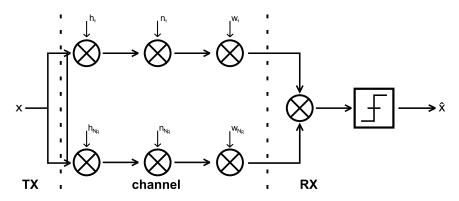
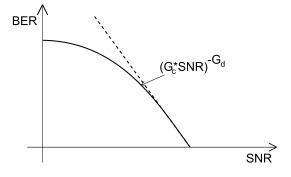
# 1 SIMO Systems

#### Remarks

- $\bullet$  In SIMO Systems only <u>coding</u> and <u>diversity</u> <u>gains</u> can be exploited (no multiplexing gains)
- To realize these gains diversity combining has to be performed
- Diversity combining schemes vary in complexity and performance
- There are many diversity combining schemes. Here we consider:
  - Maximal ratio combining (MRC)
  - Equal gain combining (EGC)
  - Selection combining (SC)
- Diversity combining problem



- how to choose combining weights  $w_n$ ?
- what performance (e.g. error rate, outage probability) is achieved?
- what diversity and coding/combining gain is achieved?



•  $G_c$ : Coding gain

•  $G_d$ : Diversity gain

## 1.1 Preliminaries

Consider an equivalent system:

$$y=hx+n;$$
 
$$\mathcal{E}\{|x^2|\}=E_s; \qquad \qquad \mathcal{E}\{|n^2|\}=\sigma_n^2; \qquad \qquad \mathcal{E}\{|h|^2\}=1$$

- Instantaneous SNR:  $\gamma_t = \frac{E_s}{\sigma_n^2} \times |h|^2$
- Average SNR:  $\bar{\gamma}_t = \mathcal{E}\{\gamma_t\} = \frac{E_s}{\sigma_n^2}$

#### Bit and Symbol Error Rate

• The Bit and Symbol Error Rate of many modulation schemes can be expressed for given  $\gamma_t$  as:

$$P_e(\gamma_t) = aQ\{\sqrt{b\gamma_t}\}$$

where:

- $Q(x) = \frac{1}{\sqrt{2\pi}} \times \int_x^\infty e^{-\frac{t^2}{2}} dt$
- $P_e(\gamma_t)$  may be exact result or approximation
- BPSK: exact with a = 1, b = 2
- M-ary QAM: tight approximation with  $a = 4\left(1 \frac{1}{\sqrt{M}}\right), b = \frac{3}{M-1}$

 $\left(Einschub:Gray-Code:BER=\frac{1}{\log_2 M}\times SER\right)$ 

• Alternative representation of Q - function:

$$Q(x) = \frac{1}{\pi} \int_0^{\frac{\pi}{2}} e^{-\frac{x^2}{2sin^2\theta}} d\theta$$

- $\rightarrow$  Integral limits are fixed and do not depend on integration variables!
- Average error probability

$$P_e = \mathcal{E}\{P_e(\gamma_t)\} = \int_0^\infty aQ(\sqrt{bx})p_{\gamma_t}(x) dx$$

- Integral may be difficult to solve analytically
- Integral has infinite support  $\rightarrow$  numerical evaluation difficult
- Using alternative representation of Q-function we get:

$$P_{e} = \int_{0}^{\infty} \frac{a}{\pi} \int_{0}^{\frac{\pi}{2}} e^{-\frac{bx}{2sin^{2}\theta}} p_{\gamma_{t}}(x) d\theta dx$$

$$= \frac{a}{\pi} \int_{0}^{\frac{\pi}{2}} \int_{0}^{\infty} p_{\gamma_{t}}(x) e^{-\frac{b}{2sin^{2}\theta}} dx d\theta \qquad = \frac{a}{\pi} \int_{0}^{\frac{\pi}{2}} M_{\gamma_{t}}(\frac{b}{2sin^{2}\theta}) d\theta$$

where:

- $-M_{\gamma_t}(s) = \int_0^\infty p_{\gamma_t}(x)e^{-sx} dx$  is the Laplace transform of  $p_{\gamma_t}$
- $-M_{\gamma_t}(-s)$  is the so called Moment Generation Function (MGF) of  $p_{\gamma_t}$
- Here, we will also refer to  $M_{\gamma_t}(s)$  as MGF
- $M_{\gamma_t}(s)$  is sometimes easier to obtain than  $p_{\gamma_t}$
- The above integral can be easily evaluated numerically because of the finite integral limits

#### Outage probability

• The outage probability is the probability that the channel cannot support a certain rate, R, i.e. (where  $\gamma_T$  is the threshold SNR):

$$C = \log_2(1 + \gamma_t) < R \quad \leftrightarrow \quad \gamma_t < 2^R - 1 \triangleq \gamma_T$$

Thus, the outage probability is given by:

$$P_{out} = P_0 \gamma_t < \gamma - T = \int_0^{\gamma_T} p_{\gamma_t}(x) \ dx$$

• Using the inverse Laplace Transform

$$p_{\gamma_t}(x) = \frac{1}{2\pi j} \int_{c-j\omega}^{c+j\omega} M_{\gamma_t}(s) e^{sx} dx$$

where c > 0 is a small constant that lies in the region of convergence of the integral, we



- 1.

$$P_{out} = \frac{1}{2\pi j} \int_{c-j\omega}^{c+j\omega} M_{\gamma_t}(s) \int_0^{\gamma_T} e^{sx} dx ds = \frac{1}{2\pi j} \int_{c-j\omega}^{c+j\omega} M_{\gamma_t}(s) e^{\gamma_T s} \frac{ds}{s}$$

(lower integral limit is 0 since  $p_{\gamma_t}(0) = 0$ )

- and 2.:

$$p_{\gamma_t}(x) = \int_0^x p_{\gamma_t}(t) \ dt = 0$$
for  $x = 0$  note:  $p_{\gamma_t}(x) \xleftarrow{Laplace}{transform} \frac{1}{s} M_{\gamma_t}(s)$ 

### General combining scheme

$$y = \left(\sum_{n=1}^{N_R} h_n w_n\right) x + \sum_{n=1}^{N_R} w_n n_n$$
$$\gamma_t = \frac{\epsilon_s \left|\sum_{n=1}^{N_R} h_n w_n\right|^2}{\sigma_n^2 \sum_{N=1}^{N_R} |w_n|^2}$$

where  $w_n$  depends on the particular combining scheme.

# 1.2 MRC (Maximum Ratio Combining)

- what weight  $w_n$  maximize  $\gamma_t$ ?
  - Cauchy-Schwarz inequality

$$\left| \sum_{n=1}^{N_R} h_n w_n \right|^2 \le \sum_{n=1}^{N_R} |h_n|^2 \cdot \sum_{n=1}^{N_R} |w_n|^2$$

where equality holds if and only if  $w_n = c \cdot h_n^*$  for some non-zero constant c.

- for  $w_n = h_n^*$ , we obtain

$$\gamma_t = \frac{\epsilon_s}{\sigma_n^2} \cdot \frac{\left(\sum_{n=1}^{N_R} |h_n|^2\right)^2}{\sum_{n=1}^{N_R} |h_n|^2} = \frac{\epsilon_s}{\sigma_n^2} \sum_{n=1}^{N_R} |h_n|^2$$

- $w_n = h_n^* \forall n$  are the MRC combining weights.
- $\bullet$  For performance analysis we assume independent identically distributed (IID) Rayleigh fading

• Error rate

$$\gamma_t = \sum_{n=1}^{N_R} \gamma_n$$

 $\rightarrow$  sum of IID random variables (r.v.s.)

