

TUTORIAL - 2

Q.1

Time Complexity of below code

```

Void fun(int n)
{

```

→ $O(1)$

```

    int j = 1; i = 0;

```

```

    while (i < n)
    {

```

```

        i = i + j;

```

```

        j++;
    }

```

 $i = 0 \rightarrow i = 0 + 1, j = 2$
 $i = 1 \rightarrow i = 1 + 2, j = 3$
 $i = 3 \rightarrow i = 1 + 2 + 3, j = 4$
 $i = 6 \rightarrow i = 6 + 4, j = 5$
 \vdots
 $i = k \rightarrow i = k + j, j = n$

Sum

or

j	1	2	3	4	...	n
---	---	---	---	---	-----	---

i	1	3	6	10	...	k
---	---	---	---	----	-----	---

$$\frac{k(k+1)}{2} > n$$

$$\frac{k^2 + k}{2} > n$$

$$k^2 = n$$

$$k = \sqrt{n}$$

$$T.C = O(k) \text{ or } O(\sqrt{n})$$

Q.2 Write recurrence relation for recursive function that prints fibonacci series.

```

int fib(int n)
{

```

```

    if (n <= 1)

```

```

        return n;

```

```

    return fib(n-1) + fib(n-2);
}

```

3

$$T(n) = \begin{cases} 1, & n \leq 1 \\ T(n-1) + T(n-2), & \text{otherwise.} \end{cases}$$

$$\begin{aligned} T(n) &= T(n-1) + T(n-2) + c \\ &= 2T(n-1) + c \end{aligned}$$

$[T(n-1) \sim T(n-2)]$ from approximation

$$\begin{aligned} T(n) &= 2[2T(n-2) + c] + c \\ &= 4T(n-2) + 3c \\ &= 8T(n-3) + 4c \end{aligned}$$

$$\vdots$$

$$2^k T(n-k) + (2^k - 1)c$$

$$\begin{aligned} T(1) &= 1 \\ T(n-k) &\sim T(1) \\ n &= k+1 \end{aligned}$$

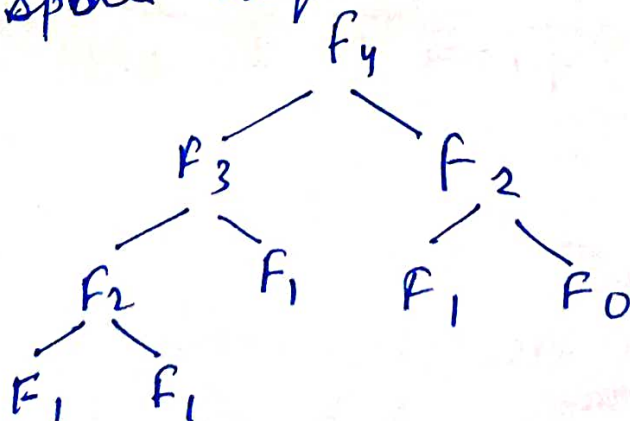
$$\boxed{k = n-1}$$

$$T(n) = 2^{n-1} T(n - (n-1)) + (2^{n-1} - 1)c$$

$$\boxed{\text{Time complexity} = O(2^n)}$$

for space complexity:

space required \propto max. dep



$\rightarrow O(4)$

\rightarrow for n elements $O(n)$

Q.3 Write programmes which have complexity ^{2.2}

1) $n \log n$

A()

```
{ int i, j;  
  for (i=1; i<=n; i++)  $\rightarrow O(n)$   
  {  
    for (j=1; j<=n; j=j/2)  $\rightarrow O(\log n)$   
    {  
      printf("#");  
    }  
  }  
}
```

Y

}

2) n^3

A()

```
{ int i, j, k;  
  for (int i=0; i<=n; i++)  $\rightarrow O(n)$   
  {  
    for (j=0; j<=n; j++)  $\rightarrow O(n)$   
    {  
      for (k=0; k<=n; k++)  $\rightarrow O(n)$   
      {  
        printf("*");  
      }  
    }  
  }  
}
```

Y
Y

Q.4 Solve the following recurrence relations.

