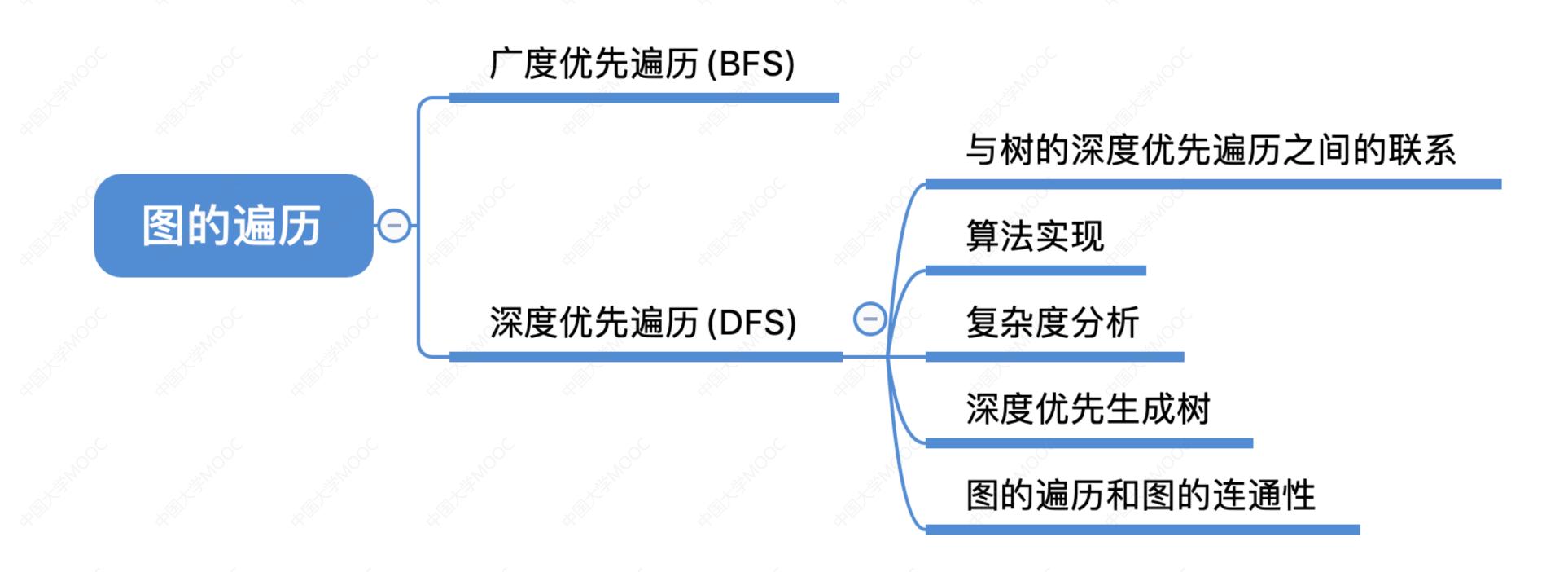
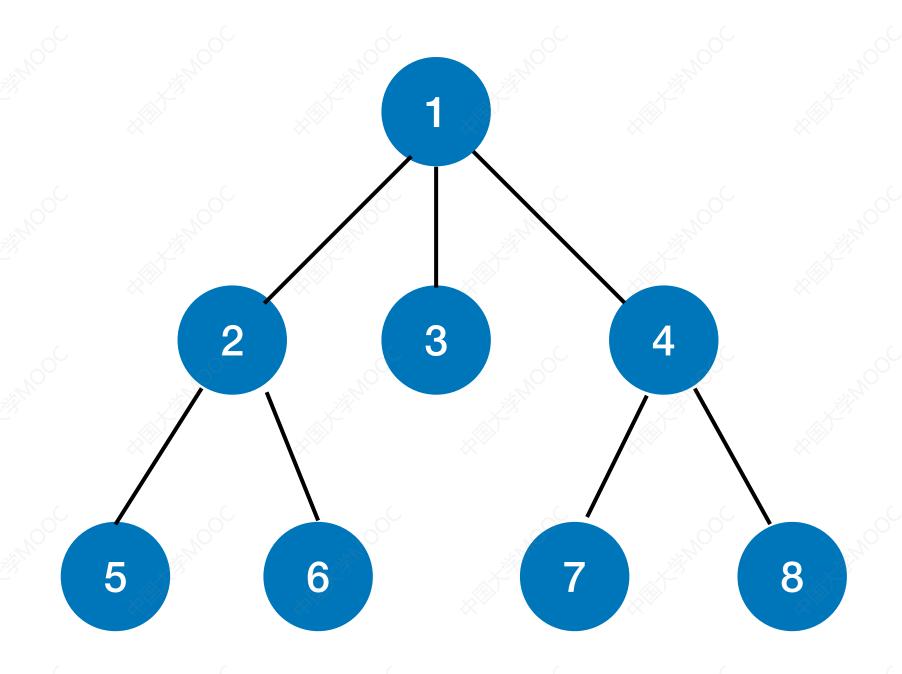
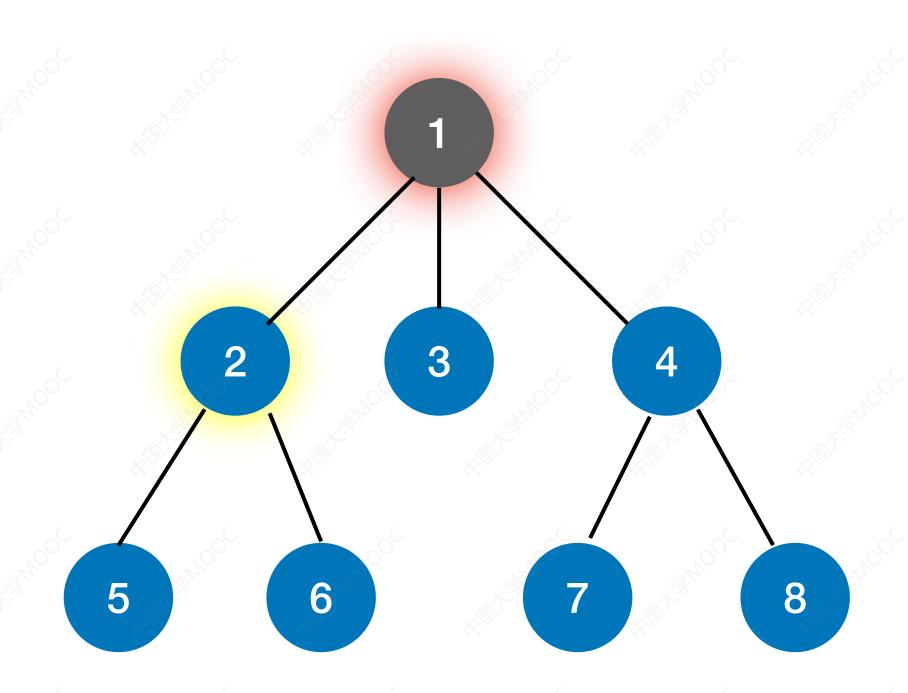
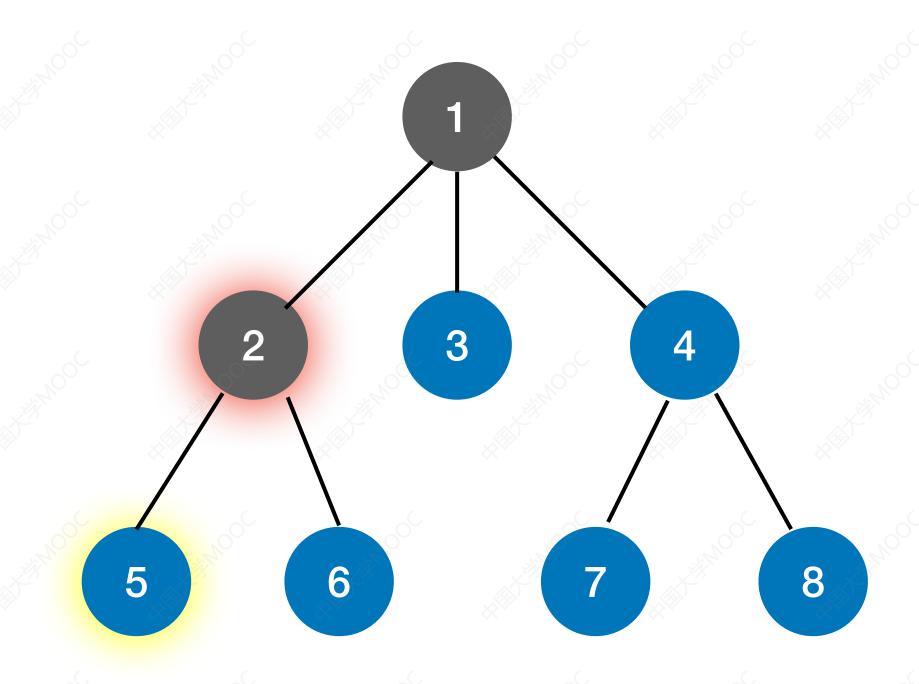
# 本节内容 图的遍历 DFS

