LP
  - On Abilene: **First-level RLT achieves optimal MLU**
- **Widely-used tunnel selection heuristics may perform poorly**
  - E.g., K-shortest (SWAN, Sigcomm '13), Shortest-Disjoint heuristics
  - More robust tunnel selection heuristic performs much better

# Synthesizing valid designs

- Validation is a stepping stone for synthesis
- Example: Optimal Capacity Augmentation
  - Incrementally add capacity to existing links
  - Minimizing cost of adding capacity
  - Ensure resulting network can handle all failure scenarios
- One can use our framework for synthesis in 2 ways:
  - 1) Get conservative solution, with a single LP
  - 2) Iterative approach, which gives a lower bound on cost at each step

# Capacity augmentation: Abilene

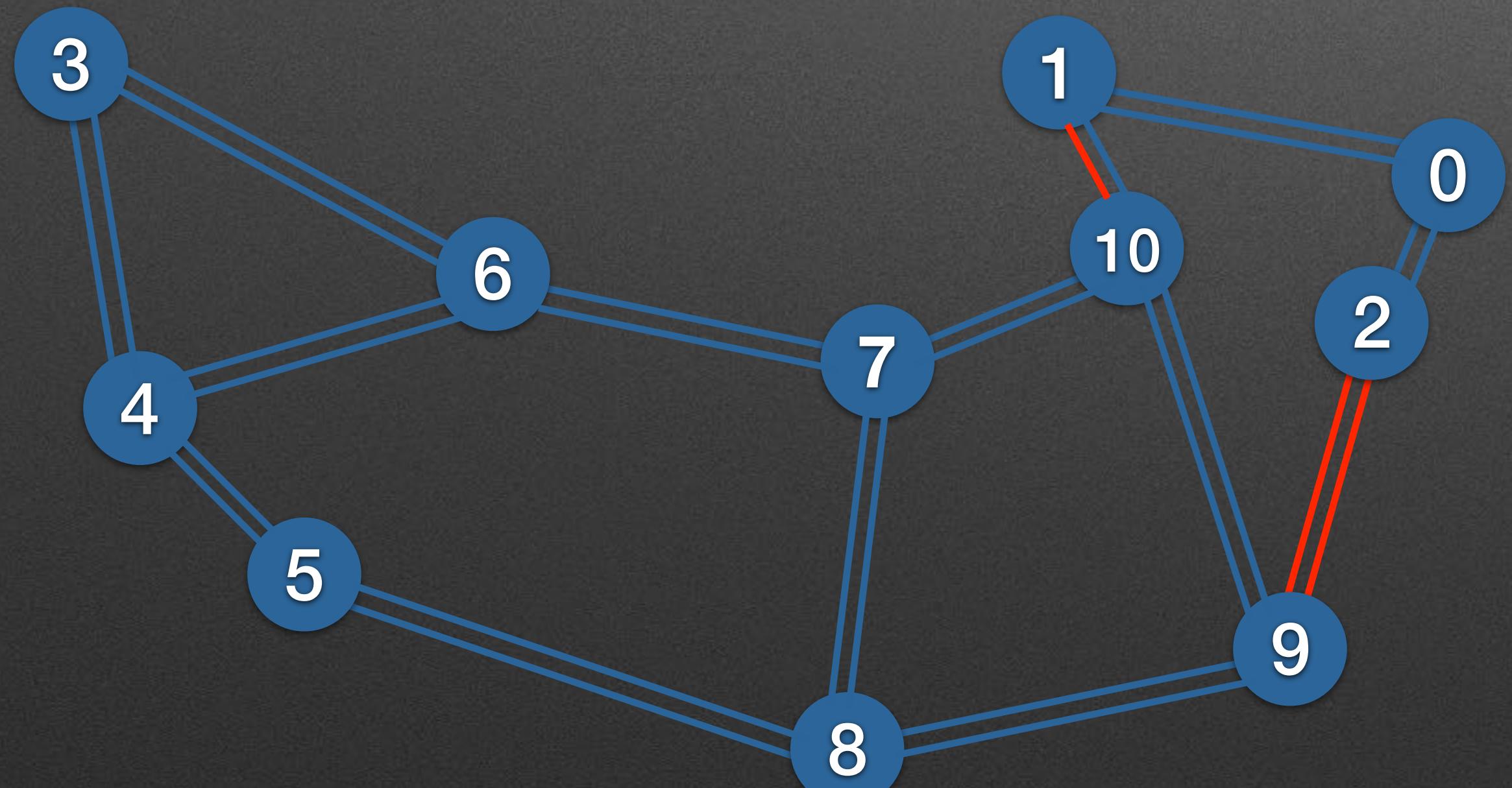
- Validate if  $MLU \leq 1$ .
- If not, run augmentation LP with counter examples



