

# Assignment #4

CSE341: Principles of Programming Languages  
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Out: Nov 19, 2020 (Thu)  
Due: Dec 2, 2020 (Wed), 23:59 (KST)

## What to submit

Submit your `Hw3.scala` file through the Blackboard.



**Info:** The directory structure of the handout is as follows.

<code>sbt/</code>	- contains the sbt program that you need to test your program.
<code>src/</code>	- where all your scala source files leave.
<code>main/scala/</code>	
<code>Hw4.scala</code>	- >>>> what you need to edit and submit. <<<<<
<code>Parser.scala</code>	- The parser driver for the languages you will interpret.
<code>main/antlr4/</code>	- where inputs to the parser generator lives. You can ignore this.
<code>test/scala/</code>	
<code>Hw4Test.scala</code>	- The tests that I wrote for you. You can edit this to further test your program.

## Rules

- You must not use the `var`, `for`, or `while` keyword.
- You must not include any additional packages or libraries besides the ones that you already have.

## Scala environment

Please refer to the instruction for the first assignment to set up the Scala environment.

# Problems

## Problem 1 (60 points)

Implement an interpreter that evaluates a language that looks like C, that adopts implicit references and procedures with multiple arguments.

### Syntax

$$\begin{aligned} P &\rightarrow E \\ E &\rightarrow \text{skip} \mid \text{true} \mid \text{false} \mid n \mid x \\ &\mid E + E \mid E - E \mid E * E \mid E / E \\ &\mid E == E \mid E <= E \mid \text{not } E \\ &\mid \text{if } E \text{ then } E \text{ else } E \\ &\mid \text{while } E \text{ } E \mid \text{let } x = E \text{ in } E \\ &\mid \text{proc } (x_1, x_2, \dots, x_n) E \\ &\mid E (E_1, E_2, \dots, E_n) \quad \text{call-by-value} \\ &\mid E \langle y_1, y_2, \dots, y_n \rangle \quad \text{call-by-reference} \\ &\mid \{ \} \mid \{ x_1 := E_1, \dots, x_n := E_n \} \\ &\mid x := E \mid E.x \mid E.x := E \mid E; E \mid \text{begin } E \text{ end} \end{aligned}$$

In scala,

```
sealed trait Program
sealed trait Expr extends Program
case object Skip extends Expr
case object False extends Expr
case object True extends Expr
case class NotExpr(expr: Expr) extends Expr
case class Const(n: Int) extends Expr
case class Var(s: String) extends Expr
case class Add(l: Expr, r: Expr) extends Expr
case class Sub(l: Expr, r: Expr) extends Expr
case class Mul(l: Expr, r: Expr) extends Expr
case class Div(l: Expr, r: Expr) extends Expr
case class LTEExpr(l: Expr, r: Expr) extends Expr
case class EQExpr(l: Expr, r: Expr) extends Expr
case class Iszero(c: Expr) extends Expr
case class Ite(c: Expr, t: Expr, f: Expr) extends Expr
case class Let(i: Var, v: Expr, body: Expr) extends Expr
case class Proc(args: List[Var], expr: Expr) extends Expr
case class Asn(v: Var, e: Expr) extends Expr
case class BeginEnd(expr: Expr) extends Expr
case class FieldAccess(record: Expr, field: Var) extends Expr
case class FieldAssign(record: Expr, field: Var, new_val: Expr) extends Expr
case class Block(f: Expr, s: Expr) extends Expr
case class PCallV(ftn: Expr, arg: List[Expr]) extends Expr
case class PCallR(ftn: Expr, arg: List[Var]) extends Expr
case class WhileExpr(cond: Expr, body: Expr) extends Expr
sealed trait RecordLike extends Expr
case object EmptyRecordExpr extends RecordLike
case class RecordExpr(field: Var, initVal: Expr, next: RecordLike)
extends RecordLike
```

The Domain on which the semantics is defined.

