Artificial Intelligence and Machine Learning (AIML)





- Last lecture: combinatorial optimization in Al
- This lecture: exact SDP methods for combinatorial optimization

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- We will concentrate on **sequential decision process (SDP)** methods as they encompass many practical AI algorithms

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- Can construct a "computational configuration graph", and special kinds of graphs arise due to the type of exact SDP algorithm: brute-force (full tree), greedy (tree with a single optimal branch at each stage), dynamic programming (incomplete tree)

- **Step 1**. *Initialization*: Start with n = 0, generate the "root" configuration(s) in the set of candidate configurations, S.
- **Step 2**. Extension: Set n = n + 1, and using input data item x_n , extend all candidate configurations in S, and append these to S.
- **Step 3**. *Reduction*: Remove any candidate configurations from *S* which cannot be extended to an optimal configuration.
- **Step 4**. *Iteration*: if n < N, go back to Step 2.
- **Step 5**. *Select best*: Select an optimal configuration *X** from the remaining candidate configurations in *S*.

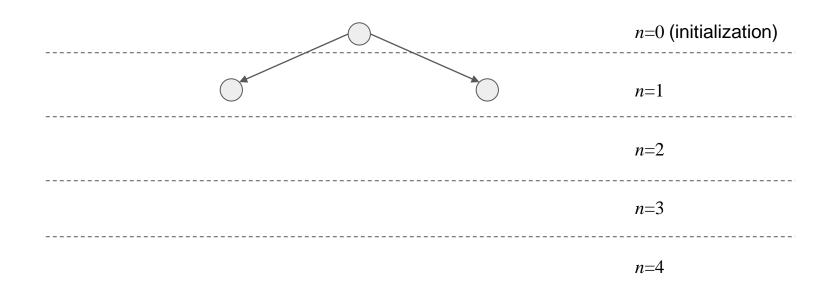
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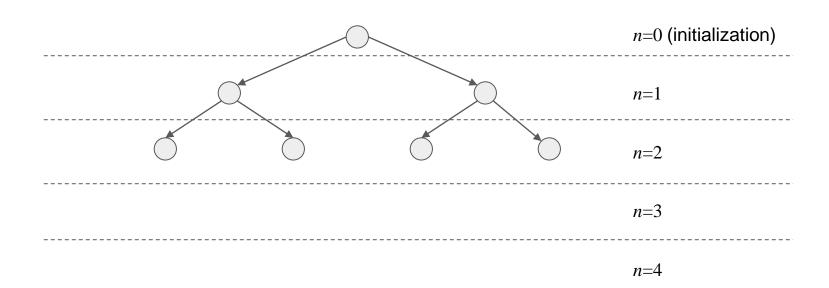
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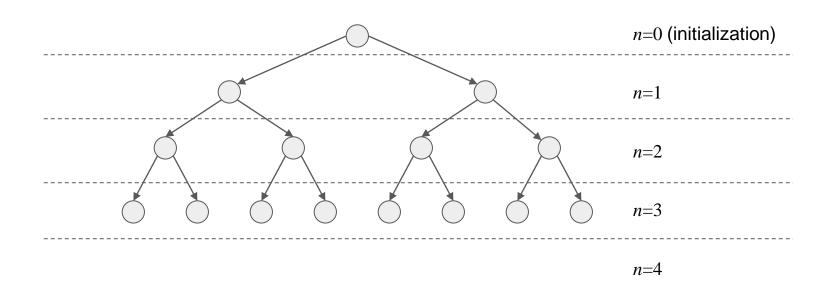
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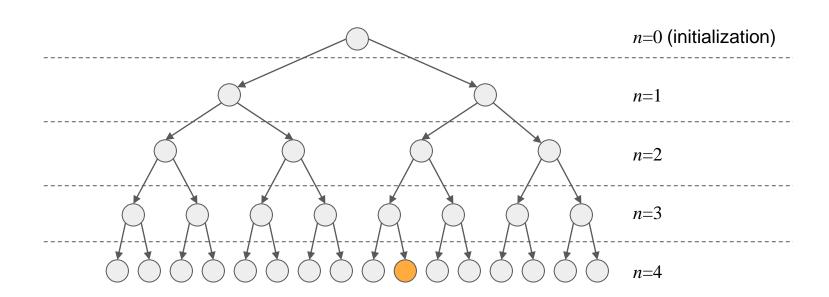
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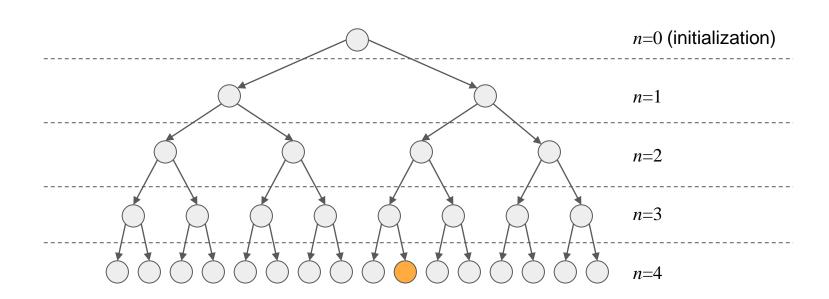
	n=0 (initialization)









