

# Roulette

Saman Saadi



## 0.1 Summary

Let's assume we want to bet for the round  $i^{th}$  in a game and our target profit is  $\$p$  for each round until we win. So far we lost in the previous rounds. We call The amount of bet in  $i^{th}$  round,  $b_i$ :

$$b_i = 2 \times b_{i-1}$$

Or we can say:

$$b_i = (2^{i-1}) \times p$$

In other words, if our previous bet was  $x$ , we should bet  $2 \times x$  this time!

If we want to afford to lose  $n - 1$  times in a row and win in the  $n^{th}$  round, our budget should be:

$$budget = (2^n - 1) \times p$$

## 0.2 Expected value

The American roulette has 37 pockets and the Canadian has 36. We use  $n$  for the number of pockets in the wheel. Let's assume we bet on  $p$  numbers. For example for betting on number 1,  $p = 1$  and for a red number  $p = 18$ . The payout is  $36 - p$  to  $p$ . That means for every  $\$p$  bet, we get  $\$36 - \$p$ . We also receive the original  $\$p$ . So our total balance for  $\$1$  bet if we would win is:

$$\begin{aligned} & \frac{36 - p + p}{p} \\ &= \frac{36}{p} \end{aligned}$$

If we lose, we lost  $\$1$ . So the expected value is:

$$\begin{aligned} E &= \frac{36 - p}{p} \times \frac{p}{n} - \frac{p}{p} \times \frac{n - p}{n} \\ &= \frac{36 - p}{n} - \frac{n - p}{n} \\ &= \frac{36 - n}{n} \\ &= \frac{36}{n} - 1 \end{aligned}$$

As you can see the expected value is not related to our choice! If we have 36 pockets, it's zero and we can call it a fair game. In Canada it's  $\frac{36}{37} - 1 = -0.027$  and in the US it's  $\frac{36}{38} - 1 = -0.053$ . In other words, in Canada on average we should lose 3 cents per dollar and in US, 5 cents per dollar.

### 0.3 General case

Let's assume we are in round  $i^{th}$  of a game. We can calculate the balance for round  $i$  that we call it  $b_i$ . We want to make  $p$  unrealized profit/loss for this game, assuming we eventually win in this round

#### 0.3.1 Mathematical induction

We define  $b_i$  as the amount of bet for round  $i$  in such a way that if we win round  $i$ , our total profit for this game would be  $p$ . We use mathematical induction and we assume we know how to solve  $b_{i-1}$ .

**When  $b_{i-1} < 0$**

Let's assume we will win at the end of round  $i$  and our profit will be  $p$ . In round  $i - 1$  we lost  $b_{i-1}$  so we just need to spend  $b_{i-1}$  to cancel it. We use mathematical hypothesis and spend another  $b_{i-1}$  to cancel losses for rounds 1 to  $i - 2$  and gain  $p$  at the end of the round (assuming we will win):

$$\begin{aligned}
 b_i &= \sum_{j=0}^{i-1} b_j - p \\
 &= \underbrace{\sum_{j=0}^{i-2} b_j - p}_{b_{i-1}} + b_{i-1} \\
 &= 2 \times b_{i-1}
 \end{aligned}$$

**When  $b_{i-1} \geq 0$**

Note that if  $b_{i-1} \geq p$ , we've already achieved our goal and no need to continue the game for the profit  $p$ . We may increase the target profit to a more ambition one later.

Let's assume  $b_{i-1} > 0 \wedge b_{i-1} < p$  and we will win at the end of round  $i$  and our profit will be  $p$ :

$$\begin{aligned}
 b_{i-1} + \underbrace{x}_{bet} - \underbrace{2 \times x}_{rewards} &= p \\
 \implies x &= b_{i-1} - p
 \end{aligned}$$

So we just need to bet on  $b_{i-1} - p$  and if we win this round our profit will be  $p$ .

### 0.3.2 General Formula

$$b_i = \begin{cases} 2 \times b_{i-1} & b_{i-1} < 0 \\ b_{i-1} - p & b_{i-1} \geq 0 \wedge b_{i-1} < p \\ b_{i-1} & b_{i-1} \geq 0 \wedge b_{i-1} \geq p \\ b_0 - p & i = 1 \end{cases}$$

Note that we define  $b_1 = b_0 - p$ . Usually  $b_0 = 0$ , but if in the round  $k$  we decide to change  $p$  to  $p'$ . Then we can initiate a new game with  $b_0 = b_k$  and profit  $p'$ .

## 0.4 Losing all previous rounds

Let's consider the worst case scenario and we are going to lose in  $n$  rounds in a row:

$$\begin{aligned} b_0 &= 0 \\ b_1 &= b_0 - p = -p \\ b_2 &= 2 \times b_1 = -2p \\ b_3 &= 2 \times b_2 = -4p \\ b_4 &= 2 \times b_3 = -8p \\ &\vdots \\ b_n &= 2 \times b_{n-1} = -2^{n-1} \times p \end{aligned}$$

## 0.5 Budget for $n$ rounds

Now let's calculate what should be our budget if we can afford to lose in  $n - 1$  rounds and win at round  $n$ . We assume  $r_i$  is our loss in round  $i^{th}$  for  $1 \leq i \leq n - 1$

$$\begin{aligned} budget &= - \sum_{i=0}^n b_i \\ &= \sum_{i=0}^n 2^{i-1} \times p \\ &= (2^n - 1) \times p \end{aligned}$$

Note that we use the following formula to get the budget:

$$\sum_{i=0}^n 2^i = 2^{n+1} - 1$$