Domain

$$\begin{aligned} Val &= \mathbb{Z} + Bool + \{\cdot\} + Procedure + Loc + Record \\ Procedure &= (Var \times Var \times \dots)E \times Env \\ r \in Record &= Field \rightarrow Loc \\ \rho \in Env &= Var \rightarrow Loc \\ \sigma \in Mem &= Loc \rightarrow Val \end{aligned}$$

In Scala,

```
sealed trait Val
case object SkipVal extends Val
case class IntVal(n: Int) extends Val
case class BoolVal(b: Boolean) extends Val
case class ProcVal(args: List[Var], expr: Expr, env: Env) extends Val
case class LocVal(l: Int) extends Val
sealed trait RecordValLike extends Val
case object EmptyRecordVal extends RecordValLike
case class RecordVal(field: Var, loc: LocVal, next: RecordValLike)
extends RecordValLike

type Env = HashMap[Var, Val]

case class Mem(m: HashMap[Loc, Val], top: Loc)
```

The semantics rules of the language is as follows.

#### Constants and Variables

$$\frac{}{\rho, \sigma \vdash \text{skip} \Rightarrow \cdot, \sigma} \quad \frac{}{\rho, \sigma \vdash \text{true} \Rightarrow \text{true}, \sigma} \quad \frac{}{\rho, \sigma \vdash \text{false} \Rightarrow \text{false}, \sigma}$$

$$\frac{}{\rho, \sigma \vdash n \Rightarrow n, \sigma} \quad \frac{}{\rho, \sigma \vdash x \Rightarrow \sigma(\rho(x)), \sigma} \quad \rho(x) \in \text{Dom}(\sigma)$$

$$\frac{}{\rho, \sigma \vdash \text{proc } (x_1, \dots, x_n) E \Rightarrow ((x_1, \dots, x_n), E, \rho), \sigma}$$

#### Unary and Binary Operations

$$\frac{\rho, \sigma_0 \vdash E_1 \Rightarrow n_1, \sigma_1 \quad \rho, \sigma_1 \vdash E_2 \Rightarrow n_2, \sigma_2}{\rho, \sigma_0 \vdash E_1 + E_2 \Rightarrow n_1 + n_2, \sigma_2} \quad \frac{\rho, \sigma_0 \vdash E_1 \Rightarrow n_1, \sigma_1 \quad \rho, \sigma_1 \vdash E_2 \Rightarrow n_2, \sigma_2}{\rho, \sigma_0 \vdash E_1 - E_2 \Rightarrow n_1 - n_2, \sigma_2}$$

$$\frac{\rho, \sigma_0 \vdash E_1 \Rightarrow n_1, \sigma_1 \quad \rho, \sigma_1 \vdash E_2 \Rightarrow n_2, \sigma_2}{\rho, \sigma_0 \vdash E_1 * E_2 \Rightarrow n_1 * n_2, \sigma_2} \quad \frac{\rho, \sigma_0 \vdash E_1 \Rightarrow n_1, \sigma_1 \quad \rho, \sigma_1 \vdash E_2 \Rightarrow n_2, \sigma_2}{\rho, \sigma_0 \vdash E_1 / E_2 \Rightarrow n_1 / n_2, \sigma_2} \quad n_2 \neq 0$$

$$\frac{\rho, \sigma_0 \vdash E_1 \Rightarrow n_1, \sigma_1 \quad \rho, \sigma_1 \vdash E_2 \Rightarrow n_2, \sigma_2}{\rho, \sigma_0 \vdash E_1 \leq E_2 \Rightarrow \text{true}, \sigma_2} \quad n_1 \leq n_2$$

$$\frac{\rho, \sigma_0 \vdash E_1 \Rightarrow n_1, \sigma_1 \quad \rho, \sigma_1 \vdash E_2 \Rightarrow n_2, \sigma_2}{\rho, \sigma_0 \vdash E_1 < E_2 \Rightarrow \text{false}, \sigma_2} \quad n_1 > n_2$$

