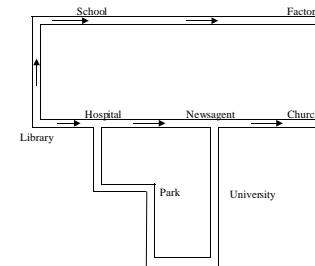


## Best First Search

- Greedy Search
  - expand the “best” successor node
  - minimise cost estimate to nearest goal:  $h(n)$ 
    - heuristic function - problem specific
    - prefers biggest local bite regardless of long term effect - hence the name!

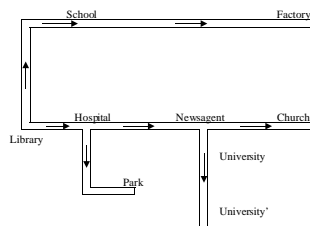
1. Start with *agenda* = [initial state]
2. While *agenda* not empty:
  - (a) remove the best node N from *agenda*
  - (b) if it is the goal then return success  
else find its successors
  - (c) assign successor nodes a score using evaluation function and add scored nodes to *agenda*

## Greedy Search cont'd

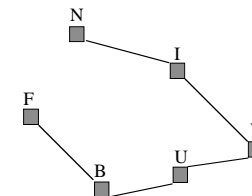


- Evaluation
  - Non-optimal finds high cost paths
  - Incomplete - gets stuck oscillating
  - Complexity - time and space exponential  $\Rightarrow b^m$

## Greedy Search cont't



## Greedy Search Again



Heuristic Functions:

$h(N) = 100$ ,  $h(I) = 150$ ,  $h(V) = 200$

Expanding Iasi gives:

(N:100, V:200)  $\leftarrow$  priority queue

Expanding Neamt (N:100) gives:

(I:150, V:200)

Expanding Iasi again (I:100) gives:

(N:100, V:200, V:200)

etc.

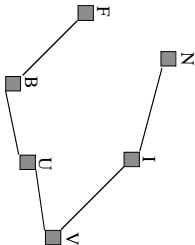
# Beam Search

- Based on Breadth first
    - Expands best *w* nodes at each level (others ignored)
    - Minimises cost estimate  $h(n)$
    - Only *w*b nodes explored at any depth
    - Only *w*d nodes stored.
1. Start with *queue* = [initial state] and *found* = FALSE
  2. While *queue* not empty and not *found* do:
    - (a) loop *w* times
      - (a1) remove first node *N* from *queue*
      - (a2) if *N* is goal state then *found* = TRUE
      - (a3) find all successor nodes of *N* and add them to the *queue*
    - (b) assign successor nodes a score and prioritise accordingly.
    - (c) remove last (*l* - *w*) nodes from *queue*

# Beam Search (2)

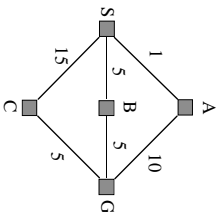
- Evaluation
  - Non-optimal
    - May still find high cost paths
      - But less likely to
  - Incomplete -
    - May still get stuck oscillating
      - But less likely to
- Complexity
  - time and space polynomial
    - *w*b and *w*d

# Beam Search Again



Heuristic Functions:  
 $h(N) = 100, h(I) = 150, h(V) = 200$   
Expanding *I*asi gives:  
(*N*:100, *V*:200)  $\Leftarrow$  priority queue  
Expanding Neamt and Vashui gives:  
(~~*I*~~:150, *U*:190, *I*:150)  
Expanding again gives:  
(*N*:100, ~~*V*~~:200, *B*:178, ~~*V*~~:200)  
etc.

# Uniform Cost Search



Expanding the start gives:  
(*A*:1 *B*:5 *C*:15)  
Expanding *A*:1 gives:  
(*B*:5 *G*:11 *C*:15)  
Expanding *B*:5 gives:  
(*G*:10 *G*:11 *C*:15)

## A\* Search

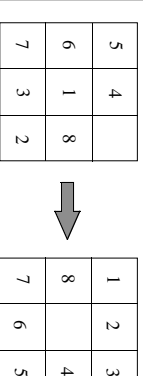
- Synthesis
  - Uniform Search minimizes  $g(n)$ 
    - optimal, complete, but inefficient
  - Greedy search minimizes  $h(n)$ 
    - non-optimal and incomplete
  - Combine and minimize  $f(n)$ 
    - $f(n) = g(n) + h(n)$
- Restrictions:
  - $h(n)$  must never overestimate the cost to reach a goal
    - Admissible heuristic - optimistic
  - $f(n)$  must be monotonic