Step	Counter examples	MLU	Links Augmented	Total new capacity (Gbps)

# Capacity augmentation: Abilene

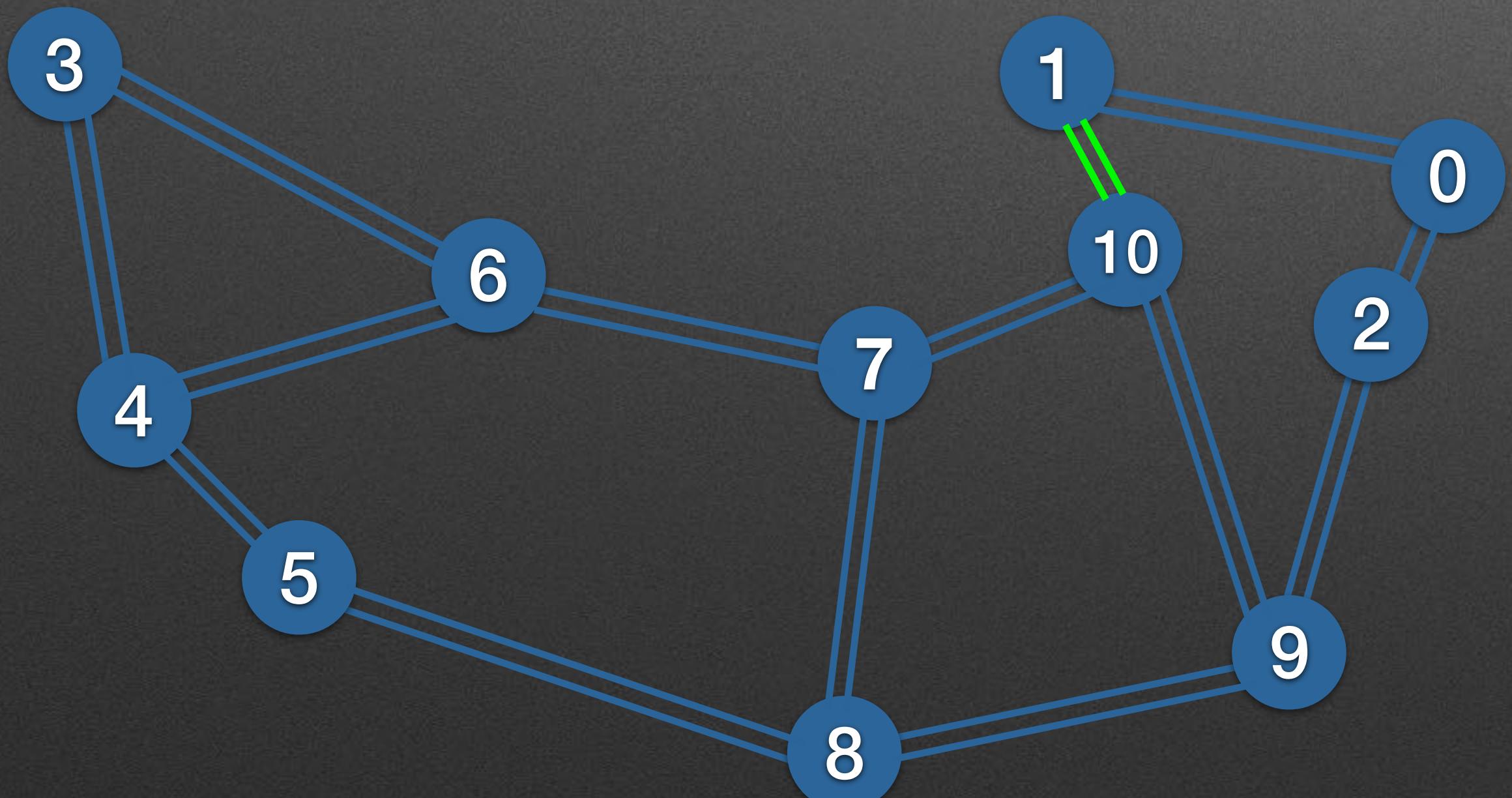
- Validate if  $MLU \leq 1$ .
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Step	Counter examples	MLU	Links Augmented	Total new capacity (Gbps)
1	(1, 10), (2, 9)	1.274		