$$\frac{\rho, \sigma_0 \vdash E_1 \Rightarrow n_1, \sigma_1 \quad \rho, \sigma_1 \vdash E_2 \Rightarrow n_2, \sigma_2}{\rho, \sigma_0 \vdash E_1 == E_2 \Rightarrow \text{true}, \sigma_2} \quad v_1 = v_2 = n \vee v_1 = v_2 = b \vee v_1 = v_2 = \cdot$$

$$\frac{\rho, \sigma_0 \vdash E_1 \Rightarrow n_1, \sigma_1 \quad \rho, \sigma_1 \vdash E_2 \Rightarrow n_2, \sigma_2}{\rho, \sigma_0 \vdash E_1 == E_2 \Rightarrow \text{false}, \sigma_2} \quad \text{otherwise}$$

$$\frac{\rho, \sigma_0 \vdash E \Rightarrow \text{true}, \sigma_1}{\rho, \sigma_0 \vdash \text{not } E \Rightarrow \text{false}, \sigma_1} \quad \frac{\rho, \sigma_0 \vdash E \Rightarrow \text{false}, \sigma_1}{\rho, \sigma_0 \vdash \text{not } E \Rightarrow \text{true}, \sigma_1}$$

Note that only integers can be compared with each other with  $>$  and  $>=$ , and the equality operation evaluates to false under type mismatches also.

#### Flow Control

$$\frac{\rho, \sigma_0 \vdash E \Rightarrow \text{false}, \sigma_1 \quad \rho, \sigma_1 \vdash E_3 \Rightarrow v, \sigma_2}{\rho, \sigma_0 \vdash \text{if } E_1 \text{ then } E_2 \text{ else } E_3 \Rightarrow v, \sigma_2}$$

$$\frac{\rho, \sigma_0 \vdash E \Rightarrow \text{true}, \sigma_1 \quad \rho, \sigma_1 \vdash E_2 \Rightarrow v, \sigma_2}{\rho, \sigma_0 \vdash \text{if } E_1 \text{ then } E_2 \text{ else } E_3 \Rightarrow v, \sigma_2}$$

$$\frac{\rho, \sigma_0 \vdash E_1 \Rightarrow \text{false}, \sigma_1}{\rho, \sigma_0 \vdash \text{while } E_1 E_2 \Rightarrow \cdot, \sigma_2}$$

$$\frac{\rho, \sigma \vdash E_1 \Rightarrow \text{true}, \sigma_0 \quad \rho, \sigma_0 \vdash E_2 \Rightarrow v_1, \sigma_1 \quad \rho, \sigma_1 \vdash \text{while } E_1 E_2 \Rightarrow v_2, \sigma_2}{\rho, \sigma \vdash \text{while } E_1 E_2 \Rightarrow \cdot, \sigma_2}$$

$$\frac{\rho, \sigma \vdash E_1 \Rightarrow v_1, \sigma_1 \quad \rho, \sigma_1 \vdash E_2 \Rightarrow v_2, \sigma_2}{\rho, \sigma \vdash E_1; E_2 \Rightarrow v_2, \sigma_2}$$

$$\frac{\rho, \sigma \vdash E \Rightarrow v, \sigma'}{\rho, \sigma \vdash \text{begin } E \text{ end} \Rightarrow v, \sigma'}$$