$$T(n) = T(n/4) + T(n/2) + cn^2$$

$$= T(n/4) \leq T(n/2)$$

$$T(n) = T(n/2) + T(n/2) + cn^2$$

$$= 2T(n/2) + cn^2$$

Comparing with master's eqⁿ we get

$$a = 2$$

$$b = 2$$

$$c = \log_2 2 = 1$$

$$n < n^2$$

$$k = 2$$

$$p = 0$$

$$a < b^k$$

$$O(n^k \log^p n)$$

$$O(n^2 (\log n)^0)$$

$$\boxed{T_c = O(n^2)}$$

Question 5:-

What is time complexity of following fn.

```

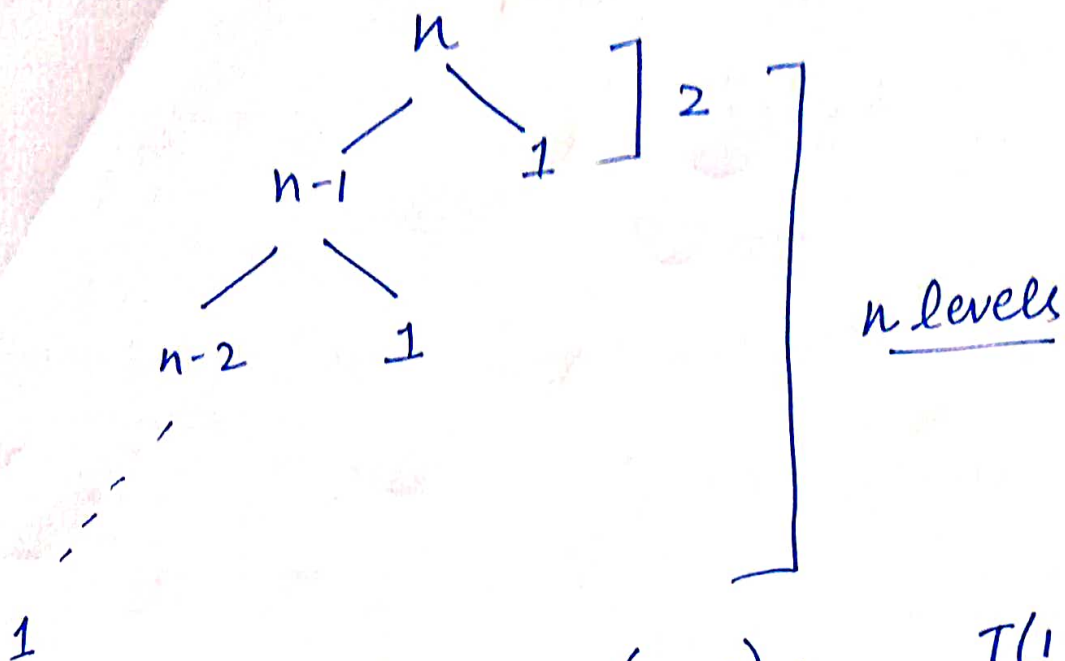
function fun ()
int fun(int n) {
    for (int i = 1; i <= n; i++) {
        for (int j = 1; j < n; j += i)
            // O(i) // task
        }
    }
}

```

Incrementation depends on i.
We'll unroll all loop.

Question 7

$$T(n) = T(n-1) + O(1)$$



$$T(n) = T(n-1) + T(n-2) + \dots + T(1) + O(1) \times n$$

$$= n \times n$$

$$\boxed{T(n) = O(n^2)}$$

lowest Height = 2
highest = n

$$\boxed{\text{Diff} = n-2}, n > 1$$

Given algo produces linear search.

•••

Question 6.

Time complexity of

```
for(int i=2; i <= n; i = pow(i, k))  
{  
    // O(1)  
}
```

$$i = 2$$

$$i = 2^k$$

$$i = (2^k)^k = 2^{k^2}$$

'2 to n times'

' 2^k to n times'

' 2^{k^2} to n times'

$$i = 2^k \log_k (\log n) = 2^{\log_2 n} = n$$

Total TC $\rightarrow O(\log \log n)$

$$2^{kl} = n$$

$$k^l = \log_2 n$$

$$k^l = \log_2 n$$

$$l = \log_k \log_2 n$$

