Data structures, Algorithms & Databases

Stacks, Queues and Heaps

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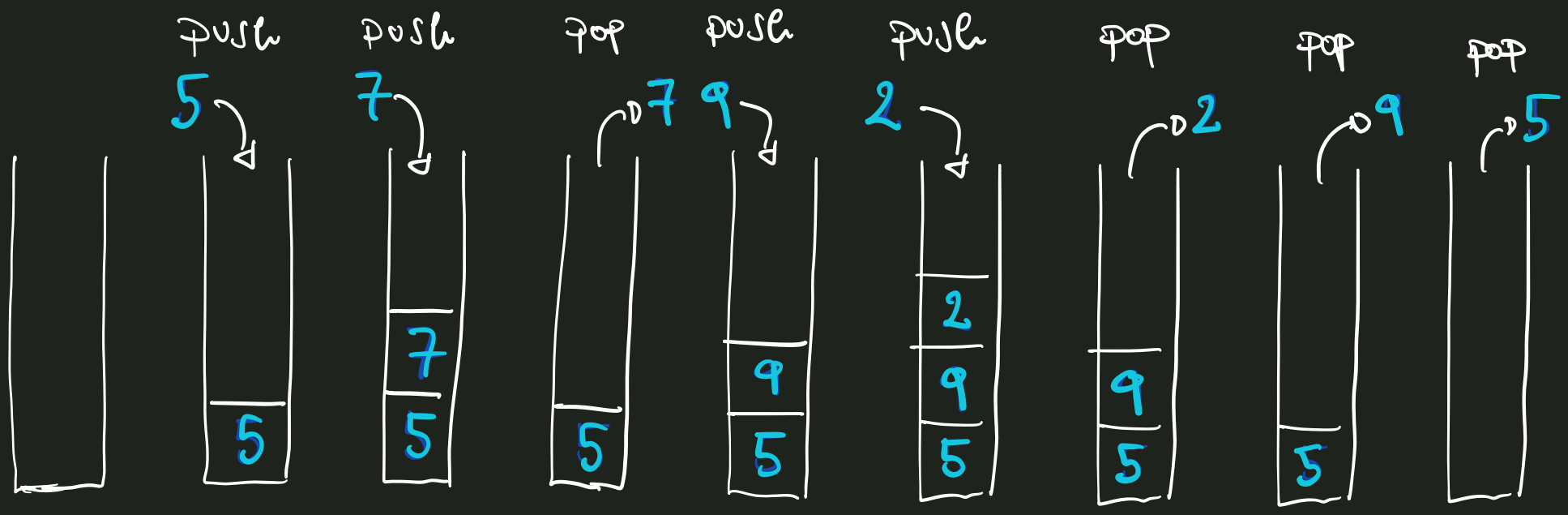
Stacks

