(a)

There will be only 4 scenarios that the index of the median of the 3 points is in the middle range of $n/4 \le i \le 3n/4$: All 3 points within [n/4, 3n/4]; 2 points within [n/4, 3n/4] and 1 point less than n/4; 2 points within [n/4, 3n/4] and 1 point greater than 3n/4; 1 point within [n/4, 3n/4], 1 point less than n/4, and 1 point greater than 3n/4.

The probabilities of each Scenario are stated as follow:

$$\Pr_1 = \left(\frac{1}{2}\right)^3 = \frac{1}{8} \tag{1}$$

$$\Pr_2 = C_3^2 \times \left(\frac{1}{2}\right)^2 \times \frac{1}{4} = \frac{3}{16} \tag{2}$$

$$\Pr_3 = C_3^2 \times \left(\frac{1}{2}\right)^2 \times \frac{1}{4} = \frac{3}{16} \tag{3}$$

$$Pr_4 = A_3^3 \times \frac{1}{2} \times \frac{1}{4} \times \frac{1}{4} = \frac{3}{16} \tag{4}$$

Hence, adding them up, we get the probability:

$$\Pr = \frac{1}{8} + \frac{3}{16} + \frac{3}{16} + \frac{3}{16} = \frac{11}{16}$$
 (5)

(b)

The algorithm is stated as follow:

Algorithm 1 Algorithm for problem 1

Guarded select (S, k)

If |S| < 4

Sort & find kth

Else i = 0

While $(n/4 \le i \le 3n/4 \text{ not true})$

Select $x \in S$ as stated in the problem

Find S-, S+; i = |S-|+1

Endwhile

If i = k: Return x

Elif i > k: Guarded select (S-, k)

Else: Guarded select (S+, k-i)

From lecture, we proved that if the probability of success in one interation is p, then the expected number of iterations is 1/p. Hence in this case, it takes 1/(11/16) = 16/11 tries to find an element

in the desired range. We need 3 comparisons when finding the median, and using such a median to compare to other (n-3) numbers will cost (n-3) comparisons. Hence the total number of comparison for a round will be 3+n-3=n, and the expected number of comparisons for an element will be 16n/11. Hence, according to what we have discussed during the class, the recurrence will be:

$$T(n) \le \frac{16}{11}n + T\left(\frac{3n}{4}\right) \tag{6}$$

The expected number of comparisons is at most:

$$\frac{16n}{11} \times \frac{1}{1 - \frac{3}{4}} = \frac{64}{11} \mathbf{n} \approx 5.818 \mathbf{n}$$
 (7)

From Equation (7) we find that the expected number of comparisons made by this method is upper-bounded by cn for a c < 6. (c = 5.818)