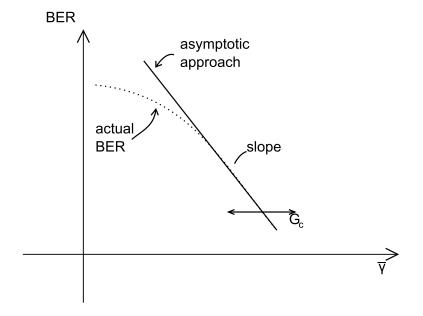
$$M_{\gamma_t}(s) = \left(M_{\gamma}(s)\right)^{N_R} = \frac{1}{(1+s\bar{\gamma})^{N_R}} = \frac{1}{\bar{\gamma}^{N_R}} \cdot \frac{1}{(s+\frac{1}{\bar{z}})^{N_R}}$$

inverse Laplace-transform (from tables)

$$p_{\gamma_t}(x) = \frac{1}{\bar{\gamma}^{N_R}} \cdot \frac{x^{N_R - 1}}{(N_R - 1)!} e^{-\frac{x}{\bar{\gamma}}}; \quad x \ge 0$$

• Direct approach

$$p_e = \int_0^\infty a \cdot Q\left(\sqrt{ax}\right) p_{\gamma_t}(x) \ dx = a \left(\frac{1-\mu}{2}\right)^{N_R} \cdot \sum_{n=0}^{N_R-1} \binom{N_R-1+n}{n} \left(\frac{1+\mu}{2}\right)^n$$
 where  $\mu = \sqrt{\frac{b\bar{\gamma}}{2+b\bar{\gamma}}}$ 



• MGF approach

$$\begin{split} p_e &= \frac{a}{\pi} \int_0^{\frac{\pi}{2}} M_{\gamma_t} \left( \frac{b}{2 \sin^2 \theta} \right) d\theta \\ &= \frac{a}{\pi} \int_0^{\frac{\pi}{2}} \frac{1}{\bar{\gamma}^{N_R} \left( \frac{b}{\sin^2 \theta} + \frac{1}{\bar{\gamma}} \right)^{N_R}} d\theta \quad \text{(numerisch berechnen!)} \end{split}$$

• high SNR:  $\bar{\gamma} \to \infty \Longleftrightarrow \frac{1}{\bar{\gamma}} \to 0$ 

$$\begin{split} p_e &= \frac{a}{\pi} \cdot \frac{1}{\bar{\gamma}^{N_R}} \cdot \left(\frac{2}{b}\right)^{N_R} \int_0^{\frac{\pi}{2}} \sin^{2N_R} \theta \ d\theta \\ \text{(from MGF approach: } \int_0^{\frac{\pi}{2}} \sin^{2N_R} \theta \ d\theta &= \frac{\pi}{2^{N_R+1}} \cdot \binom{2N_R}{N_R} \\ &= \frac{a}{2^{N_R+1} \cdot b^{N_R}} \left(2N_R \quad N_R\right) \frac{1}{\bar{\gamma}^{N_R}} \quad \text{as } \bar{\gamma} \to \infty \\ &= \frac{1}{G_c \bar{\gamma}} \end{split}$$

where: Diversity gain:  $G_d = N_R$ 

Combining/Coding gain: 
$$G_c = 2b \left(\frac{a}{2} {2N_R \choose N_R}\right)^{-\frac{1}{N_R}}$$

- MRC exploits the maximal possible diversity
- Diversity gain is not affected by correlation as the branches are not fully correlated
- Diversity gain depends on fading distribution

#### Outage probability

$$P_{out} = \int_0^{\gamma_T} p_{\gamma_t}(x) \ dx = \frac{1}{\bar{\gamma}^{N_R}} \int_0^{\gamma_T} \frac{x^{N_R - 1}}{(N_R - 1)!} e^{-\frac{x}{\bar{\gamma}}} \ dx$$
$$= 1 - e^{-\frac{\gamma_T}{\bar{\gamma}}} \cdot \sum_{n=1}^{N_R} \frac{\left(\frac{\gamma_T}{\bar{\gamma}}\right)^n}{(n-1)!}$$

• Approximation (Taylor series):  $\bar{\gamma} \to \infty$ :  $-e^{-\frac{x}{\bar{\gamma}}} = 1 - \frac{x}{\bar{\gamma}} + O(\frac{1}{\bar{\gamma}})$  where a function f(x) is O(x) if  $\lim_{x \to \infty} \frac{f(x)}{x} = 0$ .

$$\Rightarrow P_{out} = \frac{1}{\gamma^{N_R}} \int_{0}^{\gamma_T} \frac{x^{N_R - 1}}{(N_R - 1)!} \left( 1 - \frac{x}{\bar{\gamma}} + O\left(\frac{1}{\bar{\gamma}}\right) \right)$$

• Diversity and coding gain can also be defined for  $P_{out}$ 

# 1.3 EGC (Equal Gain Combining)

#### Combining Weights

- For MRC, both, the amplitudes and phases of the channel gains  $h_n = |h_n|e^{j\varphi_n}$  have to be known (or estimated in practice)
- In EGC it is assumed that only the phases are known and weights  $w_n = e^{-j\varphi_n}$  are used.

$$\Rightarrow \gamma_t = \frac{E_s}{\sigma_n^2} \frac{\left| \sum_{n=1}^{N_R} |h_n| e^{j\varphi_n} e^{-j\varphi_n} \right|^2}{\sum_{n=1}^{N_R} |e^{-j\varphi_n}|^2} = \frac{E_s}{\sigma_n^2} \frac{1}{N_R} \left( \sum_{n=1}^{N_R} |h_n| \right)^2$$
$$= \frac{1}{N_R} \left( \sum_{n=1}^{N_R} \sqrt{\gamma_n} \right)^2; \text{ with } \gamma_n = \frac{E_s}{\sigma_n^2} |h_n|^2$$

#### Performance Analysis

- IID case  $\Rightarrow \sqrt{\gamma_n}$  is Rayleigh distributed  $\Rightarrow$  Exact analysis is much more difficult than for M
  - $\Rightarrow$  Exact analysis is much more difficult than for MRC  $\Rightarrow$  see book by Simon & Alouini p.341
- Approximate result

$$P_{e} = \frac{a}{2} \left[ 1 - \sqrt{\frac{2b\bar{\gamma}}{5 + 2b\bar{\gamma}}} \sum_{n=0}^{N_{R}-1} \frac{\binom{2n}{n}}{4^{n} (1 + \frac{2}{5}b\bar{\gamma})^{n}} \right]$$

- high SNR
  - $\Rightarrow$  use high SNR analysis of Wang & Giannakis, 2003
  - $\Rightarrow$  at high SNR, only pdf of  $\gamma_n$  around 0 is relevant for performance

$$\Rightarrow \begin{array}{l} \text{Rayleigh} \\ p_{\gamma}(x) = \frac{1}{\bar{\gamma}} e^{-\frac{x}{\bar{\gamma}}} \stackrel{\text{Taylor Serie}}{=} \frac{1}{\bar{\gamma}} + O\left(\frac{1}{\bar{\gamma}}\right) \text{ as } x \to 0 \end{array}$$

