

Parsimonious Black-Box Adversarial Attacks Via Efficient Combinatorial Optimization

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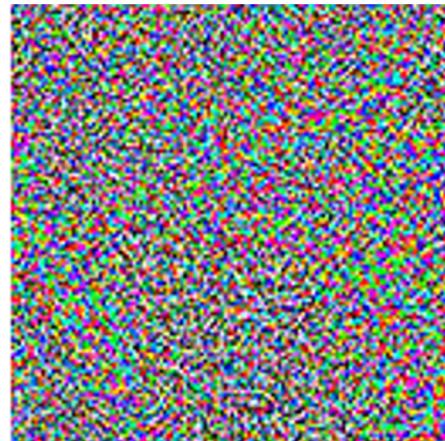


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Adversarial Perturbations



$+ \epsilon$



=



“panda”

57.7% confidence

“gibbon”

99.3% confidence

White vs Black Box Attacks

- White Box: access to parameters and therefore gradient
 - Fast Gradient Sign Method (FGSM): perturb in direction of gradient
 - Projected Gradient Descent (PGD): multiple iterations of FGSM
- Black Box: only query access
 - Substitute networks: train a network to match predictions of target network
 - Gradient estimation: directly estimate gradient via queries
 - Outperforms substitute networks

Motivation

- Focus on black-box, which is more realistic in practice
- Problems with current black-box methods
 - Substitute network attacks don't always transfer to target networks
 - Robustness of gradient estimation affected by choice of hyperparameters
 - E.g. learning rate, decay rates, update rule

Problem formulation

- Create imperceptible perturbations x_{adv} under L_∞ radius with limited query budget to maximize loss
- Attacker only has access to loss function, $l(x,y)$

$$\underset{\|x_{adv} - x\|_\infty \leq \epsilon}{\text{maximize}} \quad \ell(x_{adv})$$

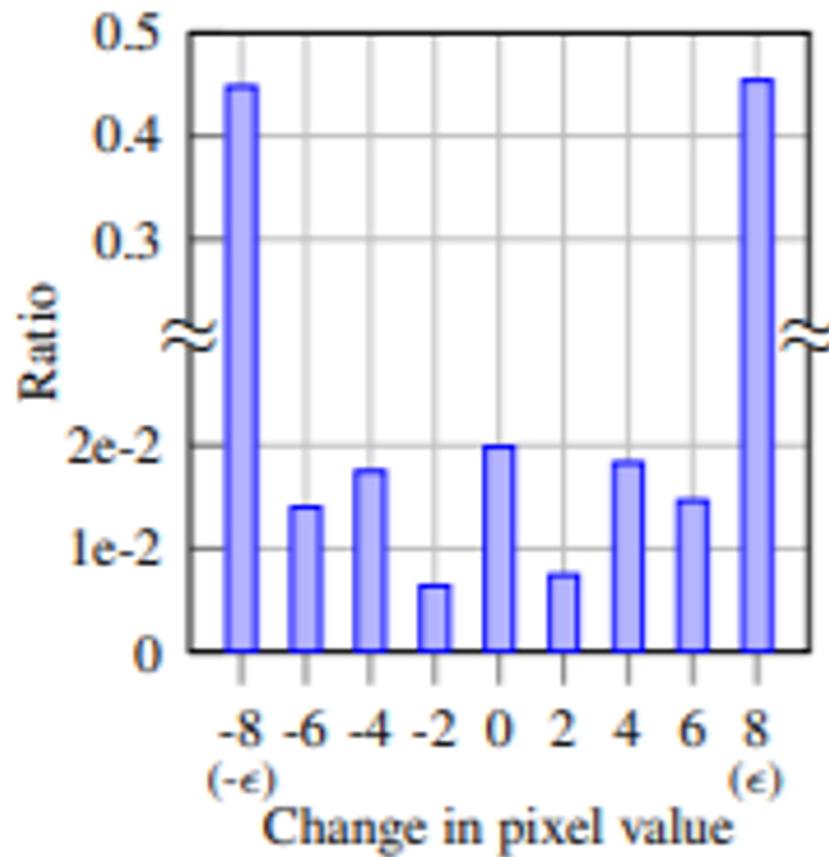
FGSM Approximation

$$\underset{\|x_{adv} - x\|_\infty \leq \epsilon}{\text{maximize}} \ell(x_{adv}) \implies \underset{x_{adv}}{\text{maximize}} x_{adv}^\top \nabla_x \ell(x, y) \quad (1)$$

subject to $-\epsilon \mathbf{1} \preceq x_{adv} - x \preceq \epsilon \mathbf{1},$

