

TUTORIAL: 3

3.1

1. Write linear search pseudocode to search an element in a sorted array with maximum comparisons,

```
for (i=0; i<length; i++)  
{  
    if (key == A[i])  
    {  
        swap (A[i], A[i-1])  
        return i-1;  
    }  
}
```

2. Pseudocode for iterative and recursive in insertion sort. Insertion sort is called online sort why?

Iterative

```
Void InsertionSort (int i=1; i<n; i++)  
{  
    j = i-1;  
    x = A[i];  
    while (j < -1 && A[j] > x)  
    {  
        A[j+1] = A[j];  
        j--;  
    }  
    A[j+1] = x;  
}
```

Recursive Insertion Sort

```
void Insertion (int arr[], int n)
```

```
{ if (n <= 1)
  return;
```

```
  Insertion (arr, n-1);
```

```
  int last = arr[n-1];
```

```
  j = n-2;
```

```
  while (j >= 0 && arr[j] > last)
```

```
  { arr[j+1] = arr[j];
    j--;
```

```
  }
```

```
  arr[j+1] = last;
```

```
}
```

In Insertion sort, we give input one by one and place each one at right order with comparison from already traced element we need the whole array simultaneously that's why it's called an online sorting algorithm.

3. Complexity of all sorting algorithms. 3.2

	<u>Best Case</u>	<u>Worst Case</u>
Bubble Sort	$O(n)$	$O(n^2)$
Selection Sort	$O(n^2)$	$O(n^2)$
Insertion Sort	$O(n)$	$O(n^2)$
Quick sort	$O(n \log n)$	$O(n^2)$
Merge Sort	$O(n \log n)$	$O(n \log n)$
Count Sort	$O(n+k)$	$O(n+k)$

4. Divide all sorting algorithms into 'inplace/stable/online sorting'.

Online sorting :- Insertion Sort (Partial)

Stable Sorting :- Merge sort, Insertion sort

Bubble sorting :-

Bubble sort, Inserting sort, Selection sort.

5. Write iterative / recursive pseudo code for binary search. What is the time and space complexity of linear and Binary search.

Iterative Binary Search

```
while (low <= high)
{
    int mid = (low + high) / 2;
    if (arr[mid] == key)
        return mid;
    else if (arr[mid] > key)
        high = mid - 1;
    else
        low = mid + 1;
}
```

Iterative Binary Search

T.C :- $O(\log n)$

S.C :- $O(1)$

Recursive Binary

Time complexity :- $O(\log n)$

Space Complexity :- $O(\log n)$

Linear search

Time complexity :- $O(n)$

space complexity :- $O(1)$

6. write recurrence relation for binary search

$$T(n) = T(n/2) + 1$$

$$T(1) = 1$$

7. Find two ~~to~~ indexed such that $A[i] + A[j] = k$ in ~~max~~ minimum time complexity.

```
int find-sumPair(int A[], int n, int k)
{
```

```
    Sort(A, n);
```

```
    i = 0;
```

```
    j = n - 1;
```

```
    while (i < j)
```

```
    { if (A[i] + A[j] == k)
```

```
        return 1;
```

```
    else if (A[i] + A[j] < k)
```

```
        i++;
```

```
    else
```

```
        j--;
```

```
    }
```

```
    return -1;
```

```
}
```