• need pdf  $\gamma_t$ : ( $\gamma_n$  bekannt,  $\rightarrow$  ges.: Wurzel, etc.) (cumulative distribution function of  $\sqrt{\gamma} \stackrel{\text{i.i.d}}{=} \sqrt{\gamma_n}$ ) (cdf))

$$\begin{split} &P_{\sqrt{\gamma}}(x) = \Pr \big\{ \sqrt{\gamma} \leq x \big\} = \Pr \big\{ \gamma \leq x^2 \big\} = P_{\gamma}(x^2) = \text{cdf of } \gamma \\ &\to p_{\sqrt{\gamma}}(x) = \frac{d}{dx} P_{\sqrt{\gamma}}(x) = 2x \cdot p_{\gamma}(x^2) = \frac{2x}{\bar{\gamma}} + O \Big( \frac{1}{\bar{\gamma}} \Big) \end{split}$$

• Laplace Transformation to MGF

$$\begin{split} & \to M_{\sqrt{\gamma}}(s) = \mathcal{L}\big\{p_{\sqrt{\gamma}}(x)\big\} = \frac{2}{\bar{\gamma}} \cdot \frac{1}{s^2} + O\big(\frac{1}{\bar{\gamma}}\big) \\ & \sqrt{\gamma_t} = \sum_{n=1}^{N_R} \frac{\sqrt{\gamma_n}}{N_R} \\ & M_{\sqrt{\gamma_t}}(s) = \mathcal{E}\Big\{\exp(-s\sqrt{\gamma_t})\Big\} = \mathcal{E}\Big\{\exp(-\frac{s}{\sqrt{N_R}} \cdot \sum_{n=1}^{N_R} \sqrt{\gamma_n})\Big\} = \Big(\mathcal{E}\Big\{\exp(-\frac{s}{\sqrt{N_R}} \cdot \sqrt{\gamma_n}\Big\}\Big)^{N_R} \\ & = \Big(M_{\sqrt{\gamma}}\big(\frac{s}{\sqrt{N_R}}\big)\Big)^{N_R} = \Big(\frac{2}{\bar{\gamma}} \cdot \frac{N_R}{s^2}\Big)^{N_R} + O\Big(\frac{1}{\bar{\gamma}^{N_R}}\Big) \end{split}$$

• inverse Laplace Transform

$$\begin{split} p_{\sqrt{\gamma_{t}}}(x) &= \mathcal{L}^{-1}\Big\{M_{\sqrt{\gamma_{t}}}(s)\Big\} = \Big(\frac{2N_{R}}{\bar{\gamma}}\Big)^{N_{R}} \cdot \frac{x^{2N_{R}-1}}{(2N_{R}-1)!} + O\Big(\frac{1}{\bar{\gamma}^{N_{R}}}\Big) \\ P_{\gamma_{t}}(x) &= \Pr\{\gamma_{t} \leq x\} = \Pr\{\sqrt{\gamma_{t}} \leq \sqrt{x}\} = P_{\sqrt{\gamma_{t}}}(\sqrt{x}) \to \operatorname{cdf of } \sqrt{\gamma_{t}} \\ p_{\gamma_{t}}(x) &= \frac{d}{dx}P_{\gamma_{t}}(x) = \frac{1}{2\sqrt{x}} \cdot p_{\gamma_{t}}(\sqrt{x}) = \frac{1}{2}\Big(\frac{2N_{R}}{\bar{\gamma}}\Big)^{N_{R}} \cdot \frac{x^{N_{R}-1}}{(2N_{R}-1)!} + O\Big(\bar{\gamma}^{-N_{R}}\Big) \\ \to M_{\gamma_{t}}(s) &= \mathcal{L}\{p_{\gamma_{t}}(x)\} = \frac{1}{2}\Big(\frac{2N_{R}}{\bar{\gamma}}\Big)^{N_{R}} \cdot \frac{(N_{R}-1)!}{(2N_{R}-1)!} + O\Big(\bar{\gamma}^{-N_{R}}\Big) \end{split}$$

• Error Probability:

$$\begin{split} P_e &= \frac{a}{\pi} \int\limits_0^{\frac{\pi}{2}} M_{\gamma_t} \left( \frac{b}{2 \sin^2(\theta)} \right) \mathrm{d}\theta \\ &= \frac{a}{\pi} \frac{1}{2} \left( \frac{2N_R}{\bar{\gamma}} \right)^{N_R} \frac{(N_R - 1)!}{(2N_R - 1)!} \frac{2^{N_R}}{b^{N_R}} \int\limits_{\frac{\pi}{2^{2N_R + 1}}}^{\frac{\pi}{2}} \sin^{2N_R}(\theta) \mathrm{d}\theta + O\left( \frac{1}{\bar{\gamma}^{N_R}} \right) \\ &= \frac{aN_R^{N_R}}{2b^{N_R}N_R!} \frac{1}{\bar{\gamma}^{N_R}} + O\left( \frac{1}{\bar{\gamma}^{N_R}} \right) \stackrel{!}{=} \left( \frac{1}{G_c} \right)^{G_d} \\ \Longrightarrow \text{Diversity gain: } G_d = N_R \\ \Longrightarrow \text{Combining gain: } G_c = \frac{b}{N_R} \left( \frac{2N_R!}{a} \right)^{\frac{1}{N_R}} \end{split}$$

vergleiche auch Blatt mit Kurven  $\overline{\mathrm{III}}$  und  $\overline{\mathrm{IV}}$ 

A similar asymptotic analysis can be conducted for the outage probability.

# 1.4 SC (Selection Combining)

### Combining weights

- only the strongest branch is chosen
- strongest branch:  $\hat{n} = \underset{n}{\operatorname{argmax}} \gamma_n \longrightarrow \gamma_t = \gamma_{\hat{n}}$
- ullet only on RF receiver chain required o saves hardware complexity

#### Performance analysis

• cdf of:  $\gamma_t$ 

$$P_{\gamma_t}(x) = \Pr\{\gamma_{\hat{n}} \le x\} = \Pr\{\gamma_1 \le x \cap \gamma_2 \le x \cap \dots \gamma_{N_R} \le x\}$$

$$\stackrel{(IID)}{=} \left(\Pr\{\gamma_n \le x\}\right)^{N_R} = \left(P_{\gamma}(x)\right)^{N_R}$$

• pdf:

$$\begin{split} p_{\gamma_t}(x) &= \frac{d}{dx} P_{\gamma_t}(x) = N_R \big( P_{\gamma}(x) \big)^{N_R - 1} \cdot p_{\gamma}(x) \\ \text{where:} \qquad p_{\gamma_t}(x) &= \frac{1}{\bar{\gamma}} e^{-\frac{x}{\bar{\gamma}}}; \quad x \geq 0 \\ P_{\gamma}(x) &= \int_0^x p_{\gamma}(x) \; dx = 1 - e^{-\frac{x}{\bar{\gamma}}}; \quad x \geq 0 \\ &\to p_{\gamma_t}(x) = \frac{N_R}{\bar{\gamma}} \big( 1 - e^{-\frac{x}{\bar{\gamma}}} \big)^{N_R - 1} e^{-\frac{x}{\bar{\gamma}}}; \quad x \geq 0 \end{split}$$

#### Error probability

- $\bullet$  direct approach  $\rightarrow$  closed-form solution possible
- MGF approach
  - Binomial expansion

$$\begin{split} p_{\gamma_t}(x) &= \frac{N_R}{\bar{\gamma}} e^{-\frac{x}{\bar{\gamma}}} \sum_{n=0}^{N_R-1} \binom{N_R-1}{n} 1^{N_R-1-n} \Big( -e^{-\frac{x}{\bar{\gamma}}} \Big)^n \\ &= \frac{N_R}{\bar{\gamma}} \sum_{n=0}^{N_R-1} \binom{N_R-1}{n} \cdot (-1)^n e^{-\frac{x(n+1)}{\bar{\gamma}}}; \quad x \geq 0 \end{split}$$