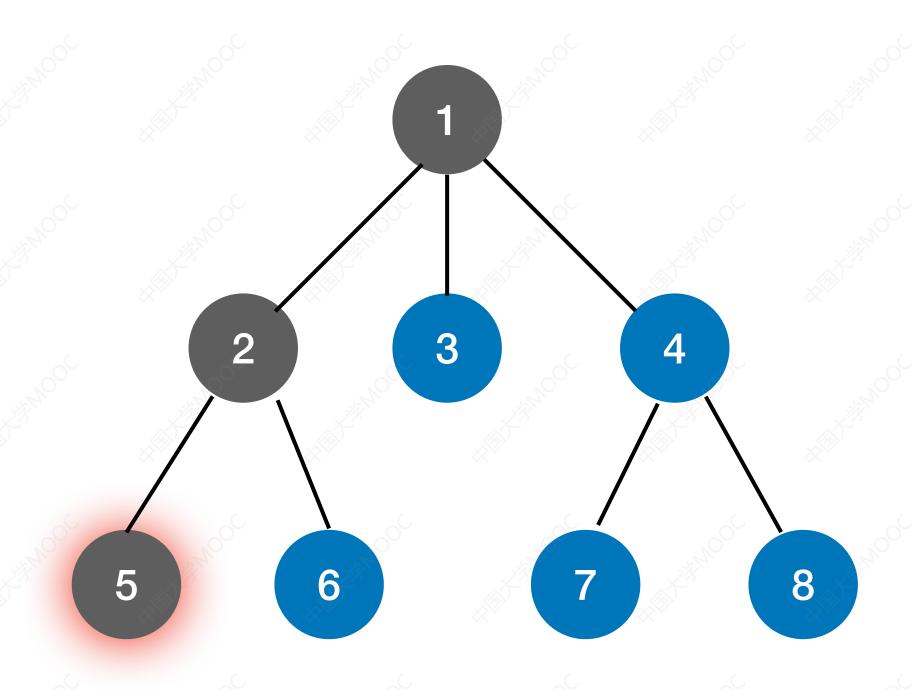
# 知识总览

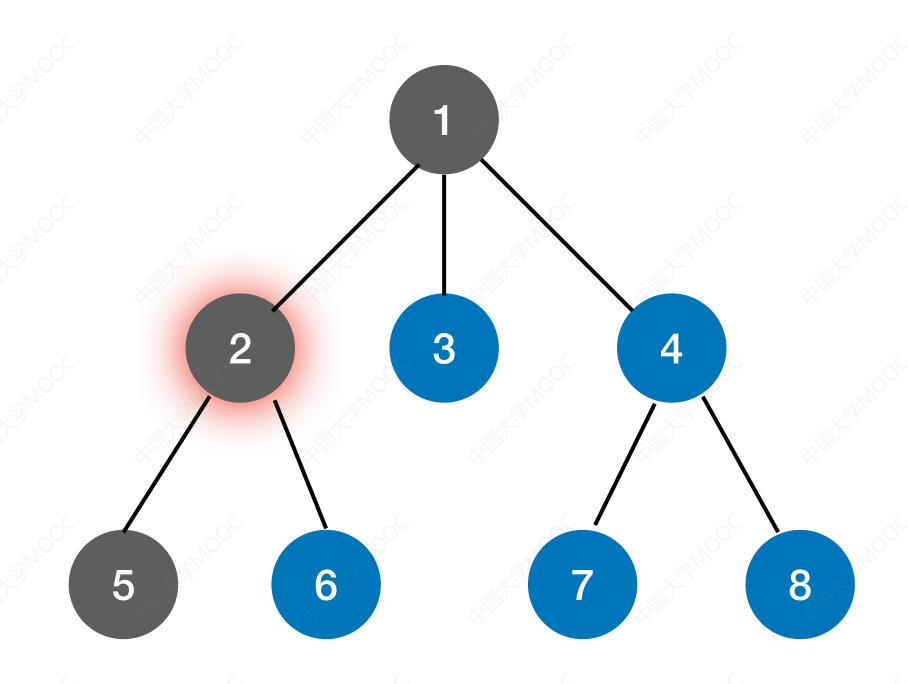


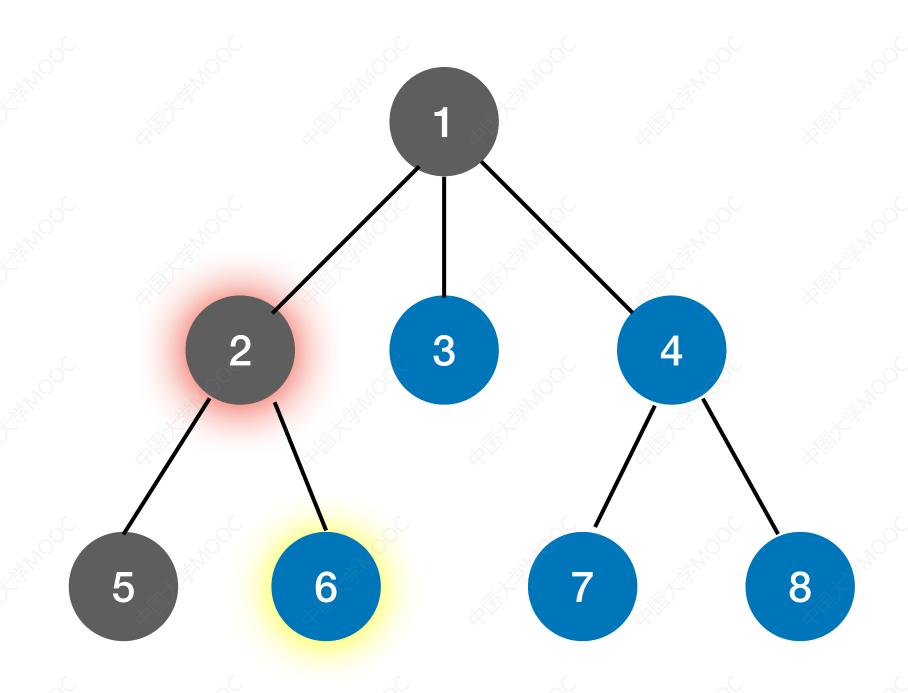


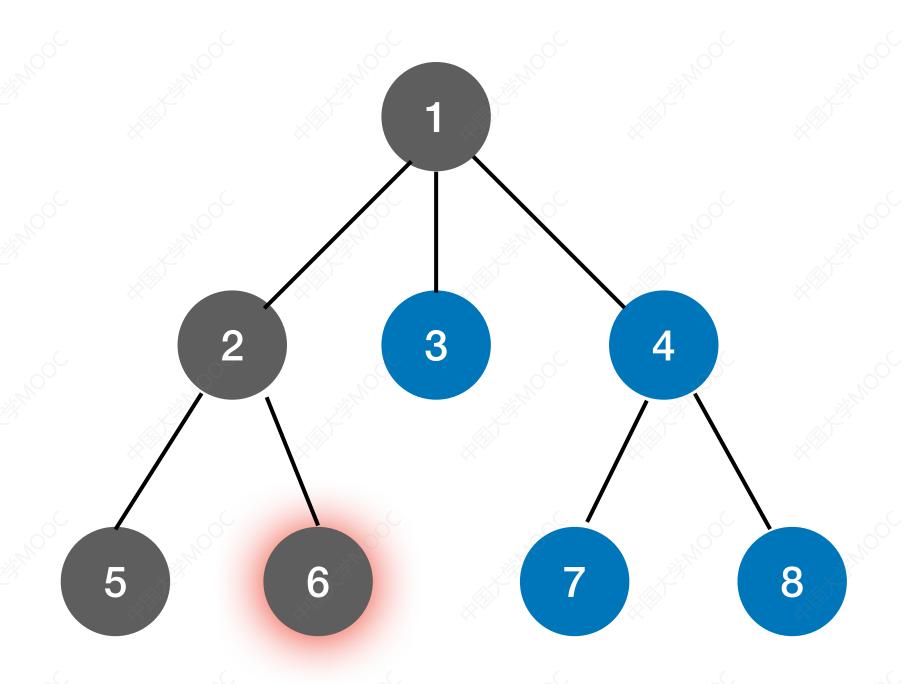


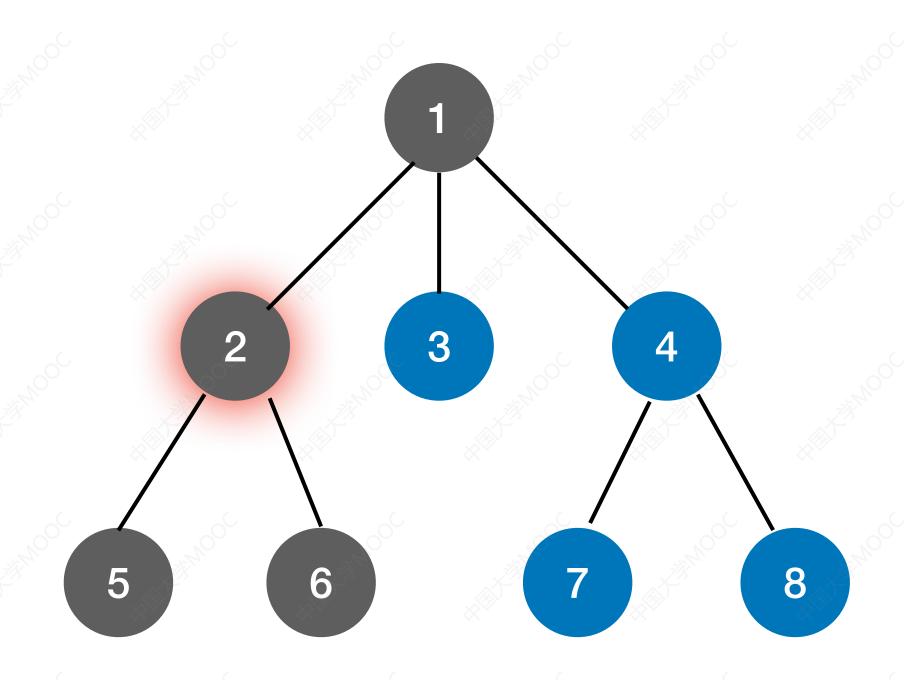


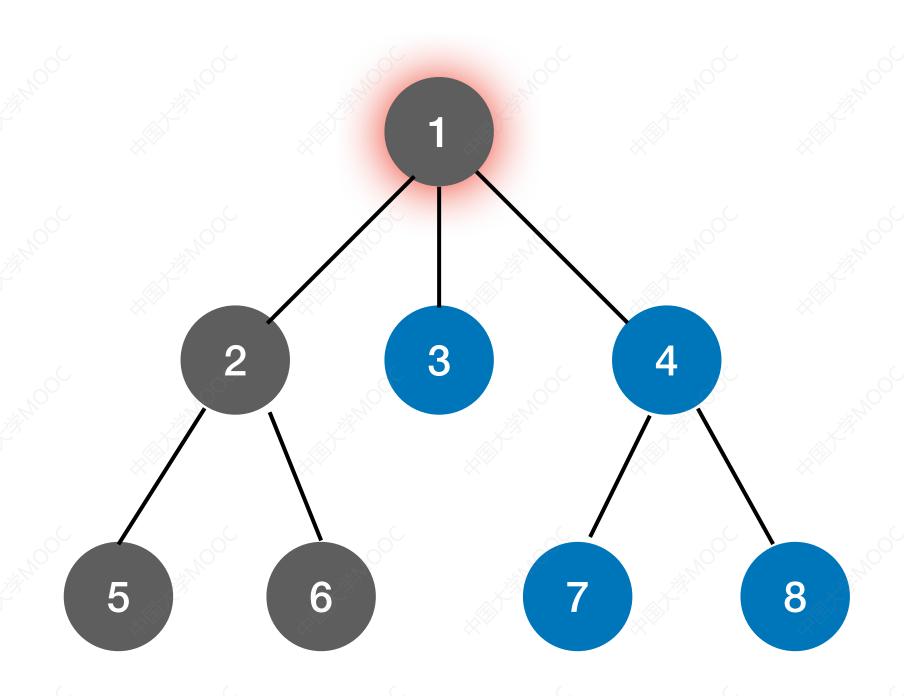


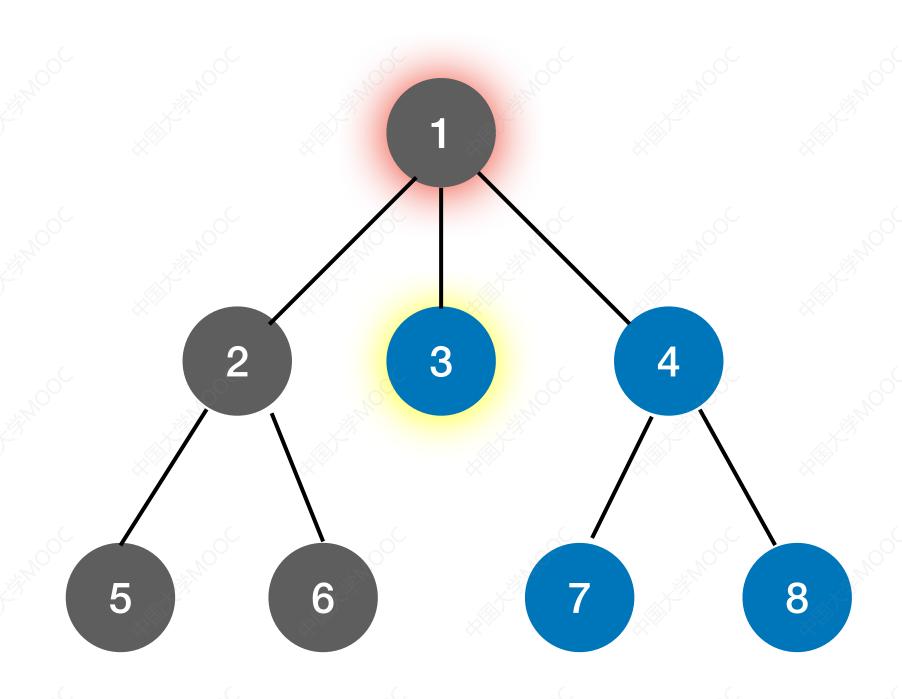


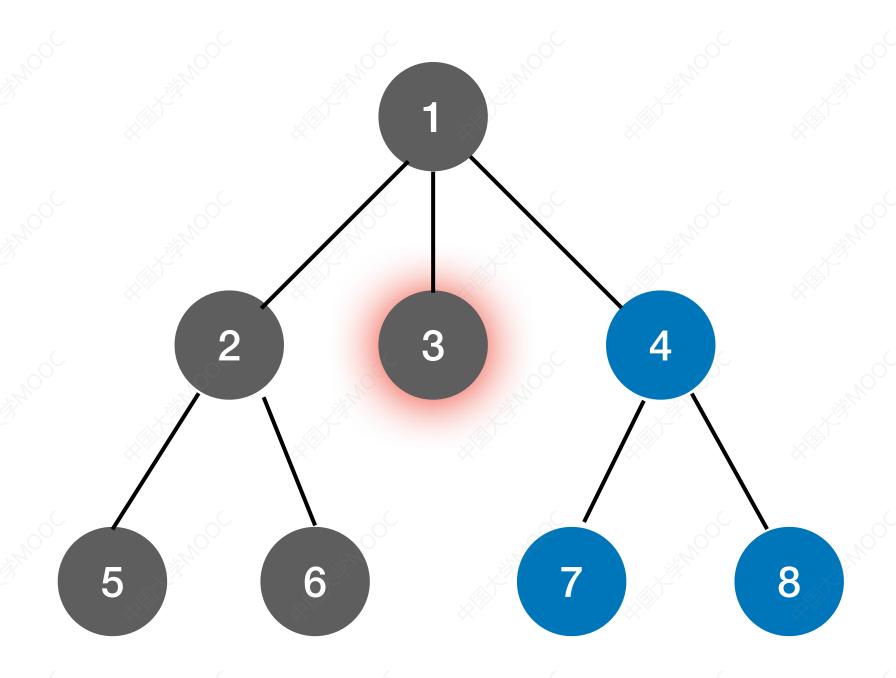


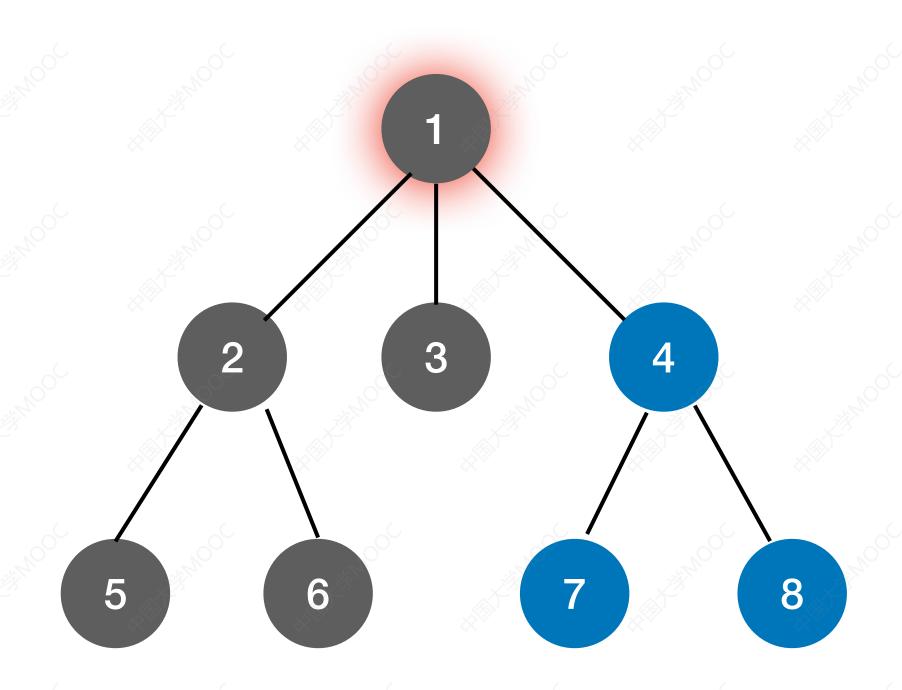


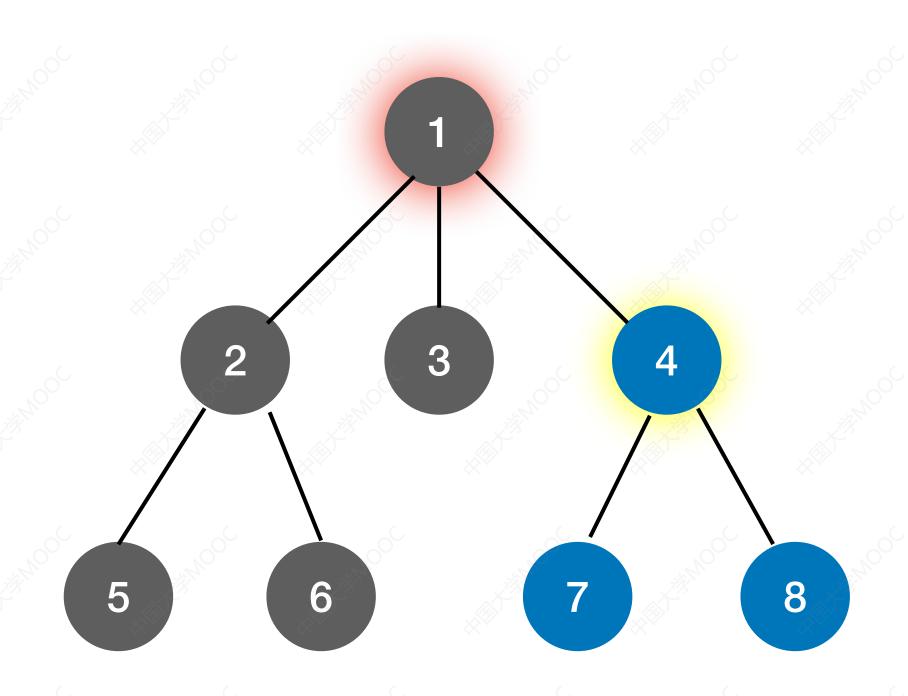


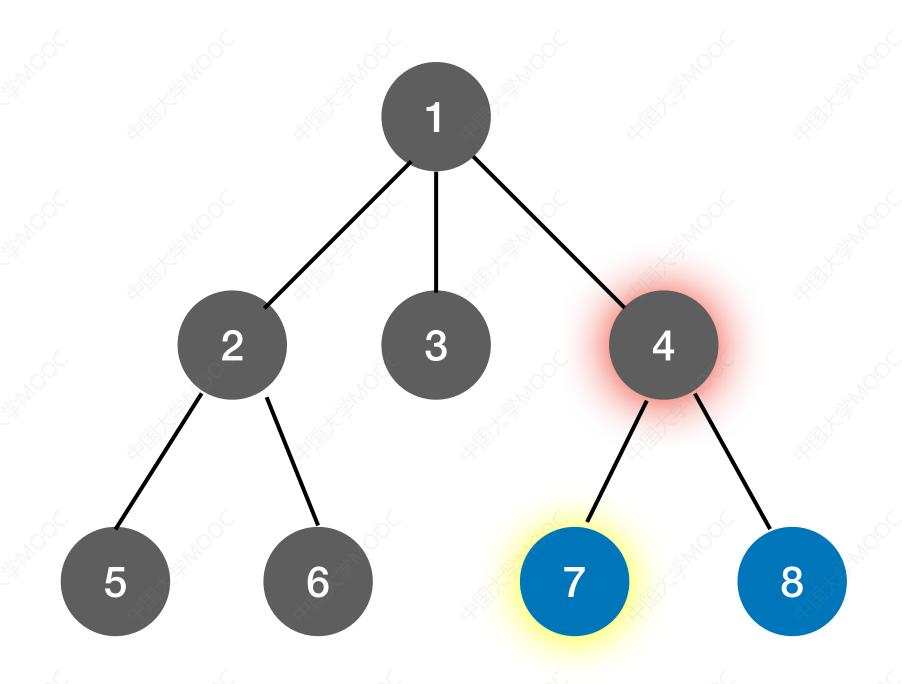


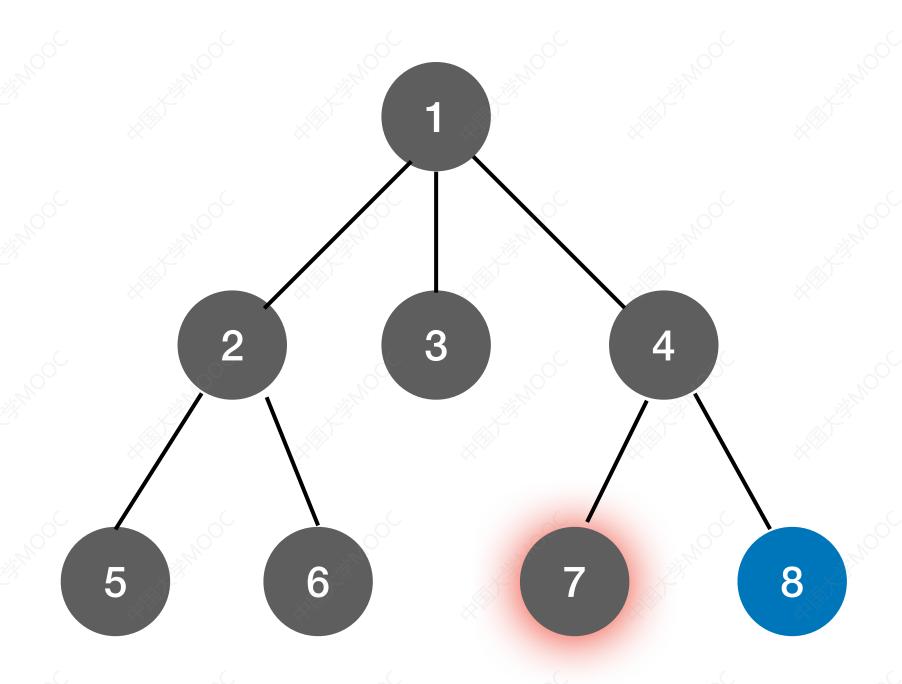


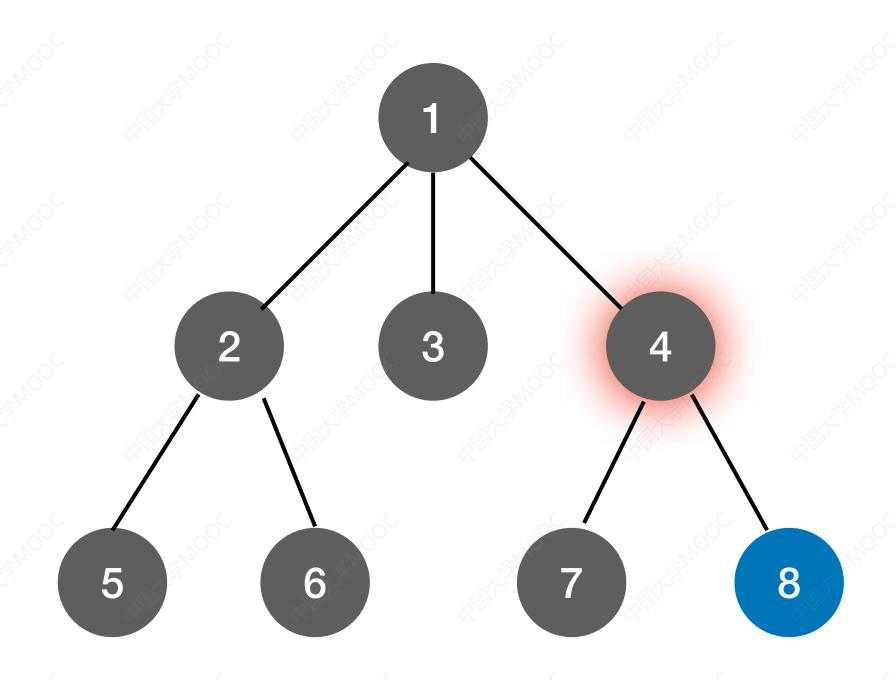


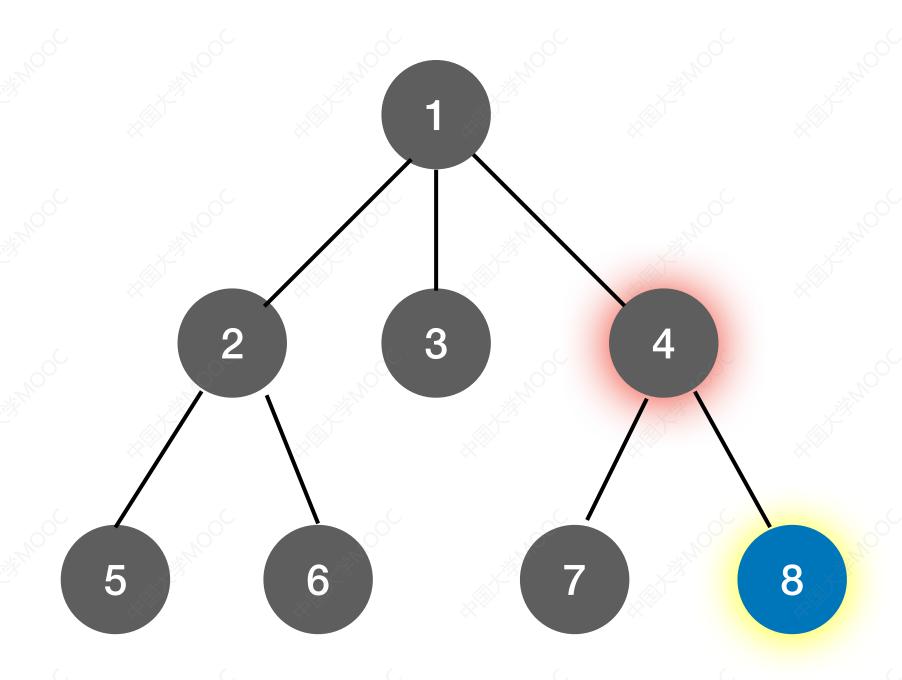


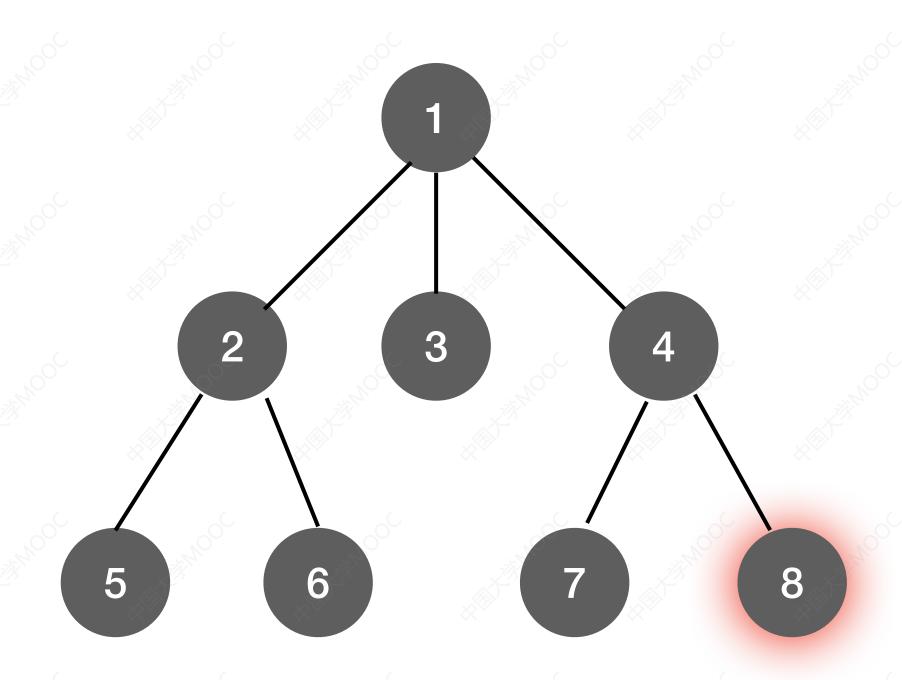


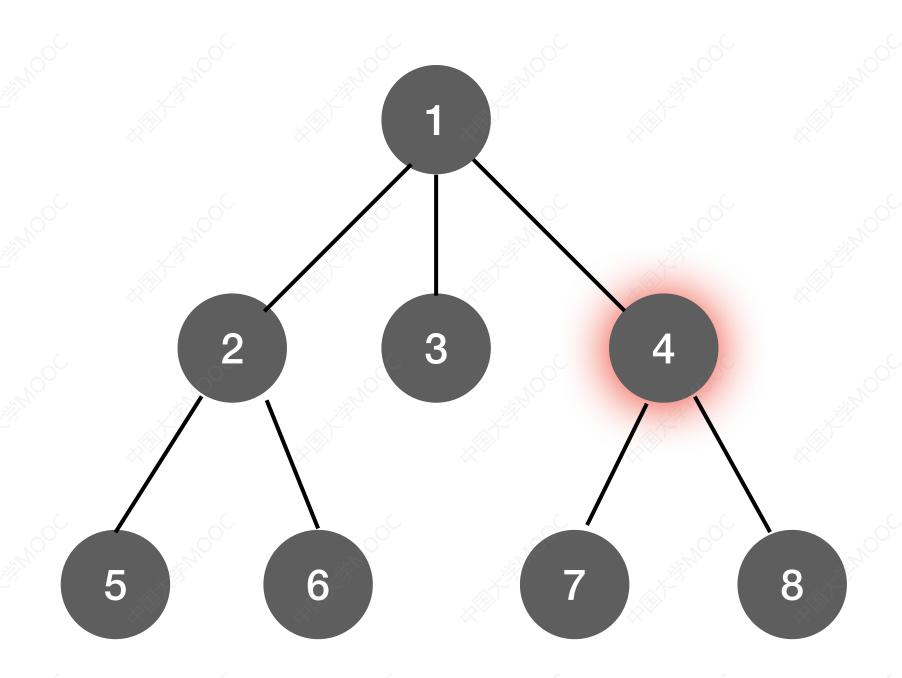


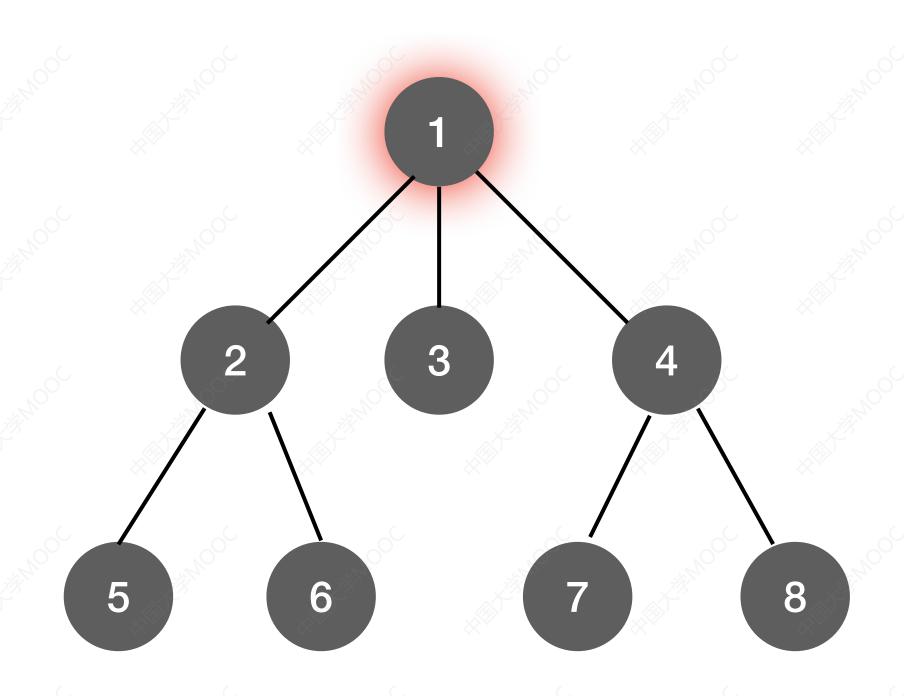


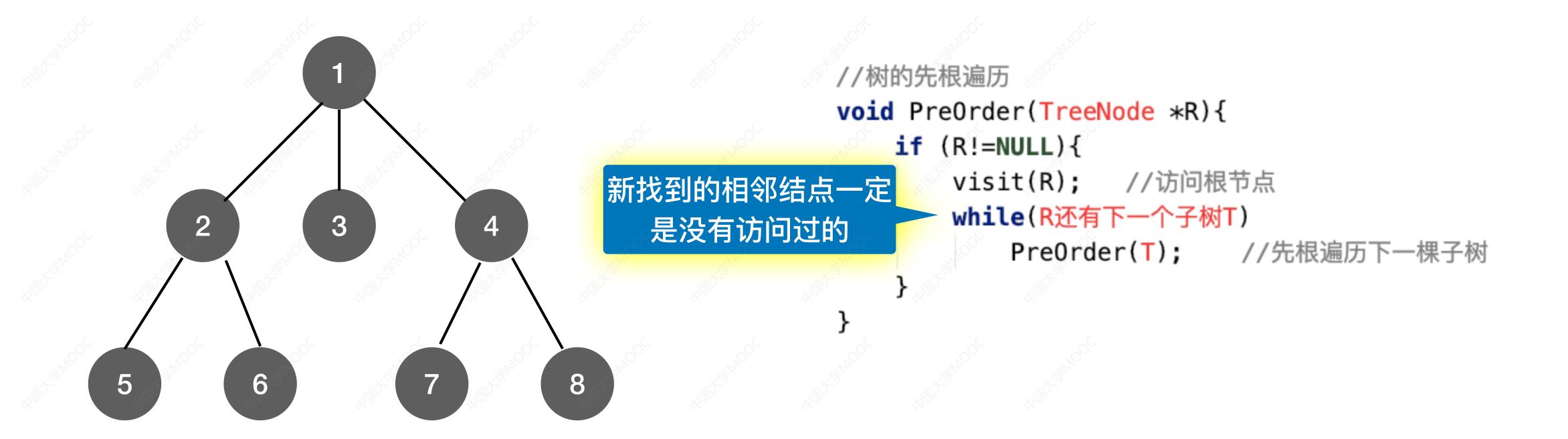






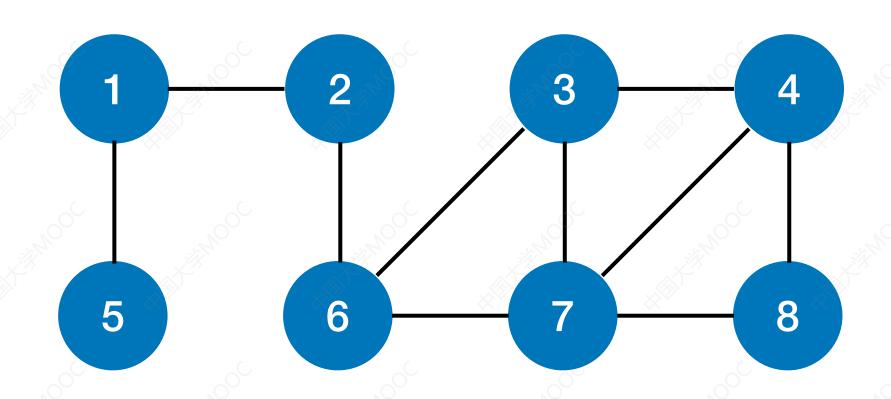






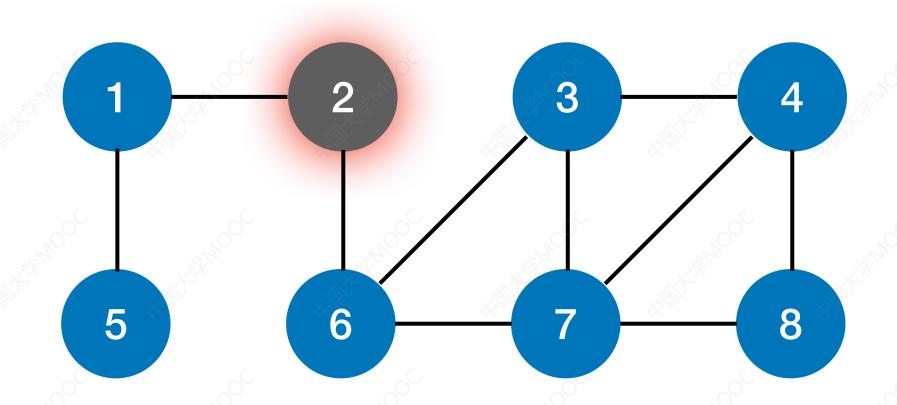
先根遍历序列: 1, 2, 5, 6, 3, 4, 7, 8

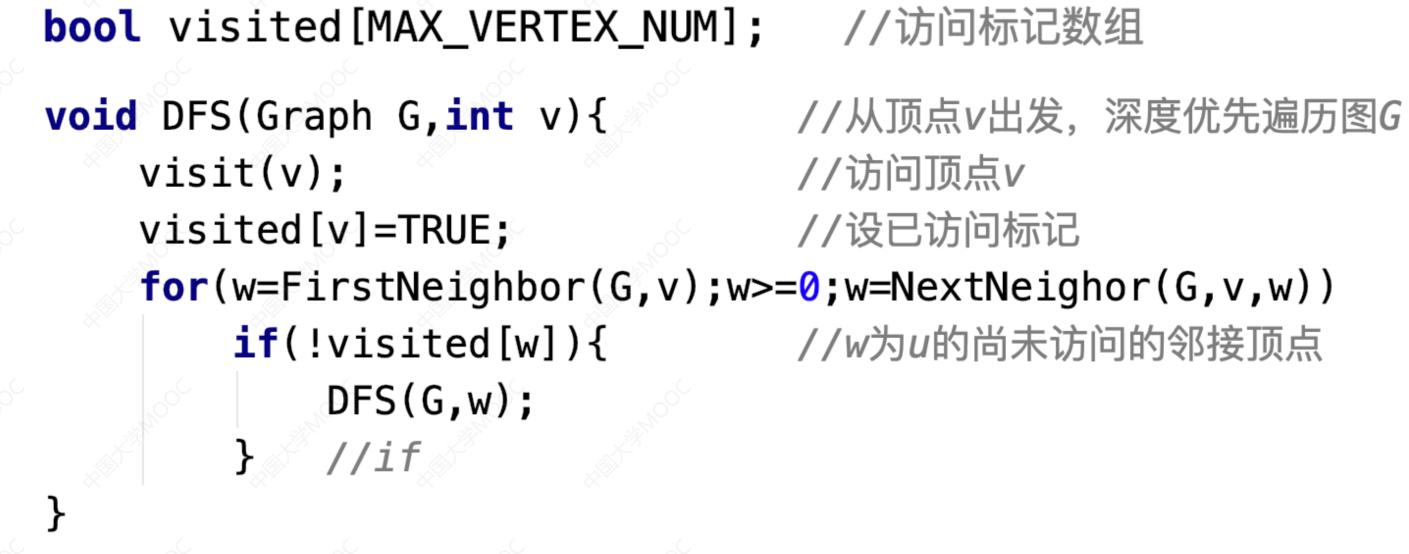
