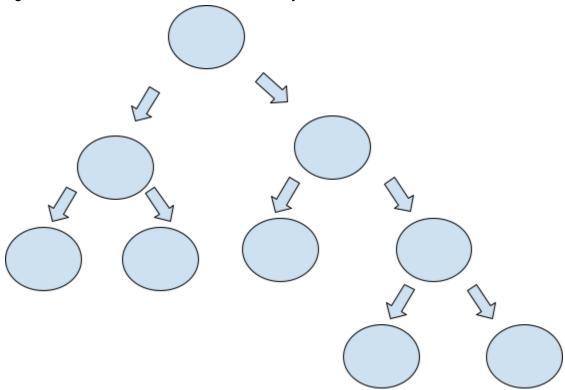
Quentin Creese Jess CS 260 Assignment 06 Due: February 26, 2024

## **Binary Search Tree**

1. Before you start coding, create a design that shows how your binary tree functions and what attributes it keeps track of to function.

Design: The top of the tree should be the middle of the tree. The right of the tree should always have a higher value. The left of the tree should always have a lower value.



Attributes to keep track of to function:

- A starting location "root"
- A value called "val"
- Current value called "current"
- Call back value to help node traversal called "callback"
- Push and pop functions to remove and call nodes
- A temporary value to help rearrange nodes called "temp"
- Left and right pointers to help with node traversal called "left" and "right".
- 2. Create some tests (at least one per function), before you start coding, that you want your Binary Search Tree (BST) to pass as evidence that it would be working correctly if it passed the tests
  - I would like to test if my BST is empty to begin
  - I would like to test that my add function is working by adding a node to my BST

- I would like to test that my remove function is working by removing a node in my BST
- I would like to test my function for pre-order traversal
- I would like to test my function for post-order traversal
- I would like to test my function for breadthFirstTraversal
- 3. Implement a binary search tree that includes:
  - 1. Nodes to store values:

2. An add function that adds a value to the appropriate location based on our ordering rules:

```
void BinarySearchTree::add(int val) {
    root = addNode(root, val); // Add a value to the binary search tree
}
```

```
// Private helper functions for the binary search tree operations
TreeNode* BinarySearchTree::addNode(TreeNode* node, int val) {
    if (node == nullptr) {
        return new TreeNode(val);
    }

    // Recursively insert the value based on the binary search tree property
    if (val <= node->val) {
        node->left = addNode(node->left, val);
    } else {
        node->right = addNode(node->right, val);
    }

    return node;
}
```

3. A remove function that finds and removes a value and then picks an appropriate replacement node:

```
void BinarySearchTree::remove(int val) {
    root = removeNode(root, val); // Remove a value from the binary search tree
}
```

```
TreeNode* BinarySearchTree::removeNode(TreeNode* node, int val) {
    if (node == nullptr) {
       return nullptr;
   if (val < node->val) {
       node->left = removeNode(node->left, val);
    } else if (val > node->val) {
       node->right = removeNode(node->right, val);
    } else {
       if (node->left == nullptr) {
           TreeNode* temp = node->right;
           delete node;
           return temp;
       } else if (node->right == nullptr) {
           TreeNode* temp = node->left;
           delete node;
           return temp;
       TreeNode* temp = findMinNode(node->right);
       node->val = temp->val;
       node->right = removeNode(node->right, temp->val);
    return node;
```

- 4. At least one tree traversal function:
  - a. The three common traversals (pre-order, post-order, in-order):

```
void BinarySearchTree::preOrder(std::function<void(TreeNode*)> callback) {
          preOrder(root, callback); // Perform pre-order traversal starting from the root
20
      void BinarySearchTree::inOrder(std::function<void(TreeNode*)> callback) {
          inOrder(root, callback); // Perform in-order traversal starting from the root
      }
      void BinarySearchTree::postOrder(std::function<void(TreeNode*)> callback) {
          postOrder(root, callback); // Perform post-order traversal starting from the root
26
100
      void BinarySearchTree::preOrder(TreeNode* node, std::function<void(TreeNode*)> callback) {
101
          if (node == nullptr) return;
102
          callback(node); // Process the current node
103
          preOrder(node->left, callback); // Traverse left subtree
104
          preOrder(node->right, callback); // Traverse right subtree
105
107
      void BinarySearchTree::inOrder(TreeNode* node, std::function<void(TreeNode*)> callback) {
108
          if (node == nullptr) return;
109
          inOrder(node->left, callback); // Traverse left subtree
110
          callback(node); // Process the current node
111
          inOrder(node->right, callback); // Traverse right subtree
112
      }
113
114
      void BinarySearchTree::postOrder(TreeNode* node, std::function<void(TreeNode*)> callback) {
115
          if (node == nullptr) return;
116
          postOrder(node->left, callback); // Traverse left subtree
117
          postOrder(node->right, callback); // Traverse right subtree
118
          callback(node); // Process the current node
119
```

b. A breadth-first traversal (sometimes called a level-order search):

```
void BinarySearchTree::breadthFirst(std::function<void(TreeNode*)> callback) {
    if (root == NULL) return; // If the tree is empty, return immediately

std::queue<TreeNode*> q;
    q.push(root);

// Perform breadth-first traversal using a queue
while (!q.empty()) {
        TreeNode* current = q.front();
        q.pop();
        callback(current); // Call the callback function for the current node

// Enqueue left and right child nodes if they exist
if (current->left != NULL)
        q.push(current->left);
if (current->right != NULL)
        q.push(current->right);
}

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```

## Main output:

```
quentin@Quentins-MacBook-Pro CS260_assignment06 % clang++ -o main main.cpp BinarySearchTree.cpp -std=c++11 -stdlib=libc++ && ./main In-order traversal: 20 20 30 40 40 20 30 40 50 60 70 80 60 60 70 80 80 Pre-order traversal: 50 30 20 40 70 60 80 30 20 40 20 40 70 60 80 60 80 Post-order traversal: 20 40 20 40 30 60 80 60 80 70 50 Breadth-first traversal: 50 30 70 20 40 60 80
```

## Test output:

```
• quentin@Quentins-MacBook-Pro CS260_assignment06 % clang++ -o BST_Test BinarySearchTree.cpp BST_Test.cpp -std=c++11 -stdlib=libc++
• quentin@Quentins-MacBook-Pro CS260_assignment06 % ./BST_Test
In-order traversal: 20 30 40 50 60 70 80
Pre-order traversal: 50 30 20 40 70 60 80
Post-order traversal: 20 40 30 60 80 70 50
Breadth-first traversal: 50 30 70 20 40 60 80
In-order traversal after removal: 40 50 60 70 80
o quentin@Quentins-MacBook-Pro CS260_assignment06 %
```

4. Analyze and compare the complexity of insert and search as compared to a binary tree without any order in its nodes (what is the run time of an unordered list?)

The efficiency of the binary search tree (BST) we created today depends on how balanced the tree remains during insertion and removal operations. In a well-balanced BST, where nodes are inserted so the tree remains balanced, the efficiency of insertion and search operations is typically O(log n) on average, where n is the number of nodes in the tree. This is because the height of the tree remains logarithmic with respect to the number of nodes, allowing for efficient traversal and search.

However, if the BST becomes unbalanced, such as in the worst-case scenario where nodes are inserted in sorted order, the efficiency of insertion and search operations degrades to O(n), where n is the number of nodes in the tree. This worst-case scenario occurs when the BST resembles a linked list, with each node having only one child.

Therefore, to ensure efficient operations in the BST we created today, it's essential to maintain balance during insertion and removal operations.