# Capacity augmentation: Abilene

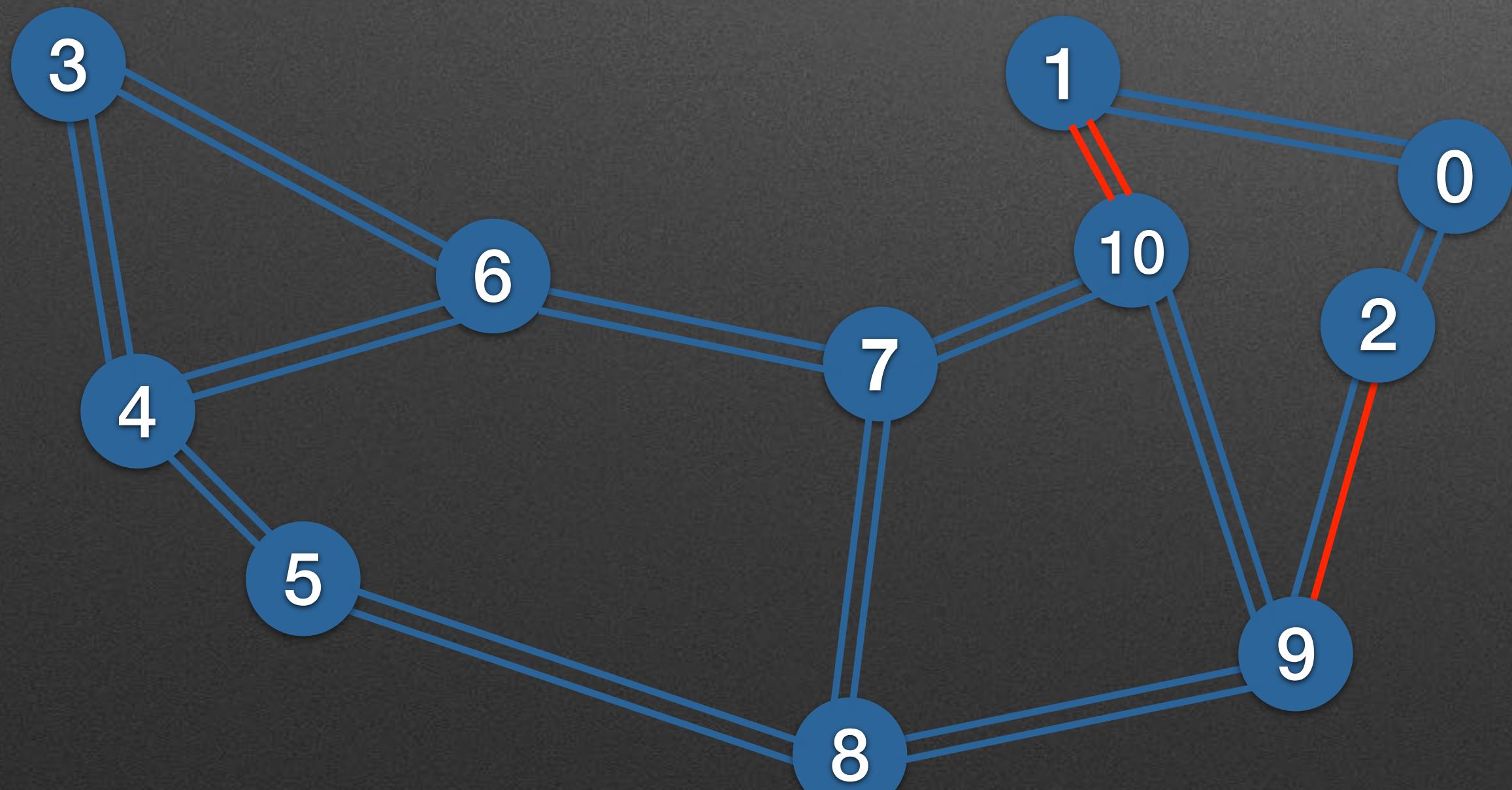
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Step	Counter examples	MLU	Links Augmented	Total new capacity (Gbps)
1	(1, 10), (2, 9)	1.274	(1, 10)	2.744

