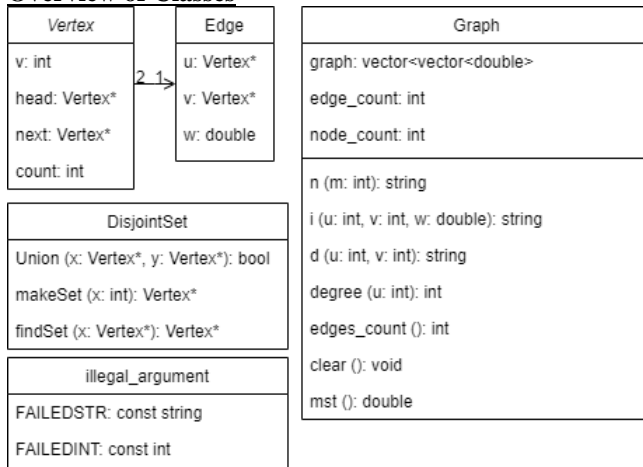


**ECE250 – Project 4**  
**MST Using Kruskal's Algorithm**  
**Design Document**  
 Qinying Wu, q227wu  
 Mar 30<sup>th</sup>, 2020

**1. Overview of Classes**



Class Name	Description	Member Variables and Functions
<b>Vertex</b>	Vertices objects in the graph	<b><i>v</i>: int</b> the integer representation of vertex <i>v</i> <b><i>head</i>: Vertex*</b> the vertex representing the disjoint set that the current vertex <i>v</i> belongs to <b><i>next</i>: Vertex*</b> the next vertex linked to the current vertex to represent a disjoint set <b><i>count</i>: int</b> if the current vertex is the representative vertex of the disjoint set, count stores the total number of vertex in the disjoint set, otherwise count=1
<b>Edge</b>	Vertices connections in the graph	<b><i>w</i>: double</b> the weight of the edge <b><i>u</i>: Vertex*</b> one of the vertices connected by the edge <b><i>v</i>: Vertex*</b> the other vertex of the edge
<b>DisjointSet</b>	Disjoint set objects to represent forest of nodes	<b><i>Union</i> (x: Vertex*, y: Vertex*): bool</b> performs a union operation on the vertex disjoint sets. Parameters x and y are vertices in the graph. Returns true if disjoint set of x is different than that of y after union. Returns false if x and y are in the same disjoint set <b><i>makeSet</i> (x: int): Vertex*</b> make a node in the graph into a vertex object. Parameter x is the integer representation of the node in the graph. Returns the newly created vertex pointer object <b><i>findSet</i> (x: Vertex*): Vertex*</b> find the disjoint set of a given vertex. Parameter x is the vertex object. Returns the head representative vertex of the disjoint set (linked list)
<b>Graph</b>	The graph object	<b><i>graph</i>: vector&lt;vector&lt;double&gt;&gt;</b> the adjacency matrix with entries as weight of the edge between two vertices in the graph <b><i>edge_count</i>: int</b> the total count of edges in the graph <b><i>node_count</i>: int</b> the total count of nodes in the graph <b><i>n</i> (m: int): string</b> initialization of the graph. Parameter m is the number of vertices in the graph. Returns success if successful in creating the graph. If m<0, throws an illegal_argument exception <b><i>i</i> (u: int, v: int, w: double): string</b> insert an edge between two vertices. Parameters u and v are integer representations of the vertex objects, w is the weight of the edge. Returns success if success in insertion or updating the weight of an existing edge. Throws an illegal_argument exception if u and/or v are outside the valid range of vertices <b><i>d</i> (u: int, v: int): string</b> deletes the edge between two vertices. Parameters u and v are integer representation of arbitrary vertices. Throws an illegal_argument exception if either u and/or v out of range. Returns “failure” if edge weight is 0 (edge does not exist). Returns “success” if the edge between vertices u and v is deleted successfully <b><i>degree</i> (u: int): int</b> returns the number of edges connected to a vertex if the vertex is in the graph. Throws an illegal_argument exception if the vertex is out of range of the graph. Parameter u is an integer representation of an arbitrary vertex. <b><i>edges_count</i> (): int</b> returns the total count of edges in the graph in string <b><i>clear</i> (): void</b> removes all the edge from the graph <b><i>mst</i> (): double</b> calculates the minimum spanning tree of the graph. Returns the minimum weight of the mst if the graph is connected, otherwise return -1 to indicate a not connected graph
<b>illegal_argument</b>	Exception class	<b><i>FAILEDSTR</i>: const string</b> indicative value for string function exception (“failure”) <b><i>FAILEDINT</i>: const int</b> indicative value for integer function exception (-1)

