



UNIVERSITI
TEKNOLOGI
MALAYSIA
www.utm.my

SCSR1013 DIGITAL LOGIC

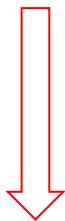
MODULE 4a: BOOLEAN ALGEBRA

FACULTY OF COMPUTING



Introduction to Boolean Algebra

- It is the **set of rules** used to simplify a given **logic expression** without changing its functionality.
- Example: $F = \bar{A}B + BC + ABC$



be minimized

$$F = \bar{A}B + BC$$

- It is used when number of variables are less, i.e., 1, 2 or 3 vars.



- **Variable**

- A symbol that represents a logical quantity
- Usually italic uppercase (A, B, C, D)
- A single variable can have a 1 or 0 value

- **Complement**

- The inverse of a variable
- Indicated by an overbar (\bar{A}) or prime (A')
- If $A = 1$, then $\bar{A} = 0$

- **Literal**

= both variable and its complement in a term

$$\bar{A} + B + C \rightarrow 3 \text{ literals}$$



Example:

$$f = x'y' + x'y = \overline{\overline{x}}\overline{y} + \overline{x}\overline{y}$$

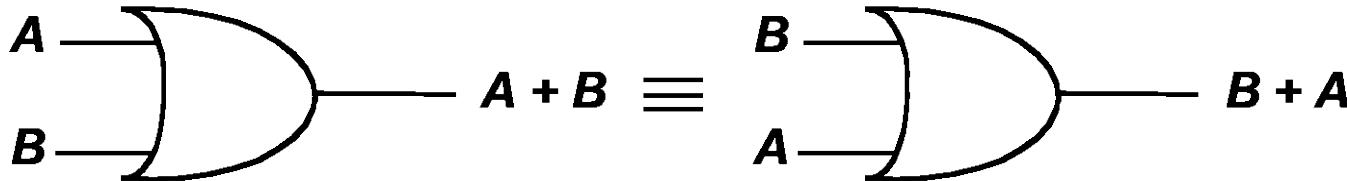
Total Number of Literals = 4 which are x' , y' , x' and y

Total Number of variables = 2 which are x and y

Total Number of terms = 2 which are $x'y'$ and $x'y$

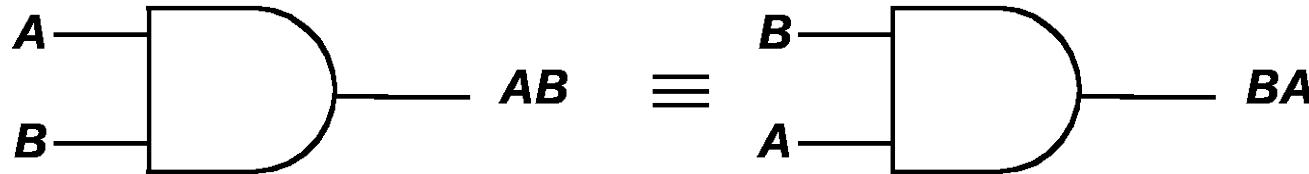


- Basic laws of BA
 - **Commutative Laws**
 - For addition and multiplication
 - **Associative Laws**
 - For addition and multiplication
 - **Distributive Laws**



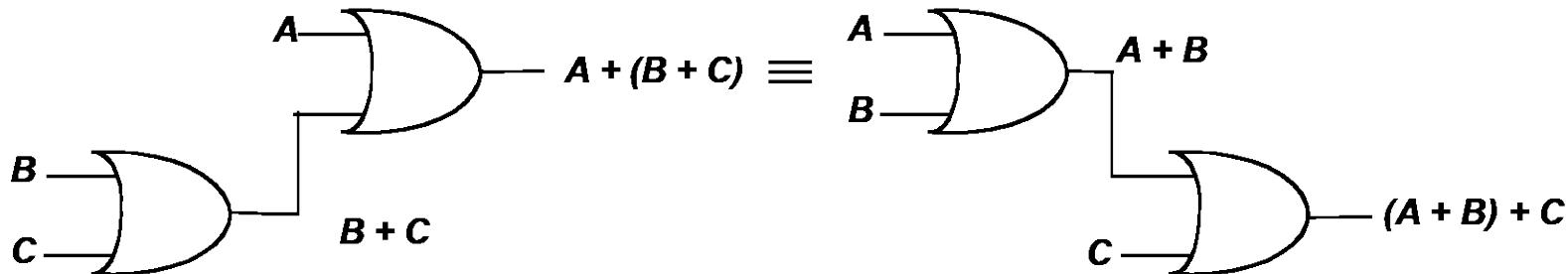
Commutative law of addition

- $A + B = B + A$
 - the **order** in which the variables are OR-ed makes no difference
 - in logic circuits, addition and the OR operation are the same.