- MGF

$$M_{\gamma_t}(s) = \frac{N_R}{\bar{\gamma}} \sum_{n=0}^{N_R-1} \binom{N_R-1}{n} (-1)^n \frac{1}{s + \frac{n+1}{\bar{\gamma}}}$$

\_

$$P_e = \frac{a}{\pi} \int_0^{\frac{\pi}{2}} M_{\gamma_t} \left( \frac{b}{2 \sin^2 \theta} \right) d\theta = \frac{a N_R}{\pi \bar{\gamma}} \sum_{n=0}^{N_R - 1} \binom{N_R - 1}{n} (-1)^n \int_0^{\frac{\pi}{2}} \frac{d\theta}{\frac{b}{2 \sin^2 \theta} + \frac{n+1}{\bar{\gamma}}}$$

 $\rightarrow$  can be evaluated numerically

– high SNR approach  $\Rightarrow \bar{\gamma} \to \infty$ 

$$\begin{split} p_{\gamma_t} &= \frac{N_R}{\bar{\gamma}} \left[ 1 - \exp\left(-\frac{x}{\bar{\gamma}}\right) \right]^{N_R - 1} \exp\left(-\frac{x}{\bar{\gamma}}\right) \\ &\stackrel{\bar{\gamma} \to \infty}{=} \frac{N_R}{\bar{\gamma}} \left[ 1 - \left(1 - \frac{x}{\bar{\gamma}} + O\left(\bar{\gamma}^{-1}\right)\right) \right]^{N_R - 1} \left(1 - \frac{x}{\bar{\gamma}} + O\left(\bar{\gamma}^{-1}\right)\right) i \\ &= \frac{N_R}{\bar{\gamma}^{N_R}} x^{N_R - 1} + o\left(\bar{\gamma}^{-N_R}\right) \end{split}$$

- MGF:

$$\begin{split} M_{\gamma_t}(s) &= \frac{N_R}{\bar{\gamma}^{N_R}} \frac{(N_R - 1)!}{s^{N_R}} + O\left(\bar{\gamma}^{-N_R}\right) \\ \left[ \to P_e &= \frac{a}{\pi} \int_0^{\frac{\pi}{2}} M_{\gamma_t} \left(\frac{b}{2\sin^2(\theta)}\right) \mathrm{d}\theta \right] \\ &= \frac{a(2N_R)!}{b^{N_R} 2^{N_R + 1} N_R!} \frac{1}{\bar{\gamma}^{N_R}} + O(\bar{\gamma}^{-N_R}) \end{split}$$

 $\implies$  Diversity gain:  $G_d = N_R$ 

$$\implies$$
 Combining gain:  $G_c = 2b \left(\frac{2N_R!}{a(2N_R)!}\right)^{\frac{1}{N_R}}$ 

- Outage Probability

$$\begin{split} P_{out} = \Pr\{\gamma_{\hat{n}} \leq \gamma_T\} = P_{\gamma_{\hat{n}}}(\gamma_T) = \left[1 - \exp\left(-\frac{\gamma_T}{\bar{\gamma}}\right)\right]^{N_R} \\ \text{high SNR: } P_{out} = \left(\frac{\gamma_T}{\bar{\gamma}}\right)^{N_R} + O\left(\bar{\gamma}^{-N_R}\right) \end{split}$$

# 1.5 Comparison

- Diversity Gain: MRC, EGC and SC all achieve the maximum possible diversity gain of  $G_d=N_R$
- Combining Gain:
  The combining gains of MRC, EGC and SC are different
  - MRC/EGC:

$$\frac{G_C^{EGC}}{G_C^{MRC}} = \frac{\frac{1}{2b} \left(\frac{a}{2} {2N_R \choose N_R}\right)^{\frac{1}{N_R}}}{\frac{N_R}{b} \left(\frac{a}{2} \frac{1}{N_R!}\right)^{\frac{1}{N_R}}} = \frac{\left[(2N_R)!\right]^{\frac{1}{N_R}}}{2N_R(N_R)^{\frac{1}{N_R}}} \le 1$$

(independent of a or b which are modulation parameters, only depends on number of antennas)

$$N_R \gg 1$$
:  $N_R! \approx \sqrt{2\pi} e^{-N_R} N_R^{N_R + \frac{1}{2}}$  (Stirling)

$$\frac{G_C^{EGC}}{G_C^{MRC}}\bigg|_{N_R \gg 1} = \frac{\left(\sqrt{2\pi}e^{-2N_R}(2N_R)^{2N_R + \frac{1}{2}}\right)^{\frac{1}{N_R}}}{2N_R\left(\sqrt{2\pi}e^{-N_R}N_R^{N_R + \frac{1}{2}}\right)^{\frac{1}{N_R}}} = \frac{2 \cdot 2^{\frac{1}{2N_R}}}{2} \stackrel{N_R \to \infty}{\longrightarrow} \frac{2}{e} \equiv -1.3\text{dB}$$

- MRC/SC:

$$\begin{split} \frac{G_C^{SC}}{G_C^{MRC}} &= \frac{2b \left(\frac{a}{2} \binom{2N_R}{N_R}\right)^{\frac{1}{N_R}}}{2b \left(\frac{a}{2} \frac{(2N_R)!}{N_R!}\right)^{\frac{1}{N_R}}} = \frac{1}{\left(N_R!\right)^{\frac{1}{N_R}}} \leq 1 \\ \frac{G_C^{SC}}{G_C^{MRC}} \bigg|_{N_R \gg 1} &= \frac{1}{\sqrt{2\pi^{\frac{1}{N_R}}}e^{-1}N_R^{1+\frac{1}{2N_R}}} N_R \xrightarrow{\rightarrow} \infty \frac{e}{N_R} \end{split}$$

 $\rightarrow$  loss increases with  $N_R$ 

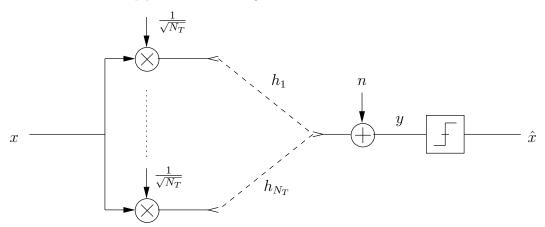
# 2 MISO Systems

#### Remarks

- Similar to SIMO systems, in MISO systems only coding and diversity gains can be obtained.
- To realize these gains, a careful transmitter design is necessary
- System design depends on whether or not channel state information (CSI) is available at transmitter

# 2.1 Naive Approach

• Assume we simply send the same signal over all  $N_T$  transmit antennas



- Transmit power:  $\mathcal{E}\left\{\left|\frac{1}{\sqrt{N_T}}x\right|^2+,\ldots,\left|\frac{1}{\sqrt{N_T}}x\right|^2\right\}=\mathcal{E}\left\{N_T\frac{1}{N_T}|x|^2\right\}=E_s$
- Received signal:  $y = \frac{1}{\sqrt{N_T}} \sum_{n=1}^{N_T} h_n \cdot x + n$
- Rayleigh fading:  $h_n$  are zero mean complex gaussian random variables  $\to h$  is also zero mean complex gaussian
- i.i.d.:

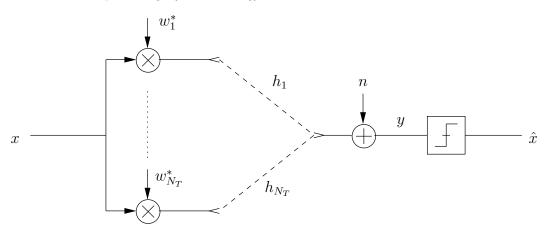
$$- \mathcal{E}\{|h_n|^2\} = 1 \ \forall n$$

$$- \mathcal{E}\{|h|^2\} = \frac{1}{N_T} \mathcal{E}\left\{ \left| \sum_{n=1}^{N_T} h_n \right|^2 \right\} = \frac{1}{N_T} \mathcal{E}\left\{ \sum_{n=1}^{N_T} |h_n|^2 \right\} = 1$$

- statistical properties of h are independent of  $N_T$
- the multiple transmit antennas have no benefit at all
- more sophisticated transmitter designs necessary

# 2.2 Full CSI Available at the Transmitter

- $h_n, n \in \{1, \dots, N_T\}$  is known at the transmitter
- Perform "precoding" (beamforming) with coefficients  $w_n$



- Transmit Power: Two constraints maybe considered
  - Average transmit power constraint

$$P_{av} = \mathcal{E}\left\{\sum_{n=1}^{N_T} |w_n^* x|^2\right\} = \sum_{n=1}^{N_T} |w_n|^2 \underbrace{\mathcal{E}\{|x|^2\}}_{E_s} = \mathcal{E}_s \Rightarrow \sum_{n=1}^{N_T} |w_n|^2 = 1$$

- Power constraint for each transmit antenna

$$\rightarrow |w_n| = \frac{1}{\sqrt{N_T}} \longrightarrow P_{av} = E_s$$

• Received signal:  $y = \underbrace{\sum_{n=1}^{N_T} w_n^* h_n}_{h} x + n$  (equivalent SISO channel)

# Maximum Ratio Transmission (MRT)

• we have only the average power constraint:  $\sum\limits_{n=1}^{N_T}|w_n|^2=1$ 

• SNR: 
$$\gamma_t = \frac{E_s |h|^2}{\sigma_n^2} = \frac{\mathcal{E}_s \left| \sum\limits_{n=1}^{N_T} w_n^* \cdot h_n \right|^2}{\sigma_n^2}$$

• Maximize SNR under constraint  $\sum_{n=1}^{N_T} |w_n|^2 = 1$ 

 $\bullet$  constraint optimization problem  $\to$  Lagrange method

$$L = \frac{E_s}{\sigma_n^2} \left| \sum_{n=1}^{N_T} w_n^* \cdot h_n \right|^2 + \lambda \left( \sum_{n=1}^{N_T} |w_n|^2 - 1 \right); \text{ where: } \lambda = \text{Lagrange Multiplier}$$

 $\Rightarrow$  Wirtinger Kalkül: treat z and  $z^*$  as independent variables for differentiation:

$$\begin{split} \frac{\partial z^*}{\partial z} &= 0; \quad \frac{\partial |z|^2}{\partial z} = \frac{\partial z \cdot z^*}{\partial z} = z^* \\ \frac{\partial x^2}{\partial x} &= 2x; \quad \frac{\partial (z^*)^2}{\partial z^*} = 2 \cdot z^*; \frac{\partial |z|^2}{\partial z} = z^* \end{split}$$

$$\frac{\partial L}{\partial w_m^*} = \frac{\epsilon_s}{\sigma_n^2} \left( \sum_{n=1}^{N_T} w_n^* \cdot h_n \right)^* h_m + \lambda w_m$$

$$\rightarrow w_m = rac{\epsilon_s}{\sigma_n^2 \cdot \lambda} \Big(\sum_{n=1}^{N_T} w_n^* h_n\Big)^* \quad h_m$$

const., independent of m := c

$$\rightarrow w_m = c \cdot h_m$$

$$\to \sum_{n=1}^{N_T} |w_n|^2 = 1 \to c^2 = \frac{1}{\sum_{n=1}^{N_T} |h_n|^2}$$

$$\rightarrow w_n = \frac{h_n}{\sqrt{\sum_{n=1}^{N_T} |h_n|^2}} \equiv \text{MRT gains}$$

$$\to \text{SNR} = \frac{E_s}{\sigma_n^2} \Big| \sum_{n=1}^{N_T} \frac{|h_n|^2}{\sqrt{\sum_{n=1}^{N_T} |h_m|^2}} \Big|^2 = \frac{\epsilon_s}{\sigma_n^2} \sum_{n=1}^{N_T} |h_n|^2$$

- ⇒ same SNR as for maximum ration combining (MRC)
- $\Rightarrow$  MRT with  $N_T$  transmit antennas achieves the same performance as MRC with  $N_T$  receive antennas
- $\Rightarrow$  MRT/MRC can be extended to  $N_T \times N_R$  MIMO systems
  - $\rightarrow$  has the same performance as MRC with  $N_T \cdot N_R$  receive antennas and one transmit antenna

# Equal Gain Transmission (EGT)

• we employ gains:  $w_n = \frac{1}{\sqrt{N_T}} \cdot \frac{h_n}{|h_m|} \to |w_n| = \frac{1}{\sqrt{N_T}}$ 

• SNR:

$$\begin{split} \gamma_t &= \frac{E_s}{\sigma_n^2} \left| \sum_{n=1}^{N_T} w_n^* h_n \right|^2 \\ &= \frac{E_s}{\sigma_n^2} \left| \sum_{n=1}^{N_T} \frac{1}{\sqrt{N_T} \cdot \frac{|h_n|^2}{|h_n|}} \right|^2 = \frac{1}{N_T} \cdot \frac{\mathcal{E}_s}{\sigma_n^2} \left| \sum_{n=1}^{N_T} |h_n| \right|^2 \\ \gamma_n &= \frac{E_s}{\sigma_n^2} |h_n|^2 \\ &\text{same SNR as for EGC} \to \gamma_t = \frac{1}{N_T} \left| \sum_{n=1}^{N_T} N_T \sqrt{\gamma_n} \right|^2 \end{split}$$

 $\rightarrow$  EGC with  $N_T$  transmit antennas achieves the same performance as EGC with  $N_T$  receive antennas

#### Transmit Antennas Selection

• select antenna with maximum channel gain for transmission:

$$w_n = \begin{cases} \frac{h_n}{|h_n|}, & \text{if } n = \hat{n} \\ 0, & \text{otherwise} \end{cases} \text{ where } \hat{n} = \underset{n}{\operatorname{argmax}} |h_n|$$

• antenna selection with  $N_T$  transmit antennas achieves the same performance as Selection Combining with  $N_T$  receive antennas

## 2.3 No CSI at Transmitter - Space-Time-Coding

- $h_n, n \in \{1, \ldots, N_T\}$ , is only known at the receiver
- "Space-time-coding" has to be employed to realize diversity gain
- $T \times N_T$  matrics **X** are transmitted in T symbol intervals over  $N_T$  antennas
- ullet X is drawn from a matrix alphabet  ${\mathcal X}$
- Example:

$$\mathbf{X} = \begin{pmatrix} x_{1,1} & x_{1,2} & \cdots & x_{1,N_T} \\ x_{2,1} & x_{2,2} & \cdots & x_{2,N_T} \\ \vdots & \vdots & \ddots & \vdots \\ x_{T,1} & x_{T,2} & \cdots & x_{T,N_T} \end{pmatrix}$$

- We distinguish:
  - Space-time-block-codes (STBCs)
    - $\to \mathbf{X}$  is obtained by mapping K scalar symbols  $s_k$ , k = 1, ..., K from a scalar alphabet  $\mathcal{A}$  to matrix  $\mathbf{X}$