#### 初始都为false



1 2 3 4 5 6 7 8 visited false false false false false false false false

#### 初始都为false

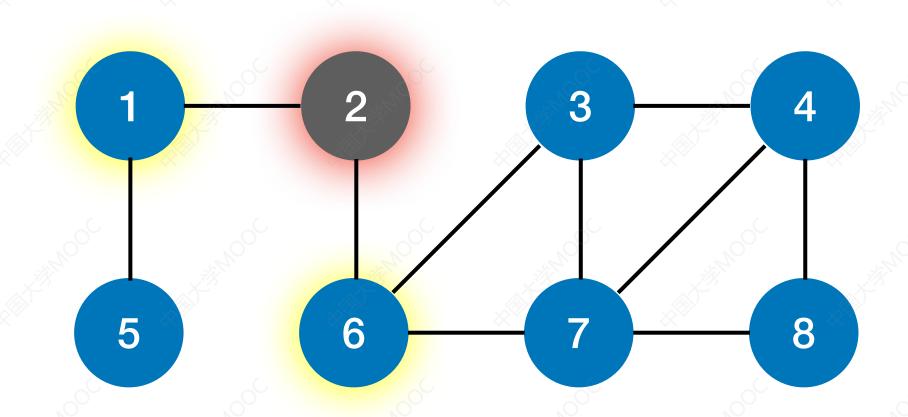




1 2 3 4 5 6 7 8 visited false true false false false false false

函数调用栈

#### 初始都为false



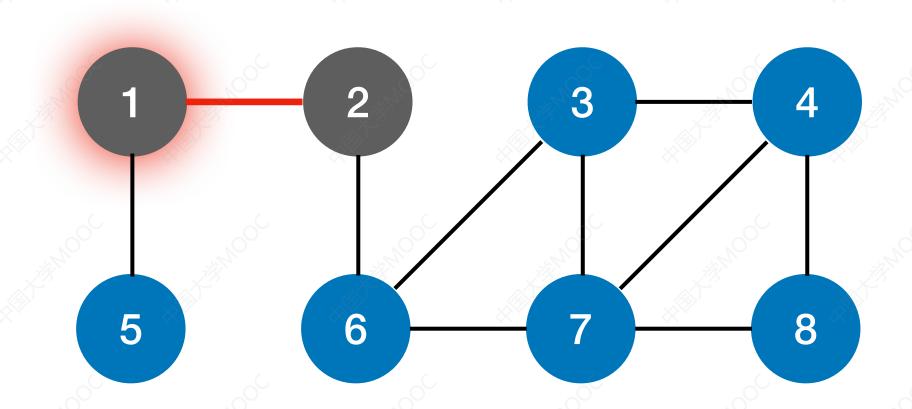
1 2 3 4 5 6 7 8 visited false true false false false false false false

函数调用栈

bool visited[MAX\_VERTEX\_NUM];

#### 初始都为false

//访问标记数组



visited true false false false false false

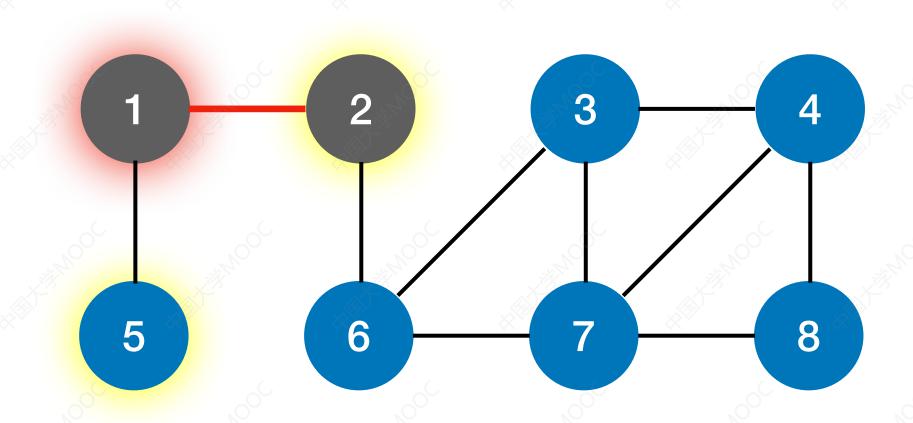
1 2, w=1

函数调用栈

bool visited[MAX\_VERTEX\_NUM];

#### 初始都为false

//访问标记数组

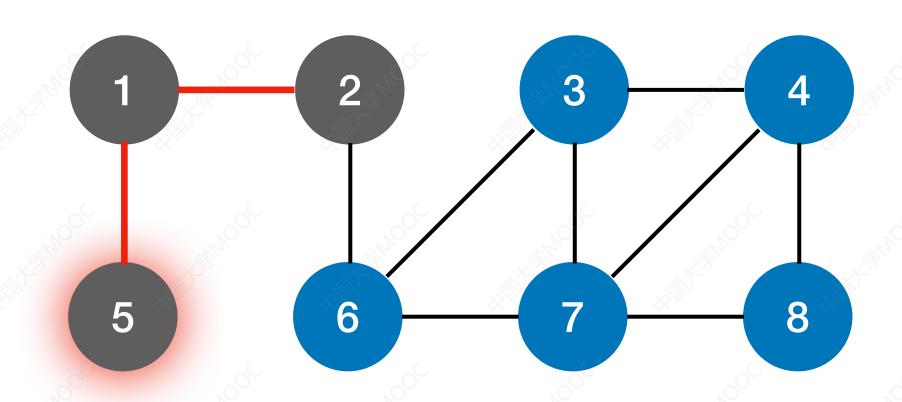


visited true false false false false false

函数调用栈

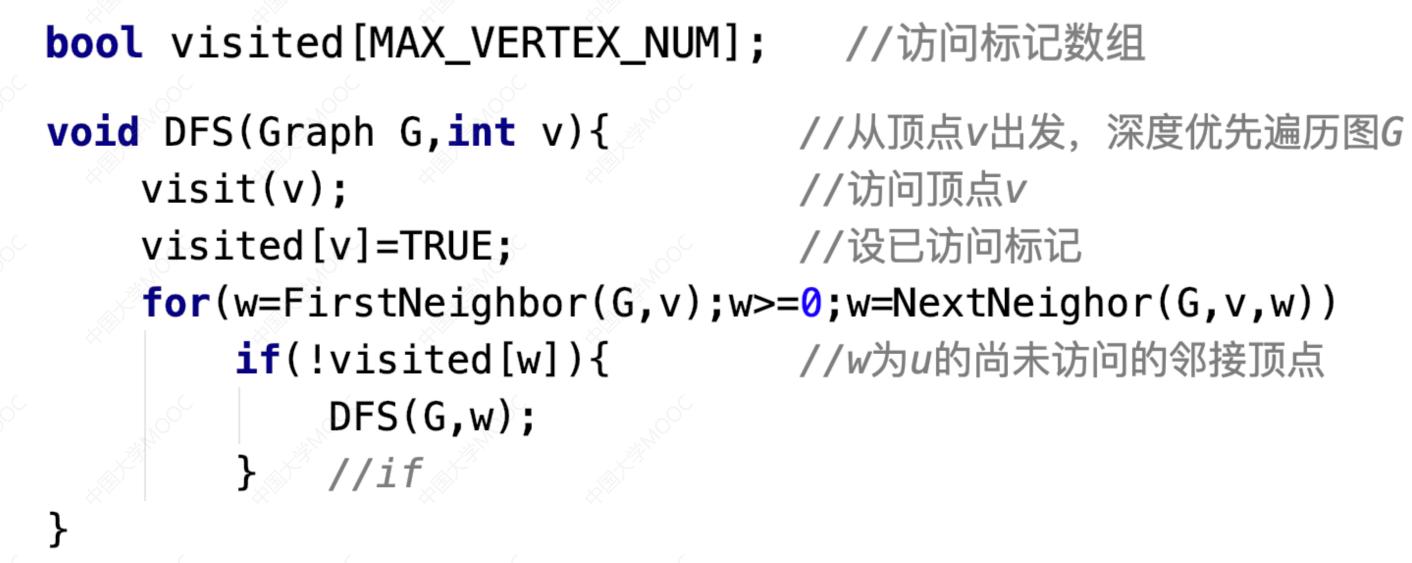
2, w=1

#### 初始都为false



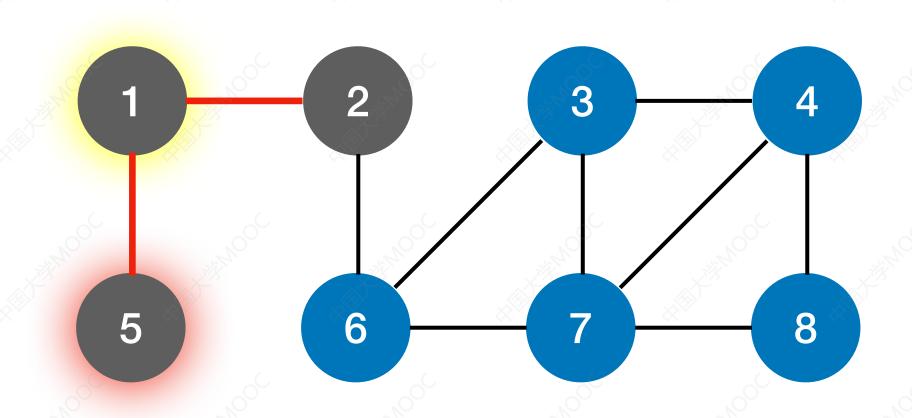
5 1, w=5 2, w=1

函数调用栈



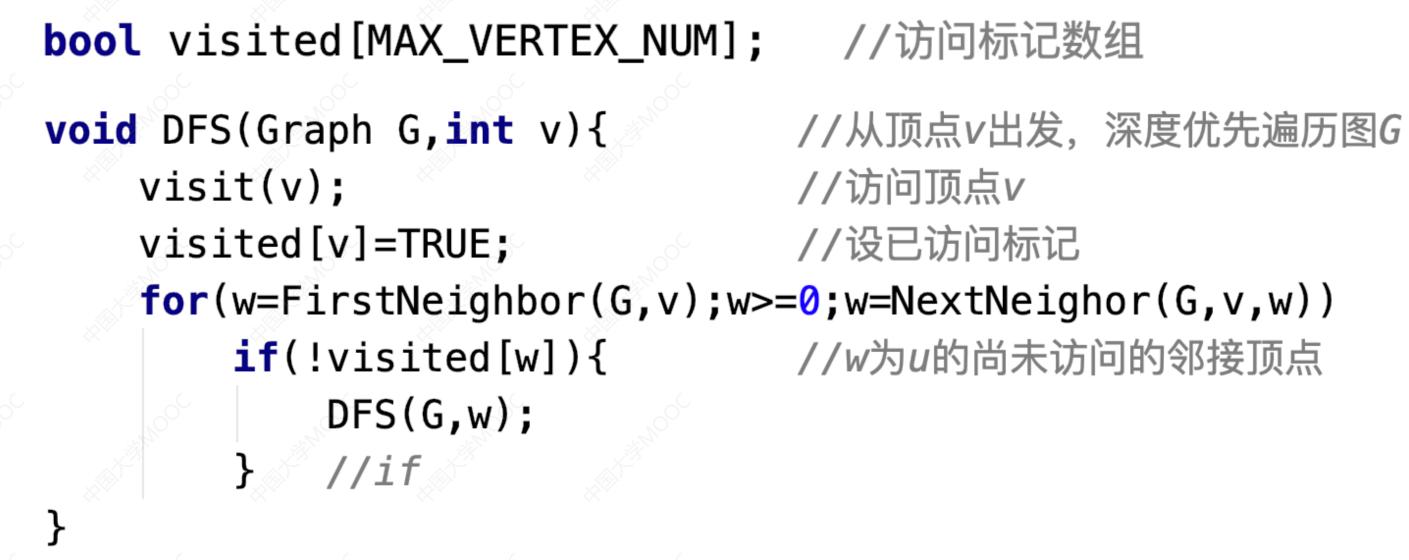
1 2 3 4 5 6 7 8 visited true true false false true false false

