Artificial Intelligence and Machine Learning (AIML)





- Last lecture: combinatorial optimization in Al
- This lecture: exact SDP methods for combinatorial optimization

 Many strategies for constructing practical exact methods: divide-and-conquer, branch-and-bound, dynamic programming

- Many strategies for constructing practical exact methods: divide-and-conquer, branch-and-bound, dynamic programming
- Usefulness/efficiency of each strategy depends upon the specific structure of the problem, no single strategy is best for all problems

- Many strategies for constructing practical exact methods: divide-and-conquer, branch-and-bound, dynamic programming
- Usefulness/efficiency of each strategy depends upon the specific structure of the problem, no single strategy is best for all problems
- Choice of strategy therefore requires understanding of the specifics of the problem itself, available computational resources and/or availability of other information about the problem

- Many strategies for constructing practical exact methods: divideand-conquer, branch-and-bound, dynamic programming
- Usefulness/efficiency of each strategy depends upon the specific structure of the problem, no single strategy is best for all problems
- Choice of strategy therefore requires understanding of the specifics of the problem itself, available computational resources and/or availability of other information about the problem
- We will concentrate on **sequential decision process (SDP)** methods as they encompass many practical AI algorithms

 An SDP algorithm is a recursive process which scans N input data items

 $x_1, ..., x_N$ in sequence, generating new candidate configurations by **extending**

- An SDP algorithm is a recursive process which scans N input data items
 - $x_1, ..., x_N$ in sequence, generating new candidate configurations by **extending** the existing configurations using each input data item in turn
- On each iteration, it **reduces** the set of candidate configurations by removing any which cannot be ultimately extended to an optimal configuration

- An SDP algorithm is a recursive process which scans N input data items
 - $x_1, ..., x_N$ in sequence, generating new candidate configurations by **extending** the existing configurations using each input data item in turn
- On each iteration, it **reduces** the set of candidate configurations by removing any which cannot be ultimately extended to an optimal configuration
- Finally, it selects an optimal configuration from the remaining candidates

- An SDP algorithm is a **recursive** process which scans N input data items
 x₁, ..., x_N in sequence, generating new candidate configurations by **extending** the existing configurations using each input data item in turn
- On each iteration, it **reduces** the set of candidate configurations by removing any which cannot be ultimately extended to an optimal configuration
- Finally, it selects an optimal configuration from the remaining candidates
- Can construct a "computational configuration graph", and special kinds of graphs arise due to the type of exact SDP algorithm: brute-force (full tree), greedy (tree with a single optimal branch at each stage), dynamic programming (incomplete tree)

- **Step 1**. *Initialization*: Start with n = 0, generate the "root" configuration(s) in the set of candidate configurations, S.
- **Step 2**. Extension: Set n = n + 1, and using input data item x_n , extend all candidate configurations in S, and append these to S.
- **Step 3**. *Reduction*: Remove any candidate configurations from *S* which cannot be extended to an optimal configuration.
- **Step 4**. *Iteration*: if n < N, go back to Step 2.
- **Step 5**. *Select best*: Select an optimal configuration *X** from the remaining candidate configurations in *S*.

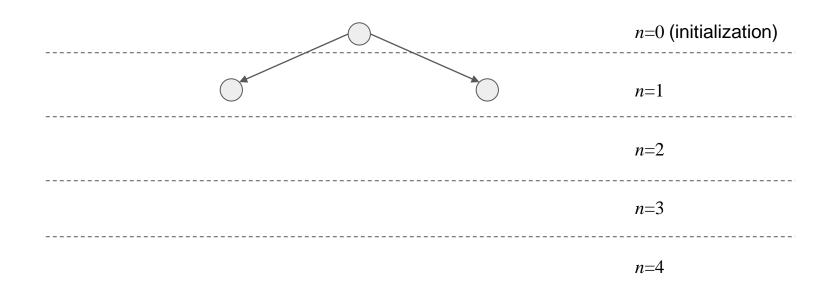
• **Step 1**. *Initialization*: Start with n = 0, generate the "root" configuration(s) in the set of candidate configurations, S.