# Capacity augmentation: Abilene

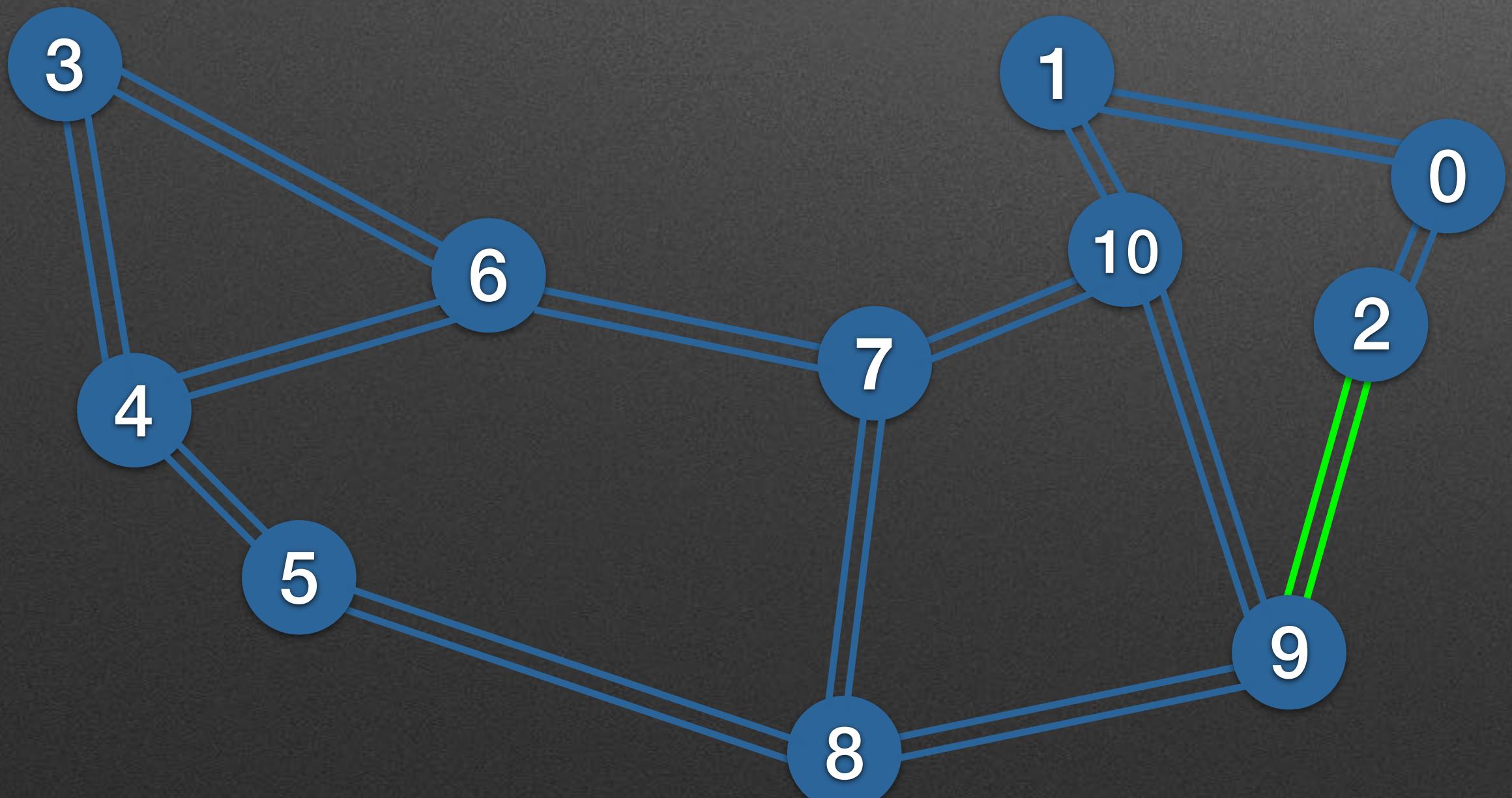
- Validate if  $MLU \leq 1$ .
- If not, run augmentation LP with counter examples