#### 初始都为false



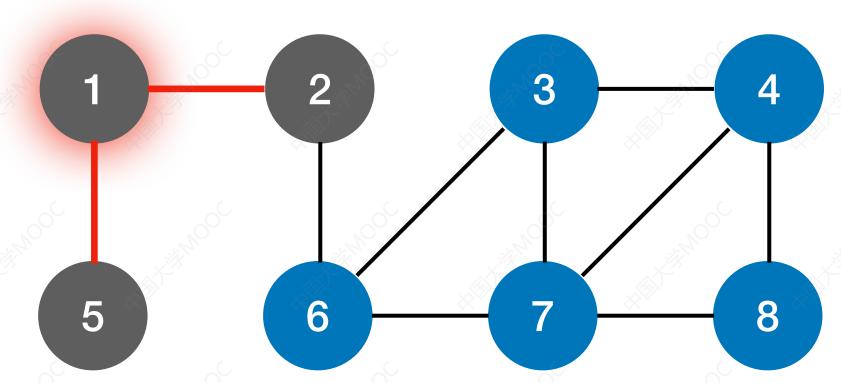
5 1, w=5 2, w=1

函数调用栈



1 2 3 4 5 6 7 8 visited true true false false true false false

#### 初始都为false



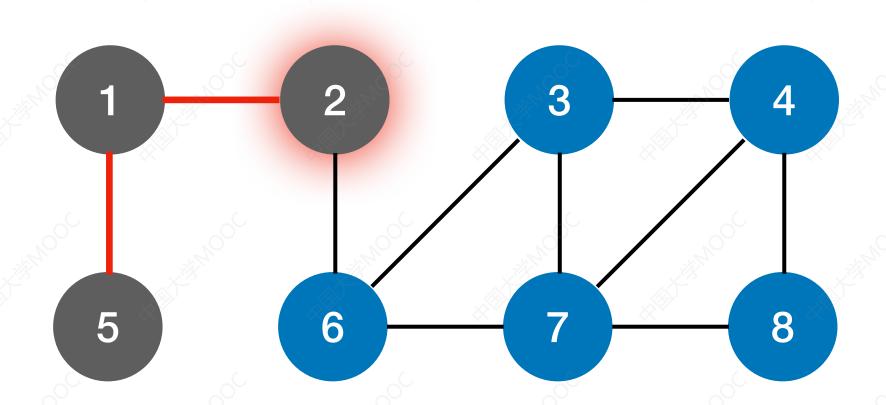
1 2 3 4 5 6 7 8 visited true true false false true false false

函数调用栈

2, w=1

w=5

#### 初始都为false

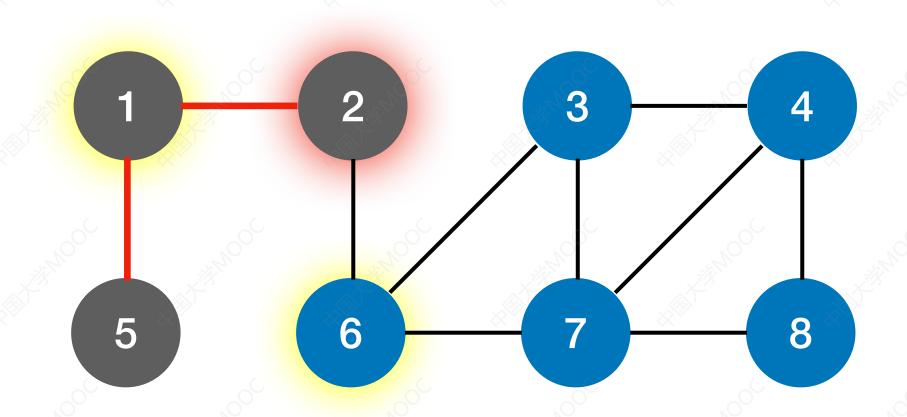


1 2 3 4 5 6 7 8 visited true true false false true false false

2, w=1

函数调用栈

#### 初始都为false

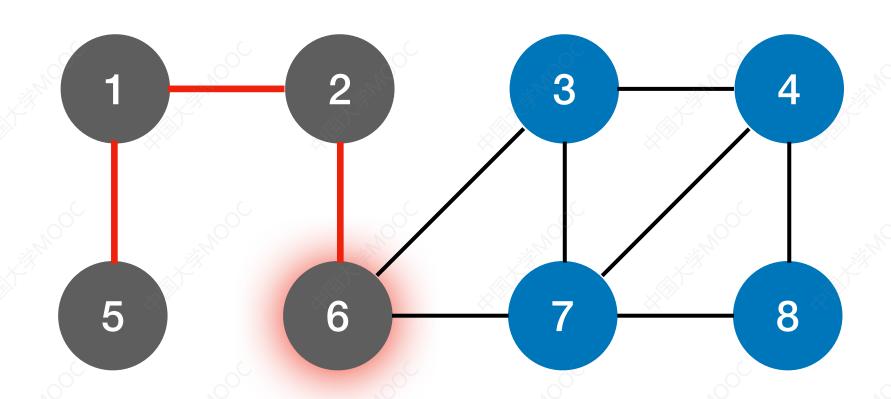


1 2 3 4 5 6 7 8 visited true true false false true false false

2, w=1

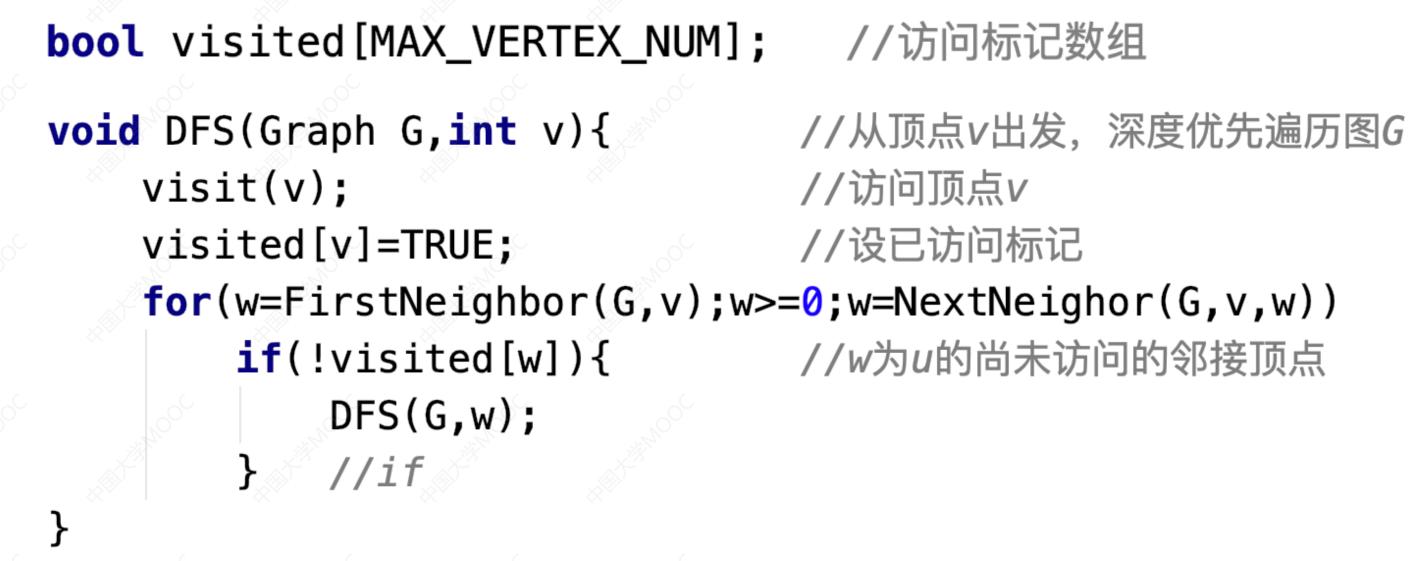
函数调用栈

#### 初始都为false



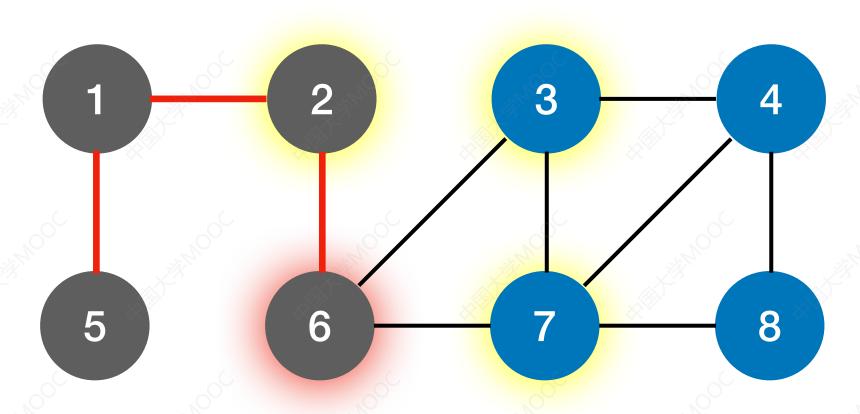
6 2, w=6

函数调用栈



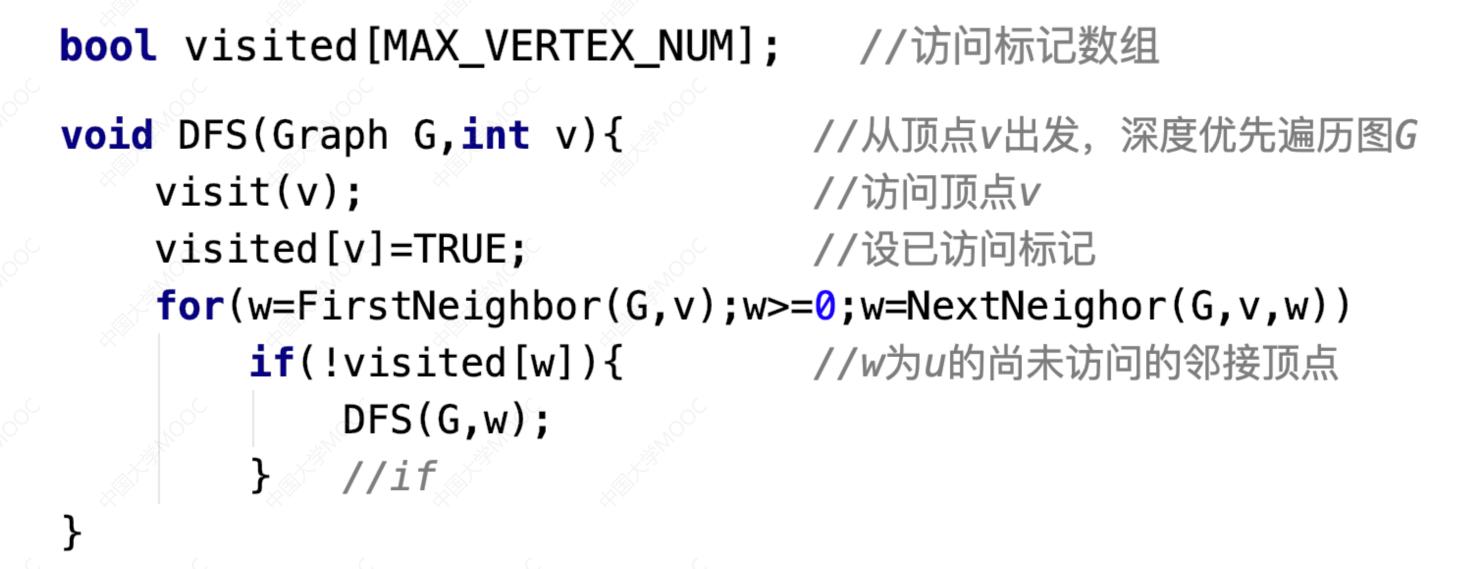
1 2 3 4 5 6 7 8 visited true true false false true true false false

#### 初始都为false



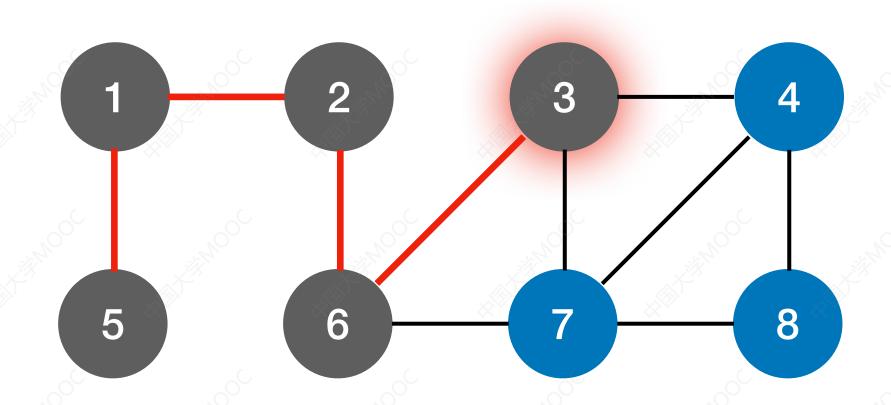
6 2, w=6

函数调用栈



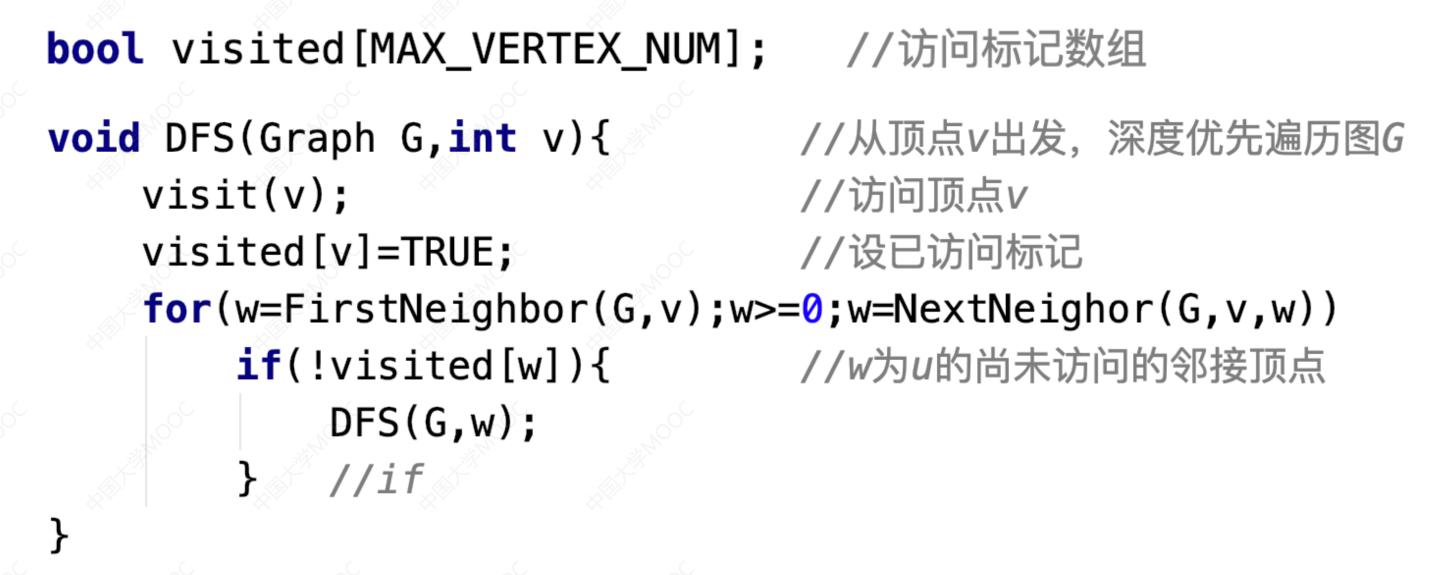
1 2 3 4 5 6 7 8 visited true true false false true true false false

#### 初始都为false



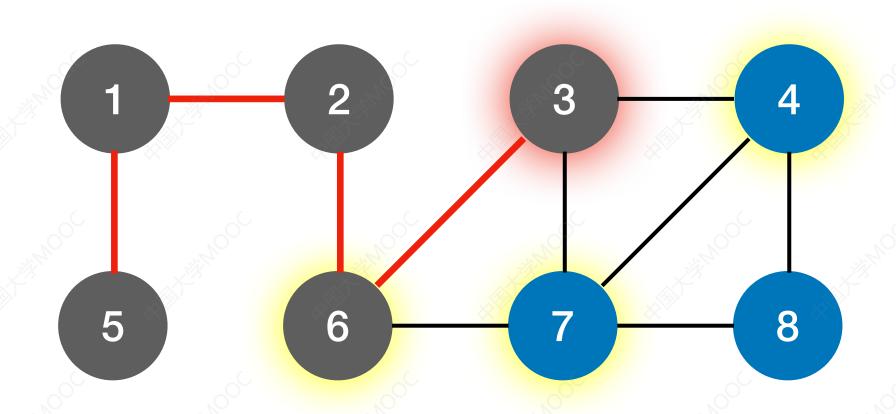
3 6, w=3 2, w=6

函数调用栈



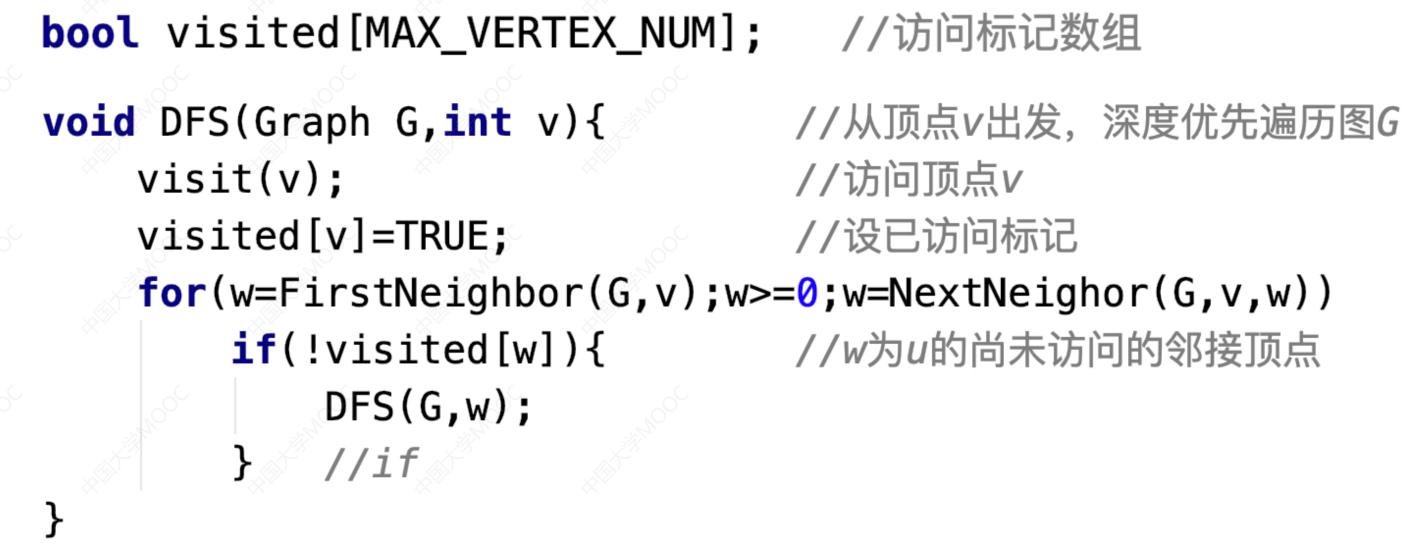
1 2 3 4 5 6 7 8 visited true true true false true true false