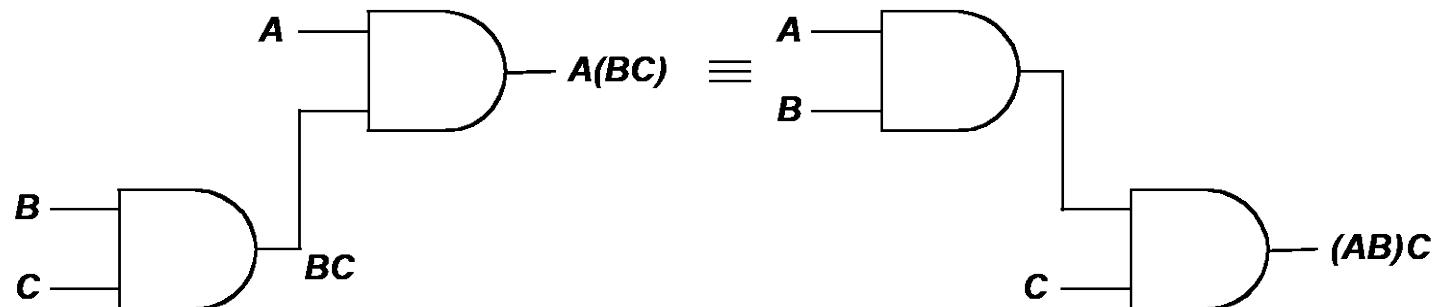
Commutative law of multiplication

- $AB = BA$
 - the **order** in which the variables are AND-ed makes no difference
 - in logic circuits, multiplication and the AND operation are the same.



Associative law of addition

- $A + (B + C) = (A + B) + C$
 - the **grouping** in which the variables are OR-ed makes no difference



Associative law of multiplication

- $A(BC) = (AB)C$
 - the **grouping** in which the variables are AND-ed makes no difference



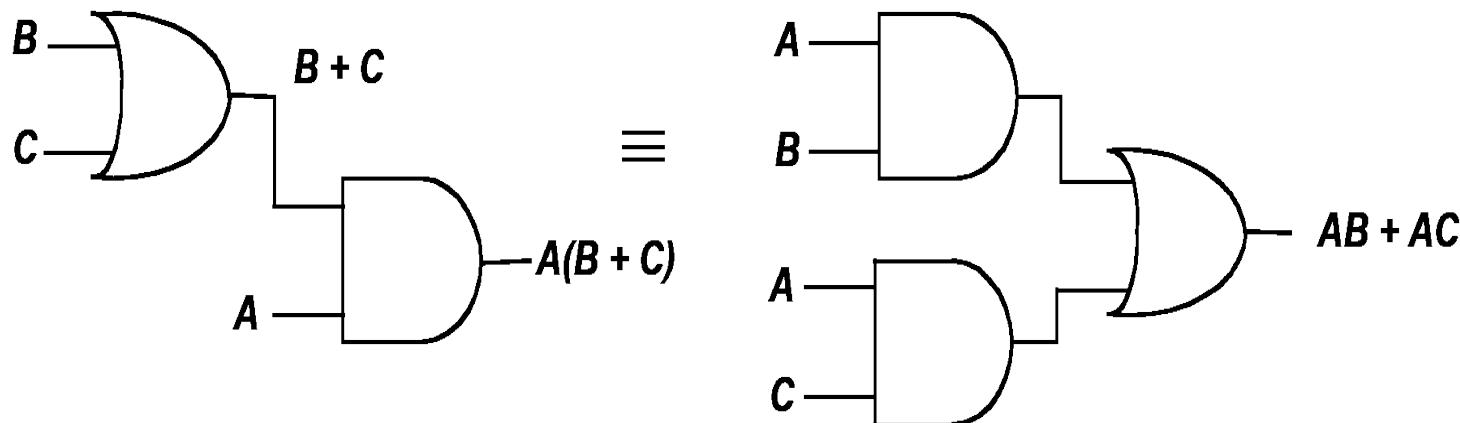
Proof

Exercise 4a.1:

Prove the Associate Law for $A(BC) = (AB)C$ using truth table.

| A | B | C | AB | BC | A(BC) | (AB)C |
|---|---|---|----|----|-------|-------|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 | 1 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 | 0 | 0 |
| 1 | 1 | 0 | 1 | 0 | 0 | 0 |
| 1 | 1 | 1 | 1 | 1 | 1 | 1 |

- $A(B + C) = AB + AC$
 - expresses the process of factoring, variable A is **factored out** of the product terms, like $AB + AC = A(B+C)$
 - ... A is factored to the $(B + C)$



Why we have this ???

- want to simplify the Boolean equation
- thus we have simple circuit
- then reduce cost of production !

1

The rules can be proven by using truth table.

| Rule 10 | | | |
|---------|---|----|----------|
| A | B | AB | $A + AB$ |
| 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 |

↑ equal ↑

Rules of Boolean Algebra

| | |
|----|---------------------------|
| 1 | $A + 0 = A$ |
| 2 | $A + 1 = 1$ |
| 3 | $A \cdot 0 = 0$ |
| 4 | $A \cdot 1 = A$ |
| 5 | $A + A = A$ |
| 6 | $A + \bar{A} = 1$ |
| 7 | $A \cdot A = A$ |
| 8 | $A \cdot \bar{A} = 0$ |
| 9 | $\bar{\bar{A}} = A$ |
| 10 | $A + AB = A$ |
| 11 | $A + \bar{A}B = A + B$ |
| 12 | $(A + B)(A + C) = A + BC$ |



2

The rules can be proven by using Boolean algebra laws and rules.

Rule 10:

$$\begin{aligned} A + AB &= A(1 + B) && \text{Factoring (distributive law)} \\ &= A \cdot 1 && \text{Rule 2: } (1 + B) = 1 \\ &= A && \text{Rule 4: } A \cdot 1 = A \end{aligned}$$

Rules of Boolean Algebra

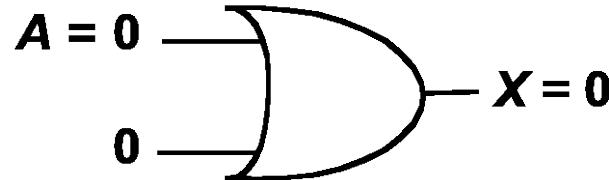
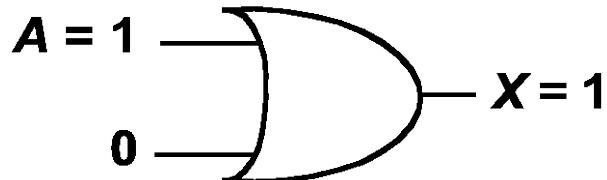
| | |
|----|---------------------------|
| 1 | $A + 0 = A$ |
| 2 | $A + 1 = 1$ |
| 3 | $A \cdot 0 = 0$ |
| 4 | $A \cdot 1 = A$ |
| 5 | $A + A = A$ |
| 6 | $A + \bar{A} = 1$ |
| 7 | $A \cdot A = A$ |
| 8 | $A \cdot \bar{A} = 0$ |
| 9 | $\bar{\bar{A}} = A$ |
| 10 | $A + AB = A$ |
| 11 | $A + \bar{A}B = A + B$ |
| 12 | $(A + B)(A + C) = A + BC$ |

Rules of Boolean Algebra: PROOF

1

$$A + 0 = A$$

Rule 1:



$$X = A + 0 = A$$

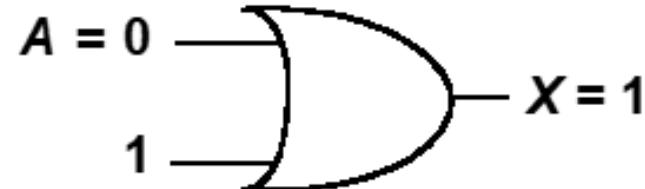
A *variable* OR-ed with **0** is always equal to the *variable*.

continue...