- Space-time-trellis-codes (STTCs)
  - $\to$  **X** is obtained from scalar symbols  $s_k$  through a trellis encoding process. [see: Tarokh, Seshadri, Calderbank: Space-time-codes for high datarate wireless communication: Performance criterions and coder construction; IEEE Trans. Inf. Theory 1998]
- here: We concentrate on space-time-block-codes (STBCs), but many results can be easily extended to space-time-trellis-codes
- STBCs:
  - K M-ary scalar symbols (e.g. M-PSK symbols) are mapped to STBC matrices  $\mathbf{X}$   $\mathbf{S} = [s_1, \dots, s_K] \to \mathbf{X}$   $s_k \in \mathcal{A} \to x \in \mathcal{X}$  with  $|\mathcal{X}| = M^K$
  - Example: "Alamouti"-Code

$$\mathbf{X} = \frac{1}{\sqrt{2}} \begin{pmatrix} s_1 & s_2 \\ -s_2^* & s_1^* \end{pmatrix}$$

[Alamouti: A simple transmit diversity technique for wireless communication, IEEE JSAC 1998]

## **Optimal Detection**

• Signal model:

$$\begin{pmatrix} y_1 \\ \vdots \\ y_T \end{pmatrix} = \mathbf{X} \begin{pmatrix} h_1 \\ \vdots \\ h_{N_T} \end{pmatrix} + \begin{pmatrix} n_1 \\ \vdots \\ n_T \end{pmatrix}$$
$$\mathbf{y} = \mathbf{X} \cdot \mathbf{h} + \mathbf{n}$$

- Optimal detection ML-detection
  - **h** is known at receiver
  - $\mathbf{n}$  is AWGN with  $\mathcal{E}\{\mathbf{n} \cdot \mathbf{n^H}\} = \sigma_{\mathbf{n}}^2 \cdots \mathbf{I}_{T \times T}$
  - $-p(\mathbf{y}|\mathbf{x})$

$$= \frac{1}{\pi^{T} |\sigma_{n}^{2} \mathbf{I}_{T \times T}|} \exp \left(-(\mathbf{y} - \mathbf{x}\mathbf{h})^{H} (\sigma_{n}^{2} \mathbf{I}_{T \times T})^{-1} (\mathbf{y} - \mathbf{x}\mathbf{h})\right)$$

$$= \frac{1}{\pi^{T} \sigma_{n}^{2T}} \exp \left(-\frac{1}{\sigma_{n}^{2}} (\mathbf{y} - \mathbf{x}\mathbf{h})^{H} (\mathbf{y} - \mathbf{x}\mathbf{h})\right) = \frac{1}{\pi^{T} \sigma_{n}^{2T}} \exp \left(||\mathbf{y} - \mathbf{x}\mathbf{h}||^{2}\right)$$

 $\rightarrow$  the optimal estimate  $\hat{\mathbf{X}}$  or equivalently the optimal estimate  $\hat{\mathbf{s}}$  can be obtained as

$$\hat{\mathbf{s}} = \underset{\mathbf{s} \in \mathcal{A}^K}{\operatorname{argmax}} \ p(\mathbf{y}|\mathbf{x}) = \underset{\mathbf{s} \in \mathcal{A}^K}{\operatorname{argmin}} ||\mathbf{y} - \mathbf{h}\mathbf{x}||^2$$

– Disadvantage: In general, metric  $||\mathbf{y} - \mathbf{h}\mathbf{x}||^2$  has to be calculated  $M^K$  times  $\to$  complexity increases exponentially with K

# Types of STBCs

- Orthogonal STBCs (OSTBCs)
  - OSTBCs are a special class of STBCs which allow independent detection of each  $s_k \to \text{only } K \cdot M$  metrics have to be evaluated
  - Rate STBCs:  $R_{STBC} = \frac{K}{T}$
  - Examples:
    - \* Alamouti Code  $(K = 2, T = 2) \rightarrow R_{STBC} = 1$

$$\mathbf{X} = \frac{1}{\sqrt{2}} \begin{pmatrix} s_1 & s_2 \\ -s_2^* & s_1^* \end{pmatrix} \updownarrow T$$

 $\rightarrow$  only "full rate" OSTBC for complex  $s_k$ 

\* 
$$N_T = 3, K = 3, T = 4$$

$$\mathbf{X} = \frac{1}{\sqrt{3}} \begin{pmatrix} s_1 & s_2 & s_3 \\ -s_2^* & s_1^* & 0 \\ s_3^* & 0 & -s_3^* \\ 0 & -s_3^* & s_2^* \end{pmatrix} \to R_{STBC} = \frac{K}{T} = \frac{3}{4}$$

- Orthogonality:  $\mathbf{X}^H \mathbf{X} = const \cdot \mathbf{I}_{N_T \times N_T}$
- Independent detection of  $s_1 \& s_2$  for Alamouti Code

$$\begin{pmatrix} y_1 \\ y_2 \end{pmatrix} = \frac{1}{\sqrt{2}} \begin{pmatrix} s_1 & s_2 \\ -s_2^* & s_1^* \end{pmatrix} \begin{pmatrix} h_1 \\ h_2 \end{pmatrix} + \begin{pmatrix} n_1 \\ n_2 \end{pmatrix}$$

$$\rightarrow \underbrace{\begin{pmatrix} y_1 \\ y_2 \end{pmatrix}}_{\tilde{y}} = \frac{1}{\sqrt{2}} \underbrace{\begin{pmatrix} h_1 & h_2 \\ h_2^* & -h_1^* \end{pmatrix}}_{\tilde{F}} \underbrace{\begin{pmatrix} s_1 \\ s_2 \end{pmatrix}}_{s} + \underbrace{\begin{pmatrix} n_1 \\ n_2^* \end{pmatrix}}_{\tilde{n}}$$

$$\mathbf{F}^H \mathbf{F} = \frac{1}{2} \begin{pmatrix} h_1^* & h_2 \\ h_2^* & -h_1 \end{pmatrix} \begin{pmatrix} h_1 & h_2 \\ h_2^* & -h_1^* \end{pmatrix} = \frac{1}{2} \begin{pmatrix} |h_1|^2 + |h_2|^2 & 0 \\ 0 & |h_1|^2 + |h_2|^2 \end{pmatrix}$$

- $\rightarrow \frac{\sqrt{2}}{\sqrt{|h_1|^2+|h_2|^2}} \cdot \mathbf{F}$  is unitary matrix
- $\rightarrow \text{ ML decision: } \hat{\mathbf{s}} = \underset{\mathbf{s}}{\operatorname{argmin}} \left| \left| \frac{2}{|h_1|^2 + |h_2|^2} \cdot \mathbf{F}^H \cdot \tilde{\mathbf{y}} \mathbf{s} \right| \right|^2$
- $\rightarrow$  independent ML decoding

$$\hat{s}_1 = \underset{s_1}{\operatorname{argmin}} \left| s_1 - \frac{h_1^* y_1 + h_2 y_2^*}{\frac{1}{\sqrt{2}} (|h_1|^2 + |h_2|^2)} \right|$$

$$\hat{s}_2 = \underset{s_2}{\operatorname{argmin}} \left| s_2 - \frac{h_1^* y_1 - h_2 y_2^*}{\frac{1}{\sqrt{2}} (|h_1|^2 + |h_2|^2)} \right|$$

