#### **CS 228 : Logic in Computer Science**

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- ▶ A disjunction of literals  $L_1 \lor L_2 \lor ... L_n$  is valid iff ...

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- A disjunction of literals L₁ ∨ L₂ ∨ ... Ln is valid iff<sup>Any</sup> one is true for all assignment.
  A conjunction of literals L₁ ∧ L₂ ∧ ... Ln is satisfiable iff ...

all are true for any one assignment

## **Normal Forms: CNF Validity**

Let  $\varphi = C_1 \wedge C_2 \wedge \cdots \wedge C_n$  be in CNF.

- ▶ Checking if  $\varphi$  is satisfiable is NP-complete.
- $\blacktriangleright$  Checking if  $\varphi$  is valid is polynomial time. Why?
- Question raised in class: If validity is polytime, so should be satisfiability. Is this true?

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