Rules of Boolean Algebra: PROOF

2 $A + 1 = 1$

Rule 2:



$$X = A + 1 = 1$$

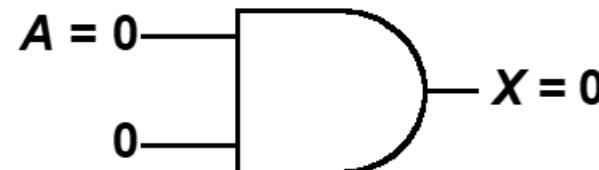
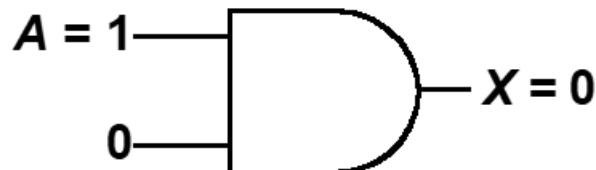
A *variable* OR-ed with 1 is always equal to the 1.

continue...

3

$A \cdot 0 = 0$

Rule 3:



$$X = A \cdot 0 = 0$$

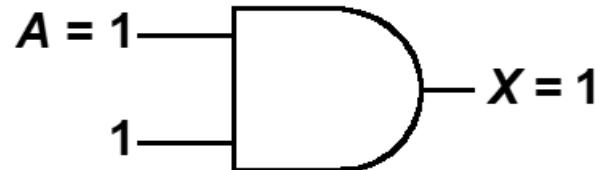
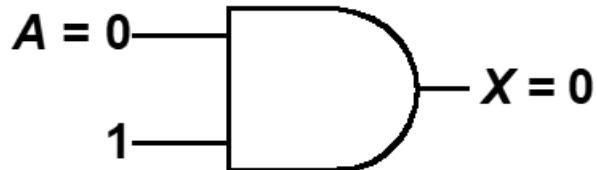
A *variable* AND-ed with 0 is always equal to 0.

continue...

4

$A \cdot 1 = A$

Rule 4:



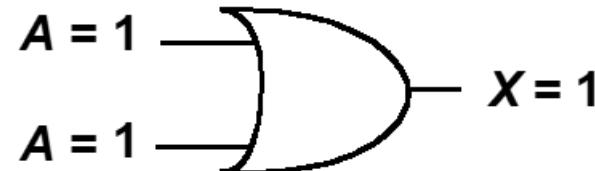
$$X = A \cdot 1 = A$$

A variable AND-ed with 1 is always equal to variable.

continue...

5 | $A + A = A$

Rule 5:



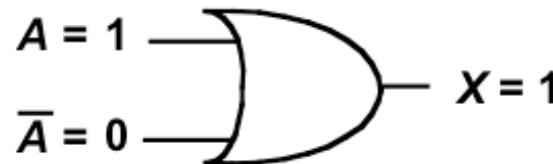
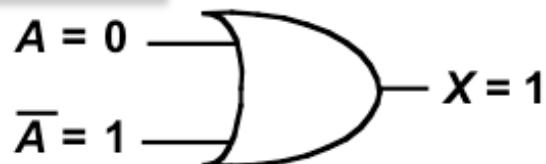
$$X = A + A = A$$

A variable OR-ed with itself is always equal to the variable.

continue...

6 $A + \bar{A} = 1$

Rule 6:



$$X = A + \bar{A} = 1$$

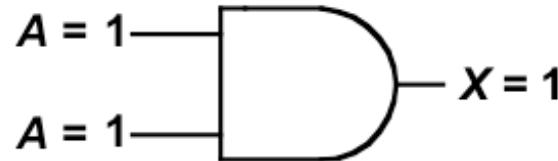
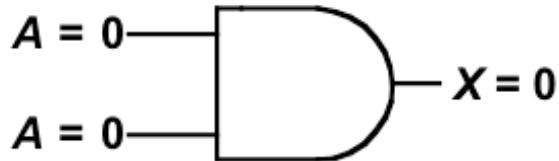
A *variable* OR-ed with its *complement* is always equal to 1.

continue...

