01-D1

绪论

复杂度分析：级数

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谁校对时间，谁就会突然老去。

算法分析

❖ 两个主要任务 = 正确性（不变性 x 单调性） + 复杂度

❖ 为确定后者，真地需要将算法描述为RAM的基本指令，再统计累计的执行次数？不必！

❖ C++等高级语言的**基本指令**，均等效于常数条RAM的**基本指令**；在渐进意义下，二者大体相当

- 分支转向：goto //算法的灵魂；出于结构化考虑，被隐藏了
- 迭代循环：for()、while()、... //本质上就是“if + goto”
- 调用 + 递归（自我调用） //本质上也是goto

❖ 复杂度分析的主要方法

- 迭代：级数求和
- 递归：递归跟踪 + 递推方程
- 猜测 + 验证

级数

- ❖ **算术级数：与末项平方同阶** $T(n) = 1 + 2 + \ldots + n = \binom{n+1}{2} = \frac{n(n+1)}{2} = \mathcal{O}(n^2)$

- ❖ **幂方级数：比幂次高出一阶** $\sum_{k=0}^{n} k^d \approx \int_0^n x^d dx = \left.\frac{x^{d+1}}{d+1}\right|_0^n = \frac{n^{d+1}}{d+1} = \mathcal{O}(n^{d+1})$

$$T_2(n) = \sum_{k=1}^{n} k^2 = 1^2 + 2^2 + 3^2 + \ldots + n^2 = n(n+1)(2n+1)/6 = \mathcal{O}(n^3)$$

$$T_3(n) = \sum_{k=1}^{n} k^3 = 1^3 + 2^3 + 3^3 + \ldots + n^3 = n^2(n+1)^2/4 = \mathcal{O}(n^4)$$

$$T_4(n) = \sum_{k=1}^{n} k^4 = 1^4 + 2^4 + 3^4 + \ldots + n^4 = n(n+1)(2n+1)(3n^2+3n-1)/30 = \mathcal{O}(n^5)$$

- ❖ **几何级数：与末项同阶**

$$T_a(n) = \sum_{k=0}^{n} a^k = a^0 + a^1 + a^2 + a^3 + \ldots + a^n = \frac{a^{n+1}-1}{a-1} = \mathcal{O}(a^n), \quad 1 < a$$

$$T_2(n) = \sum_{k=0}^{n} 2^k = 1 + 2 + 4 + 8 + \ldots + 2^n = 2^{n+1} - 1 = \mathcal{O}(2^{n+1}) = \mathcal{O}(2^n)$$

收敛级数

❖ $$\sum_{k=2}^{n} \frac{1}{(k-1)\cdot k} = \frac{1}{1\cdot 2} + \frac{1}{2\cdot 3} + \frac{1}{3\cdot 4} + \ldots + \frac{1}{(n-1)\cdot n} = 1 - \frac{1}{n} = \mathcal{O}(1)$$

$$\sum_{k=1}^{n} \frac{1}{k^2} = 1 + \frac{1}{2^2} + \frac{1}{3^2} + \ldots + \frac{1}{n^2} < 1 + \frac{1}{2^2} + \frac{1}{3^2} + \ldots = \frac{\pi^2}{6} = \mathcal{O}(1)$$

$$\sum_{k\ is\ a\ perfect\ power} \frac{1}{k-1} = \frac{1}{3} + \frac{1}{7} + \frac{1}{8} + \frac{1}{15} + \frac{1}{24} + \frac{1}{26} + \frac{1}{31} + \frac{1}{35} + \ldots = 1 = \mathcal{O}(1)$$

❖ **几何分布：** $(1-\lambda)\cdot[1+2\lambda+3\lambda^2+4\lambda^3+\ldots] = 1/(1-\lambda) = \mathcal{O}(1), \quad 0<\lambda<1$

❖ **有必要讨论这类级数吗？**

难道，基本操作次数、存储单元数可能是分数？

是的，某种意义上！

不收敛，但有限

❖ **调和级数：** $h(n) = \sum_{k=1}^{n} \frac{1}{k} = 1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \ldots + \frac{1}{n} = \ln n + \gamma + \mathcal{O}(\frac{1}{2n}) = \Theta(\log n)$

❖ **对数级数：** $\sum_{k=1}^{n} \ln k = \ln \prod_{k=1}^{n} k = \ln n! \approx (n + 0.5) \cdot \ln n - n = \Theta(n \cdot \log n)$

❖ **对数 + 线性 + 指数：** $\sum_{k=1}^{n} k \cdot \log k \approx \int_{1}^{n} x \ln x dx = \left. \frac{x^2 \cdot (2 \cdot \ln x - 1)}{4} \right|_{1}^{n} = \mathcal{O}(n^2 \log n)$

$$\sum_{k=1}^{n} k \cdot 2^k = \sum_{k=1}^{n} k \cdot 2^{k+1} - \sum_{k=1}^{n} k \cdot 2^k = \sum_{k=1}^{n+1} (k-1) \cdot 2^k - \sum_{k=1}^{n} k \cdot 2^k$$

$$= n \cdot 2^{n+1} - \sum_{k=1}^{n} 2^k = n \cdot 2^{n+1} - (2^{n+1} - 2) = (n-1) \cdot 2^{n+1} + 2 = \mathcal{O}(n \cdot 2^n)$$

❖ **如有兴趣，不妨读读：**Concrete Mathematics　　//ex-2.35, Goldbach Theorem