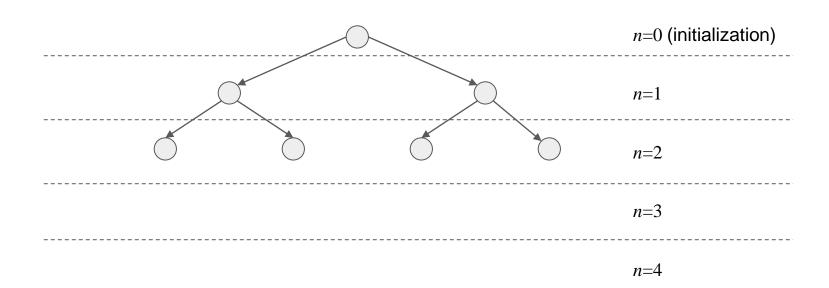
- **Step 1**. *Initialization*: Start with n = 0, generate the "root" configuration(s) in the set of candidate configurations, S.
- **Step 2**. Extension: Set n = n + 1, and using input data item x_n , extend all candidate configurations in S, and append these to S.

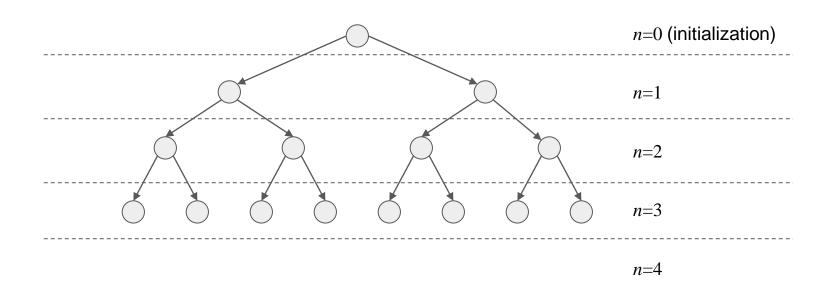
- **Step 1**. *Initialization*: Start with n = 0, generate the "root" configuration(s) in the set of candidate configurations, S.
- **Step 2**. Extension: Set n = n + 1, and using input data item x_n , extend all candidate configurations in S, and append these to S.
- **Step 3**. *Reduction*: Remove any candidate configurations from *S* which cannot be extended to an optimal configuration.

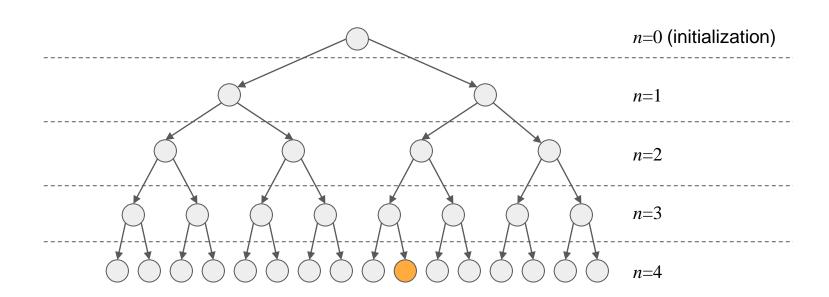
- **Step 1**. *Initialization*: Start with n = 0, generate the "root" configuration(s) in the set of candidate configurations, S.
- **Step 2**. Extension: Set n = n + 1, and using input data item x_n , extend all candidate configurations in S, and append these to S.
- **Step 3**. *Reduction*: Remove any candidate configurations from *S* which cannot be extended to an optimal configuration.
- **Step 4**. *Iteration*: if n < N, go back to Step 2.
- **Step 5**. *Select best*: Select an optimal configuration *X** from the remaining candidate configurations in *S*.