7

$$A \bullet A = A$$

Rule 7:



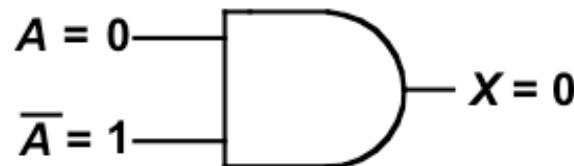
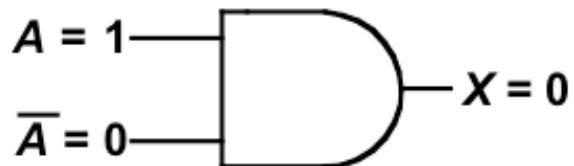
$$X = A \bullet A = A$$

A *variable* AND-ed with itself is always equal to the *variable*.

continue...

8 $A \bullet \bar{A} = 0$

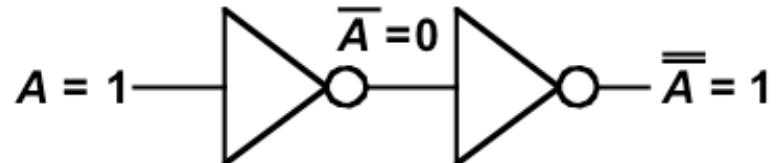
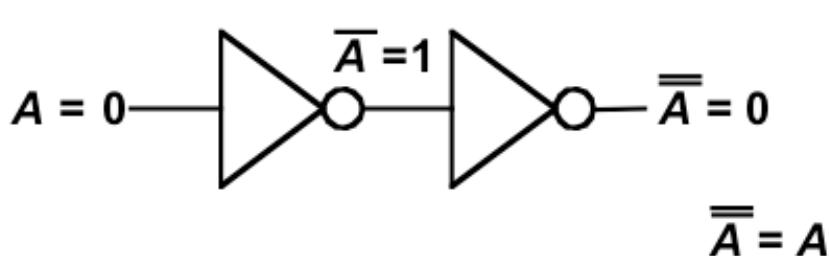
Rule 8:



$$X = A \bullet \bar{A} = 0$$

A variable AND-ed with its complement is always equal to 0

continue...

9 | $\overline{\overline{A}} = A$ **Rule 9:**

The *double complement* of a variable is always equal to the *variable*

continue...



Rules 10, 11, and 12 can be proven by using Boolean algebra laws.

10 $A + AB = A$

Rule 10:

$$A + AB = A(1 + B) \quad \text{Factoring (distributive law)}$$

$$= A \cdot 1 \quad \text{Rule 2: } (1 + B) = 1$$

$$= A \quad \text{Rule 4: } A \cdot 1 = A$$

continue...



Rules 10, 11, and 12 can be proven by using Boolean algebra laws.

11 $A + \bar{A}B = A + B$

Rule 11:

$$\begin{aligned} A + \bar{A}B &= (A + AB) + \bar{A}B \\ &= (AA + AB) + \bar{A}B \\ &= AA + AB + A\bar{A} + \bar{A}B \\ &= (A + \bar{A})(A + B) \\ &= 1 \cdot (A + B) \\ &= A + B \end{aligned}$$

Rule 10: $A = A + AB$

Rule 7: $A = AA$

Rule 8: adding $\bar{A}\bar{A} = 0$

Factoring

Rule 6: $A + \bar{A} = 1$

Rule 4: drop the 1

continue...



Rules 10, 11, and 12 can be proven by using Boolean algebra laws.

12

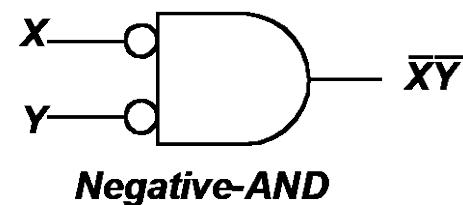
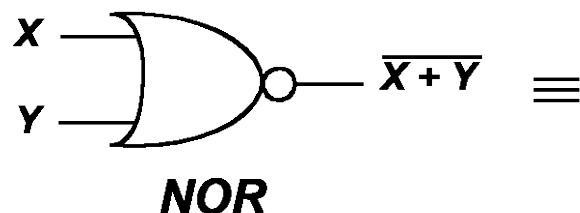
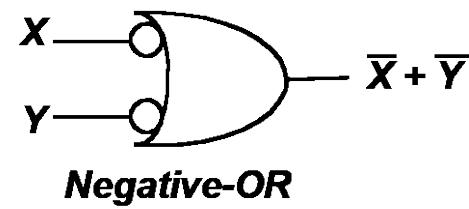
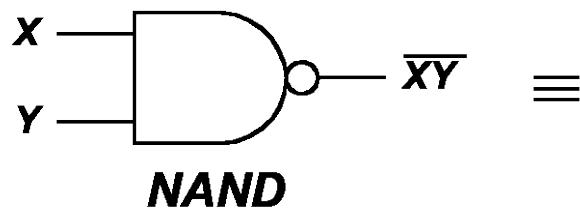
$(A + B)(A + C) = A + BC$