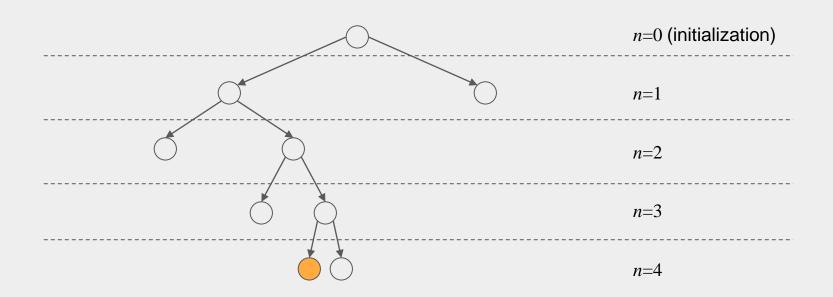


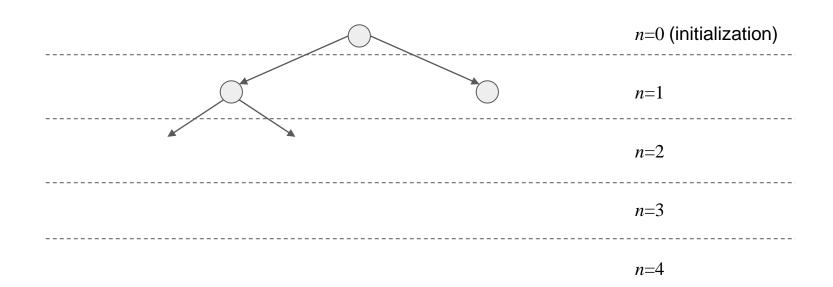
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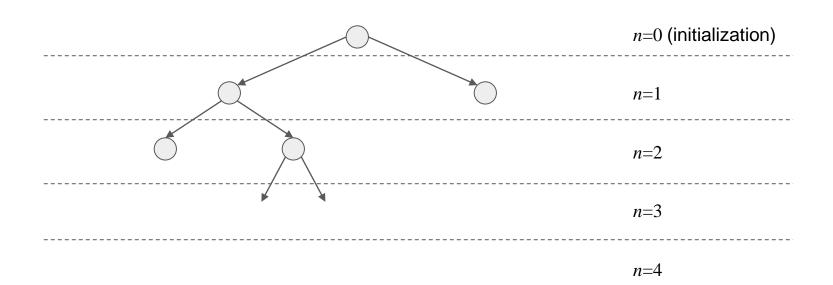
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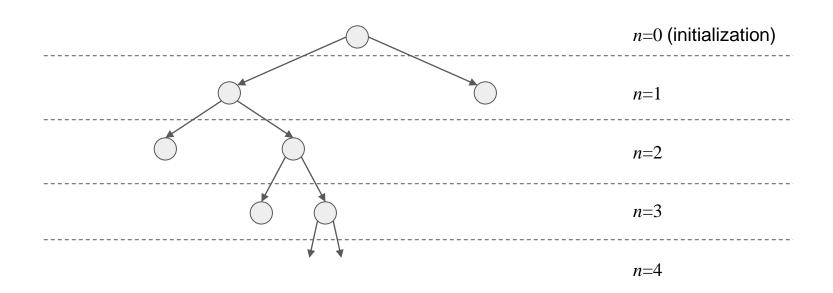
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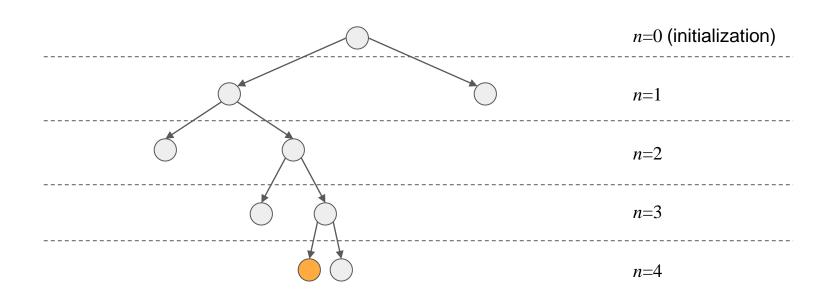
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- **Complexity**: typically O(Nk)