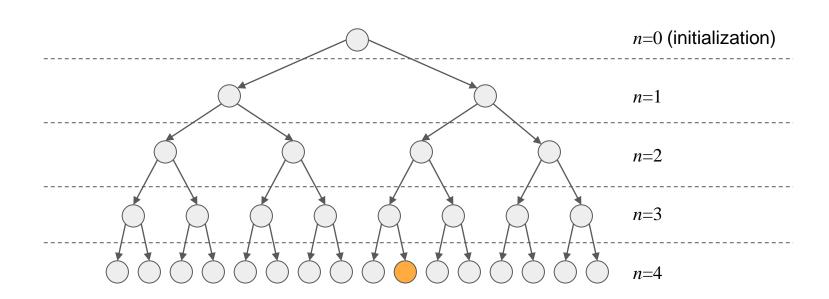
	n=0 (initialization)









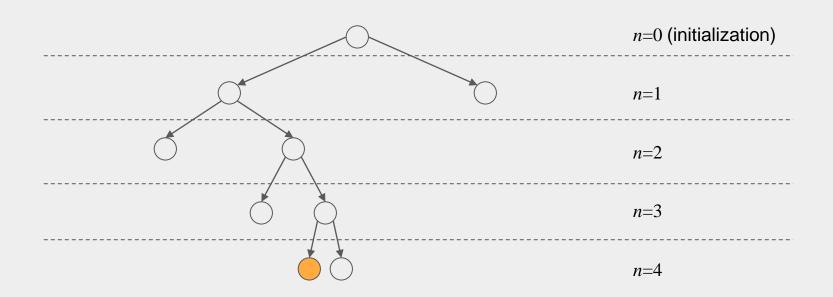


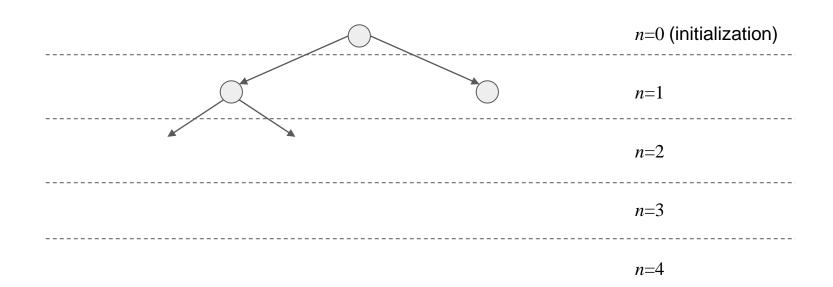
 From every exact configuration at stage n < N, there is exactly one possible optimal configuration obtained by extending the current exact configuration

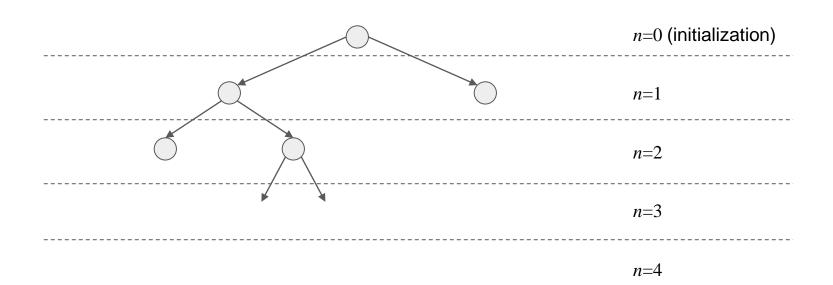
- From every exact configuration at stage n < N, there is exactly one possible optimal configuration obtained by extending the current exact configuration
- Generate all possible extensions from stage n to stage n + 1, compute objective function, retain only the best one

