## Homework on Robust Optimization

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In this homework we investigate  $\Gamma$ -uncertainty in the context of linear programs (LPs). In particular we consider LPs of the following form:

where  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$ , and  $x, \ell, u \in \mathbb{R}^n$ . We assume without loss of generality that the objective function is certain. On the other hand, for each row  $A_i$  of A we have a set  $J_i \subseteq \{1, 2, \ldots, n\}$  of indices of the elements of  $A_i$  that are uncertain. The elements  $\{1, 2, \ldots, n\} \setminus J_i$  of  $A_i$  are certain. Furthermore, for  $1 \leq i \leq m$ , we have a value  $\Gamma_i \in [0, |J_i|]$  that controls how much the uncertainty may affect the elements of  $A_i$  indexed by  $J_i$ . Suppose that the element  $a_{ij}$  of A is uncertain (i.e.,  $j \in J_i$ ). Then the value of  $a_{ij}$  will be in the interval  $[a_{ij} - \xi_{ij}, a_{ij} + \xi_{ij}]$  for some given  $\xi_{ij} \geq 0$ . Hence the uncertainty is coordinate-wise, that is,  $\mathcal{U} = \mathcal{U}_1 \times \ldots \times \mathcal{U}_m$ . For a given solution  $x \in \mathbb{R}^n$ , we quantify the amount of uncertainty that can hit constraint i by

$$\beta_i(x, \Gamma_i) := \max_{\{S_i \cup \{t_i\} | S_i \subseteq J_j, |S_i| = \lfloor \Gamma_i \rfloor, t_i \in J_i \setminus S_i\}} \left\{ \sum_{j \in S_i} \xi_{ij} |x_j| + (\Gamma_i - \lfloor \Gamma_i \rfloor) \xi_{it_i} |x_{t_i}| \right\}$$

We obtain the following non-linear robust counterpart of (LP):

$$\max_{s.t.} c^T x$$
s.t.  $A_i x + \beta_i(x, \Gamma_i) \le b$   $(1 \le i \le m)$  (RLP)
$$\ell \le x \le u$$

Exercise 1. Our first goal is to show that the non-linear optimization problem (RLP) can be written as a linear program.

(a) Show that for a feasible solution x of (RLP) and  $1 \le i \le m$ , the function  $\beta_i(x, \Gamma_i)$  is equal to the optimal value of

$$\max \sum_{j \in J_i} \xi_{ij} | x_j | z_{ij}$$
s.t. 
$$\sum_{j \in J_i} z_{ij} \le \Gamma_i$$

$$0 \le z_{ij} \le 1 \qquad (j \in J_i)$$

$$(P_i)$$

(b) By using the LP dual of  $(P_i)$  show that (RLP) can be written as a linear program of the following form

$$\max \quad c^{T}x$$
s.t. 
$$A_{i}x + z_{i}\Gamma_{i} + \sum_{j \in J_{i}} p_{ij} \leq b_{i} \qquad (1 \leq i \leq m)$$

$$z_{i} + p_{ij} \geq \xi_{ij}y_{j} \qquad (1 \leq i \leq m, j \in J_{i})$$

$$p_{ij} \geq 0 \qquad (1 \leq i \leq m, j \in J_{i})$$

$$-y \leq x \leq y$$

$$\ell \leq x \leq u$$

$$y \geq 0, z, p \geq 0$$
(RLP<sub>lin</sub>)

where  $x, y, u, \ell \in \mathbb{R}^n$  and  $z \in \mathbb{R}^m$ .

Exercise 2. The parameters  $\Gamma_i$  control the amount of uncertainty we allow per constraint. We would like to investigate the impact of a very small change  $\Delta\Gamma_i$  of  $\Gamma_i$  on the objective function. For this purpose we assume that (RLP<sub>lin</sub>) has a unique optimal primal solution  $z^*$  and a unique optimal dual solution  $q^*$ . Furthermore, let B be the unique optimal basis corresponding to  $z^*$ . Our goal is to show that the derivative of the objective function with respect to  $\Gamma_i$  is equal to  $z_i^*q_i^*$ .

(a) We consider (RLP<sub>lin</sub>) written in standard form and depending on the parameter  $\Gamma_i$ :

$$G(\Gamma_i) := \begin{cases} \max & c^T x \\ \text{s.t.} & Ax + \Gamma_i z_i e_i = b \end{cases} \quad (1 \le i \le m)$$

where  $e_i$  is a unit vector whose *i*th element is 1. Show that if  $z_i$  is a non-basic variable with respect to B then the derivative of the objective function given above is correct.

(b) To conclude assume that  $z_i$  is a basic variable with respect to B. First, write the basis associated to  $G(\Gamma_i + \Delta \Gamma_i)$  as the sum of two matrices and apply the Sherman-Morrison identity (that can be found on wikipedia) to obtain its inverse. Use the result to show that

$$\lim_{\Delta\Gamma_i \to 0} \frac{G(\Gamma_i + \Delta\Gamma_i) - G(\Gamma_i)}{\Delta\Gamma_i} = z_i^* q_i^* .$$