#### 初始都为false



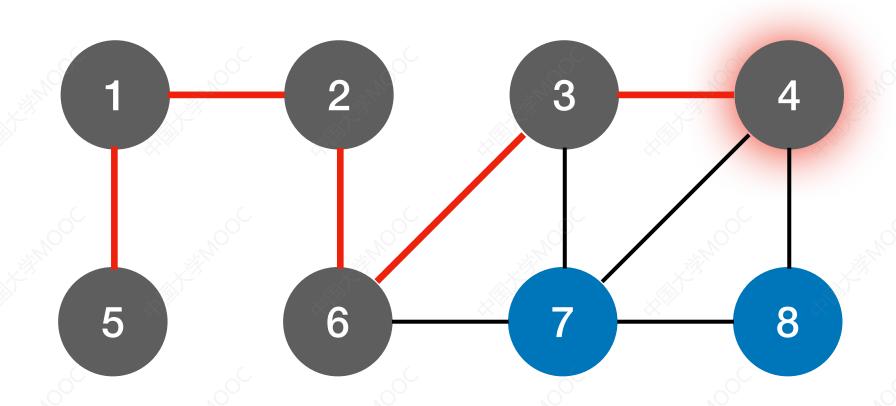
3 6, w=3 2, w=6

函数调用栈

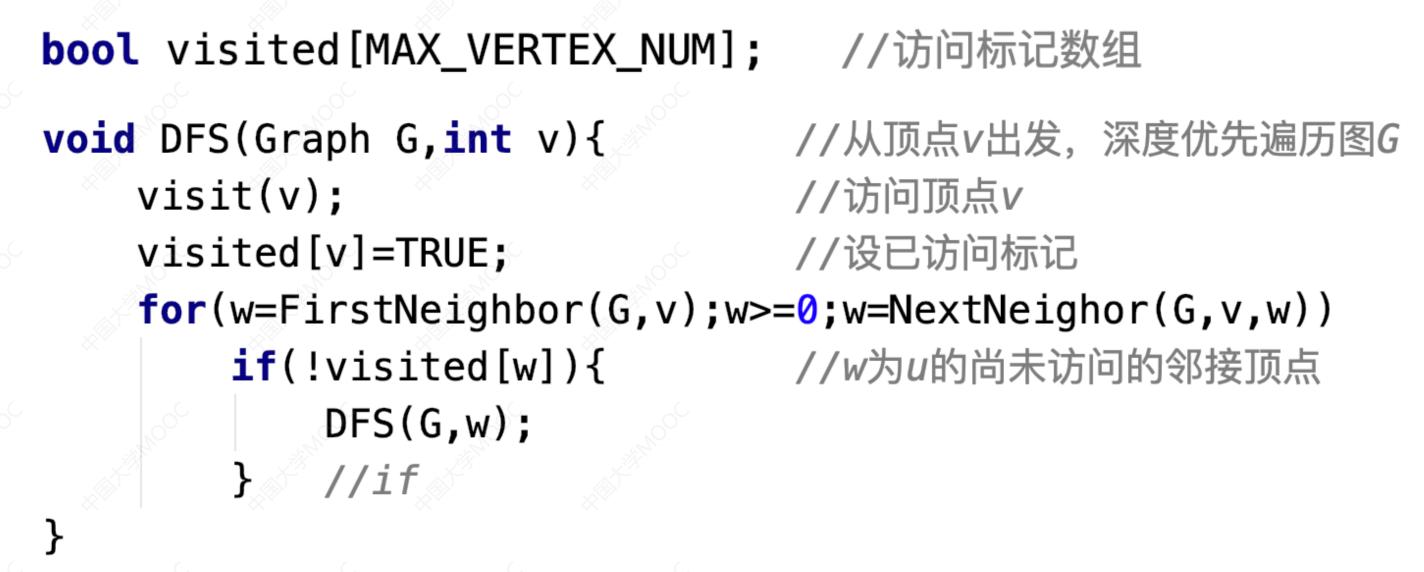


<i>**</i> **	1	2	3	4	5	6	7	8
visited	true	true	true	false	true	true	false	false

#### 初始都为false

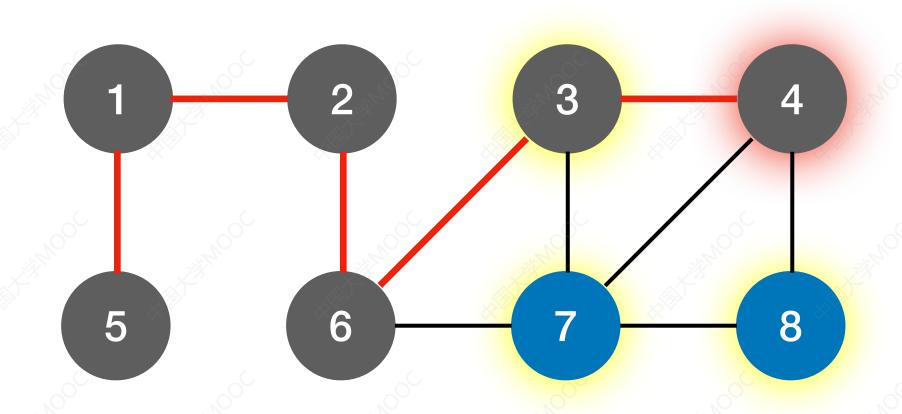


4 3, w=4 6, w=3 2, w=6

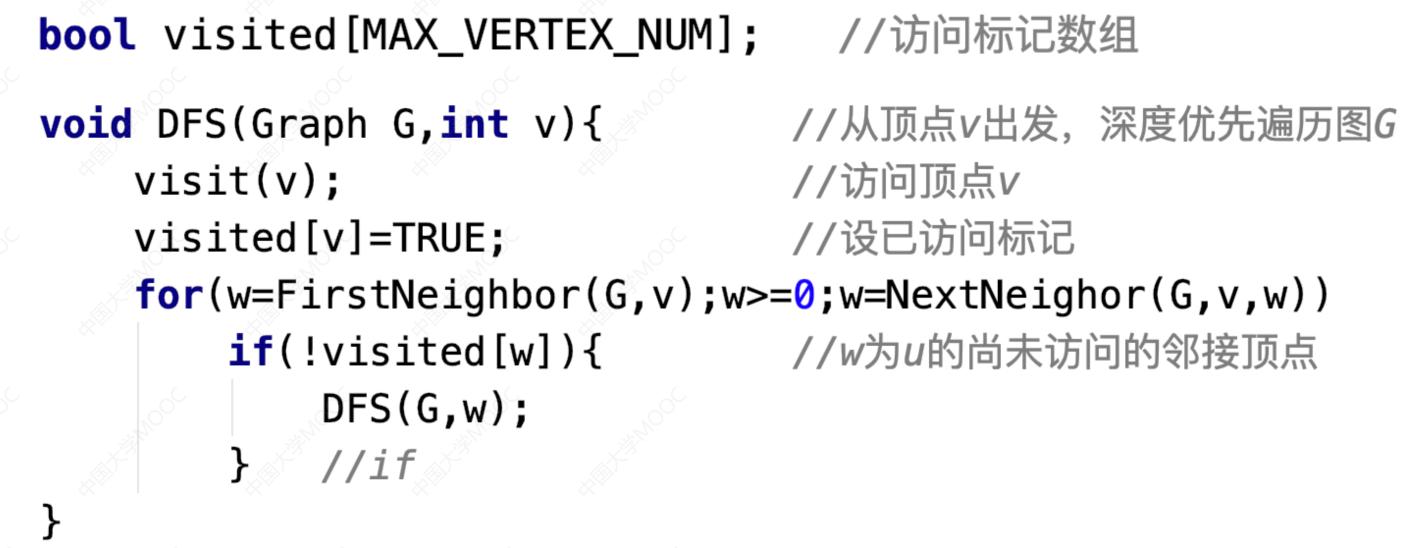


	1	2	3	4	5	6	7	8
visited	true	true	true	true	true	true	false	false

#### 初始都为false

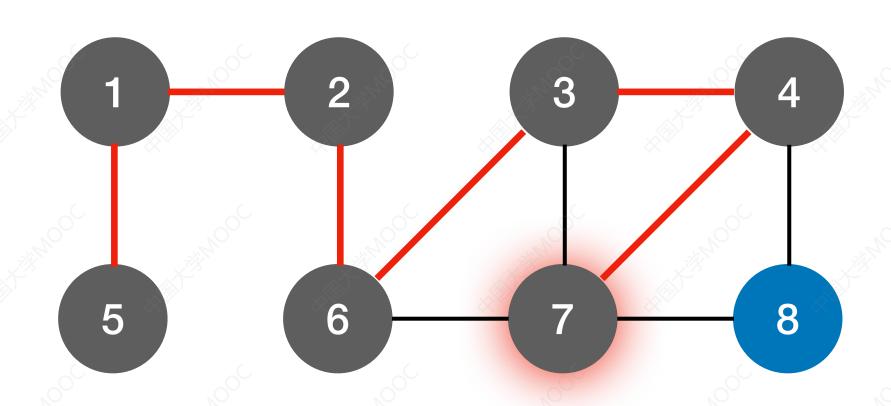


4
3, w=4
6, w=3
2, w=6

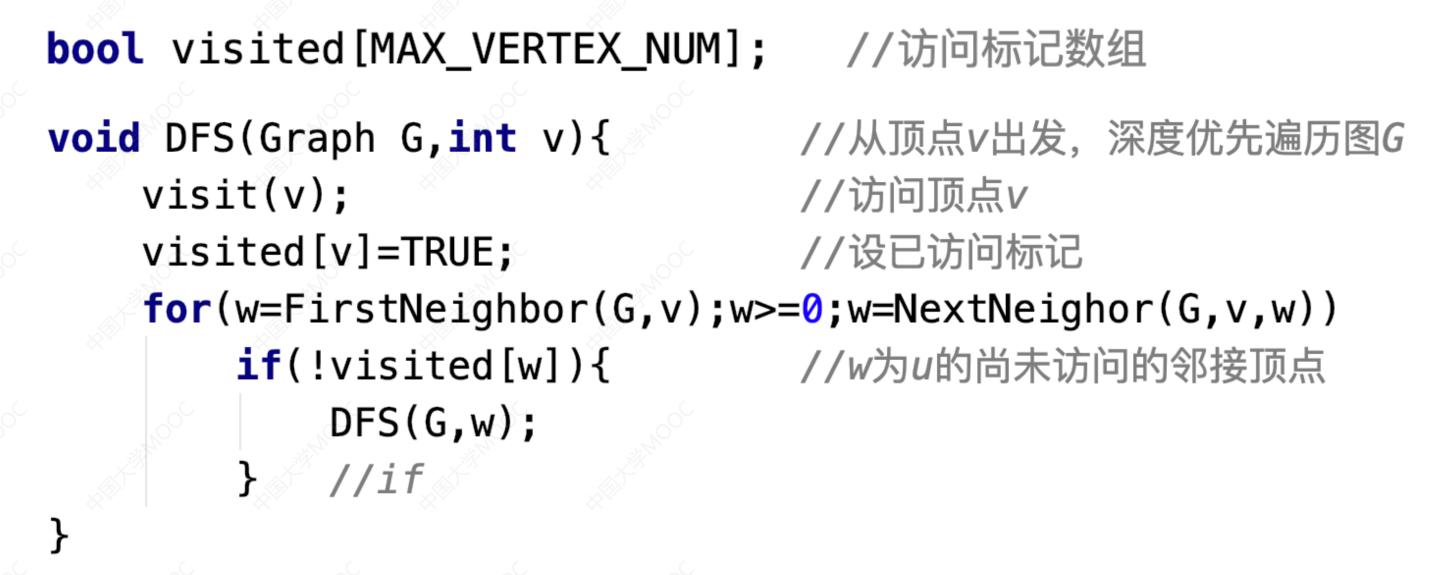


<i>**</i> **********************************	1	2	3	4	5	6	7	8
visited	true	true	true	true	true	true	false	false