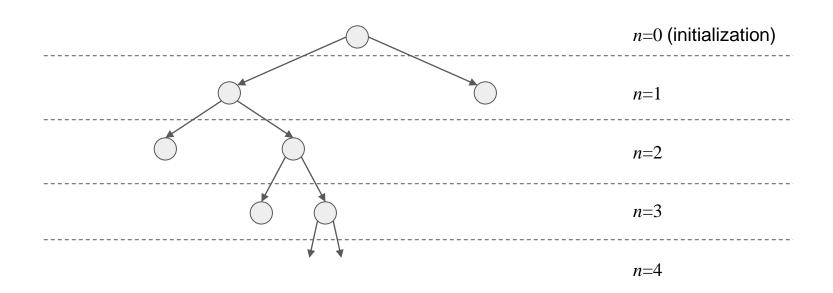
- From every exact configuration at stage n < N, there is exactly one possible optimal configuration obtained by extending the current exact configuration
- Generate all possible extensions from stage n to stage n + 1,
 compute objective function, retain only the best one
- Efficiency: avoid computing all configurations before selecting the best one and discarding sup-optimal solutions, only a small number of extensions at each stage; only one solution remaining (no selection required)

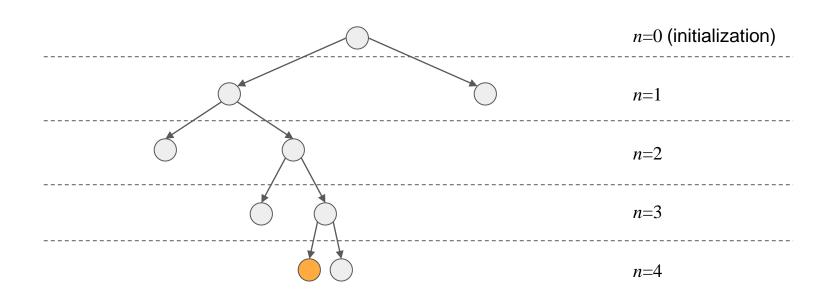
- From every exact configuration at stage n < N, there is exactly one possible optimal configuration obtained by extending the current exact configuration
- Generate all possible extensions from stage n to stage n + 1,
 compute objective function, retain only the best one
- Efficiency: avoid computing all configurations before selecting the best one and discarding sup-optimal solutions, only a small number of extensions at each stage; only one solution remaining (no selection required)
- **Complexity**: typically O(Nk)

