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1.

$A ::= [B, A] \mid B.$

$B ::= C \mid (A; C).$

$C ::= \{C\} \mid D.$

$D ::= a \mid b \mid c.$

a).  $[C, (b; a)]$

$A \rightarrow [B, A].$

$\rightarrow [C, B].$

$\rightarrow [D, (A; C)].$

$\rightarrow [C, (B; D)].$

$\rightarrow [C, (D; D)].$

$\rightarrow [C, (b; a)].$

A can generate.

b).  $([C, b], C; a)$

$A \rightarrow B$

$\rightarrow (A; C).$

$\rightarrow ([C, A], C).$

$\rightarrow ([C, b], C; a).$

$\rightarrow ([C, b], C; a).$

$\rightarrow ([C, b], C; a).$

$\rightarrow ([C, b], C; a).$

A, B can generate.

c).  $[C, ([C, b], C; a)]$

$A \rightarrow [B, A].$

$\rightarrow [C, ([C, b], C; a)].$

$= [C, ([C, b], C; a)].$

since there is no way for C to derive  $[a, C]$ . Thus, there is no way we can derive this based on given grammar [none].

2.

a).

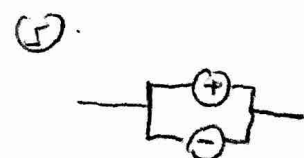
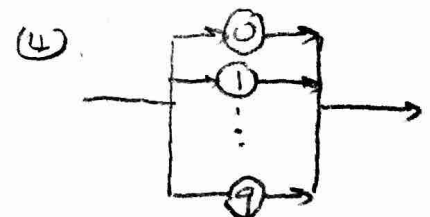
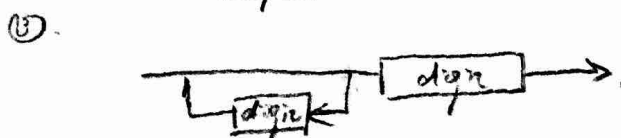
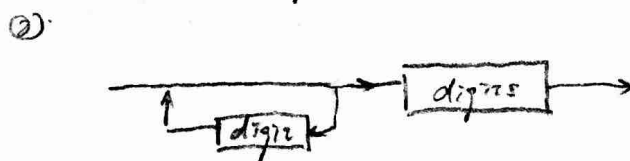
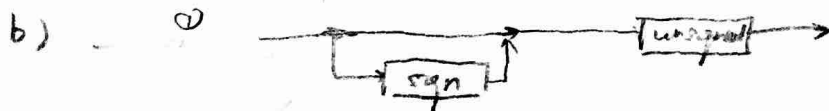
$\langle \text{integer} \rangle ::= [\langle \text{sign} \rangle] \langle \text{unsigned} \rangle$

$\langle \text{unsigned} \rangle ::= \{ \langle \text{digits} \rangle \} \langle \text{digits} \rangle$

$\langle \text{digits} \rangle ::= \{ \langle \text{digit} \rangle \} \langle \text{digit} \rangle$

$\langle \text{digit} \rangle ::= 0 \mid 1 \mid \dots \mid 9.$

$\langle \text{sign} \rangle ::= + \mid -$



3. What language is generated by each of BNF grammars below.

a).

$\langle S \rangle ::= 0 \langle S \rangle 1111 \mid \text{empty}$

$$L = \{ 0^k (1)^{4k} \mid k \geq 0 \}$$

b)

$\langle S \rangle ::= \langle x \rangle \mid \langle y \rangle \mid \text{empty}$

$\langle x \rangle ::= 1 \langle x \rangle 00 \mid \text{empty}$

$\langle y \rangle ::= 0 \langle y \rangle 1 \mid \text{empty}$

$$L_x = \{ 1^k 0^{2k} \mid k \geq 0 \}$$

$$L_y = \{ 0^k 1^k \mid k \geq 0 \}$$

$$L = \{ L_x \cup L_y \}$$

c)

$\langle S \rangle ::= \langle x \rangle \mid \langle y \rangle$

$\langle x \rangle ::= 0 \langle x \rangle 1 \mid \langle x1 \rangle$

$\langle x1 \rangle ::= 0 \langle x1 \rangle \mid 0$

$\langle y \rangle ::= 0 \langle y \rangle 11 \mid \langle y1 \rangle$

$\langle y1 \rangle ::= \langle y1 \rangle 111$

$$L_{x1} = \{ 0^k \mid k \geq 1 \}$$

$$L_{y1} = \{ 1^k \mid k \geq 1 \}$$

$$L_x = \{ 0^x 1^y \text{ or } 0^k \mid k \geq 1 \& x \geq 1 \}$$

$$L_y = \{ 0^x 1^{2k} \text{ or } 1^k \mid k \geq 1 \& x \geq 1 \}$$

$$L = \{ L_x \cup L_y \}$$

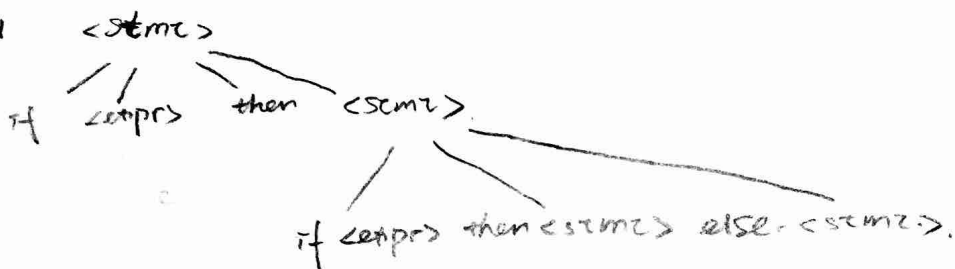
4.