#### 初始都为false

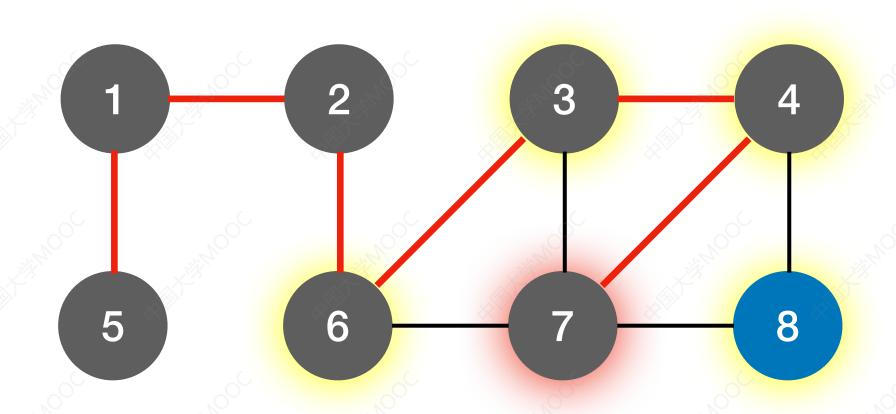


7
4, w=7
3, w=4
6, w=3
2, w=6

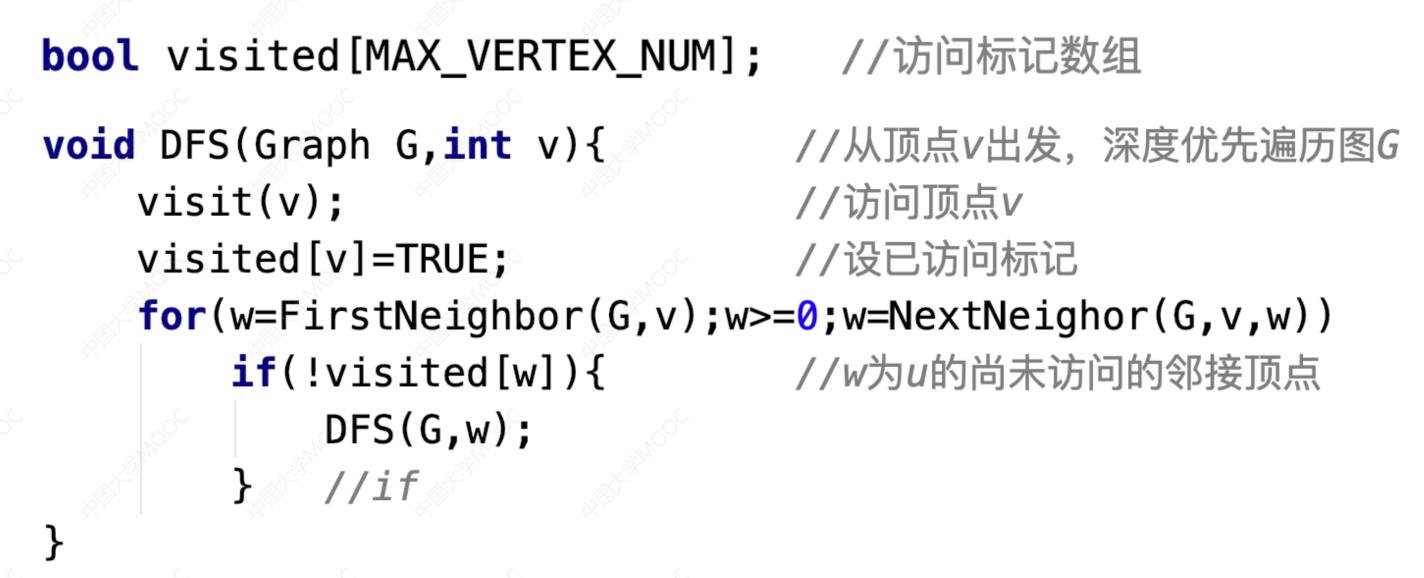




#### 初始都为false

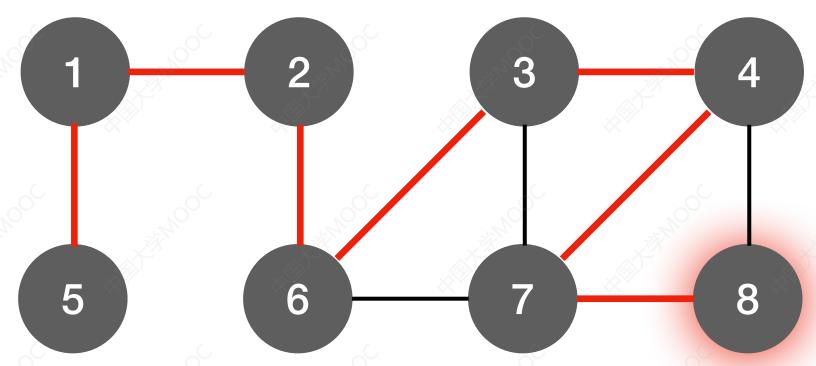


7
4, w=7
3, w=4
6, w=3
2, w=6



	1 2		3	4	5 6		7	8
visited	true	false						

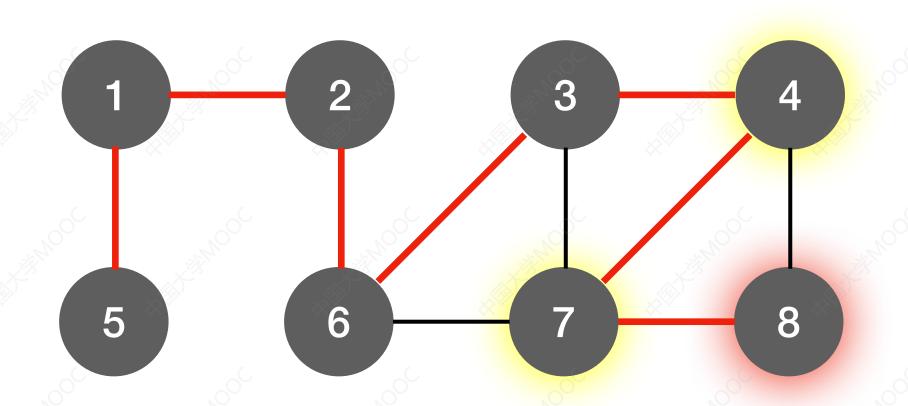
#### 初始都为false



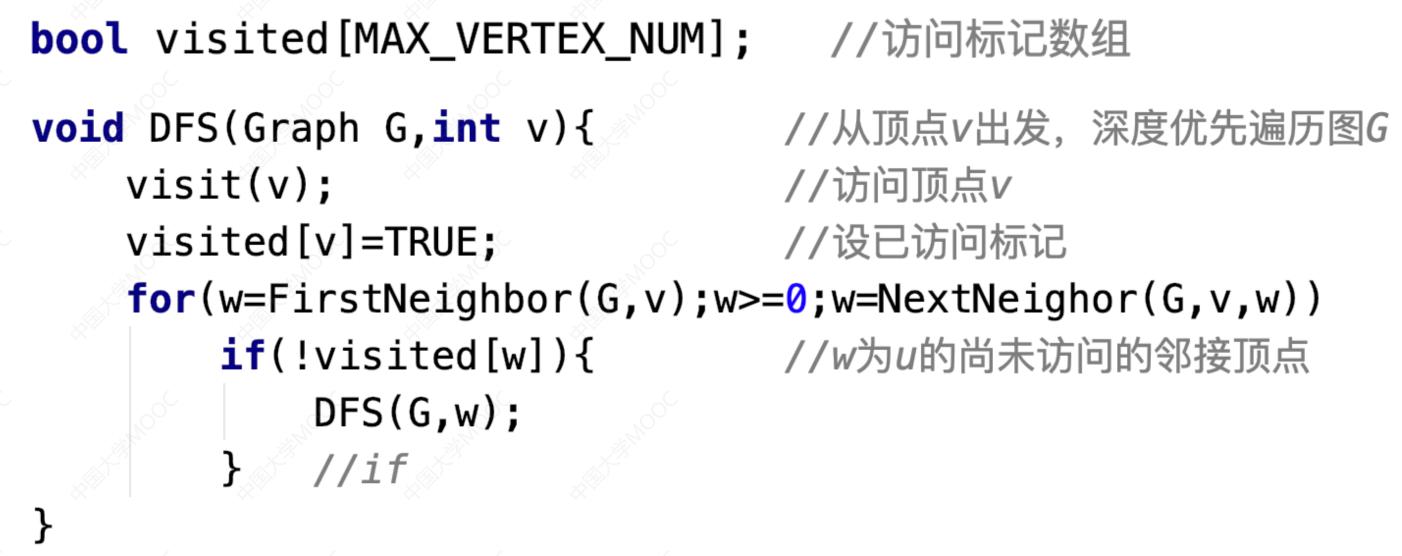
8
7, w=8
4, w=7
3, w=4
6, w=3
2, w=6

	1	2	3	4	5	6	7	8
visited	true							

#### 初始都为false

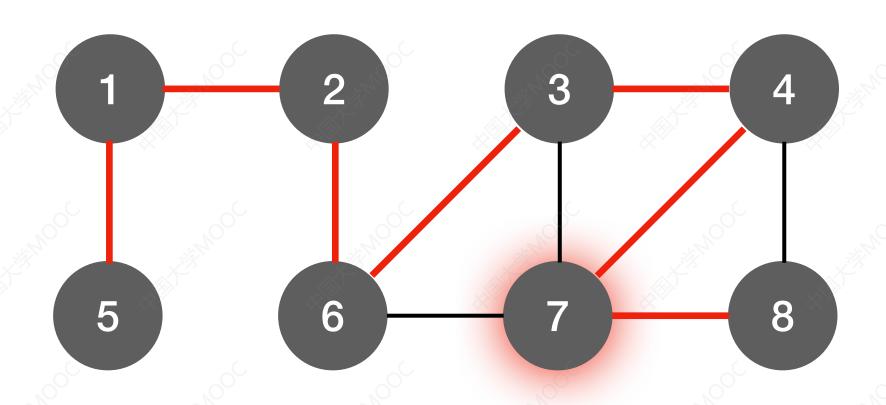


8
7, w=8
4, w=7
3, w=4
6, w=3
2, w=6





#### 初始都为false

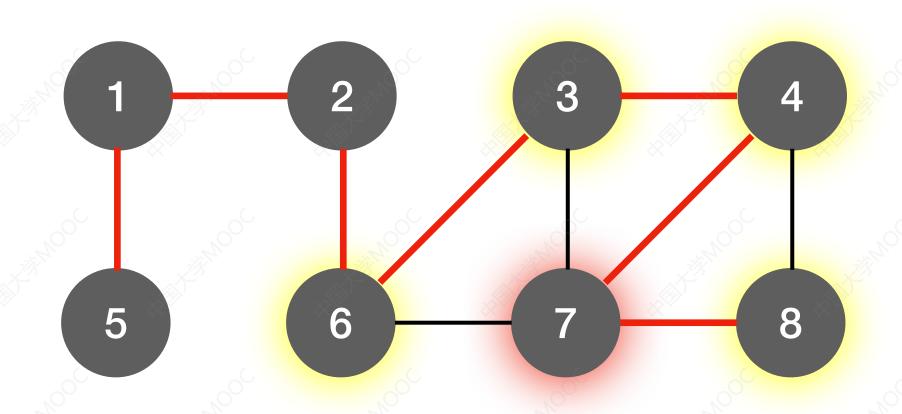


7, w=84, w=7 3, w=46, w=3 2, w=6

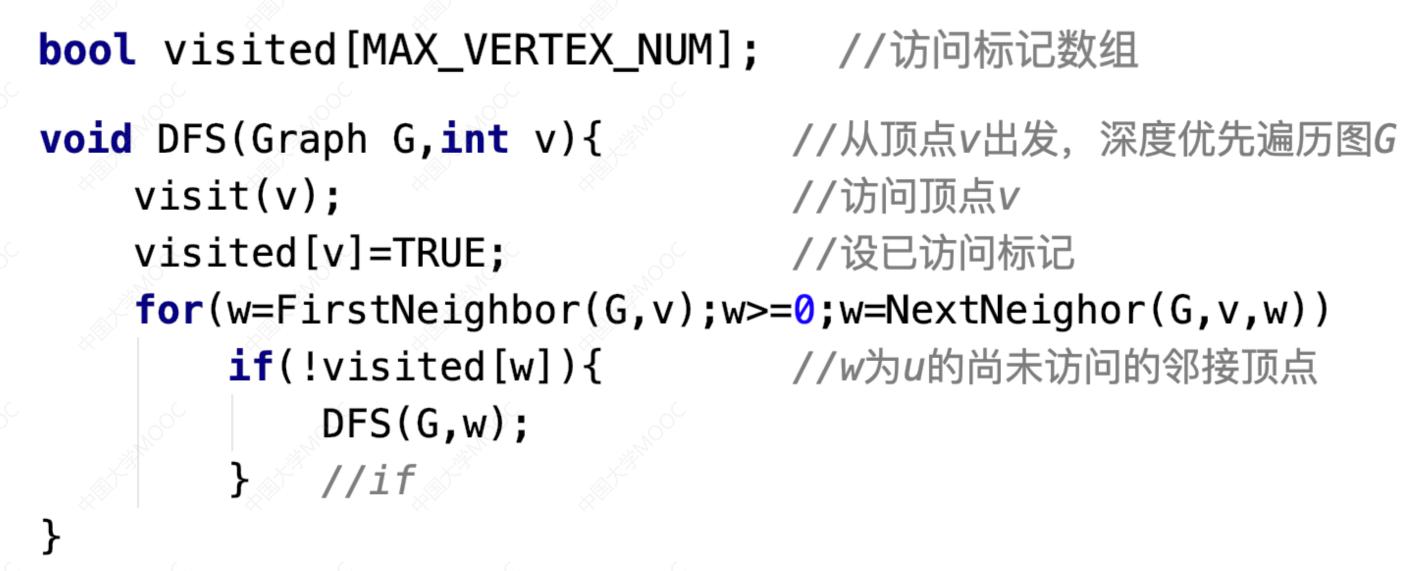
```
bool visited[MAX_VERTEX_NUM];
                              //访问标记数组
void DFS(Graph G,int v){
                            //从顶点v出发,深度优先遍历图G
   visit(v);
                            //访问顶点v
   visited[v]=TRUE;
                            //设已访问标记
   for(w=FirstNeighbor(G,v);w>=0;w=NextNeighor(G,v,w))
       if(!visited[w]){
                          //w为u的尚未访问的邻接顶点
          DFS(G,w);
          //if
```

visited true true true true true true true

#### 初始都为false

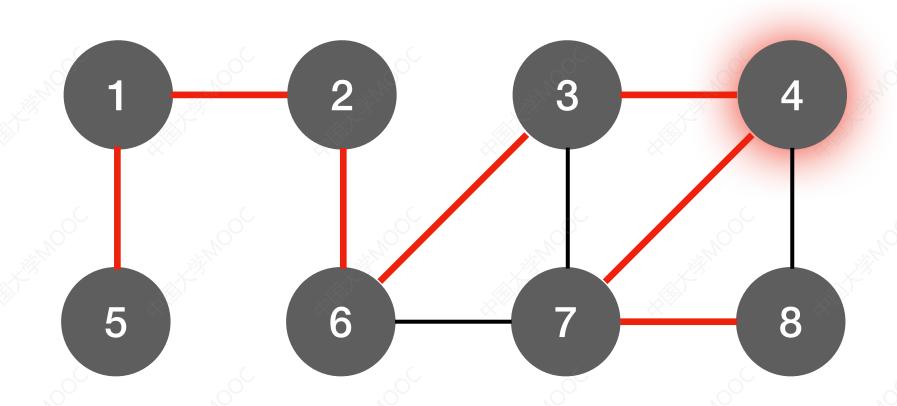


7, w=8
4, w=7
3, w=4
6, w=3
2, w=6

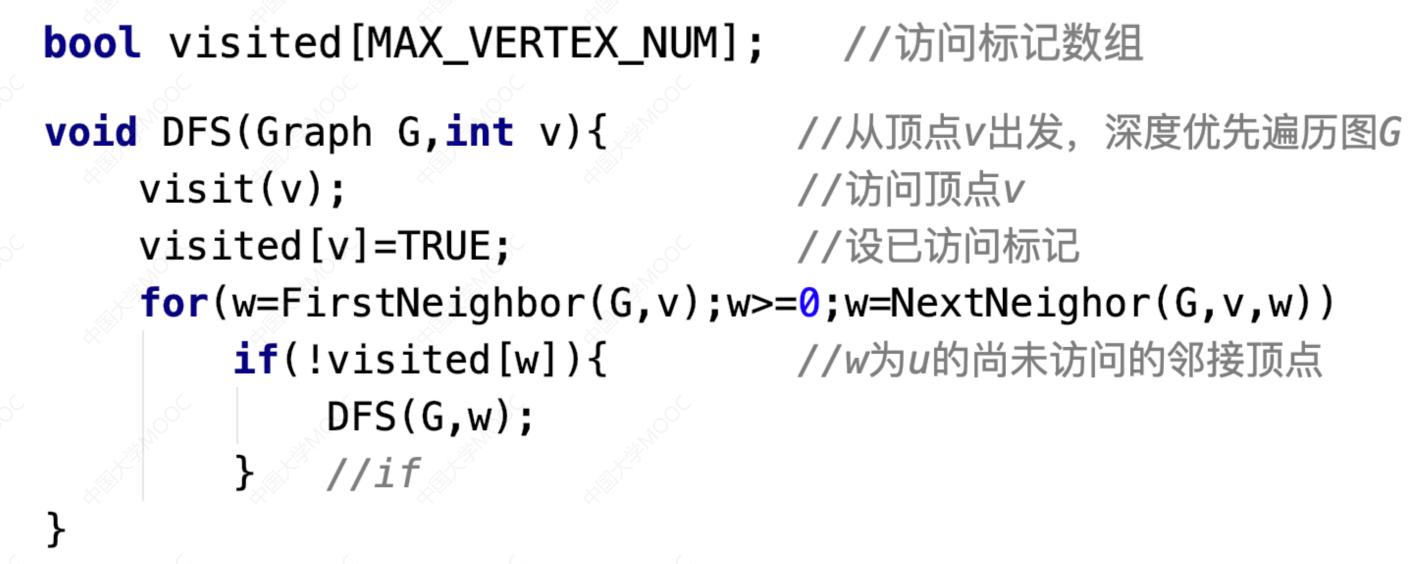


	1	2	3	4	5	6	7	8
visited	true							

#### 初始都为false

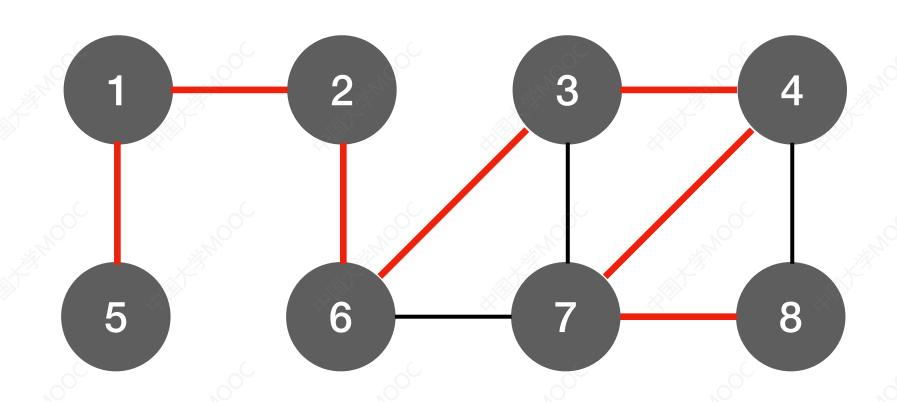


4, w=7
3, w=4
6, w=3
2, w=6



	1	2	3 4		5 6		7 8	
visited	true							

#### 初始都为false

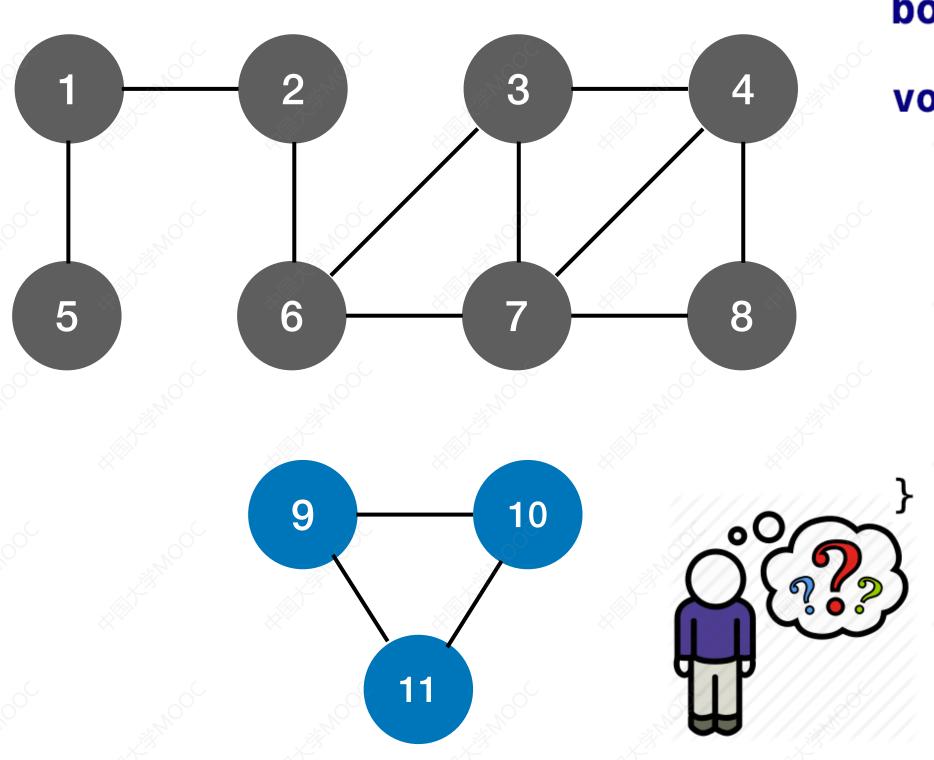


	1 2		3	4	5	6	7	8
visited	true							

从2出发的深度遍历序列: 2, 1, 5, 6, 3, 4, 7, 8

### 算法存在的问题

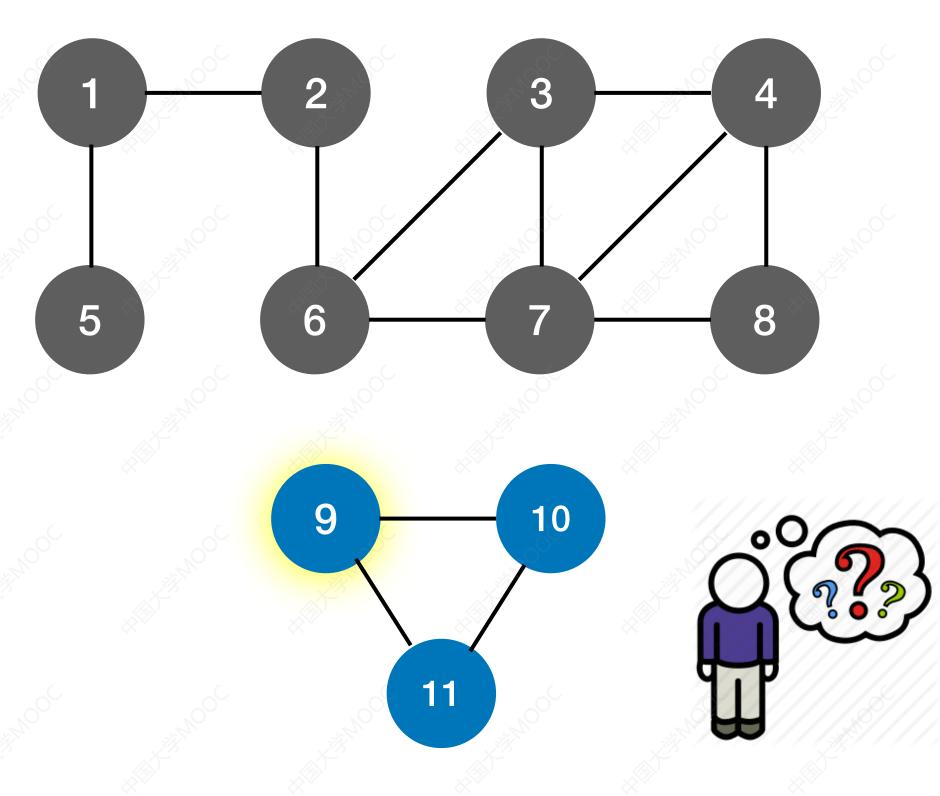
#### 初始都为false



	1	2	3	4	5	6	7	8	9	10	11
visited	true	false	false	false							

如果是非连通图,则无法遍历完所有结点

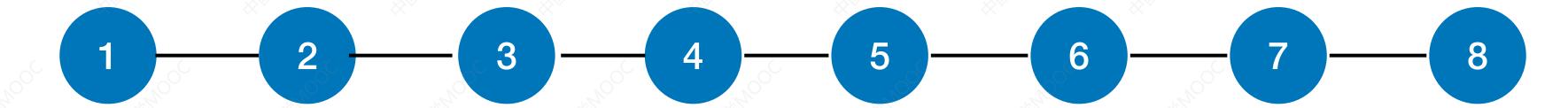
### DFS算法 (Final版)



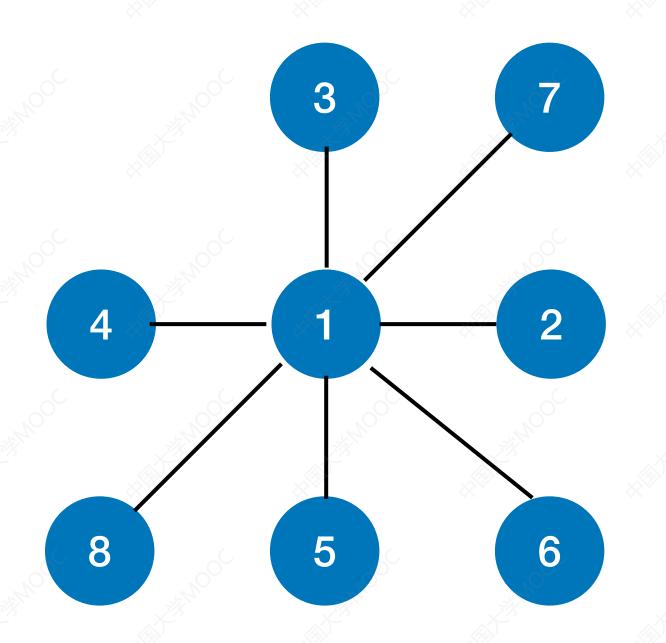
如果是非连通图,则无法遍历完所有结点

```
bool visited[MAX_VERTEX_NUM];
                                //访问标记数组
void DFSTraverse(Graph G){
                              //对图G进行深度优先遍历
    for(v=0; v<G.vexnum; ++v)</pre>
       visited[v]=FALSE;
                              //初始化已访问标记数据
    for(v=0; v<G.vexnum; ++v)</pre>
                              //本代码中是从v=0开始遍历
       if(!visited[v])
           DFS(G,v);
void DFS(Graph G,int v){
                             //从顶点v出发,深度优先遍历图G
   visit(v);
                             //访问顶点v
   visited[v]=TRUE;
                             //设已访问标记
   for(w=FirstNeighbor(G,v);w>=0;w=NextNeighor(G,v,w))
       if(!visited[w]){
                             //w为u的尚未访问的邻接顶点
           DFS(G,w);
           //if
```