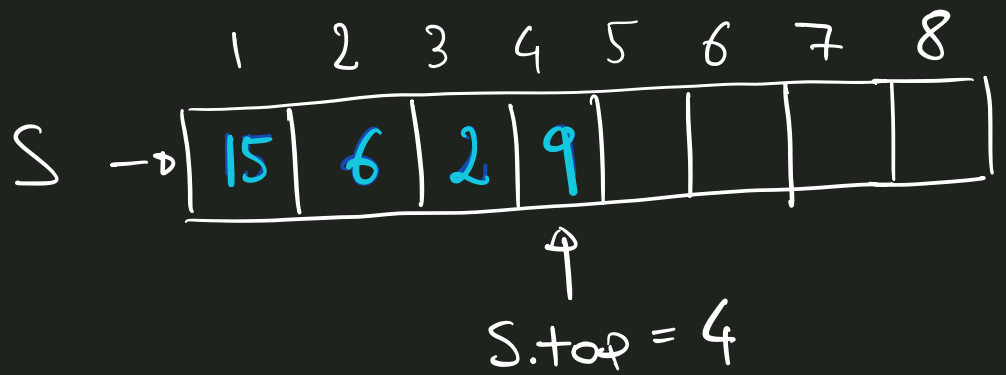


LIFO : Last in, first out

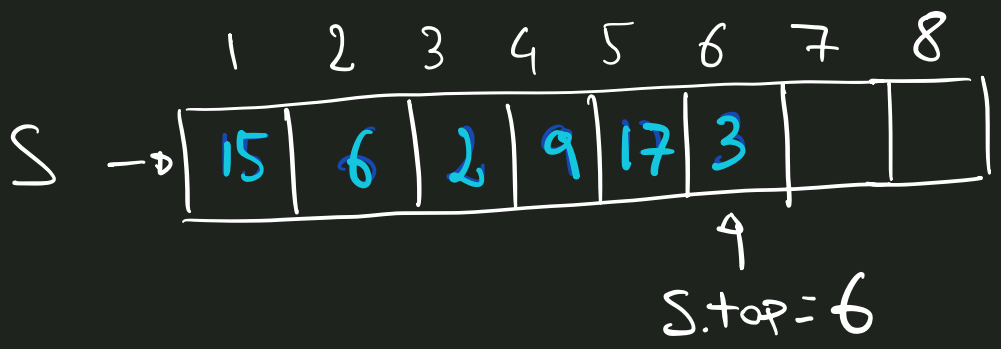
Stock operations

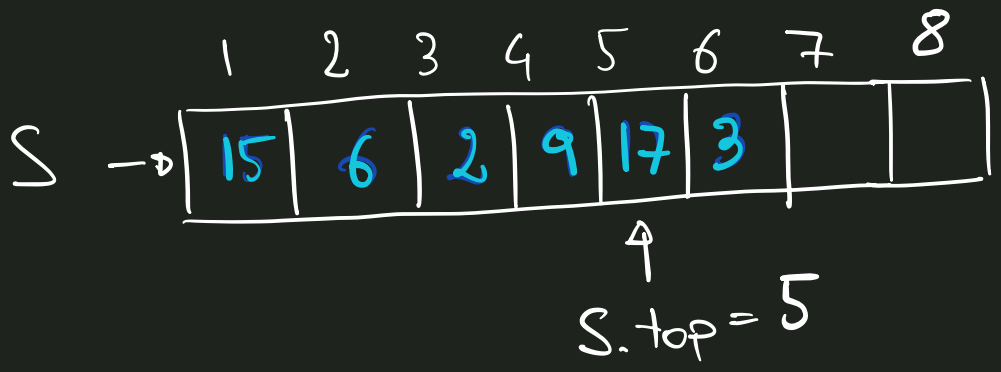
- Stack-Empty(S): returns true when empty
- Push(S, x): insert element
- Pop(S): remove and return last element

Stack Implementation Using Array



S.Size = 8

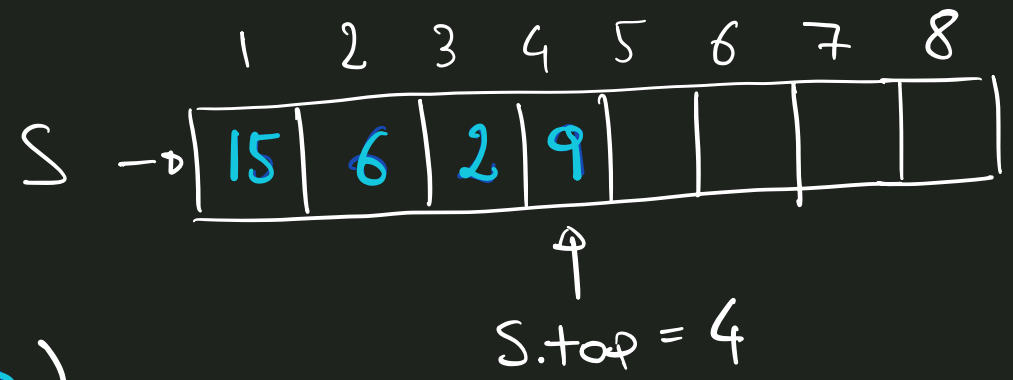




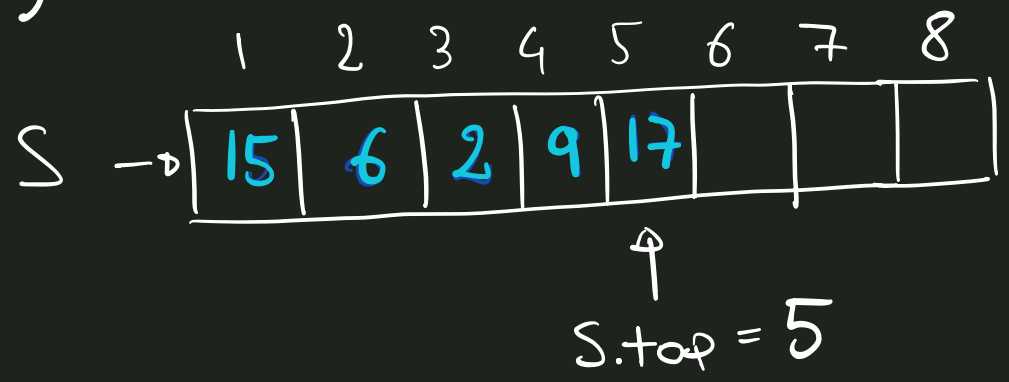
```
Push (S, x)
  if S.top == S.size
     error "overflow"
  else
     S.top = S.top + 1
     S[S.top] = x
```

Complexity ?