### Records

$$\begin{array}{c}
\frac{}{\rho, \sigma \vdash \{ \} \Rightarrow \cdot, \sigma} \\
\frac{\rho, \sigma_0 \vdash E_1 \Rightarrow v_1, \sigma_1 \quad \dots \quad \rho, \sigma_n \vdash E_n \Rightarrow v_n, \sigma_n \quad l_1, \dots, l_n \notin \text{Dom}(\sigma_n)}{\rho, \sigma_0 \vdash \{ x_1 := E_1, \dots, x_n := E_n \} \Rightarrow \{ x_1 \mapsto l_1, \dots, x_n \mapsto l_n \}, [l_1 \mapsto v_1, \dots, l_n \mapsto v_n] \sigma_n} \\
\frac{\rho, \sigma_0 \vdash E \Rightarrow r, \sigma_1}{\rho, \sigma_0 \vdash E.x \Rightarrow \sigma_1(r(x)), \sigma_1} \\
\frac{\rho, \sigma_0 \vdash E_1 \Rightarrow r, \sigma_1 \quad \rho, \sigma_1 \vdash E_2 \Rightarrow v, \sigma_2}{\rho, \sigma_0 \vdash E_1.x := E_2 \Rightarrow v, [r(x) \mapsto v] \sigma_2}
\end{array}$$

### Assignments

$$\begin{array}{c}
\frac{\rho, \sigma \vdash E \Rightarrow v, \sigma_1}{\rho, \sigma \vdash x := E \Rightarrow v, [\rho(x) \mapsto v] \sigma_1} \\
\frac{\rho, \sigma_0 \vdash E_1 \Rightarrow v_1, \sigma_1 \quad [x \mapsto l] \rho, [l \mapsto v_1] \sigma_1 \vdash E_2 \Rightarrow v, \sigma_2 \quad l \notin \text{Dom}(\sigma_1)}{\rho, \sigma_0 \vdash \text{let } x = E_1 \text{ in } E_2 \Rightarrow v, \sigma_2}
\end{array}$$

### Procedure Calls

$$\begin{array}{c}
\frac{\rho, \sigma \vdash E_0 \Rightarrow ((x_1, \dots, x_n), E, \rho') \quad \rho, \sigma_0 \vdash E_1 \Rightarrow v_1, \sigma_1 \quad \dots \quad \rho, \sigma_{n-1} \vdash E_n \Rightarrow v_n, \sigma_n}{\frac{[x_1 \mapsto l_1, \dots, x_n \mapsto l_n] \rho', [l_1 \mapsto v_1, \dots, l_n \mapsto v_n] \sigma_n \vdash E \Rightarrow v, \sigma'}{\rho, \sigma \vdash E_0 (E_1, \dots, E_n) \Rightarrow v, \sigma'} \quad l_1, \dots, l_n \notin \text{Dom}(\sigma_n)} \\
\frac{\rho, \sigma \vdash E_0 \Rightarrow ((x_1, \dots, x_n), E, \rho'), \sigma_0 \quad [x_1 \mapsto \rho(y_1), \dots, x_n \mapsto \rho(y_n)] \rho', \sigma_0 \vdash E \Rightarrow v, \sigma'}{\rho, \sigma \vdash E_0 \langle y_1, \dots, y_n \rangle \Rightarrow v, \sigma'}
\end{array}$$

In the skeleton, you can find the [MiniCInterpreter](#) object whose apply method looks like:

```
def apply(program: String): (Val, Mem)
```

and calls the parser and the interpreter for you. Your job is to fill out the body of this method.

```
def eval(env: Env, mem: Mem, expr: Expr): Result = Result(SkipVal, mem)
```

As noted in class, a valid program that passes the parser may not have its semantics. If this is the case, this time, you have to throw a particular exception that is defined in the object as follows.

```
case class UndefinedSemantics(msg: String = "", cause: Throwable = None.orNull)
extends Exception("Undefined Semantics: " ++ msg, cause)
```

You can throw the exception as follows.

```
throw new UndefinedSemantics(s"message ${variable}")
```

Problem 2 (40 points)

In this language, we allocate new memory locations in `let`, `call`, and `record` expressions, but they never get freed, leading to memory exhaustion. Fill out another method `gc` in the Interpreter and call it. Please note that your interpreter does not need to call this `gc`. It will be evaluated separately as being done tested with the embedded test.

```
def gc(env: Env, mem: Mem): Mem
```