- independent decoding property can be proved for all OSTBCs
- low complexity is at the expense of a rate-loss compared to other STBCs for  $N_T > 2$ 
  - $\rightarrow$  Frequenzhopping
  - $\rightarrow$  keine Kanalinformation aus vorher empfangenen Symbolen möglich  $\Rightarrow$  Kanal ändert sich ständig: nur Entscheidung, ob Rauschen oder Signal + Rauschen
- Performance Analysis of Alamouti Code
  - Decision-variables after combining

$$r_1 = \sqrt{2} \frac{h_1^* y_1 + h_2 y_2^*}{|h_1|^2 + |h_2|^2}$$
$$r_2 = \sqrt{2} \frac{h_1^* y_1 - h_2 y_2^*}{|h_1|^2 + |h_2|^2}$$

because of symmetry it suffices to consider  $r_1$ 

$$\begin{split} r_1 &= \sqrt{2} \frac{h_1^* \left(\frac{1}{\sqrt{2}} s_1 h_1 + \frac{1}{\sqrt{2}} h_2 s_2 + n_1\right) + h_2 \left(-\frac{1}{\sqrt{2}} h_2 s_1^* + \frac{1}{\sqrt{2}} h_1 s_2^* + n_2\right)^*}{|h_1|^2 + |h_2|^2} \\ &= \sqrt{2} \frac{\frac{1}{\sqrt{2}} \left(|h_1|^2 + |h_2|^2\right) s_1 + h_1^* n_1 + h_2 n_2^*}{|h_1|^2 + |h_2|^2} \\ &= 1 \cdot s_1 + n_{eq} \end{split}$$

where

$$n_{eq} = \sqrt{2} \frac{h_1^* n_1 + h_2 n_2^*}{|h_1|^2 + |h_2|^2}$$

$$SNR \to \gamma_t = \frac{E_s \cdot 1^2}{\sigma_{eq}^2} \quad \text{with} \quad \mathcal{E}\{|s_1|^2\} = \mathcal{E}_s$$

$$\sigma_{eq}^2 = 2 \frac{|h_1|^2 \sigma_n^2 + |h_2|^2 \sigma_{eq}}{\left(|h_1|^2 + |h_2|^2\right)^2} = \frac{2\sigma_n^2}{|h_1|^2 + |h_2|^2}$$

$$\rightarrow \gamma_t = \frac{1}{2} \frac{E_s}{\sigma_n} (|h_1|^2 + |h_2|^2)$$

- $\rightarrow SNR_{Alamouti} = \frac{1}{2}SNR_{MRC} = \frac{1}{2}SNR_{MRT}$
- $\rightarrow$  Alamouti code has diversity gain  $G_d = 2$
- $\rightarrow$  Transmission with Alamouti STBC requires 3dB higher SNR to achieve same performance as MRT  $\rightarrow$  3dB loss in coding gain  $G_c$
- → Lack of CSI knowledge at transmitter "costs" 3dB in power efficiency
- $\rightarrow$  General:
  - · OSTBCs achieve a diversity gain of  $G_d = N_T$  if only one receive antenna is available
  - · if  $N_R$  receive antennas are available, MRC can be used at the receiver to yield a diversity gain of  $G_d = N_T N_R$
- Other STBCs:

- Quasi orthogonal STBCs
  - \* higher rate than OSTBCs
  - \* only subset of symbols have to be decoded jointly
  - \* Example:  $K = N_T = T = 4$

$$\mathbf{X} = \frac{1}{2} \begin{pmatrix} s_1 & s_2 & s_3 & s_4 \\ -s_2^* & s_1^* & -s_4^* & s_3^* \\ -s_3^* & -s_4^* & s_1^* & s_2^* \\ s_4 & -s_3 & -s_2 & s_1 \end{pmatrix}$$

- $\ast$  Anmerkung 1: X ist ähnlich zu Alamouti Code
- \* Anmerkung 2:  $\mathbf{X}^H \mathbf{X}$ : viele Nicht-diagonal Elemente sind Null; die, die ungleich Null sind, zeigen, welche Symbole gemeinsam entschlüsselt werden müssen
- Golden Code for  $N_T=N_R=2$ : achieves a rate of  $R_{STBC}=2$  and full diversity of  $G_d=N_T,N_R=4$
- Differential STBCs:  $\mathbf{X}_k = \mathbf{X}_{k-1} \cdot \mathbf{D}_k$ .  $\mathbf{X}_k$  is transmitted,  $\mathbf{D}_k$  is transmitted
- Linear dispersion codes: designed to achieve high mutual information
- noncoherent STBCs (On-Off-Keying)

# Space Time Code Design

Given:

- Code  $\mathscr{X} = \{\mathbf{X}_1, \dots, \mathbf{X}_{|\mathscr{X}|}\}$
- Channel: IID Rayleigh-fading:
  - $-h_n \sim C\mathcal{N}(0,1); \quad n \in \{1, 2, \dots, N_T\}$
  - AWGN  $n \sim C\mathcal{N}(0, \sigma_n^2)$

Problem: How should we design codebook  $\mathcal{X}$ ?

- Need to derive error rate for general codebooks  $\mathcal{X}$ !
  - Codeword error rate

$$P_e = \frac{1}{|\mathcal{X}|} \sum_{i=1}^{|\mathcal{X}|} \Pr{\{\mathbf{x}_i \neq \hat{\mathbf{x}}_i\}}$$

where  $\hat{\mathbf{x}}_i$  is the detected codeword and we assume that all codewords are equally likely

Problem:  $\Pr{\{\mathbf{x}_i \neq \hat{\mathbf{x}}_i\}}$  is not tractable in general

• Use union bound to upper bound  $\Pr\{\mathbf{x}_i \neq \hat{\mathbf{x}}_i\}$  as upper sum over pairwise error probabilities (PEP)  $\Pr\{\mathbf{x}_i \to \mathbf{x}_j\}$  where it is assumed that  $\mathbf{x}_i$  was transmitted and  $\overline{\mathbf{x}}_i$  and  $\overline{\mathbf{x}}_j$  are the only codewords in the codebook

$$P_e \leq \frac{1}{|\mathcal{X}|} \sum_{i=1}^{|\mathcal{X}|} \sum_{j=1}^{|\mathcal{X}|} \Pr{\{\mathbf{x}_i \to \mathbf{x}_j\} \text{ where } j \neq i}$$

#### Calculation of PEPs

Recall:  $\hat{\mathbf{x}} = \underset{\mathbf{x} \in \mathscr{X}}{\operatorname{argmin}} ||\mathbf{y} - \mathbf{x}\mathbf{h}||^2$ 

Now,  $\mathbf{x}_i$  and  $\mathbf{x}_j$  are the only alternatives and an error is made if  $||\mathbf{y} - \mathbf{x}_i \mathbf{h}||^2 > ||\mathbf{y} - \mathbf{x}_j \mathbf{h}||^2$  since  $\mathbf{x}_i$  was sent but  $\mathbf{x}_j$  was detected

$$\begin{aligned} & \rightarrow ||\mathbf{x}_{i}\mathbf{h} + \mathbf{n} - \mathbf{x}_{i}\mathbf{h}||^{2} > ||\mathbf{x}_{i}\mathbf{h} + \mathbf{n} - \mathbf{x}_{j}\mathbf{h}||^{2} \\ & \quad ||\mathbf{n}|| > ||(\mathbf{x}_{i} - \mathbf{x}_{j})\mathbf{h} + \mathbf{n}||^{2} \\ & \rightarrow ||\mathbf{n}|| > \underbrace{\mathbf{h}^{H}(\mathbf{x}_{i} - \mathbf{x}_{j})^{H}(\mathbf{x}_{i} - \mathbf{x}_{j})\mathbf{h}}_{\Delta} + \mathbf{h}^{H}(\mathbf{x}_{i} - \mathbf{x}_{j})\mathbf{n} + \mathbf{n}^{H}(\mathbf{x}_{i} - \mathbf{x}_{j})\mathbf{h} + ||\mathbf{n}||^{2} \end{aligned}$$

