Analysis of Algorithms - CS 323 Lecture #14 – May 18, 2016

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#### **Announcements**

The final exam will start at 6:30 p.m. on Wednesday, May 25th.

You can submit today's lecture notes for extra credit (2 points or 3 if it's really good) by 11 a.m. on Sunday, May  $22^{nd}$ .

## **Lecture Notes**

Countable  $\rightarrow$  there is a mapping of every element in a set to the natural numbers  $(\mathbb{N})$ 

Uncountable  $\rightarrow$  you will miss some numbers in the set ( $\mathbb{R}$ ), Cantor's Diagonalization Argument

Inclusive (up to and including)  $\rightarrow$  [0, 1)  $\leftarrow$  Exclusive (up to but not including)

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\begin{array}{c} 0 \cdot \boldsymbol{b_{1,1}} \ b_{1,2} \ b_{1,3} \ b_{1,4} \dots \\ 0 \cdot b_{2,1} \ \boldsymbol{b_{2,2}} \ b_{2,3} \ b_{2,4} \dots \end{array}
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 $0.b_{3,1}b_{3,2}b_{3,3}b_{3,4}...$ 

Look at the diagonal values and take the complement

$$\overline{b}_{\mathrm{i,i}} = 1 - \overline{b}_{\mathrm{i,i}}$$

 $0 \rightarrow 1$ 

 $1 \rightarrow 0$ 

$$r^* = 0 . \overline{b}_{1,1} \overline{b}_{2,2} \overline{b}_{3,3} \overline{b}_{4,4} ...$$

$$f(x) \rightarrow y$$

 $x \in \mathbb{N}$ 

 $y \in \mathbb{N}$ 

X	f <sub>0</sub> (x)	f <sub>1</sub> (x)	f <sub>2</sub> (x)	f <sub>3</sub> (x)	 f*(x)
0	$f_0(0)$	$f_1(0)$	$f_2(0)$	f <sub>3</sub> (0)	$f_0(0) + 1$
1	f <sub>0</sub> (1)	f <sub>1</sub> (1)	f <sub>2</sub> (1)	f <sub>3</sub> (1)	$f_1(1) + 1$
2	f <sub>0</sub> (2)	f <sub>1</sub> (2)	f <sub>2</sub> (2)	f <sub>3</sub> (2)	$f_2(2) + 1$
3	f <sub>0</sub> (3)	f <sub>1</sub> (3)	f <sub>2</sub> (3)	f <sub>3</sub> (3)	$f_3(3) + 1$
4	f <sub>0</sub> (4)	f <sub>1</sub> (4)	f <sub>2</sub> (4)	f <sub>3</sub> (4)	
					$f_{i}(i) + 1$

There are countably infinite programs

There are uncountably infinite simple integer functions

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The Halting Problem
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The Median Problem

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NP Complete Problems IS \ (Independent \ Set) \ Problem \\ Subgraph \ S \ such \ that \ no \ edges \ in \ V_s \ are \ adjacent \\ VC \ (Vertex \ Cover) \ Problem \\ Subgraph \ S \ such \ that \ each \ v_i \in V_s \ is \ adjacent \ to \ another \ vertex \ v_i \in V_s
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IS and VC are complements of each other

## **Final Exam Review**

Time complexities Names of main algorithms

1 question from the first exam and 1 question from the second exam (more general)

For example:

- 1) Sorting algorithms
- 2) Graph algorithms
- 3) Pseudocode for something ← e.g., Halting problem

The rest are from the last 1.5 lectures

Sequences (know the most basic ones, like these)

$$1 + 2 + 3 + ... + n = n(n+1)/2$$

$$r + r^2 + r^3 + ... + r^n = (r^{n+1} - 1)/(r - 1)$$

$$f(n) = f(n) + f(n-2) \leftarrow$$
 Characteristic equation (be able to identify what this is)

$$f(n) - f(n-1) - f(n-2) = 0$$

$$x^2 - x - 1 = 0$$

$$n = 2^k$$

$$T(n) = T(n/2) + f(n) \leftarrow$$
 Domain transformation (change the input)  $S(k) = S(k-1)$ 

$$T(n) = 2T(n/2) + ... \leftarrow$$
 Range transformation (change the output)

Master theorem ← What theorem helps algorithms estimate time complexity

Dynamic programming technique

Floyd

Diikstra

Prim

Bellman Ford  $\leftarrow$  Handles negative weight complexity, O(V \* E)

Weighted triangulation ← Not on the final

## Dynamic problem solution

Knapsack problem  $\leftarrow$  Classic NP complete problem, O(n \* k) where k is the upper bound on weight capacity of knapsack

Classic knapsack problem → You want to maximize the value

0-1 knapsack problem

0-1 knapsack problem with integer capacity  $\rightarrow$  0(n \* k) but k will be very large to n, more like exponential

Fractional knapsack problem  $\rightarrow$  O(n), sort by: (value/weight)

# Basic matrix multiplication problem

Divide and conquer approach, divide into 4 quadrants

Now you have 8 subproblems, solve those

$$T(n) = 8T(n/2) + O(n^2)$$

$$T(n) = O(n^3)$$

#### Strassen

$$T(n) = 7T(n/2) + O(n^2)$$

$$T(n) = O(n^{\log_2 7})$$