$O(1)$



Push (S, 17)

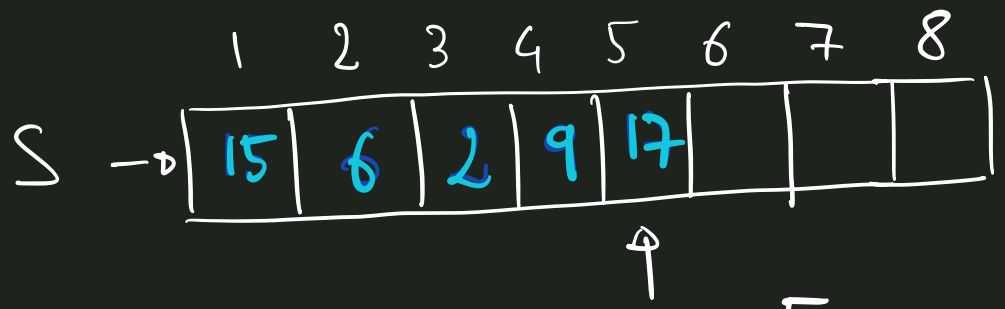


Pop (S):

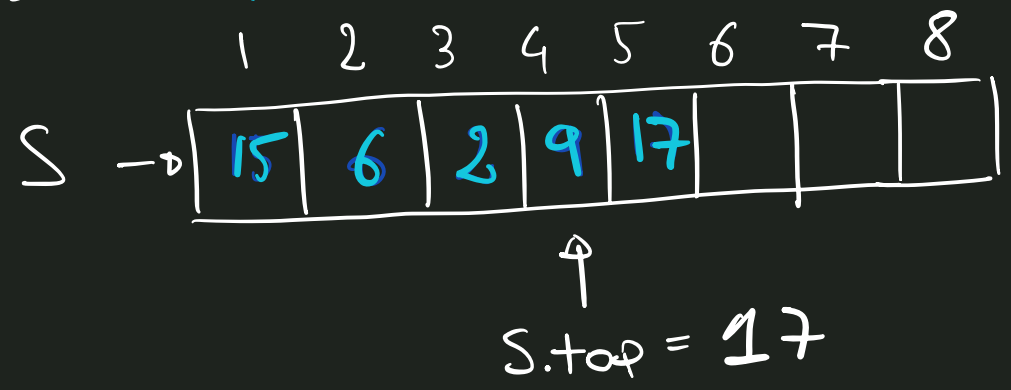
Complexity ?

```
if S.top == 0
    error "underflow"
else
    S.top = S.top - 1
    return S[S.top + 1]
```

