

9.3

语句  $Q$  不合法。语句  $\exists x \text{ AsHighAs}(x, \text{Everest})$  为  $x$  与 Everest 等高, 显然将  $x$  实例化为 Everest 无实际意义。

语句  $b, c$  合法。谓语句  $b, c$  分别将  $x$  实例化了 1 次、2 次, 都是合法实例化。

9.4

a.  $\sigma_1 = \{A/x\}$   $W_1 = \{p(A, B, B), p(x, y, z)\}$   $\sigma_1 = \{p(A, B, B), p(A, y, z)\}$   
 $\sigma_2 = \{B/y\}$   $W_2 = W_1$   $\sigma_2 = \{p(A, B, B), p(A, B, z)\}$   
 $\sigma_3 = \{B/z\}$   $W_3 = W_2$   $\sigma_3 = \{p(A, B, B)\}$   
 $\sigma = \sigma_1 \circ \sigma_2 \circ \sigma_3 = \{A/x, B/y, B/z\}$  为最一般合-置换

b.  $\sigma_1 = \{G(x, x)/y\}$   $W_1 = \{Q(y, G(A, B)), Q(G(x, x), y)\}$   $\sigma_1 = \{Q(G(x, x), G(A, B)), Q(G(x, x), G(x, x))\}$   
 对于  $G(A, B), G(x, x)$ , 不存在  $x$  使得不存在置换使其合一  
 因此无最一般合-置换

c.  $\sigma_1 = \{y/x\}$   $W_1 = \{older(father(y), y), older(father(x), John)\}$   $\sigma_1 = \{older(father(y), y), older(father(y), John)\}$   
 $\sigma_2 = \{John/y\}$   $W_2 = W_1$   $\sigma_2 = \{older(father(John), John)\}$   
 $\sigma = \sigma_1 \circ \sigma_2 = \{John/x, John/y\}$   
 $\sigma$  为最一般合-置换

d.  $\sigma_1 = \{Father(y)/x\}$   $W_1 = \{knows(f(y), y), knows(x, x)\}$   $\sigma_1 = \{knows(f(y), y), knows(f(y), f(y))\}$   
 对于  $y, f(y)$ , 不存在  $y = f(y)$   $Father(y)$   
 因此无最一般合-置换

9.7

a. 设  $P(x, y)$  表示  $x > y$   
 则  $\forall x \exists y P(x, y)$  成立;  $\exists q P(q, q)$  不成立.

b. 将前提 skolem化:  $P(x, f(x))$   
 目标取否定:  $\forall y \sim P(\frac{P, P}{y, y})$ .  
 如果可以合一, 则将归结为 NULL

c. 前提 skolem化:  $P(x, sko)$   
 目标取否, 在置换  $\{q/sko, x/sko\}$  作用下, 归结为 NULL

d. 对于某假设  $\exists x P(x)$ , 若证明  $P(A)$   
 将前提 skolem化:  $P(sko)$   
 目标取否, 在置换  $\{sko/A\}$  作用下, 归结为 NULL.

9.19

- a. (i) 有解  $\{John/y\}$   
 (ii) 有解  $\{John/y\}$ .  
 (iii) 有解  $\{ \}$   
 (iv) 从不终止

b. 不能证明. KB 仅包含了通过母亲来判断祖先的归结规则, 并不包含通过其他亲属关系判断是否为祖先的规则.

c. 仍无法证明

9.23.

a. 前提  $\forall x \text{ horse}(x) \Rightarrow \text{Animal}(x)$

结论  $\forall x, h \text{ horse}(x) \wedge \text{headof}(h, x) \Rightarrow \exists y \text{ Animal}(y) \wedge \text{headof}(h, y)$

b.

$$\neg(\forall x \text{ horse}(x) \Rightarrow \text{Animal}(x)) \wedge$$

$$(\forall x \text{ horse}(x) \Rightarrow \text{Animal}(x)) \wedge \neg(\forall x, h \text{ horse}(x) \wedge \text{headof}(h, x) \Rightarrow \exists y \text{ Animal}(y) \wedge \text{headof}(h, y))$$

$$(\forall x \neg \text{horse}(x) \vee \text{Animal}(x)) \wedge \neg(\neg \forall x, h \text{ horse}(x) \wedge \text{headof}(h, x) \vee \neg \exists y \text{ Animal}(y) \wedge \text{headof}(h, y))$$