## A\* Search 2

- PATHMAX
  - maintains monotonicity
    - if  $n'$  is a child of  $n$  and  $f(n') < f(n)$  then:  
 $f(n') = \max(f(n), g(n') + f(n'))$
- Proof of A\* optimality
  - Russel and Norvig page 99
- Proof of A\* completeness
  - Russel and Norvig page 100
- A\* complexity
  - exponential in the solution length
    - subexponential condition:  
 $|h(n) - h^*(n)| \leq O(\log h^*(n))$
    - *memory will run out first!!!*

## Heuristic Functions

- The Eight Puzzle



- Possible heuristics
  - No of tiles in wrong position:  $h1$
  - Sum of distances from goal:  $h2$ 
    - City block or **Manhattan Distance**

## Accuracy & Performance

- Effective Branching Factor
  - uniform tree equivalent with N nodes if A\* expands N nodes
  - Domination:
    - if one heuristic function has a higher value than another.

It is always better to use heuristic functions with higher values as long as they do not overestimate!

## Accuracy & Performance

$d$	Search Cost			Effective Branching Factor			
	IDS	$A^*(h)$	$A^*(b_2)$	IDS	$A^*(h)$	$A^*(b_2)$	
2	10	6	6	2.45	1.79	1.79	
4	112	13	12	2.87	1.48	1.45	
6	680	20	18	2.73	1.34	1.30	
8	6384	39	25	2.80	1.33	1.24	
10	47127	93	39	2.79	1.38	1.22	
12	364404	227	73	2.78	1.42	1.24	
14	3473941	539	113	2.83	1.44	1.23	
16	-	1301	211	-	1.45	1.25	
18	-	3056	363	-	1.46	1.26	
20	-	7276	676	-	1.47	1.27	
22	-	18094	1219	-	1.48	1.28	
24	-	39135	1641	-	1.48	1.26	

## Inventing Heuristics

- Relaxed problem
  - less restrictions on operators
  - exact sol'n to RP  $\Rightarrow$  good heuristic
- Best heuristic?
  - not always clear
    - $h(n) = \max(h_1(n), \dots, h_m(n))$
- Statistical Information
  - random tests to gather statistics
  - can lose admissibility
- Pick out “Features”
  - relevant info for heuristics from present state
- Machine Learning

Good Heuristic Functions must be efficient as well as accurate

## Memory Bounded Search

- Memory is usually the first thing to give!
- IDA\*
  - extension of ID search to use heuristic information
- SMA\*
  - restricts A\* agenda size to fit available memory.

## IDA\* Search

- Depth First Iterations
  - f-cost replaces depth limit
  - complete search inside f-contour
  - next contour min of successors f-cost
- Space complexity
  - polynomial (aprox.  $bd$ )
- Time complexity
  - no priority queue saves time
  - proportional to heuristic values
  - $\epsilon$ -Admissible
    - fixed step increase for each iteration
    - sub-optimal by at most  $\epsilon$

## SMA\* Search

- Problems with IDA\*
  - uses too little memory
  - only keeps current f-cost
  - forgets history -> repeats it
- Simplified Memory-bounded A\*
  - uses all available memory for search
  - leads to improved search efficiency
- Other Properties
  - Avoids repeated states
    - as far as memory allows
  - Complete
    - if memory stores shallowest sol'n path
  - Optimal
    - if shallowest optimal path can be stored
  - Optimally Efficient
    - if memory can store entire search tree

## SMA\* Search cont'd

- Forgotten Nodes
  - nodes dropped if high f-cost
    - ancestor retains info of best cost in forgotten sub-tree
  - only reconstructed if all other paths look worse.

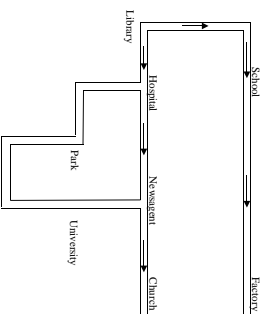
## Heuristics for CSP

- Most constrained variable heuristic
  - variable with *fewest* possible values chosen
- Most constraining variable heuristic
  - variable involved in largest no. of constraints on other variables chosen.
- Least constraining value heuristic
  - choose a value that rules out the least no. of values in variables connected to the current variable by constraints.

## Iterative Improvement Algorithms

- Hill Climbing (Gradient Descent)
  - Head towards "best" successor node that is better than present one
- 1. Start with *current-state* = initial state.
- 2. Until *current-state* = goal-state or not there is no change in *current-state* do:
  - (a) Get successors of *current-state* and use evaluation function to assign score to each successor.
  - (b) If one of successors has a better score than *current-state* then set new *current-state* to be the successor with the best score
- Problems
  - Local Maxima (Minima)
  - Plateaux
  - Ridges

## Hill Climbing cont'd



- Random Restart Hill Climbing
  - improves probability of finding global solution.