$O(1)$

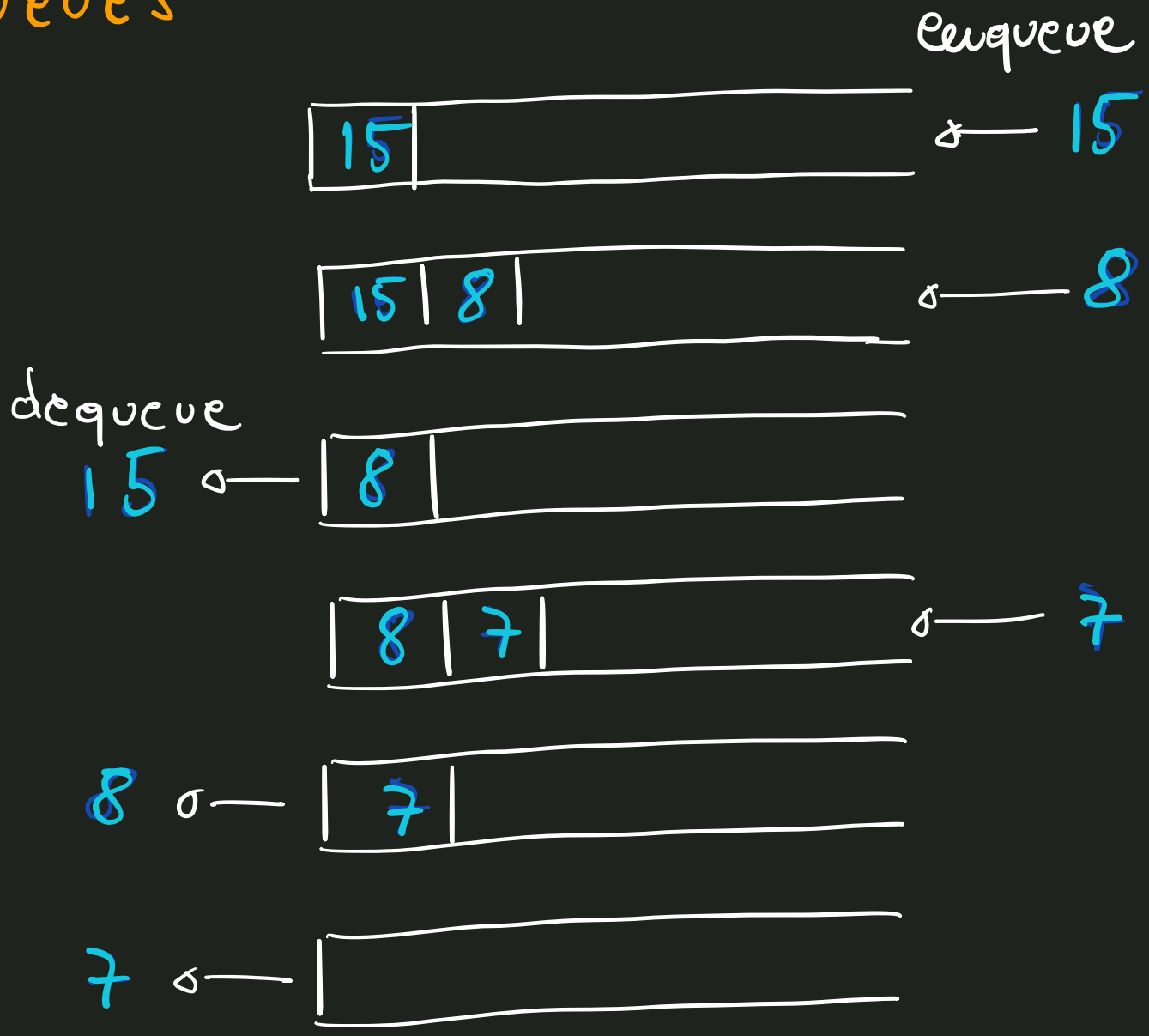


Pop(S) → 17



push	pop	stack-empty
$O(1)$	$O(1)$	$O(1)$

Queues



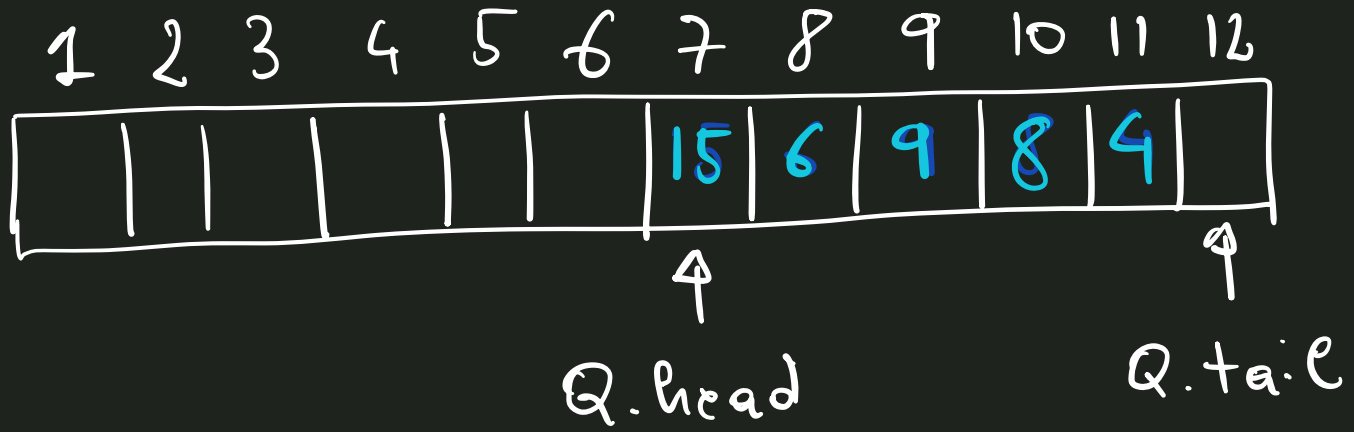
FIFO: first in, first out

Queues Operations

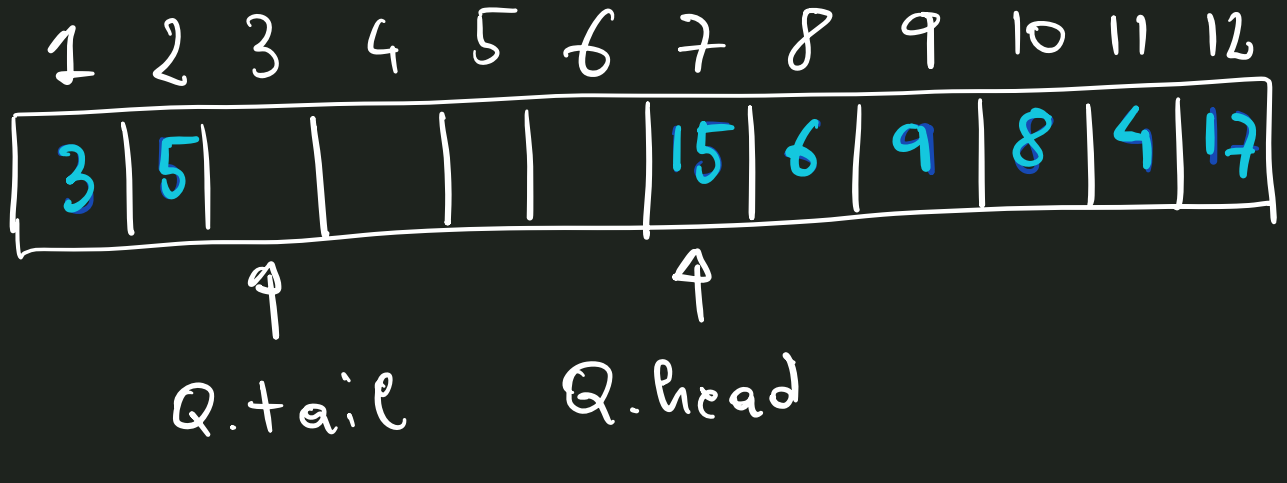
- Enqueue(Q, x) : insert x at the tail of the queue Q

