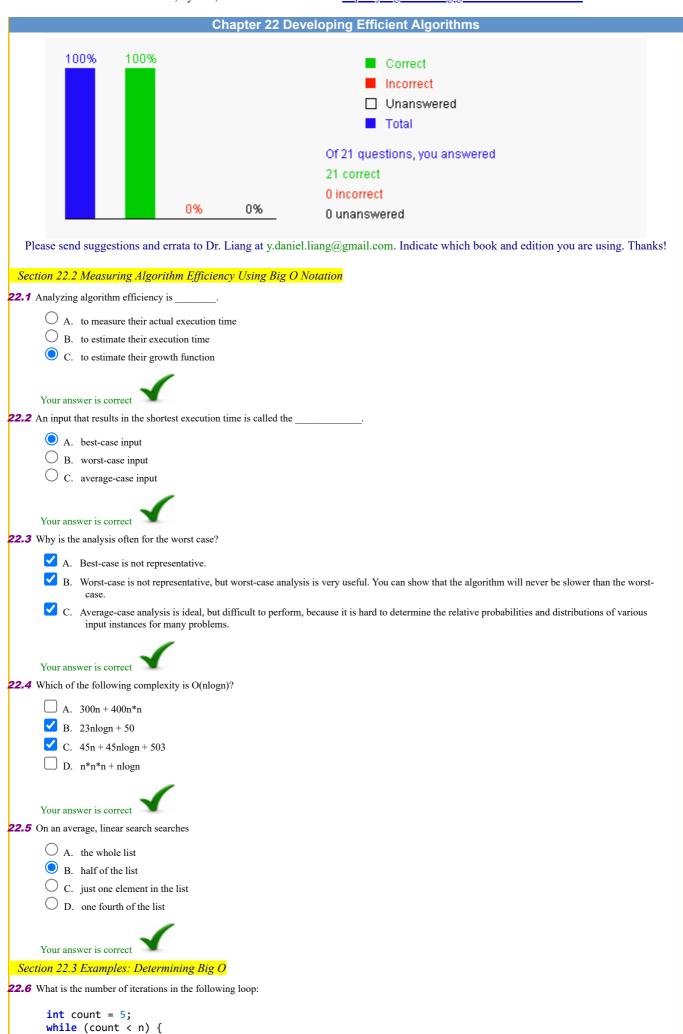
Introduction to Java Programming, Includes Data Structures, Eleventh Edition, Y. Daniel Liang

This quiz is for students to practice. A large number of additional quiz is available for instructors using Quiz Generator from the Instructor's Resource Website.

Videos for Java, Python, and C++ can be found at https://yongdanielliang.github.io/revelvideos.html.



| count = count + 3; |
|---|
| } |
| O A. n-5 |
| O B. n-3 |
| O C. n/3-1 |
| \bigcirc D. $(n-5)/3$ |
| \bullet E. the ceiling of (n - 5) / 3 |
| Your answer is correct |
| Section 22.4 Analyzing Algorithm Time Complexity |
| 22.7 For a sorted list of 1024 elements, a binary search takes at most comparisons. |
| ● A. 11 |
| |
| O B. 100 |
| O C. 512 |
| O D. 6 |
| Your answer is correct |
| 22.8 O(1) is |
| • A. constant time |
| |
| B. logarithmic time |
| C. linear time |
| O D. log-linear time |
| v · · · · · · · · · · · · · · · · · · · |
| Your answer is correct |
| 22.9 The time complexity for the Towers of Hanoi algorithm in the text is |
| \bigcirc A. $O(n)$ |
| O B. O(n^2) |
| O C. O(n^3) |
| D. O(2^n) |
| |
| v · · · · · · · · · · · · · · · · · · · |
| Your answer is correct |
| 22.10 The time complexity for the selection sort algorithm in the text is |
| O A. O(nlogn) |
| ⊕ B. O(n^2) |
| C. O(logn) |
| O D. O(2^n) |
| |
| Your answer is correct |
| 22.11 The time complexity for the insertion sort algorithm in the text is |
| |
| A. O(nlogn) |
| B. O(n^2) |
| C. O(logn) |
| O D. O(2^n) |
| |
| Your answer is correct |
| Section 22.5 Finding Fibonacci Numbers using Dynamic Programming |
| 22.12 approach is the process of solving subproblems, then combining the solutions of the subproblems to obtain an overall solution. This |
| naturally leads to a recursive solution. However, it would be inefficient to use recursion, because the subproblems overlap. The key idea behind dynamic programming is to solve each subproblem only once and store the results for subproblems for later use to avoid redundant computing of the subproblems. |
| A. Divide-and-conquer |
| B. Dynamic programming |
| C. Brutal-force |
| D. Backtracking |
| D. Duvatuvanig |
| |
| Your answer is correct |
| 22.13 The time complexity for the recursive Fibonacci algorithm in the text is |
| A. O(nlogn) |