$$\rightarrow \underbrace{-\mathbf{h}^{H}(\mathbf{x}_{i} - \mathbf{x}_{j})^{H}\mathbf{n} - \mathbf{n}^{H}(\mathbf{x}_{i} - \mathbf{x}_{j})\mathbf{h}}_{\mathbf{r}} > \Delta$$

for given  $\mathbf{h}$ , z is a gaussian random variable

$$\sigma_z^2 = \mathcal{E}\{|z|^2\} = \mathcal{E}\{2\mathbf{h}^H(\mathbf{x}_i - \mathbf{x}_j) \overbrace{\mathbf{n}\mathbf{n}^H}^{\sigma_n^2 \mathbf{I}} (\mathbf{x}_i - \mathbf{x}_j)\mathbf{h} + 2\mathbf{h}^H(\mathbf{x}_i - \mathbf{x}_j)^H \overbrace{\mathbf{n}\mathbf{n}^T}^{=0} (\mathbf{x}_i - \mathbf{x}_j)^* \mathbf{h}^*\}$$

$$= 2\sigma_n^2 \Delta + 0$$

$$\Pr\{\mathbf{x}_i \to \mathbf{x}_j\} = \int_{\Delta}^{\infty} \frac{1}{\sqrt{2\pi}\sigma_z} \exp\left(-\frac{z^2}{2\sigma_z^2}\right) dz, \ t = \frac{z}{\sigma_z}$$
$$= \frac{1}{\sqrt{2\pi}} \int_{\frac{\Delta}{\sigma_z}}^{\infty} e^{-\frac{t^2}{2}} dt = Q\left(\frac{\Delta}{\sigma_z}\right) = Q\left(\frac{\Delta}{\sqrt{2\sigma_n^2 \Delta}}\right)$$
$$= Q\left(\sqrt{\frac{\Delta}{2\sigma_n^2}}\right)$$

- $\Pr{\mathbf{x}_i \to \mathbf{x}_j} = \mathcal{E}\left{Q\left(\sqrt{\frac{\Delta}{2\sigma_n^2}}\right)\right}$ 
  - to avoid cumbersome Q-function we use Chernoff bound:

$$Q(x) \le \frac{1}{2}e^{-\frac{x^2}{2}}$$

$$\Pr{\mathbf{x}_i \to \mathbf{x}_j} \le \frac{1}{2} \mathcal{E}_h \left\{ \exp\left(-\frac{\mathbf{h}^H \mathbf{Q} \mathbf{h}}{4\sigma_n^2}\right) \right\}$$
where  $\mathbf{Q} = (\mathbf{x}_i - \mathbf{x}_j)^H (\mathbf{x}_i - \mathbf{x}_j)$ 

- Eigendecomposition:  $\mathbf{Q} = \mathbf{U}^H \mathbf{\Lambda} \mathbf{U}$  with  $\mathbf{\Lambda} = \text{diag}\{\lambda_1, \dots, \lambda_r, 0, \dots, 0\}$   $r = \text{rank}\{\mathbf{Q}\}$
- Elements h are i.i.d. Gaussian
  - $-\underline{\beta} = \mathbf{U}\mathbf{h}$  has also i.i.d. Gaussian random variables as elements since  $\mathbf{U}$  is unitary matrix

$$- \mathbf{h}^{H}\mathbf{Q}\mathbf{h} = \underbrace{\mathbf{h}^{H}\mathbf{U}^{H}}_{\beta^{*}}\mathbf{\Lambda}\underbrace{\mathbf{U}\mathbf{h}}_{\beta} = \sum_{i=1}^{r} \lambda_{i}|\beta_{i}|^{2} \text{ with } \underline{\beta} = [\beta_{1}, \dots, \beta_{N_{T}}]$$

$$\begin{aligned} \Pr\{\mathbf{x}_i \to \mathbf{x}_j\} &= \frac{1}{2} \mathcal{E}_{\underline{\beta}} \left\{ \exp\left(-\frac{\sum\limits_{i=1}^r \lambda_i |\beta_i|^2}{4\sigma_n^2}\right) \right\} \\ &= \frac{1}{2} \mathcal{E}_{\underline{\beta}} \left\{ \prod\limits_{i=1}^r e^{-\frac{\lambda_i}{4\sigma_n^2} |\beta_i|^2} \right\} \\ &= \frac{1}{2} \prod\limits_{i=1}^r \mathcal{E}_{\beta_i} \left\{ e^{-\frac{\lambda_i}{4\sigma_n^2} |\beta_i|^2} \right\} \\ &= \frac{1}{2} \prod\limits_{i=1}^r \mathcal{E}_{|\beta_i|^2} \left\{ e^{-\frac{\lambda_i}{4\sigma_n^2} |\beta_i|^2} \right\} \stackrel{\frown}{=} \text{MGF of exponentially distributed variable } \alpha_i = |\beta_i|^2 \end{aligned}$$

$$\to P_{\alpha_i}(x) = e^{-x}, \ x \ge 0$$

• upper bound on  $P_e$ :

$$\begin{split} & \lambda_n(i,j) = n \text{th eigenvalue of } (\mathbf{x}_i - \mathbf{x}_j)^H (\mathbf{x}_i - \mathbf{x}_j) \\ & r(i,j) = \text{ rank of } (\mathbf{x}_i - \mathbf{x}_j)^H (\mathbf{x}_i - \mathbf{x}_j) \end{split}$$

$$\rightarrow P_e \leq \frac{1}{|\mathcal{X}|} \sum_{i=1}^{|\mathcal{X}|} \sum_{j=1}^{|\mathcal{X}|} 2^{2r(i,j)-1} \frac{1}{\prod\limits_{i=1}^{r(i,j)} \lambda_n(i,j)} \left(\frac{1}{\sigma_n^2}\right)^{-r(i,j)}$$

• generally loose bound but offers significant insight for code design

#### Two criteria:

**Rank criterion:** The diversity gain of a ST code is given by  $G_d = \min_{i,j} (r(i,j)) = \min_{i,j} \operatorname{rank} \left( (\mathbf{x}_i - \mathbf{x}_j)^H (\mathbf{x}_i - \mathbf{x}_j) \right)$   $\rightarrow$  Design code such that minimum rank of all possible matrices  $(\mathbf{x}_i - \mathbf{x}_j)^H (\mathbf{x}_i - \mathbf{x}_j)$  is maximized

$$T \updownarrow \overset{\stackrel{N_T}{\longleftrightarrow}}{\mathbf{X}_i} \Rightarrow r(i,j) = N_T \qquad \forall i \neq j$$

**<u>Determinant criterion:</u>** To maximize the coding gain among all codes with  $r(i,j) = N_T$ , we need to maximize  $\max \min_{i,j} \prod_{n=1}^{N_T} \lambda_n(i,j) = \max \min_{i,j} \left| (\mathbf{x}_i - \mathbf{x}_j)^H (\mathbf{x}_i - \mathbf{x}_j) \right| \quad \forall i \neq j$ 

- Rank and determinant criterion can be used for the search for good space-time block codes and space-time trellis codes. These two criteria were first derived by Tarokh, et. al. 1998.
- ullet diversity increases to  $N_T N_R$  if  $N_K$  receive antennas are available
- Example: see Bäro, Bauch, Hansmann: Improved codes for space-time trellis coded modulation. IEEE Comm. Letters, 2000.