# 复杂度分析

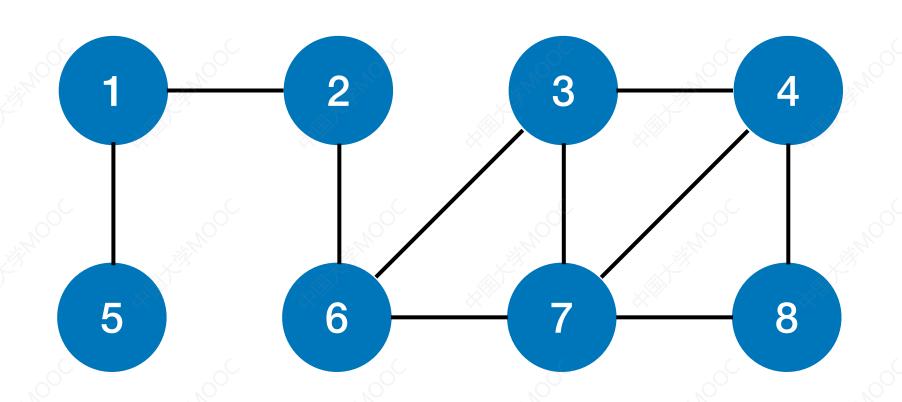


空间复杂度:来自函数调用栈,最坏情况,递归深度为O(|V|)

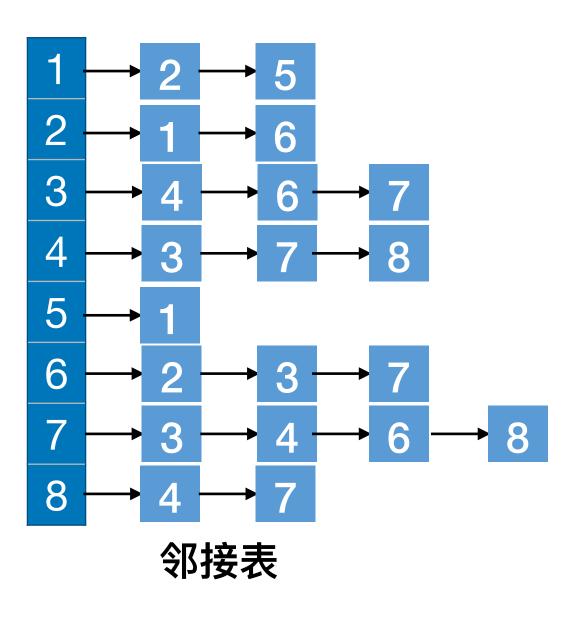


空间复杂度:最好情况,O(1)

### 复杂度分析







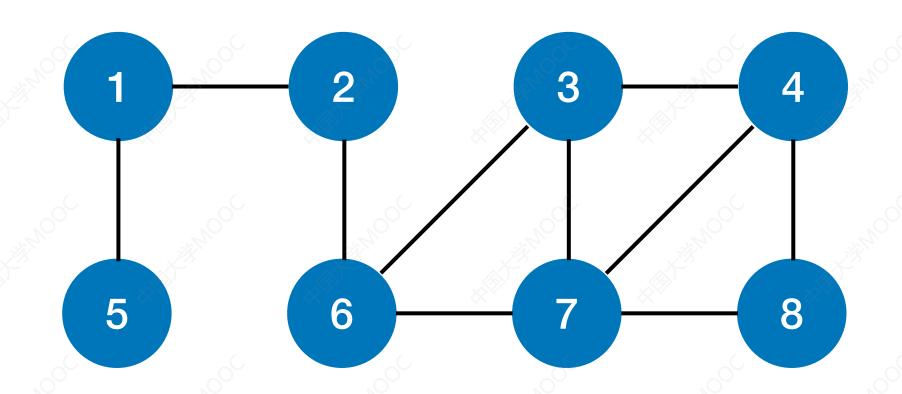
时间复杂度=访问各结点所需时间+探索各条边所需时间

#### 邻接矩阵存储的图:

访问 |V| 个顶点需要O(|V|)的时间 查找每个顶点的邻接点都需要O(|V|)的时间,而总共有|V|个顶点 时间复杂度= O(|V|<sup>2</sup>)

#### 邻接表存储的图:

访问 |V| 个顶点需要O(|V|)的时间 查找各个顶点的邻接点共需要O(|E|)的时间, 时间复杂度= O(|V|+|E|)





 1
 2
 5

 2
 1
 6

 3
 4
 6
 7

 4
 3
 7
 8

 5
 1
 8

 6
 2
 3
 7

 7
 3
 4
 6
 8

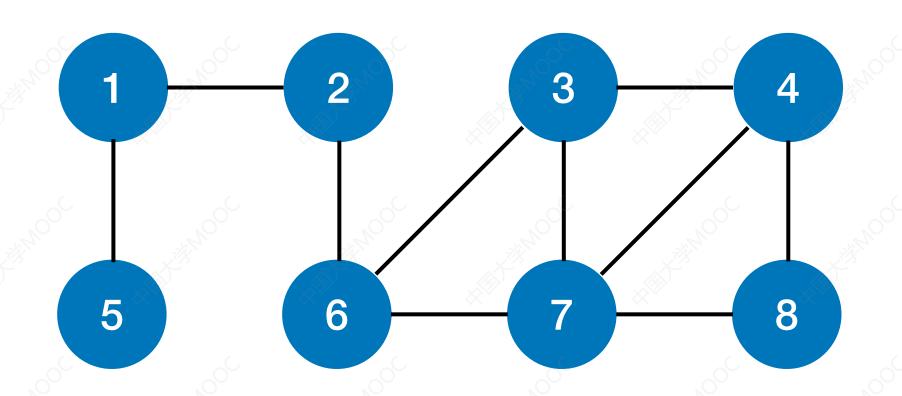
 8
 4
 7

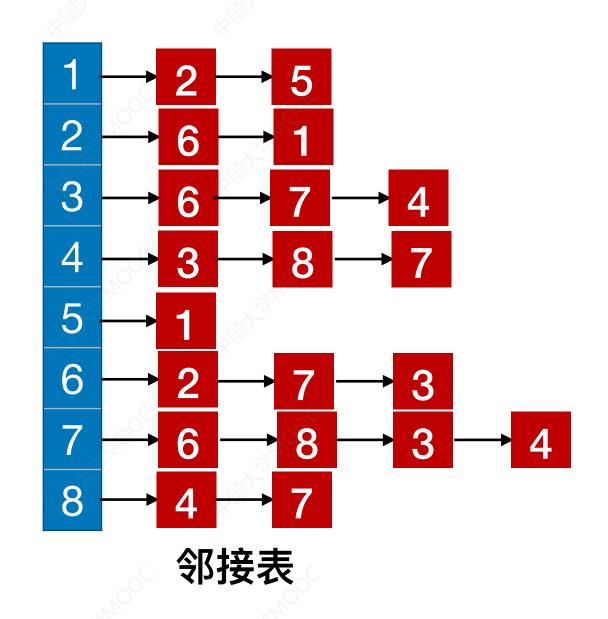
 \$\$i\$\$\frac{1}{2}\$\$
 \$\$i\$\$\frac{1}{2}\$\$
 \$\$i\$\$\$\frac{1}{2}\$\$

从2出发的深度优先遍历序列: 2, 1, 5, 6, 3, 4, 7, 8

从3出发的深度优先遍历序列: 3, 4, 7, 6, 2, 1, 5, 8

从1出发的深度优先遍历序列: 1, 2, 6, 3, 4, 7, 8, 5

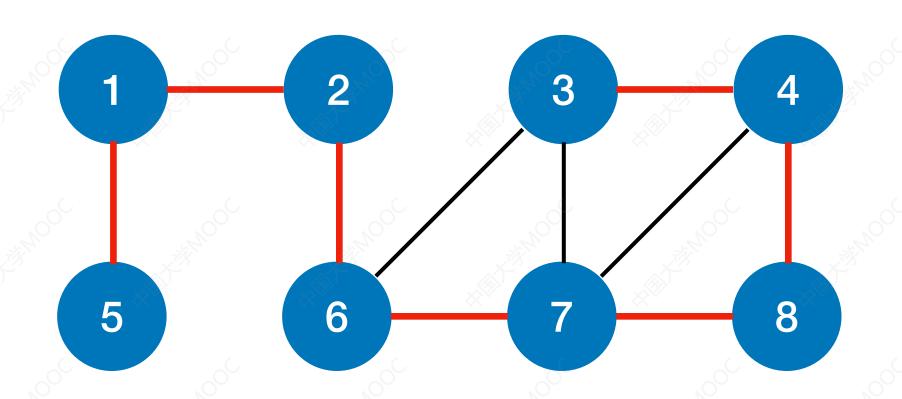


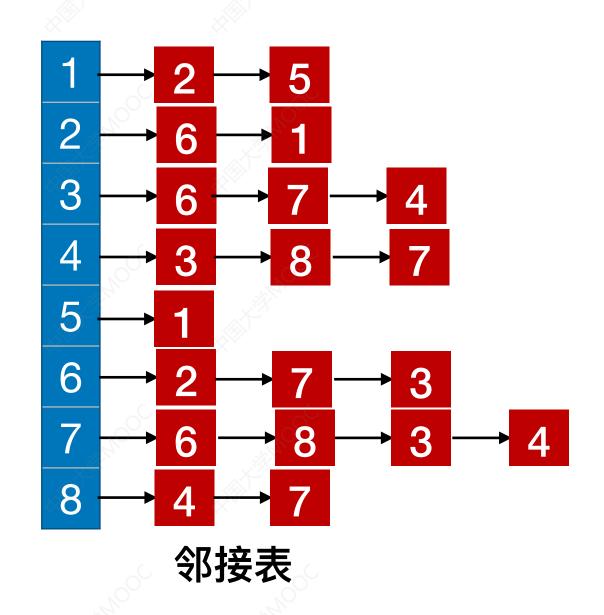


从2出发的深度优先遍历序列: 2, 6, 7, 8, 4, 3, 1, 5

从3出发的深度优先遍历序列?

从1出发的深度优先遍历序列?

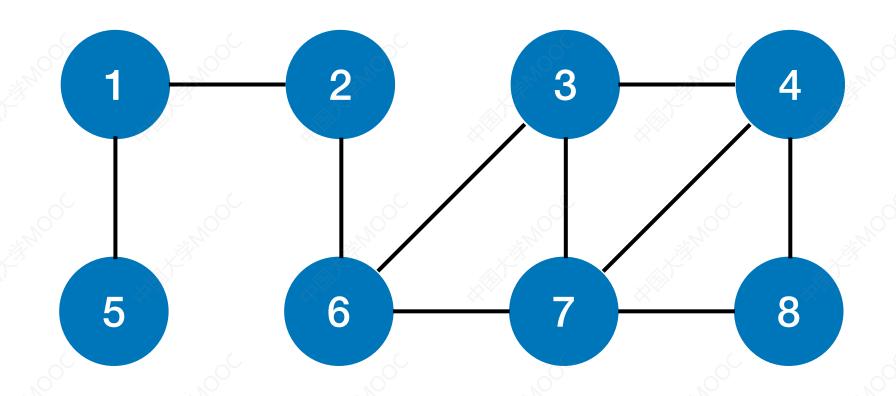




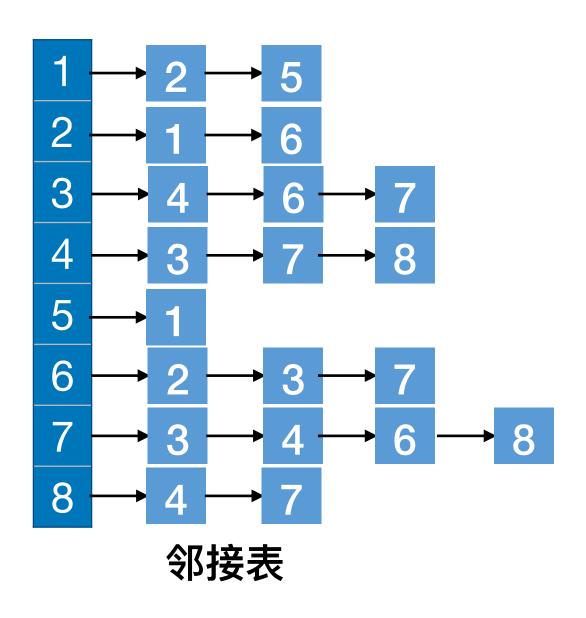
从2出发的深度优先遍历序列: 2, 6, 7, 8, 4, 3, 1, 5

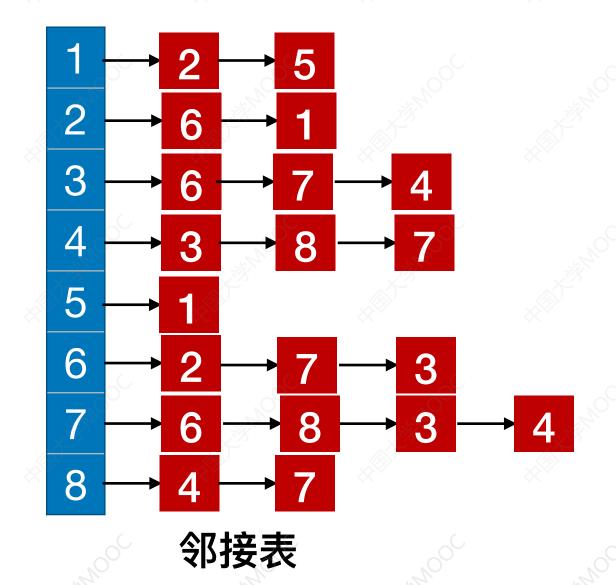
从3出发的深度优先遍历序列?

从1出发的深度优先遍历序列?



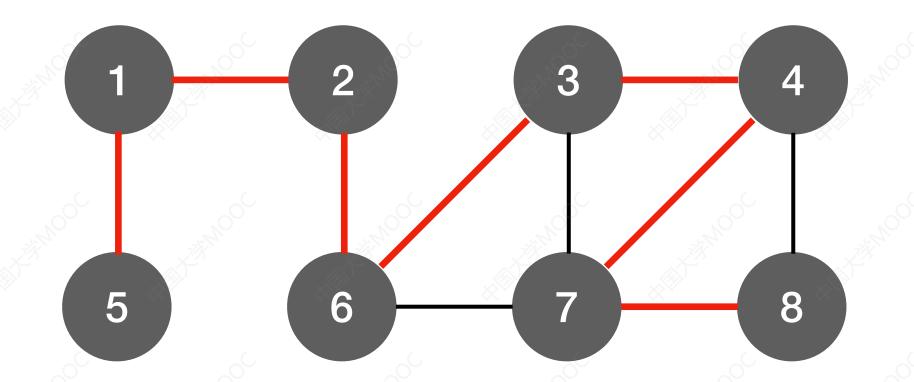


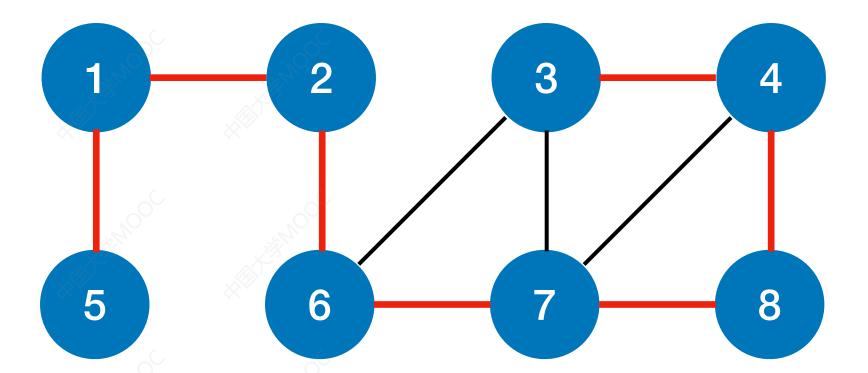




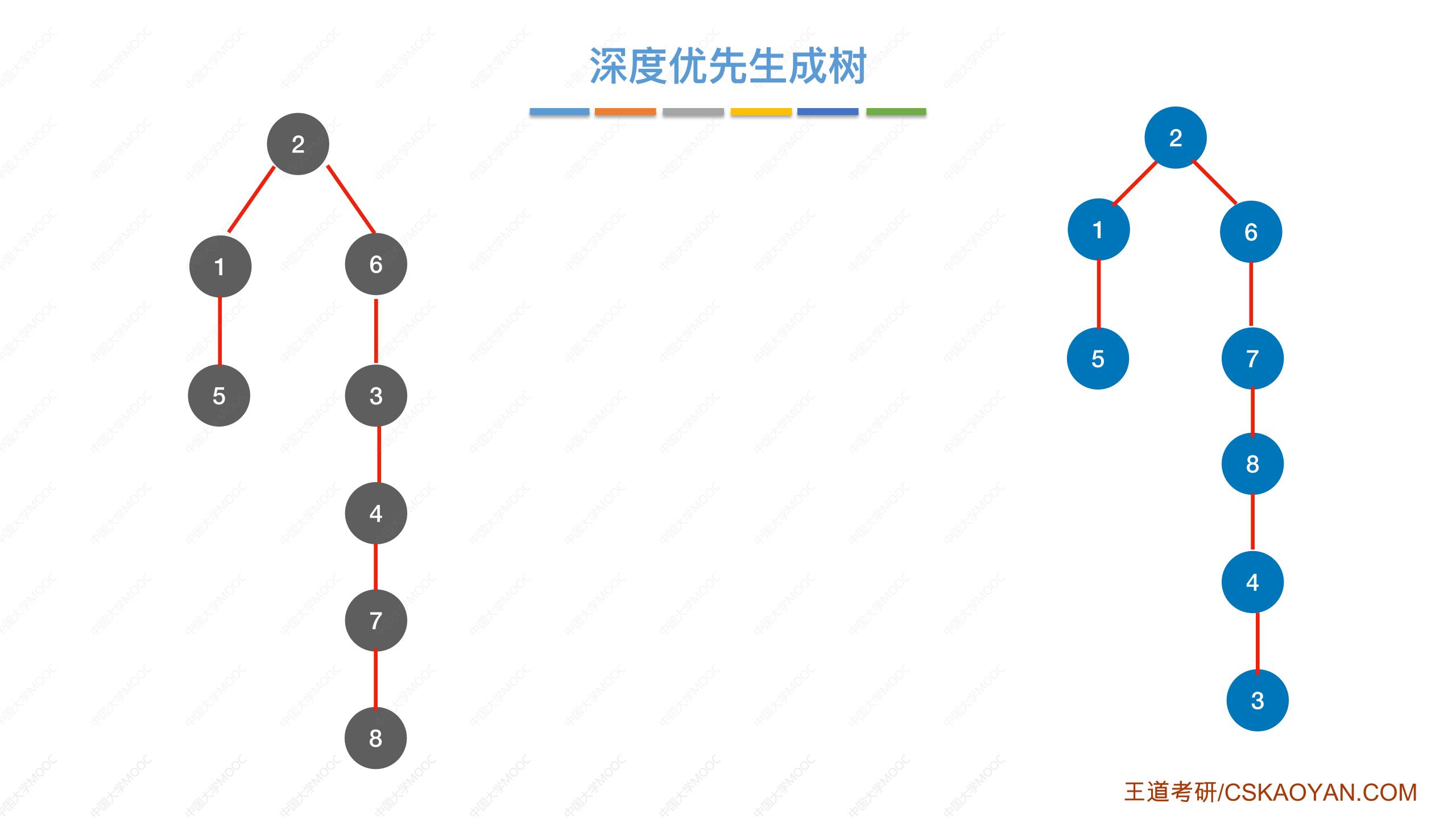
同一个图的邻接矩阵表示方式唯一,因此深度优先遍历序列唯一同一个图邻接表表示方式不唯一,因此深度优先遍历序列不唯一