| | \bigcirc B | . O(n^2) |
|---------------------------------|--|--|
| | \bigcirc c | . O(logn) |
| | D | o. O(2^n) |
| | | |
| | Your a | nswer is correct |
| 22.14 | | ne complexity for the algorithm using the dynamic programming approach is |
| | | . O(n) |
| | _ | . O(n) . O(n^2) |
| | _ | |
| | _ | O(logn) O(2^n) |
| | O L | $O(2^n)$ |
| | | |
| | Your ar | iswer is correct |
| Sect | ion 22. | 6 Finding Greatest Common Divisors Using Euclid?s Algorithm |
| 22.15 | The tir | ne complexity for the Euclid?s algorithm is |
| | O A | . O(n) |
| | _ | . O(n^2) |
| | _ | . O(logn) |
| | _ | O. O(2^n) |
| | O L | . 0(2 ii) |
| | | |
| | Your ar | swer is correct |
| Sect | ion 22. | 7 Efficient Algorithms for Finding Prime Numbers |
| 22.16 | The tir | ne complexity for the Sieve of Eratosthenes algorithm is |
| | \bigcirc \triangle | . O(n) |
| | _ | $O(n^{(1.5)/\log n})$ |
| | _ | . O(logn) |
| | _ | O. O(2^n) |
| | O L | . 0(2 11) |
| | | |
| | | swer is correct |
| Sact | ion 22. | 8 Finding the Closest Pair of Points Using Divide-and-Conquer |
| Seci | | |
| | The tir | ne complexity for the the closest pair of points problem using divide-and-conquer is |
| | | ne complexity for the the closest pair of points problem using divide-and-conquer is |
| | O A | ne complexity for the the closest pair of points problem using divide-and-conquer is O(n) |
| | ○ A | ne complexity for the the closest pair of points problem using divide-and-conquer is O(n) O(nlogn) |
| | ○ A ○ B ○ C | ne complexity for the the closest pair of points problem using divide-and-conquer is O(n) O(nlogn) O(logn) |
| | ○ A ○ B ○ C | ne complexity for the the closest pair of points problem using divide-and-conquer is O(n) O(nlogn) |
| | ○ A○ B○ C○ D | ne complexity for the the closest pair of points problem using divide-and-conquer is O(n) O(nlogn) O(logn) O(2^n) |
| 22.17 | A B C C C T | ne complexity for the the closest pair of points problem using divide-and-conquer is O(n) O(nlogn) O(logn) O(2^n) |
| 22.17 22.18 | A B C C C T | ne complexity for the the closest pair of points problem using divide-and-conquer is O(n) O(nlogn) O(logn) O(2^n) swer is correct approach divides the problem into subproblems, solves the subproblems, then combines the solutions of the subproblems to obtain |
| 22.17 22.18 | A B C C C T | ne complexity for the the closest pair of points problem using divide-and-conquer is O(n) O(nlogn) O(logn) O(2^n) |
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| 22.17 22.18 Sect | Your and the solution A A B B C C C C C C C C C C C C C C C C | ne complexity for the the closest pair of points problem using divide-and-conquer is O(n) O(nlogn) O(2^n) swer is correct approach divides the problem into subproblems, solves the subproblems, then combines the solutions of the subproblems to obtain approach divides the problem. Unlike the approach, the subproblems in the divide-and-conquer approach don?t overlap. A subproblem is original problem with a smaller size, so you can apply recursion to solve the problem. Divide-and-conquer/dynamic programming Dynamic programming/divide-and-conquer Brutal-force/divide-and-conquer Backtracking/dynamic programming swer is correct 9 Solving the Eight Queens Problem Using Backtracking approach searches for a candidate solution incrementally, abandoning that option as soon as it determines that the candidate cannot |
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| O B. O(nlogn) | | | |
|---|--|--|--|
| C. O(logn) | | | |
| D. O(n^2) | | | |
| Your answer is correct | | | |
| 22.21 The Graham?s algorithm for finding a convex hull takes time. | | | |
| \bigcirc A. O(n) | | | |
| B. O(nlogn) | | | |
| O C. O(logn) | | | |
| O D. O(n^2) | | | |
| Your answer is correct | | | |