8 Which sorting is best for practical use?
Explain.

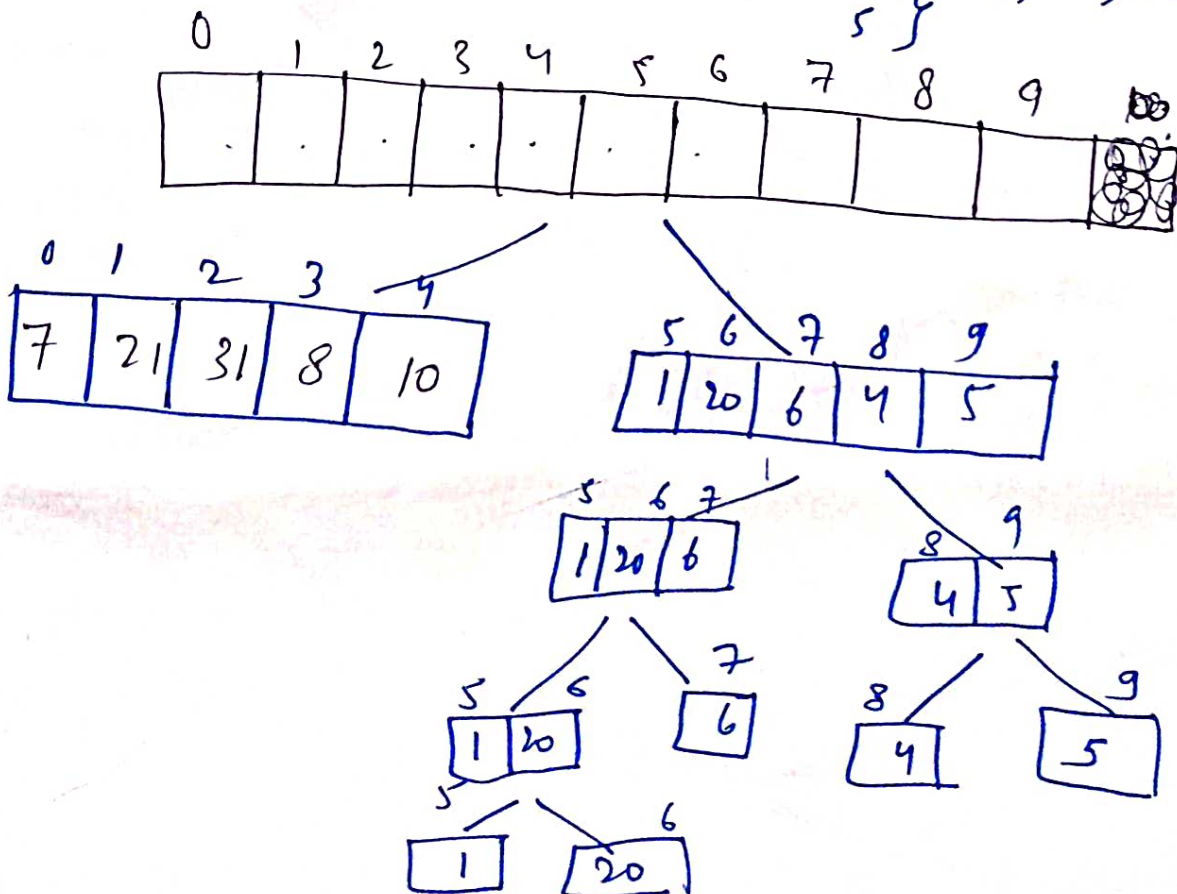
Quick sort is the fastest general purpose sort.
In most practical situations, quick sort is a method of choice.

If stability is important and space is available, merge sort might be best.

q. What do you mean by inversions?

Inversion count of an array suggests how close array is from being sorted.

Given :- Array $arr[] = \{7, 21, 31, 8, 10, 1, 20, 6, 4, 5\}$



Q. In which cases Quick sort will give best and worst time complexity. 3.4

Worst case: Time complexity of Quick sort is $O(n^2)$ it occurs when picked pivot is always an extreme (smallest or largest) element. This happens when input array is sorted or reverse sorted and either first or last element is picked.

Best case: - Best case time complexity of Quick sort is $O(n \log n)$ it occurs when pivot element is middle element always for the partitioning process.

11. Quick SORT

Worst Case :-

$$\begin{aligned} T(0) &:- \\ T(0) &= T(1) = 0 \\ T(N) &= N + T(N-1) \\ &\rightarrow O(n^2) \end{aligned}$$

$$\begin{aligned} \text{Best Case :-} \\ T(0) &= T(1) = 0 \\ T(N) &= 2T(N/2) + N \\ &\rightarrow O(n \log n) \end{aligned}$$

Merge SORT

$$T(n) = 2T\left(\frac{n}{2}\right) + n$$

<u>Basis</u>	Quick Sort	Merge Sort
Portion works well on	splitting is done in any relation. Smaller array	Array is partitioned in just 2 halves Fine on any size of array.
Addition & space	less (inplace)	More (not inplace)
Sorting method	Internal	External
Stability	Unstable	Stable

12 Selection sort isn't stable by default but can you write a stable ~~sorted~~ version of selection sort.

Selection sort can be made stable if instead of swapping, minimum element is placed in its position without swapping i.e. by placing the number in its position by pushing every element one step forward.

13 Bubble sort scans whole array even when array is sorted. Can you modify bubble sort.

Modified Bubble sort:-

3.5

```
void bubble sort (int A[], int n)
{
    int flag;
    for (int i = 0; i < n - 1; i++) {
        flag = 0;
        for (j = 0; j < n - 1 - i; j++)
        {
            if (A[j] > A[j + 1])
            {
                swap (A[j], A[j + 1]);
                flag = 1;
            }
        }
        if (flag == 0)
            break;
    }
}
```