- Dequeue(Q) : remove and return the head of the queue Q

Implementation using Arrays

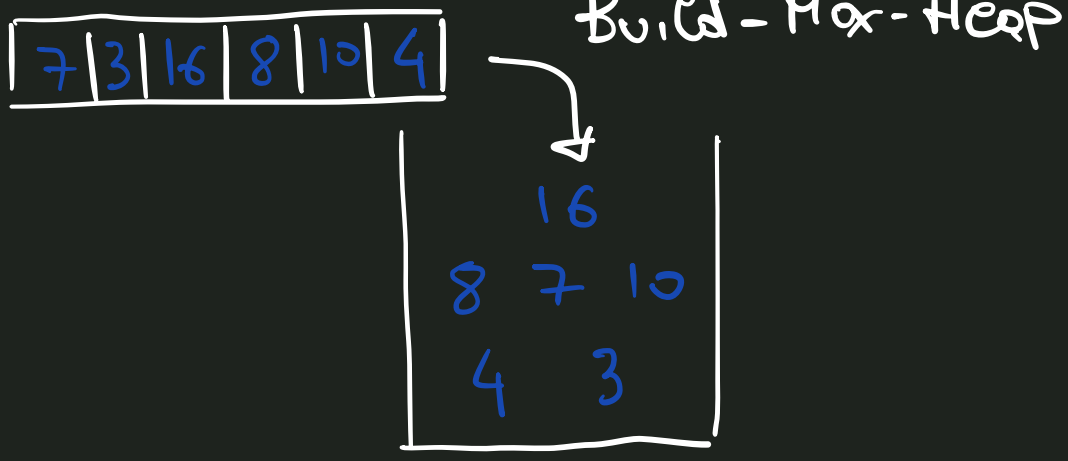


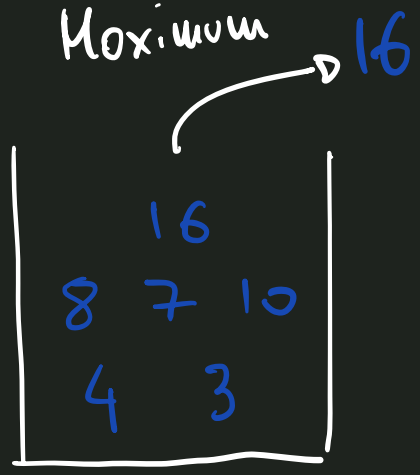
Q.size = 12

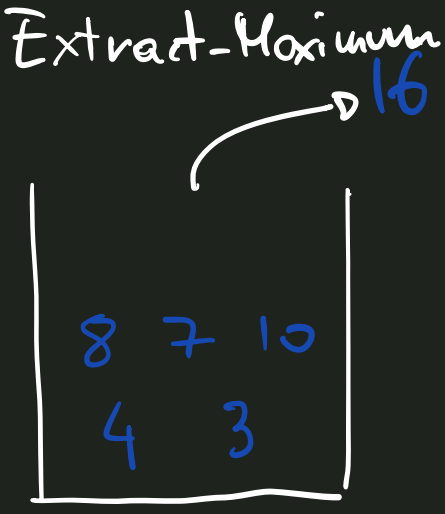


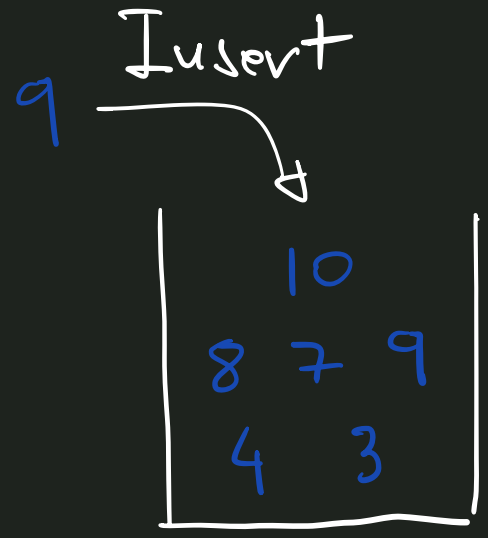
Enqueue	Dequeue
O(1)	O(1)