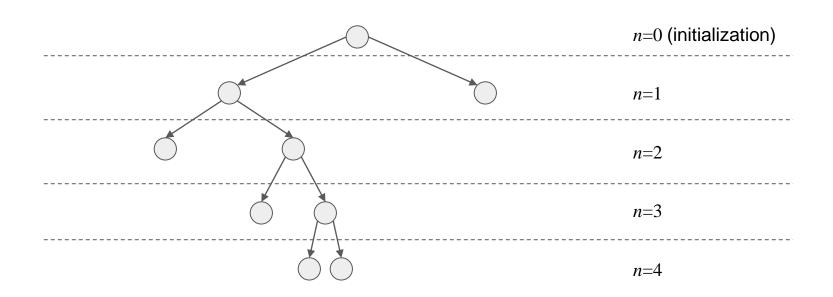






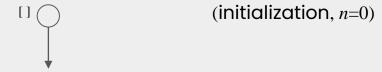


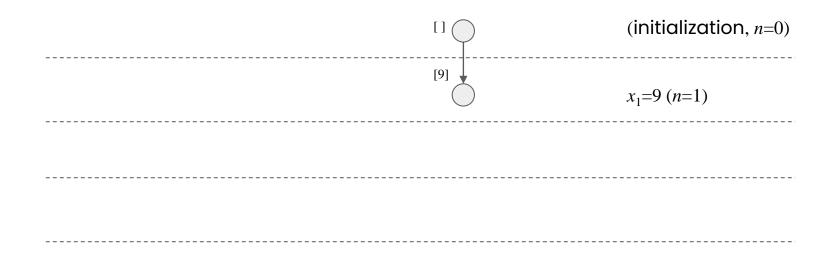


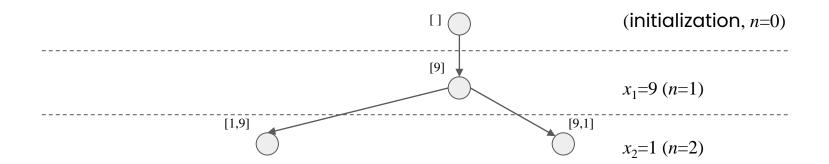


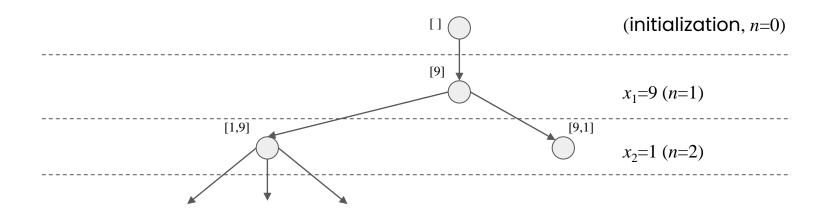
Insertion sort

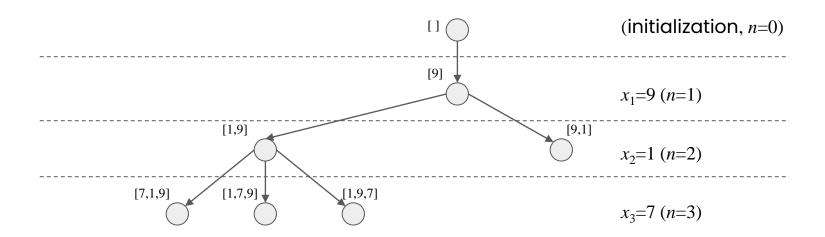
- A sorting algorithm, like selection sort
- Let's say you want to sort an array A with four elements
- \bullet A = [9, 1, 7, 3]

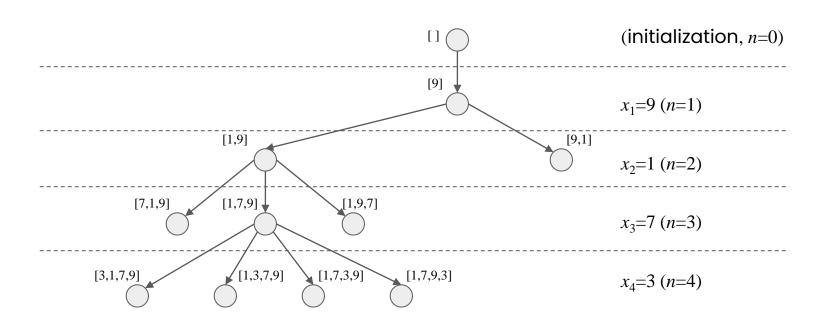


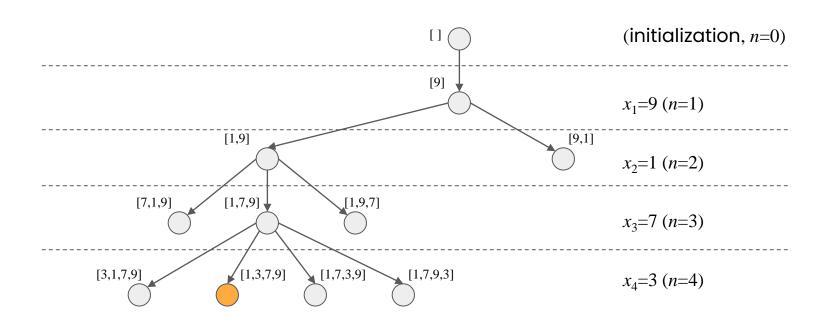












References and further reading

- MLSP, Section 2.6
- CLRS, Chapter 21