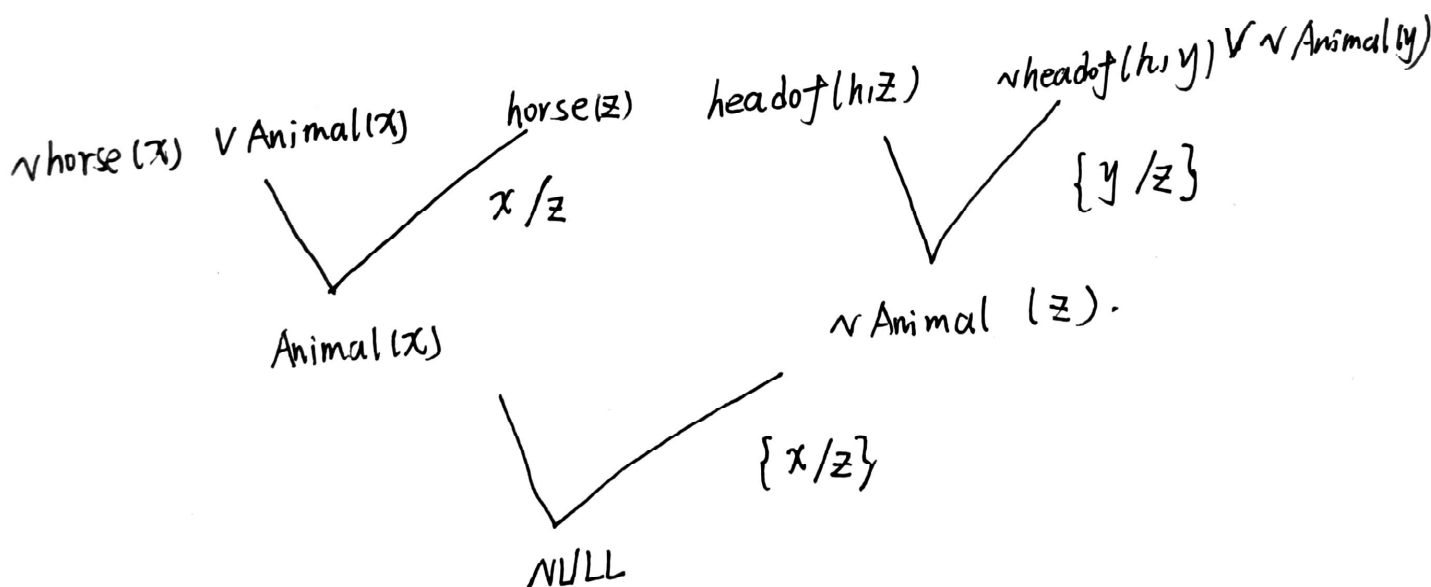
$$(\forall x \neg \text{horse}(x) \vee \text{Animal}(x)) \wedge \neg(\neg \forall x, h \text{ horse}(x) \wedge \text{headof}(h, x) \vee (\exists y \text{ Animal}(y) \wedge \text{headof}(h, y)))$$

$$(\forall x \neg \text{horse}(x) \vee \text{Animal}(x)) \wedge (\forall z, h \text{ horse}(z) \wedge \text{headof}(h, z) \vee (\neg \exists y \text{ Animal}(y) \wedge \text{headof}(h, y)))$$

$$(\neg \text{horse}(x) \vee \text{Animal}(x)) \wedge \text{horse}(z) \wedge \text{headof}(h, z) \wedge (\forall y \neg \text{Animal}(y) \vee \neg \text{headof}(h, y))$$

$$(\neg \text{horse}(x) \vee \text{Animal}(x)) \wedge \text{horse}(z) \wedge \text{headof}(h, z) \wedge (\neg \text{Animal}(y) \vee \neg \text{headof}(h, y))$$

c.



因此可由前提归结推导出结论.