- Where the \leq is element-wise inequality and $\mathbf{1}$ is a vector of ones
- Optimal solution will be obtained at extreme point of feasible set, or a vertex of the L_∞ ball
- PGD on Cifar-10 does give solutions close to vertices of L_∞ ball

PGD Pixel-level Perturbations



Discrete formulation

- Only consider pixel perturbations of $+\/-\epsilon$

$$\underset{x_{adv} \in \mathbb{R}^p}{\text{maximize}} \ f(x_{adv}) \quad \Rightarrow \quad \underset{x_{adv}}{\text{maximize}} \ f(x_{adv}) \quad (2)$$

subject to $\|x_{adv} - x\|_\infty \leq \epsilon$ subject to $x_{adv} - x \in \{\epsilon, -\epsilon\}^p$,

- $f(x) = I(x, y_{gt})$ for untargeted attacks, $-I(x, y_{target})$ for targeted attacks
- Set maximization problem in which we choose from all pixels \vee a set S with $+\epsilon$ perturbations, with the rest having $-\epsilon$ perturbations

Submodularity

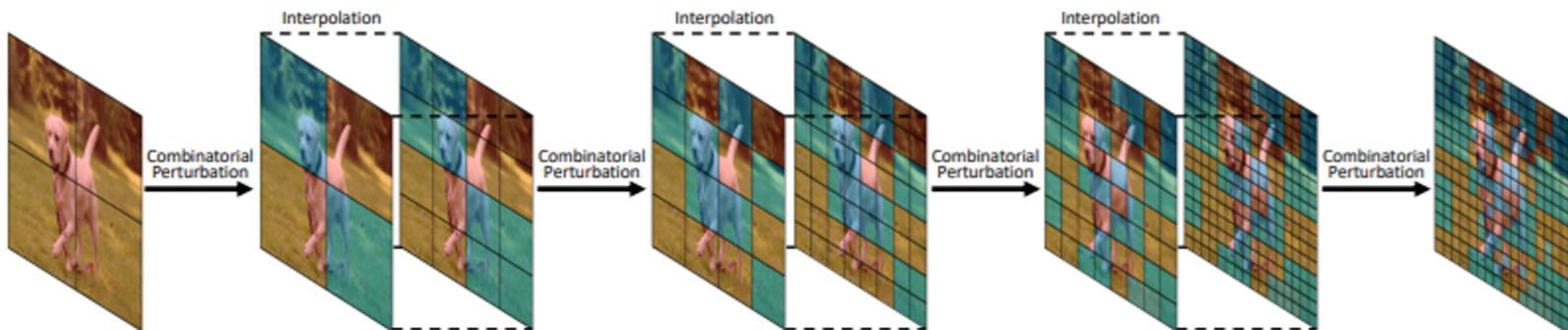
- Define $F(S \cup \{e\}) - F(S)$ as the marginal gain from adding pixel e to S
 - Submodularity implies that this marginal gain will be smaller when S has more elements
 - “Diminishing returns”
- This is not completely true but algorithm assumes it is approximately true to cut save queries

Lazy Greedy Insertion

- First, query marginal gain for all elements not in S , and insert these elements into a max heap
 - this is treated as an upper bound because of submodularity
- While the heap isn't empty:
 - Pop the top element, update its upper bound
 - If it's greater than the new top element
 - If it's > 0 , add it to S
 - else, end
 - If it's less than the new top element
 - Add it back to the heap

Implementation

- Exploit locally regular structure to do hierarchical evaluation
 - At each level, do one iteration of lazy insertion then lazy deletion
 - Terminate when converges or query limit reached



Diminishing gains

Submodular set functions are set functions who exhibit **diminishing returns**

for all $A \subseteq B$

$$\underline{F(A \cup e) - F(A)} \geq \underline{F(B \cup e) - F(B)}$$

Basically: as the size of the input set increases, the value that a single element adds decreases

Our problem: approximate submodularity

- As it turns out, our problem is not technically submodular
- However, as long as submodularity is not “severely deteriorated” (Zhou & Spanos, 2016), submodular maximization algorithms still work very well
- This means that we can compute an approximately optimal solution with a greedy algorithm!