**2. Constructors and Destructor Decisions**

Class Name	Constructor	Destructor
<b>Vertex</b>	One integer parameter – vt is assigned to v to represent the current vertex. Member variables head and next are each assigned to nullptr. Member variable count is assigned to 1 since each vertex is in its own disjoint set upon initialization	Member variables head and next are each dereferenced to nullptr
<b>Edge</b>	Three parameters – vertex objects u and v are assigned to member variables u and v, respectively. Double w is assigned to the member variable w	Vertex u and v are dereferenced to nullptr
<b>DisjointSet</b> <b>illegal_argument</b>	Does not pass in any parameters, simply an initialization of objects	Does not pass in any parameters, simply destructs the object
<b>Graph</b>	Two integer parameters – e is assigned to edge_count, v is assigned to node_count	Clears the graph adjacency matrix

### 3. Asymptotic Upper Bounds

<b>Union</b>	$O(\min(\text{findSet}(x).\text{size}, \text{findSet}(y).\text{size}))$ first check the disjoint set size of each x and y, and add the smaller set to the larger set
<b>makeSet</b> <b>findSet</b> <b>edges_count</b> <b>i, d</b>	$O(1)$ makeSet – converts one integer representation of the vertex into a pointer vertex object and returns it findSet, edges_count – accessing the corresponding member variable in their respective ADTs is constant time i, d – updates the entry in the adjacency matrix by directly accessing the given matrix row and column indexes
<b>degree</b>	$O(V)$ traverse through the row of the corresponding square adjacency matrix index. The number of entries visited is equivalent to the number of vertices in the graph represented by V here
<b>n</b>	$O(V^2)$ V is the total count of vertices in the graph. The graph is initialized with V rows and V columns which is equivalent to the square of V
<b>clear</b>	$O(V^2)$ traverse through every entry of the adjacency matrix. Since it is a square matrix with same row and column counts equal to the number of vertices in the graph, the time complexity would thus be the square of the vertex count
<b>mst</b>	$O(V+E+E\log(E))$ Make each vertex in the graph into a disjoint set requires $V * T_{\text{makeSet}} = O(V)$ time. Make each edge in the graph into an Edge object requires $E * (O(1)) = O(E)$ time. Sort the edges in ascending order using <code>std::sort()</code> requires $O(E\log(E))$ time. Finally traverse through the sorted list of edge objects to find the mst is $E * T_{\text{Union}} = O(E)$ .

### 4. Test Cases

a) i u;v;w, d u;v, degree u <ul style="list-style-type: none"> <li>a. u=v</li> <li>b. u!=v <ul style="list-style-type: none"> <li>i. both u and v in range</li> <li>ii. either u or v, or both u and v out of range</li> </ul> </li> <li>c. <math>w \leq 0</math> and <math>w &gt; 0</math> (for i)</li> <li>d. after calling the clear function</li> </ul> b) edges_count <ul style="list-style-type: none"> <li>a. before/after inserting a valid/invalid edge</li> <li>b. before/after deleting a valid/invalid edge</li> <li>c. before/after calling the clear function</li> </ul> c) mst <ul style="list-style-type: none"> <li>a. disconnected graph</li> <li>b. connected graph</li> <li>c. after calling the clear function</li> </ul> d) Union (u,v) <ul style="list-style-type: none"> <li>a. <math>u.\text{DSet} = v.\text{DSet}</math></li> <li>b. <math>u.\text{DSet} \neq v.\text{DSet}</math></li> </ul> e) findSet (u) <ul style="list-style-type: none"> <li>a. u.DSet contains itself only</li> <li>b. u.DSet contains other nodes <ul style="list-style-type: none"> <li>i. u is the representative node of the set</li> <li>ii. u is not the representative node of the set</li> </ul> </li> </ul>	Example: (expected output) n 4 (success) d 0;2 (failure) d 1;1 (failure) degree 2 (failure) edge_count (edge count is 0) mst (not connected) clear (success) mst (not connected) i 3;3;2 (failure) i 91;2;33 (failure) i 2;3;1 (success) i 0;1;1 (success) i 0;2;2 (success) i 1;2;1 (success) i 1;3;44 (success) i 3;1;2 (success) degree 1 (degree of 1 is 3) degree 0 (degree of 0 is 2) degree 2 (degree of 2 is 3) degree 4 (failure) degree 3 (degree of 3 is 2) mst (3) edge_count (edge count is 5) clear (success) edge_count (edge count is 0)
--	---