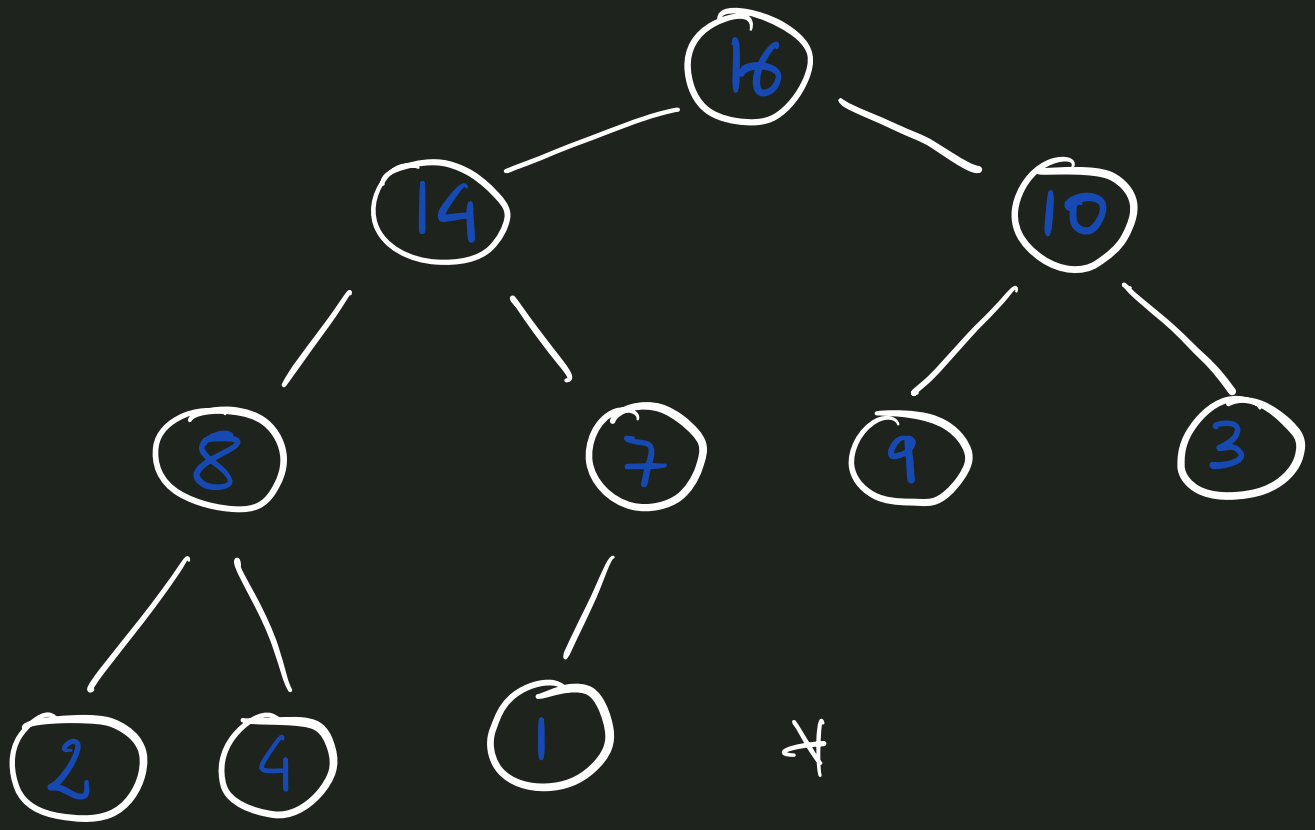
Max-Heaps


7 3 16 8 10 4
Build-Max-Heap
16
8 7 10
4 3


Maximum
16
16
8 7 10
4 3


Extract-Maximum
16
8 7 10
4 3


9
Insert
10
8 7 9
4 3



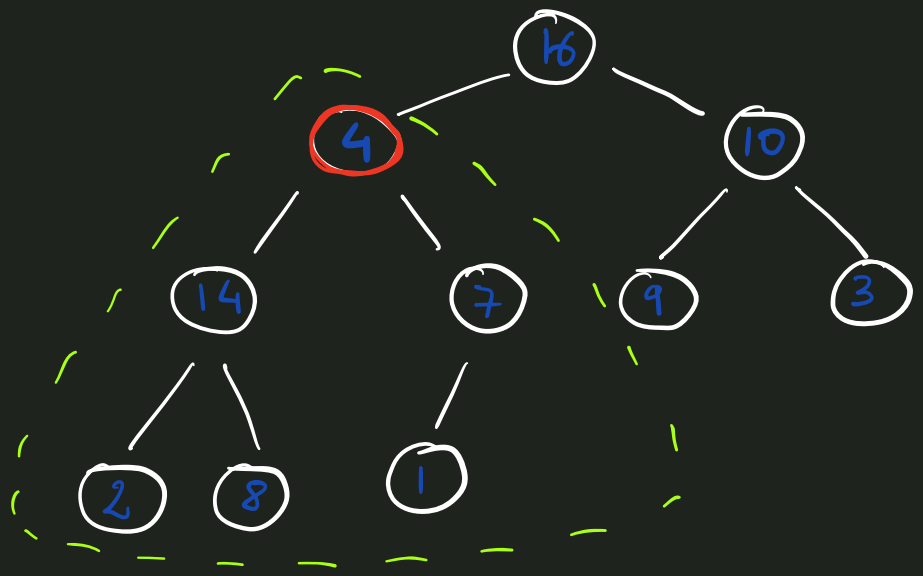
Max-Heap Property:

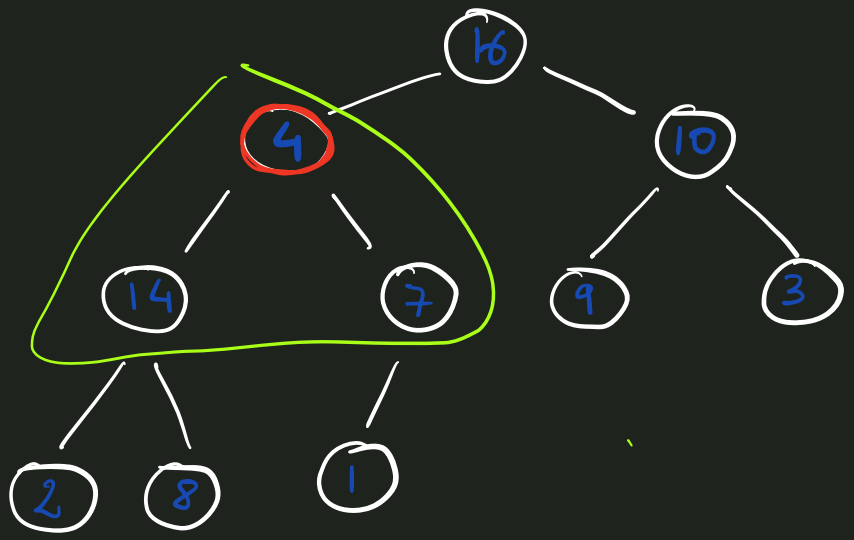
value of parent ≥ value of child

Max-Heapify Procedure

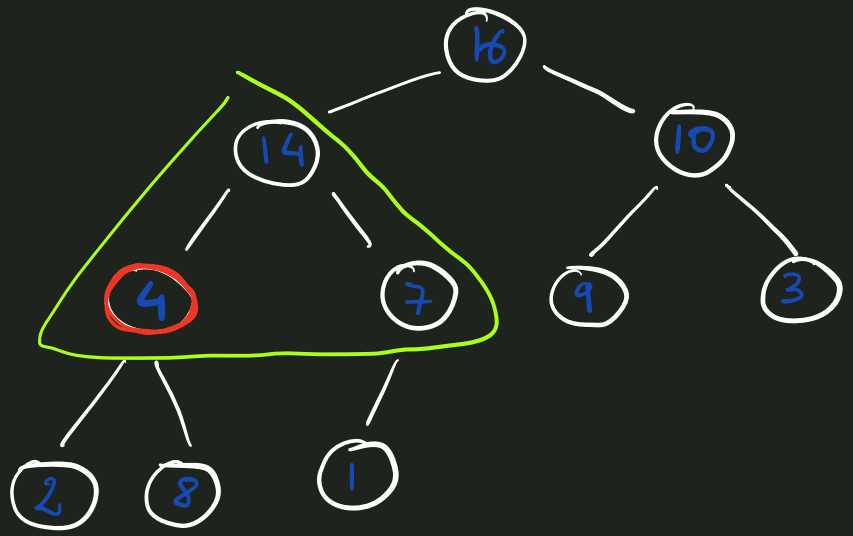
- Enforces the max-heap property of a node N