Step	Counter examples	MLU	Links Augmented	Total new capacity (Gbps)
1	(1, 10), (2, 9)	1.274	(1, 10)	2.744
2	(2, 9), (10, 1)	1.274		

# Capacity augmentation: Abilene

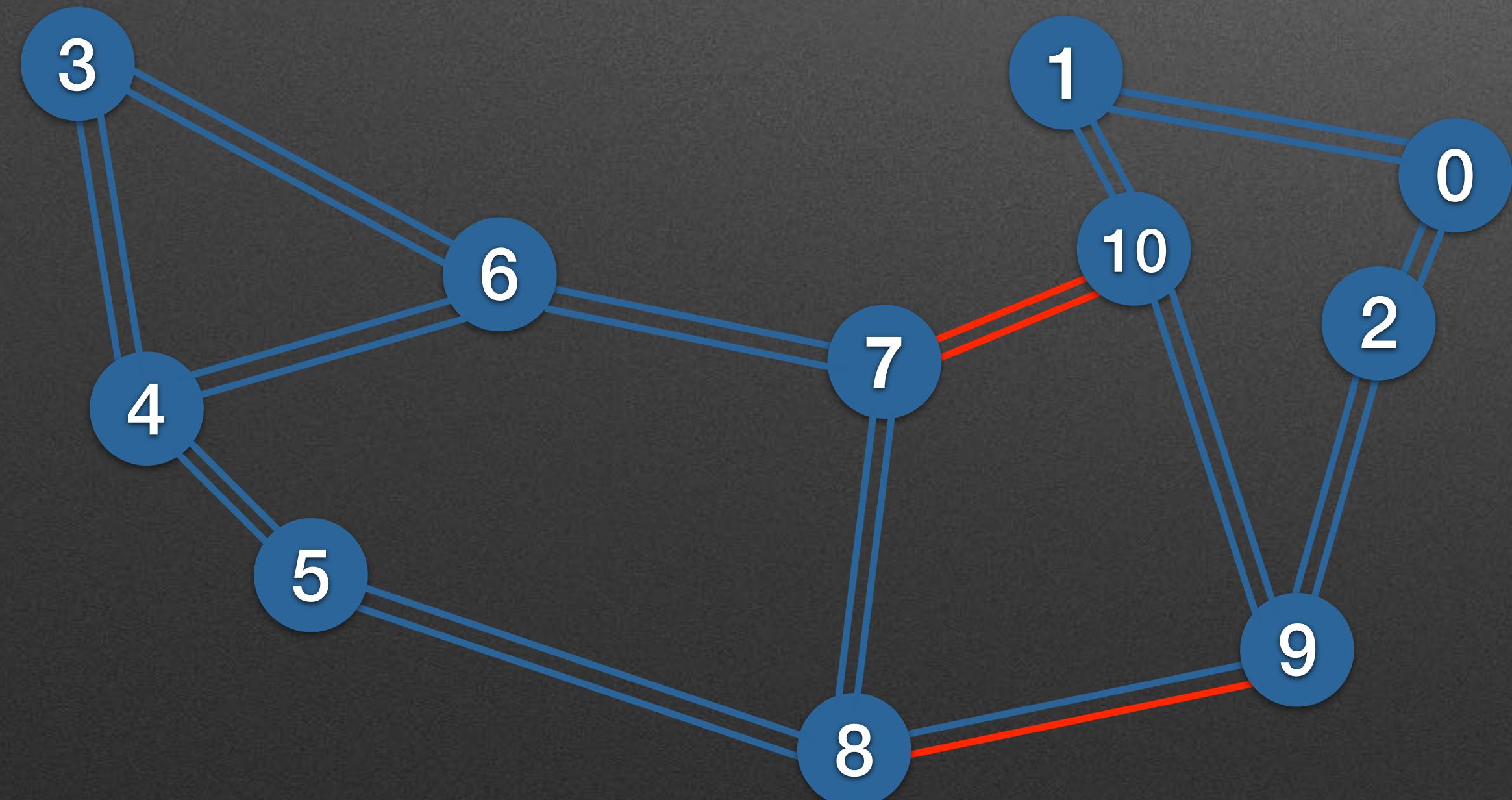
- Validate if  $MLU \leq 1$ .
- If not, run augmentation LP with counter examples