# 深度优先生成树

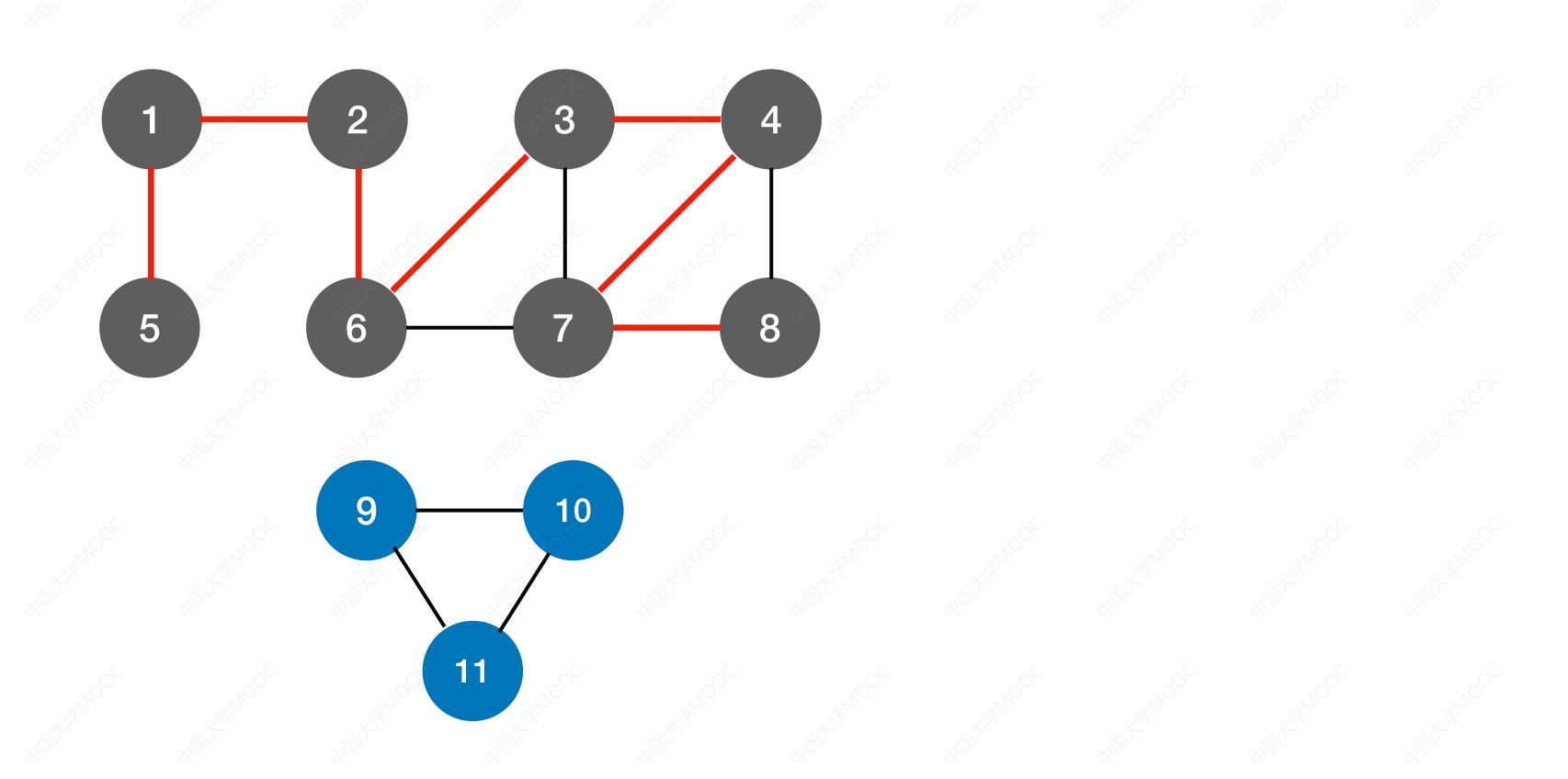




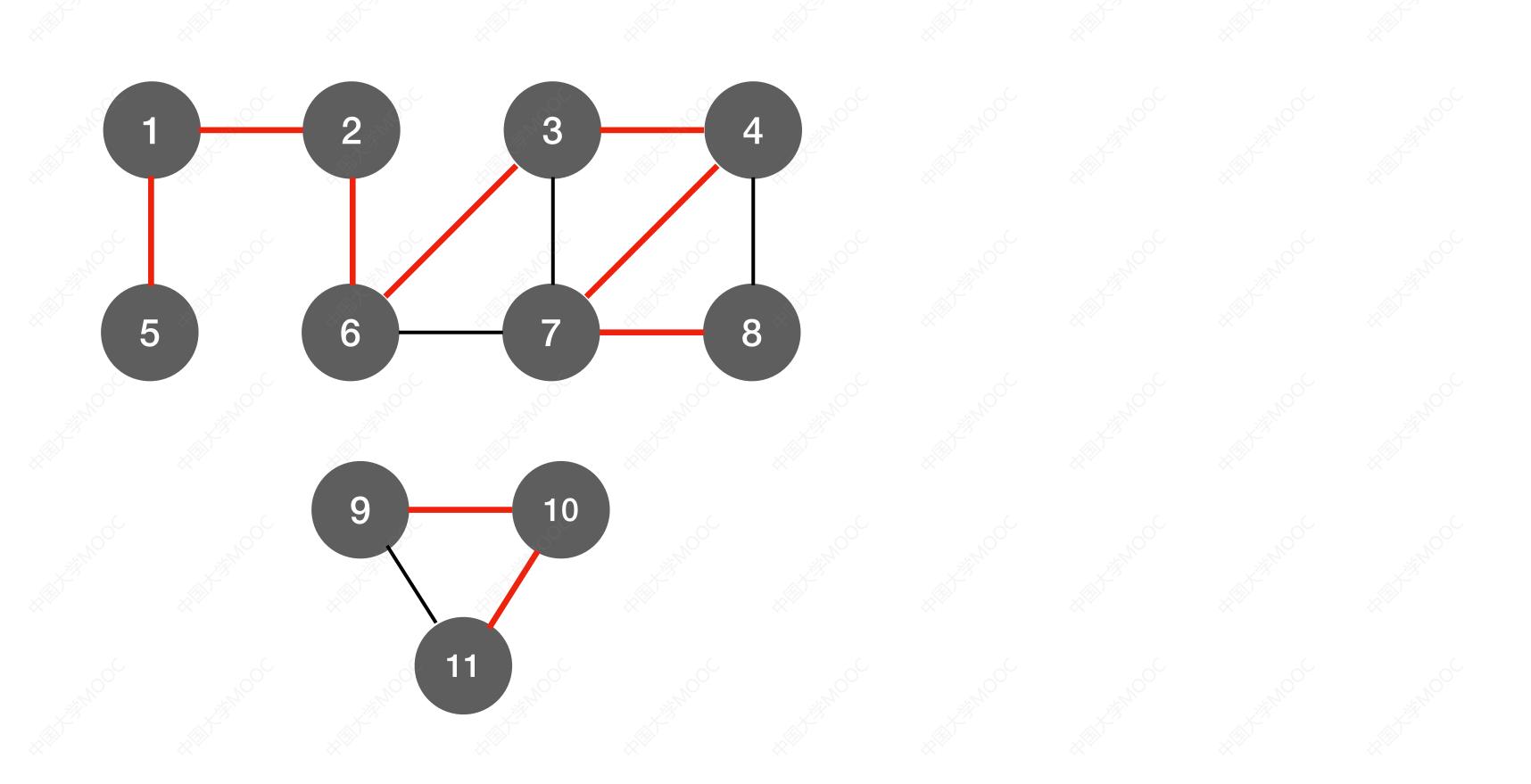
同一个图的邻接矩阵表示方式唯一,因此深度优先遍历序列唯一,深度优先生成树也唯一同一个图邻接表表示方式不唯一,因此深度优先遍历序列不唯一,深度优先生成树也不唯一

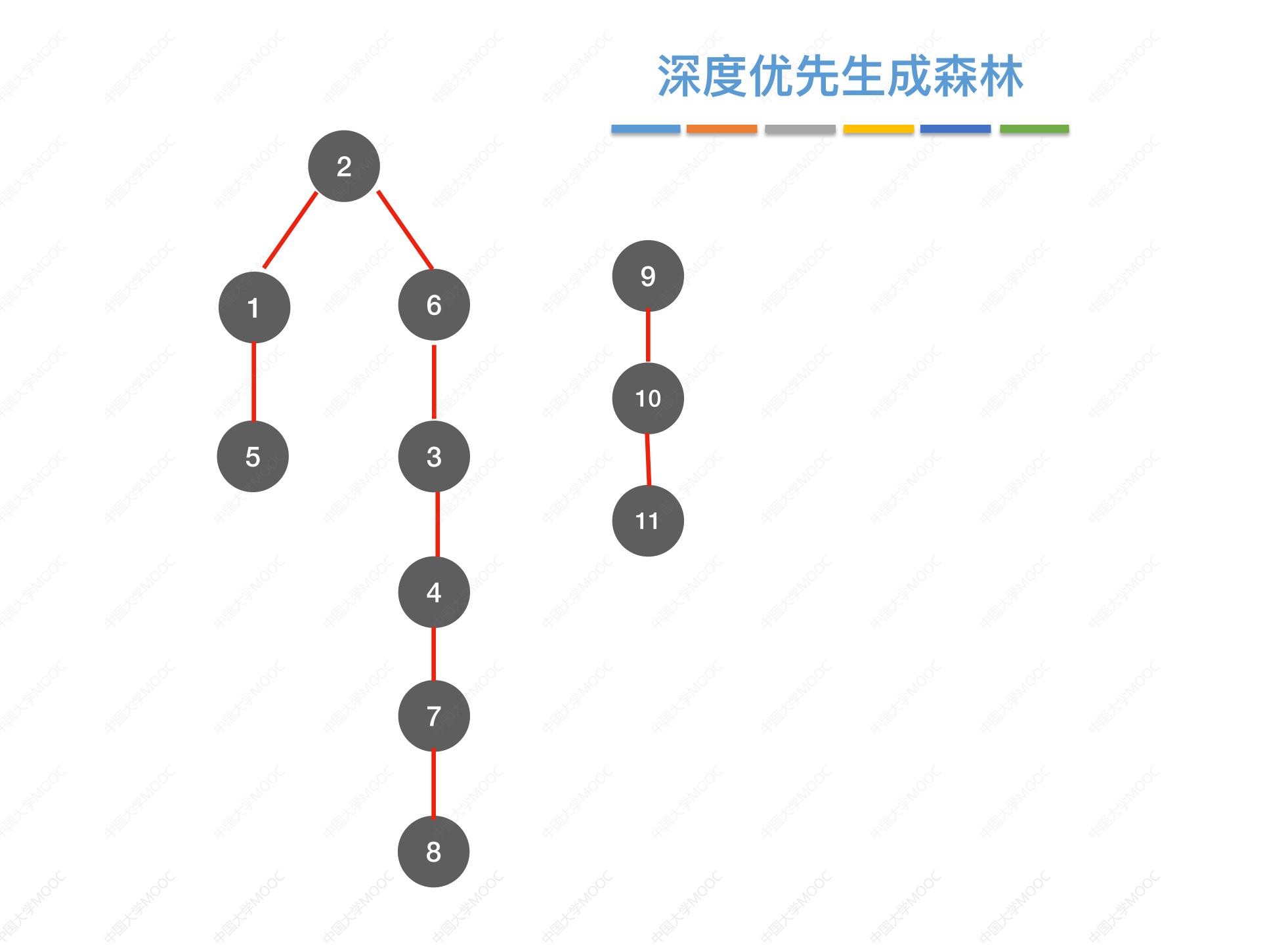


# 深度优先生成森林

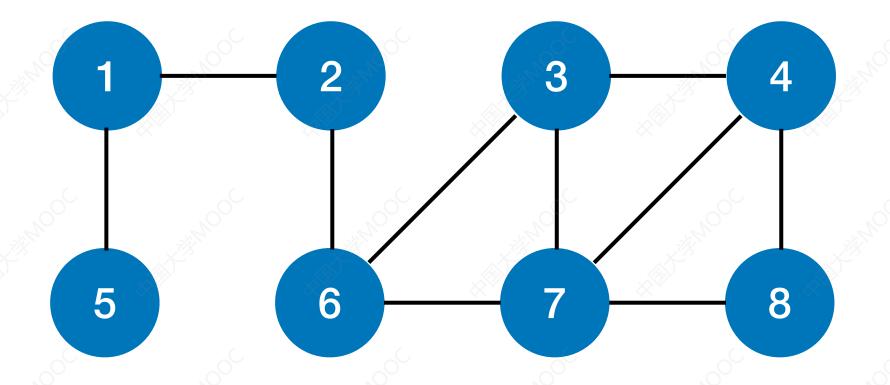


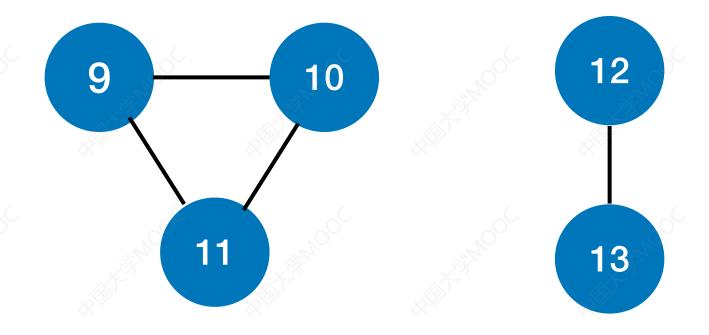
# 深度优先生成森林





# 图的遍历与图的连通性

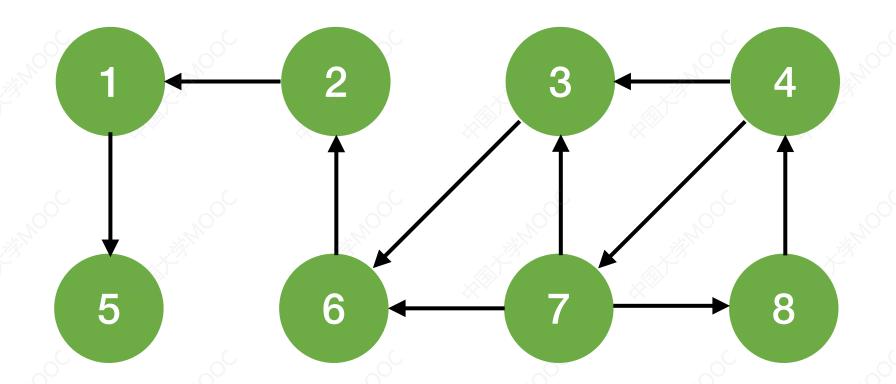




对无向图进行BFS/DFS遍历 调用BFS/DFS函数的次数=连通分量数

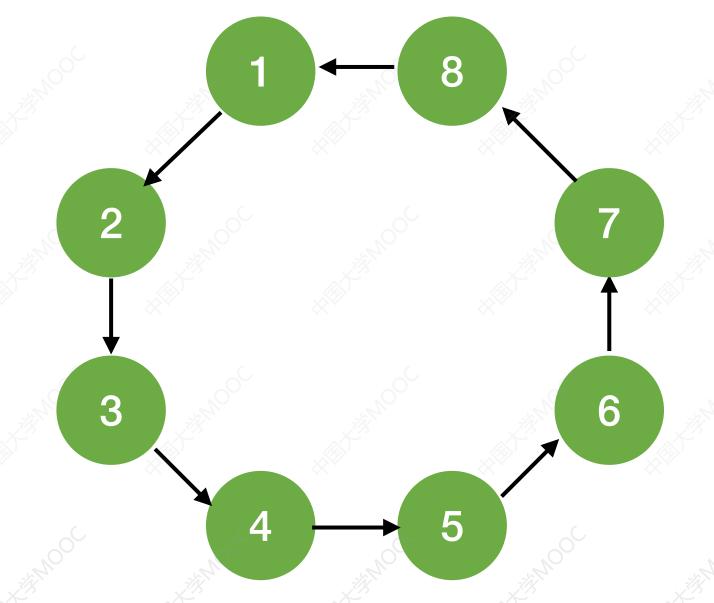
对于连通图,只需调用1次 BFS/DFS

# 图的遍历与图的连通性



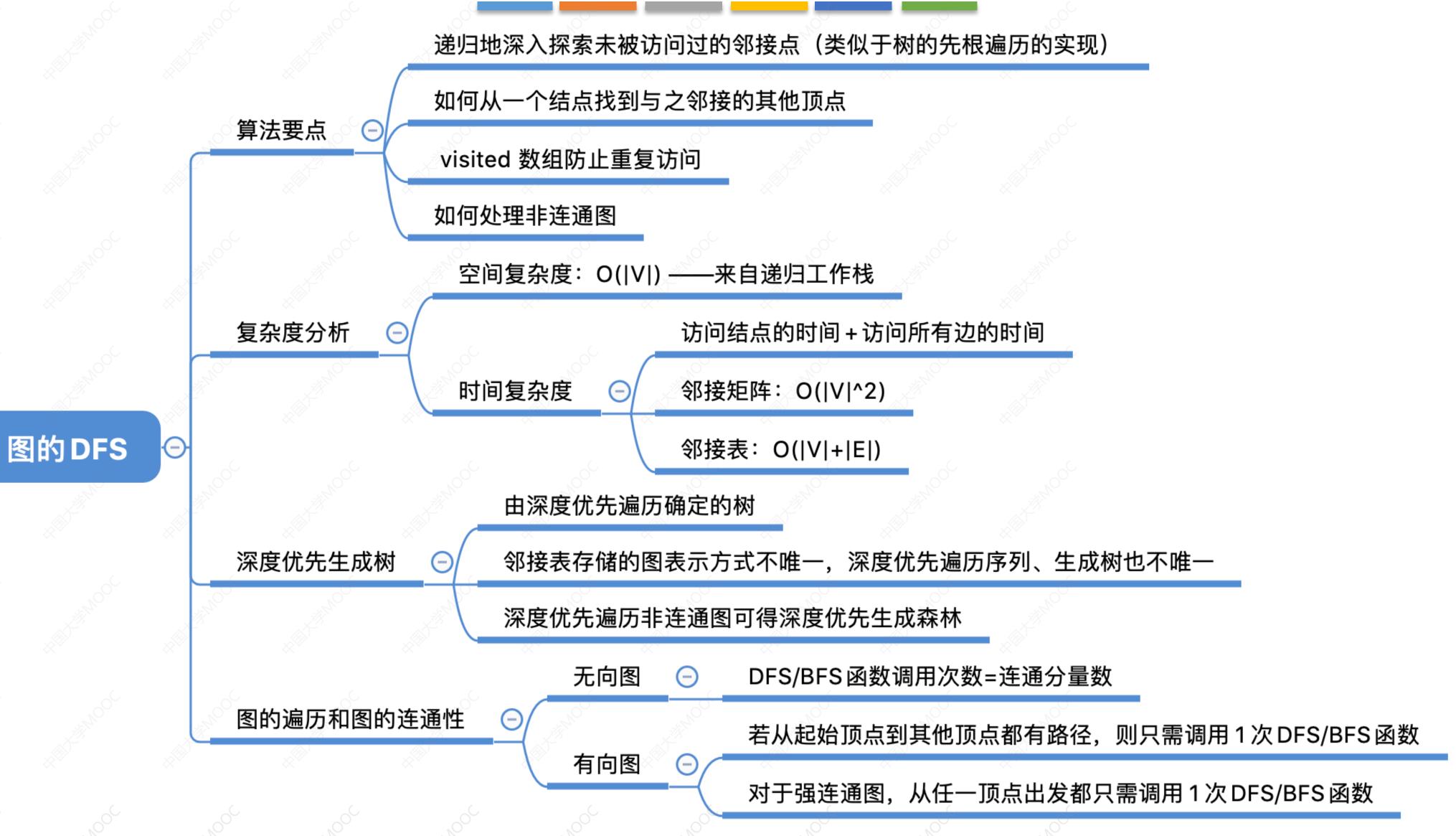
对有向图进行BFS/DFS遍历 调用BFS/DFS函数的次数要具体问题具体分析

若起始顶点到其他各顶点都有路径,则只需调用1次 BFS/DFS 函数



对于强连通图,从任一结点出发都只需调用1次 BFS/DFS

### 知识回顾与重要考点



# 欢迎大家对本节视频进行评价~



学员评分: 6.3.2 图的...





公众号: 王道在线



b站: 王道计算机教育



抖音: 王道计算机考研