Given the following grammar.

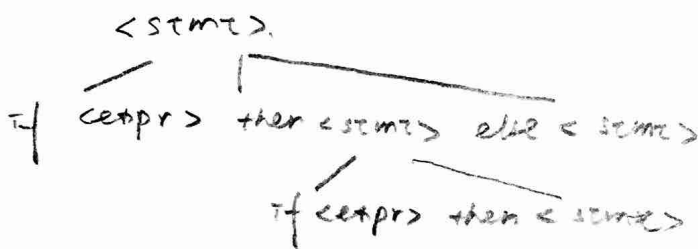
$\langle \text{stmt} \rangle ::= \text{if } \langle \text{expr} \rangle \text{ then } \langle \text{stmt} \rangle$   
 $\quad \quad \quad | \text{if } \langle \text{expr} \rangle \text{ then } \langle \text{stmt} \rangle \text{ else } \langle \text{stmt} \rangle.$   
 $\quad \quad \quad | \text{other}$   
 $\langle \text{expr} \rangle ::= \text{true} \mid \text{false}.$

a).

Tree 1



Tree 2

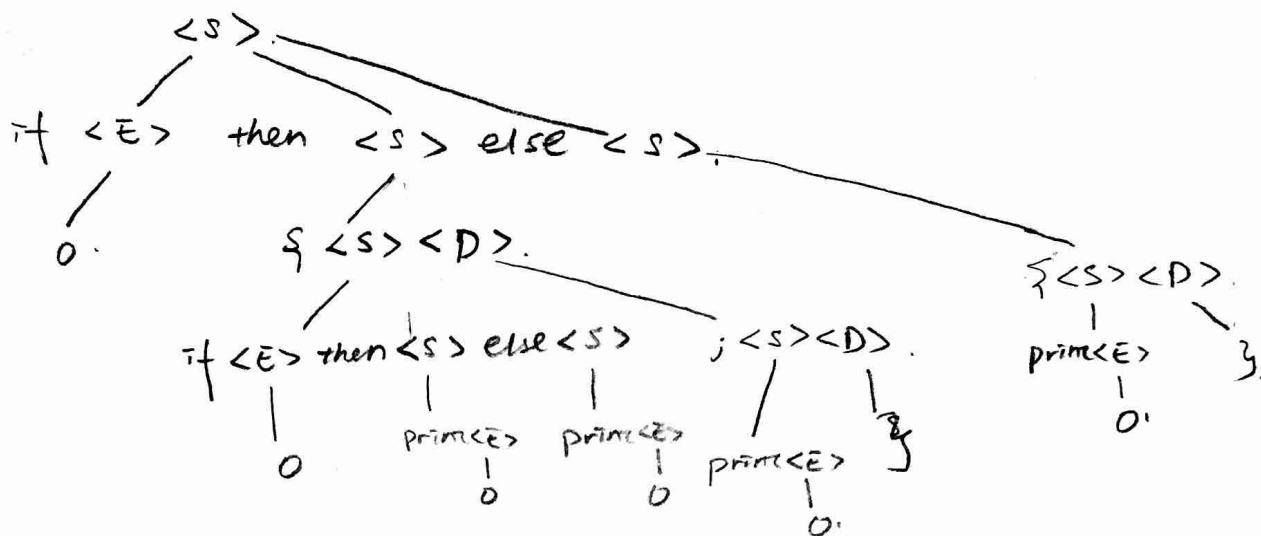


string: if <expr> then if <expr> then <stmt> else <stmt>

b).

$\langle \text{stmt} \rangle ::= \text{if } \langle \text{expr} \rangle \text{ then } \langle \text{NT} \rangle \mid \langle \text{NT} \rangle \mid \text{other}.$   
 $\langle \text{NT} \rangle ::= \text{if } \langle \text{expr} \rangle \text{ then } \langle \text{NT} \rangle \text{ else } \langle \text{stmt} \rangle \mid \text{other}.$   
 $\langle \text{expr} \rangle ::= \text{true} \mid \text{false}.$

5.



6.

a. Algorithm:

Base Case : if the nonterminal can be derive into any terminal.  
then remove the nonterminal symbol from the.  
nonterminal array. Return the nonterminal array.

Recursion: for (each nonterminal in the nonterminal array) {  
if (it can be derive into terminal) {  
Remove itself from the nonterminal array

} else if (it can only be derive into nonterminals (doesn't contain itself))  
pass that nonterminals into the function again

} else if (it can only be derive into nonterminals  
that contain itself) {  
keep reading & don't recurse.

b. Read the first line " $\langle S \rangle ::= 0 \mid \langle A \rangle \mid \langle C \rangle$ " since  $\langle S \rangle$  can be  
derive into terminal 0. Remove  $\langle S \rangle$  from nonterminal array.

The second line " $\langle A \rangle ::= \langle A \rangle \langle B \rangle$ ". Since  $\langle A \rangle$  can only be derive into nonterminals  
 $\langle A \rangle \langle B \rangle$ , which contains  $\langle A \rangle$  itself, so it stays in the nonterminal array and we move on.

After we read the 3rd and 4th line, we see that both  $\langle B \rangle$  and  
 $\langle C \rangle$  can be removed from nonterminal array.

Finally we return the array nonterminal, which now only contains  
the useless nonterminal  $\langle A \rangle$