Step	Counter examples	MLU	Links Augmented	Total new capacity (Gbps)
1	(1, 10), (2, 9)	1.274	(1, 10)	2.744
2	(2, 9), (10, 1)	1.274	(2, 9)	5.488

# Capacity augmentation: Abilene

- Validate if  $MLU \leq 1$ .
- If not, run augmentation LP with counter examples



Step	Counter examples	MLU	Links Augmented	Total new capacity (Gbps)
1	(1, 10), (2, 9)	1.274	(1, 10)	2.744
2	(2, 9), (10, 1)	1.274	(2, 9)	5.488
3	(9, 8), (10, 7)	1.217		

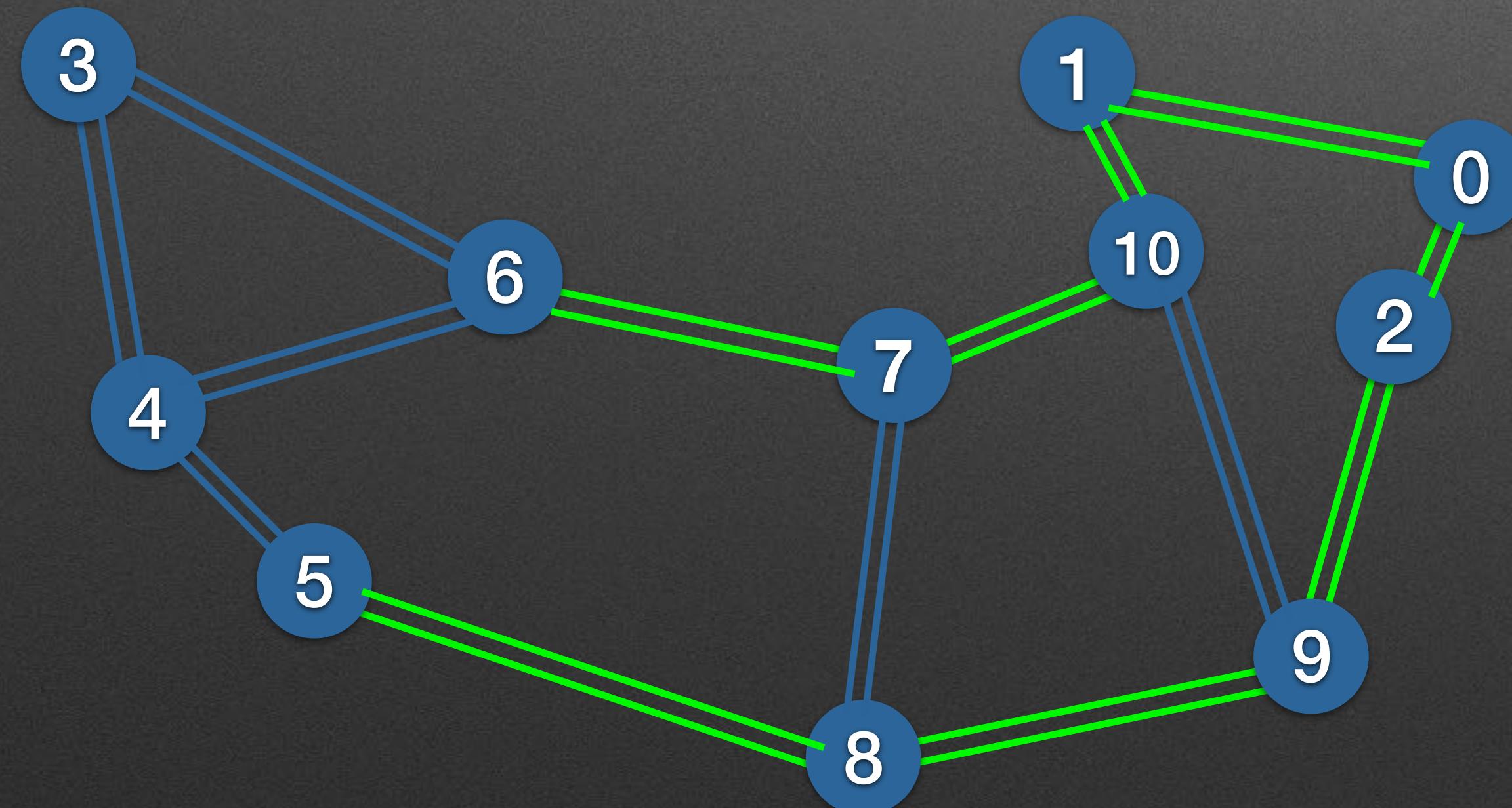
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Step	Counter examples	MLU	Links Augmented	Total new capacity (Gbps)
1	(1, 10), (2, 9)	1.274	(1, 10)	2.744
2	(2, 9), (10, 1)	1.274	(2, 9)	5.488
3	(9, 8), (10, 7)	1.217	(9, 8)	7.653

# Capacity augmentation: Abilene

- Validate if  $MLU \leq 1$ .
- If not, run augmentation LP with counter examples



Step	Counter examples	MLU	Links Augmented	Total new capacity (Gbps)
1	(1, 10), (2, 9)	1.274	(1, 10)	2.744
2	(2, 9), (1, 10)	1.274	(2, 9)	5.488
3	(9, 8), (10, 7)	1.217	(9, 8)	7.653
4	(10, 7), (9, 8)	1.217	(10, 7)	9.818
5	(0, 2), (1, 10)	1.192	(0, 2)	11.743
6	(1, 0), (1, 10)	1.071	(1, 0)	12.452
7	(7, 6), (8, 5)	1.006	(7, 6)	12.509
8	(8, 5), (7, 6)	1.006	(8, 5)	12.566
9	—	1.000	—	—

# Conclusions

- Early effort at formally verifying quantitative network properties under uncertainty
- Generic framework for a wide class of network validation problems
- Modeling adaptivity results in intractable problems
  - RLT relaxations promising
  - Tighter bounds than oblivious
  - Exact in multiple failures case and predicted demand case
- Validation framework enables network synthesis

**Thanks!  
Questions?**