Local-search optimization

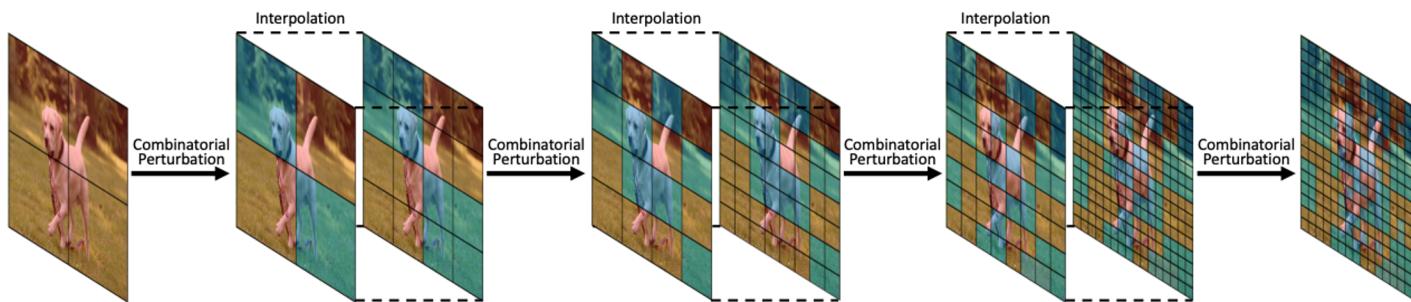
- Notation
 - \mathbf{P} is all pixels,
 - \mathbf{S} is pixels to add $+\epsilon$ to
 - $\mathbf{P} \setminus \mathbf{S}$ are pixels to add $-\epsilon$
- Basically, we can greedily choose to insert a pixel into S if the marginal gain is strictly positive, and remove it from S if the marginal gain is strictly negative
 - Then once the algorithm converges, it will converge to a **local optimum**
- End up with a set of pixels S to perturb the input image with $+\epsilon$, and $\mathbf{P} \setminus \mathbf{S}$ to perturb with $-\epsilon$

Speedup #1: Acceleration with lazy evaluations

- At each step we have to find the element that maximizes the marginal gain
 - Therefore, our greedy algorithm has to make $O(|P| |S|)$ queries
 - This may be impractical for query-limited black-box attacks
- Speed this up: use the Lazy-Greedy algorithm (Minoux, 1978)
 - Instead of re-computing the marginal gain for each pixel at each iteration, keep the upper bounds on the marginal gains in a **max-heap**
 - Theoretically has the same worst-case number of function evaluations but provides a speedup of several orders of magnitude in practice!
 - Why? Because of submodularity! (wow)

Speedup #2: Hierarchical lazy evaluation

- Exploit the locally regular structure of most images and do this on a hierarchical scale for another speed boost



Blue squares are in S , red squares are in $P \setminus S$

Experimental Results

Method	Success rate	Avg. queries	Med. queries	Avg. queries (NES success)
PGD (white-box)	47.2%	20	-	-
NES	29.5%	2872	900	2872
Bandits	38.6%	1877	459	520
Ours	48.0%	1261	356	247

Table 1. Results for ℓ_∞ untargeted attacks on Cifar-10. Maximum number of queries set to 20,000.

Method	Success rate	Avg. queries	Med. queries	Avg. queries (NES success)
PGD (white-box)	99.9 %	20	-	-
NES [†]	77.8%	1735	-	1735
NES	80.3%	1660	900	1660
Bandits [†]	95.4%	1117	-	703
Bandits	94.9%	1030	286	603
Ours	98.5%	722	237	376

Table 2. Results for ℓ_∞ untargeted attacks on ImageNet. Maximum number of queries set to 10,000.

Method	Success rate	Avg. queries	Med. queries	Avg. queries (NES success)
PGD (white-box)	100%	200	-	-
NES [†]	99.2%	-	11550	-
NES	99.7%	16284	12650	16284
Bandits	92.3%	26421	18642	26421
Ours	99.9%	7485	5373	7371

Table 3. Results for ℓ_∞ targeted attacks on ImageNet. Maximum number of queries set to 100,000.

Conclusion

- Practical method for black-box adversarial attacks
- No gradient estimation required
 - No update hyperparameters
- State of the art success and query rates for both targeted and untargeted attacks