- Assumes that the children of N are valid max-heaps

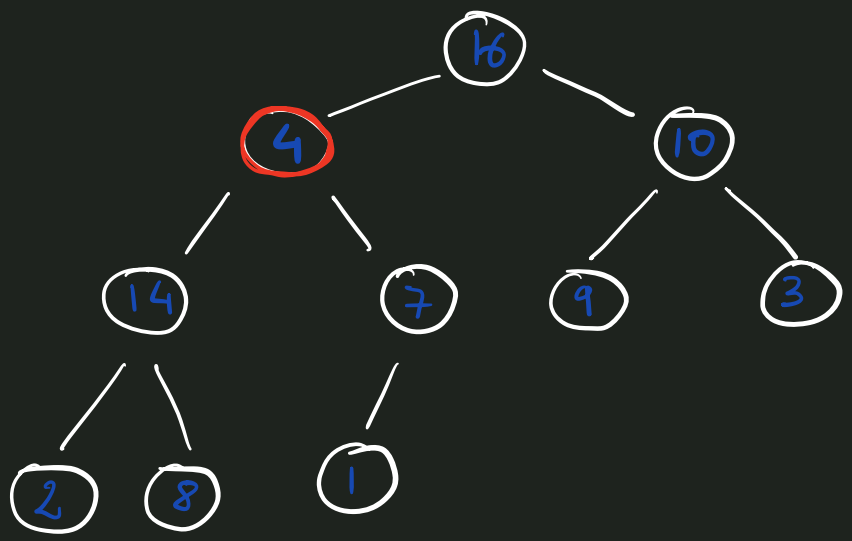

16
4
10
14
7
9
3
2
8
1

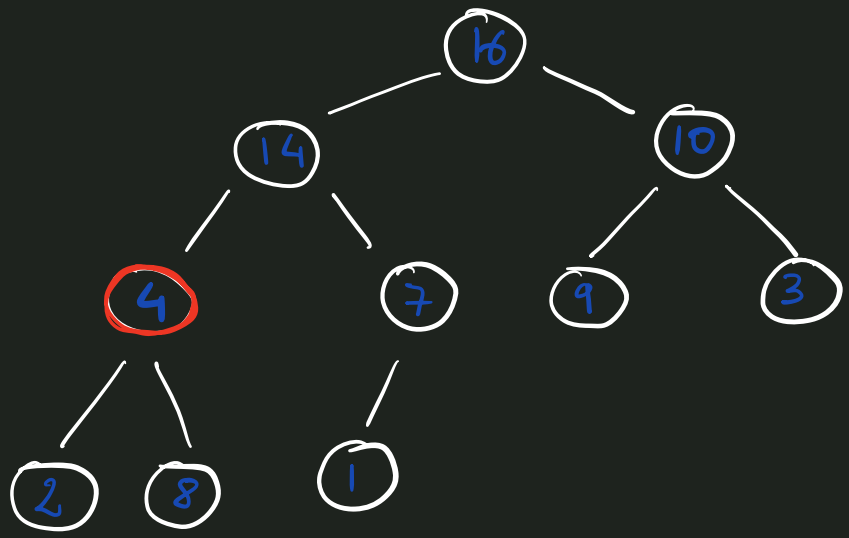


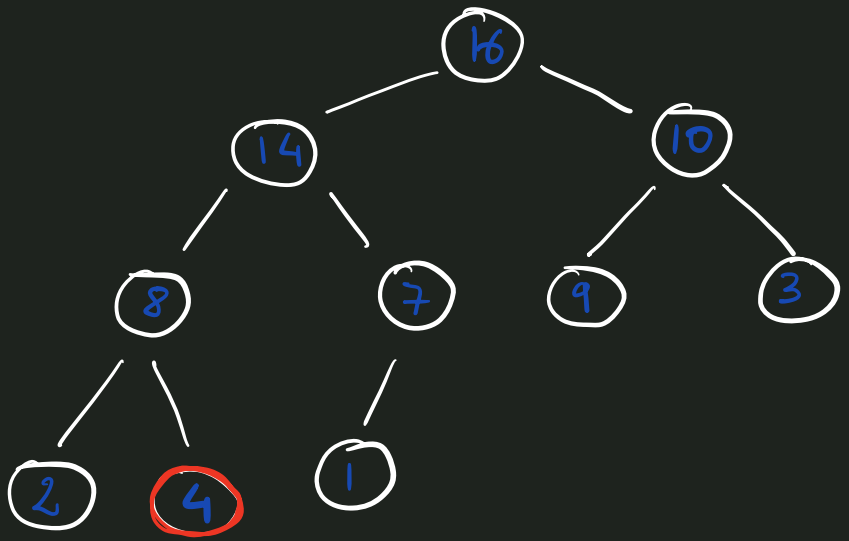
pick the maximum



swap


16
4
10
14
7
9
3
2
8
1


16
14
10
4
7
9
3
2
8
1


16
14
10
8
7
9
3
2
4
1

Time Complexity of Max-Heapify

Let h be the height of the tree

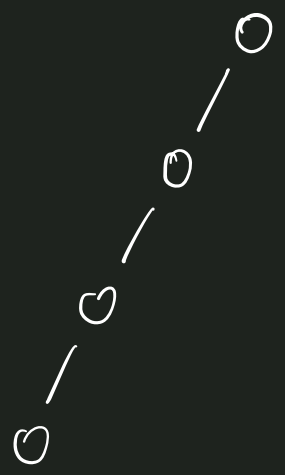
- How many swaps in the worst case, as a function of h?

$O(h)$

How many nodes does a binary tree of height h contain?

How many nodes does a binary tree of height h contain?

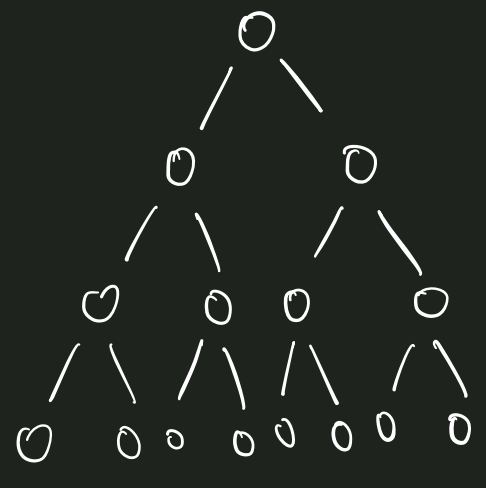
h	$\#n$
1	1
2	2
3	3
4	4



at least h nodes

How many nodes does a binary tree of height h contain?

h	$\#n$
1	1
2	3
3	7
4	15



at most $2^h - 1$ nodes

number of nodes

$h \leq n \leq 2^h$

height

when completely skewed

when perfectly balanced

Logarithm in base 2

Definition: $\log_2 n$ is the number k such that $2^k = n$

n	1	2	4	8	16	...	131072	262144
$\log_2 n$	0	1	2	3	4		17	18

$\log 250000$ → somewhere in between

height

=

worse case time complexity of Max-Heapify

$$\log n \leq h \leq n$$

when perfectly balanced

when completely skewed

Max-Heapify
$O(\log n)$