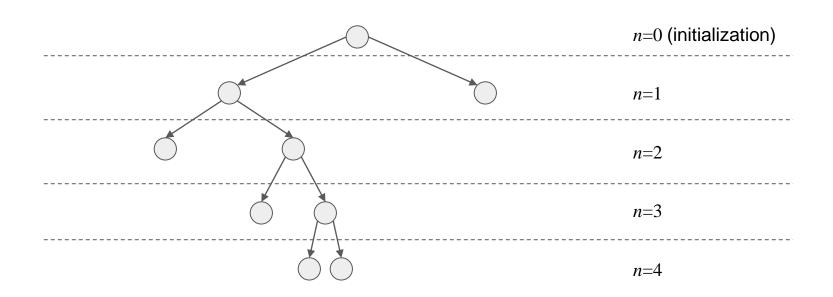






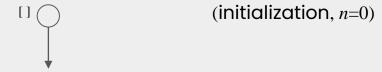


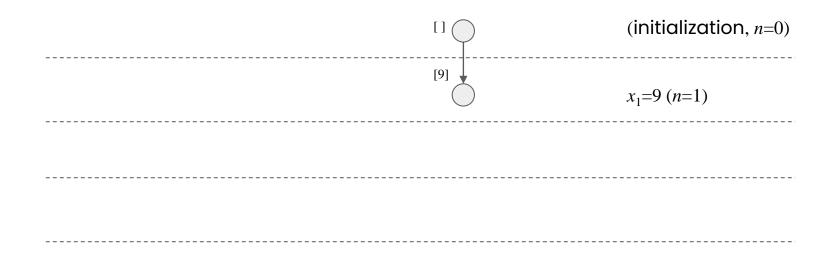


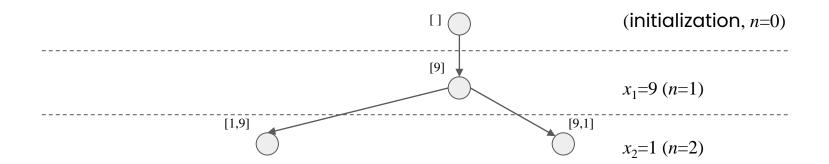


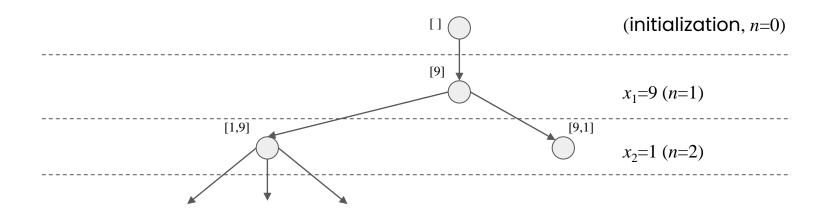
Insertion sort

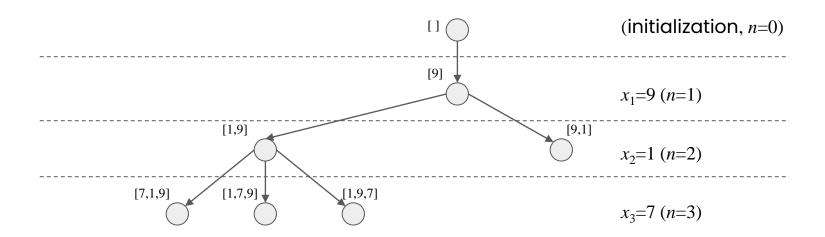
- A sorting algorithm, like selection sort
- Let's say you want to sort an array A with four elements
- \bullet A = [9, 1, 7, 3]

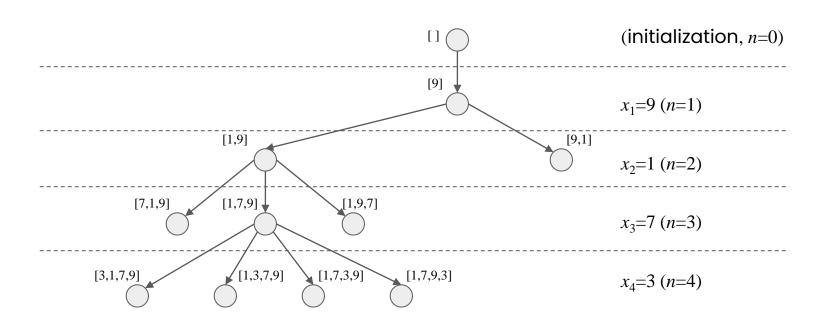


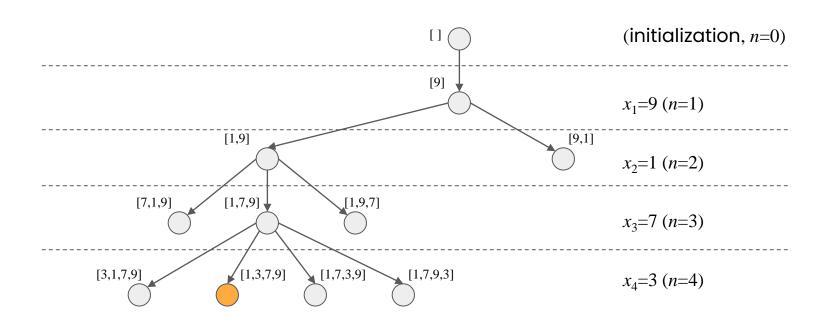












References and further reading

- MLSP, Section 2.6
- CLRS, Chapter 21