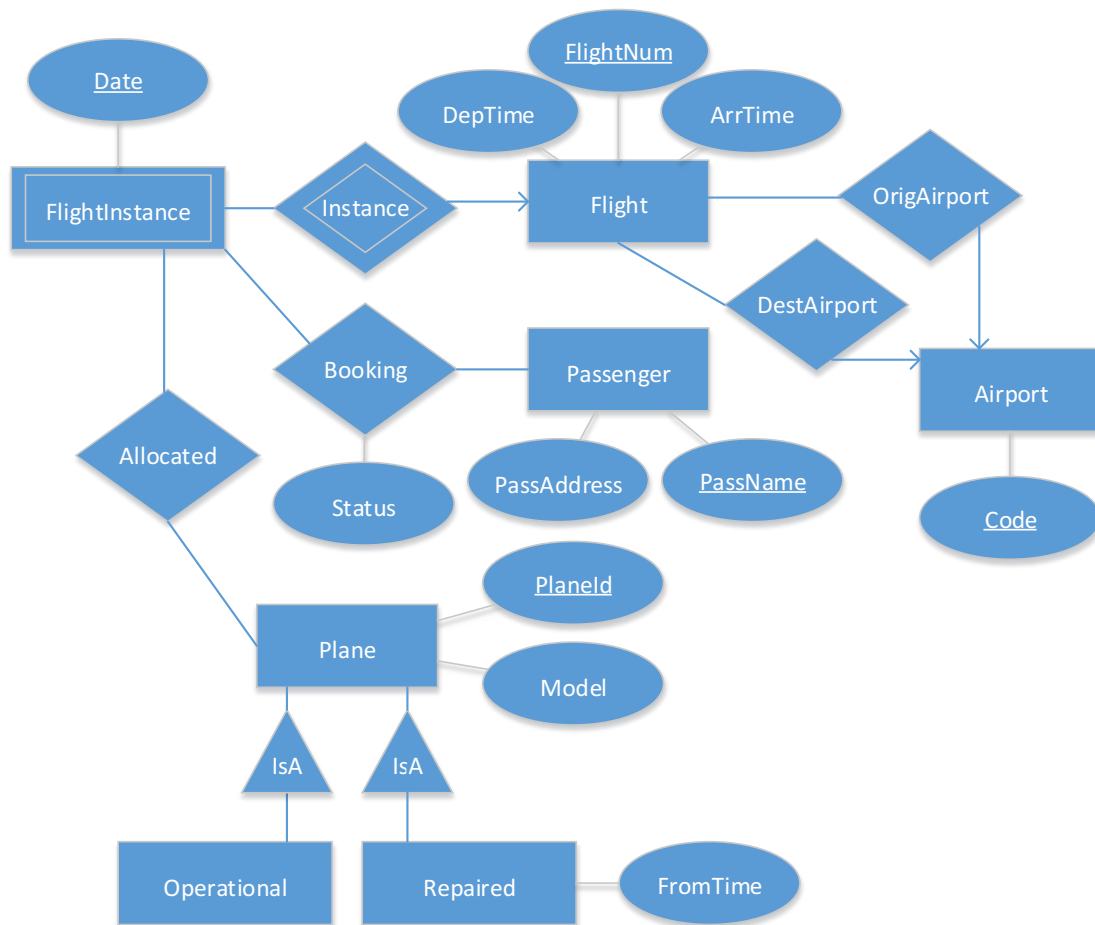


Problem 1



Problem 2

According to the textbook, the definition of “Satisfying the BCNF rule” is:

The left side of every non-trivial FD must contain a key.

1.

This schema is not in BCNF.

We rename Course as “C”, rename Teacher as “T”, and rename Hour as “H”, rename Room as “R”, rename Student as “S”. The original relation is {C T H R S}

From $HT \rightarrow R$, $HR \rightarrow C$, we get $HST \rightarrow HSCTR$, so {HST} is the key of this schema. Consider the FD $C \rightarrow T$, the lhs of the FD is not a superkey (It does not contain the key) of this schema, so this schema is not in BCNF.

2.

Lossless join decomposition into BCNF:

For {C, T, H, R, S}

The key is {H,S,T}

$C \rightarrow T$ violates the BCNF rule,

The $R_1 = C$ deviation = {C, T};

The $R_2 = C \text{ UNION } (R - R_1) = C \text{ UNION } (H, R, S)$.

For R_1 , FD is $C \rightarrow T$;

According to the textbook, the definition of “Satisfying the BCNF rule” is:

The left side of every non-trivial FD must contain a key.

The key of R1 is C.

R1 is in BCNF rule;

For R2,

Since FD $C \rightarrow T$ is in R1, so there is no T in the R2;

But it is lossless join, T is still existed in the other FDs, so if R1 is {C, T}

R2 is {C, H, R, S, T}.

The Key of R2 is HST.

FD $HR \rightarrow C$ and $HS \rightarrow R$ and $HT \rightarrow R$ is R2's FD.

From the FD $HR \rightarrow C$,

HR is not a superkey, so $HR \rightarrow C$ violates the BCNF.

Decompose for the $HR \rightarrow C$;

So, $R_{21} = \{H, R, C\}$, $R_{22} = \{H, T, S, R\}$

For R21, The FD is $HR \rightarrow C$

The key is HR, R1 satisfies the BCNF rule;

For $R_{22} = \{H, R, S\}$,

The key is HS, the FD is $HS \rightarrow R$ and $HT \rightarrow R$;

R2 satisfies the BCNF rule.

So, in the end, it is:

$R_1 = \{C, T\}$

$R_{21} = \{H, R, C\}$

$R_{22} = \{H, T, S, R\}$

With the FD $C \rightarrow T$, $H, R \rightarrow C$, $H, S \rightarrow R$, $H, T \rightarrow R$;

It preserves the given functional dependencies.

Problem 3

1.

No.

From $B \rightarrow A$, we know $A^+ = \{A\}$.

$B \notin A^+$, so $A \rightarrow B$ doesn't follow from $B \rightarrow A$.

2.

Yes.

From $C \rightarrow A$, $C^+ = \{C, A\}$; from $AC \rightarrow B$, $C^+ = \{C, A, B\}$.

$B \in C^+$, so $C \rightarrow B$ follow from $AC \rightarrow B$ and $C \rightarrow A$.

3.

No.

From $AB \rightarrow C$, we know $A^+ = \{A\}$.

$C \notin A^+$, so $A \rightarrow C$ doesn't follow from $AB \rightarrow C$.

4.

No.

From $A \rightarrow C$ and $A \rightarrow B$, we know A^+ is $\{A B C\}$

$D \notin A^+$, so $A \rightarrow D$ doesn't follow from $A \rightarrow C$ and $A \rightarrow B$.