Rule 12:

$$\begin{aligned}(A + B)(A + C) &= AA + AC + AB + BC && \text{Distributive law} \\&= A + AC + AB + BC && \text{Rule 7: } AA = A \\&= A(1 + C) + AB + BC && \text{Factoring (distributive law)} \\&= A \cdot 1 + AB + BC && \text{Rule 2: } 1 + C = 1 \\&= A(1 + B) + BC && \text{Factoring (distributive law)} \\&= A \cdot 1 + BC && \text{Rule 2: } 1 + B = 1 \\&= A + BC && \text{Rule 4: } A \cdot 1 = A\end{aligned}$$

DeMorgan's Theorems

- To minimize the **variety** and **number of logic gates** IC.
- Provides mathematical verification for:
 - NAND \equiv negative-OR
 - NOR \equiv negative-AND





- **DM theorem 1:**

- The complement of a product of variables is equal to the sum of the complements of the variables

$$\overline{XY} = \overline{X} + \overline{Y}$$

- **DM theorem 2:**

- The complement of a sum of variables is equal to the product of the complements of the variables

$$\overline{X + Y} = \overline{X}\overline{Y}$$



Example of applying DeMorgan's Theorems.

DM theorem 1:

The complement of a product of variables is equal to the sum of the complements of the variables

$$\overline{XYZ} = \overline{X} + \overline{Y} + \overline{Z}$$

$$\overline{WXYZ} = \overline{W} + \overline{X} + \overline{Y} + \overline{Z}$$

DM theorem 2:

The complement of a sum of variables is equal to the product of the complements of the variables

DeMorgan's Theorem II

$$\overline{X + Y + Z} = \overline{XYZ}$$

$$\overline{W + X + Y + Z} = \overline{W}\overline{X}\overline{Y}\overline{Z}$$



Example of applying DeMorgan's Theorems.

$$\overline{(AB + C)(A + BC)} = \overline{(AB + C)} + \overline{(A + BC)} \quad (\text{Theorem I})$$

$$\overline{(AB + C)} + \overline{(A + BC)} = \overline{(AB)}\overline{C} + \overline{A}\overline{(BC)} \quad (\text{Theorem II})$$

$$= (\bar{A} + \bar{B})\bar{C} + \bar{A}(\bar{B} + \bar{C}) \quad (\text{Theorem I})$$

(Theorem I)

$$\overline{XY} = \overline{X} + \overline{Y}$$

(Theorem II)

$$\overline{X+Y} = \overline{X}\overline{Y}$$



Extra

Self-Test:

Prove that \overline{AB} is equal or not equal with $\overline{A}\overline{B}$ by using the truth table.



Extra

Self-Test:

Prove that \overline{AB} is **equal** or **not equal** with $\overline{A}\overline{B}$ by using the truth table.

Solution:

| A | B | \overline{A} | \overline{B} | AB | \overline{AB} | $\overline{A}\overline{B}$ |
|---|---|----------------|----------------|----|-----------------|----------------------------|
| | | | | | | |
| | | | | | | |
| | | | | | | |
| | | | | | | |



Extra

Self-Test:

Prove that \overline{AB} is equal or not equal with $\overline{\overline{A}\overline{B}}$ by using the truth table.

Solution:

| A | B | \overline{A} | \overline{B} | AB | $\overline{A}\overline{B}$ | $\overline{\overline{A}\overline{B}}$ |
|---|---|----------------|----------------|------|----------------------------|---------------------------------------|
| 0 | 0 | 1 | 1 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | 0 | 1 | 0 | 0 |

$$\overline{AB} \neq \overline{\overline{A}\overline{B}}$$



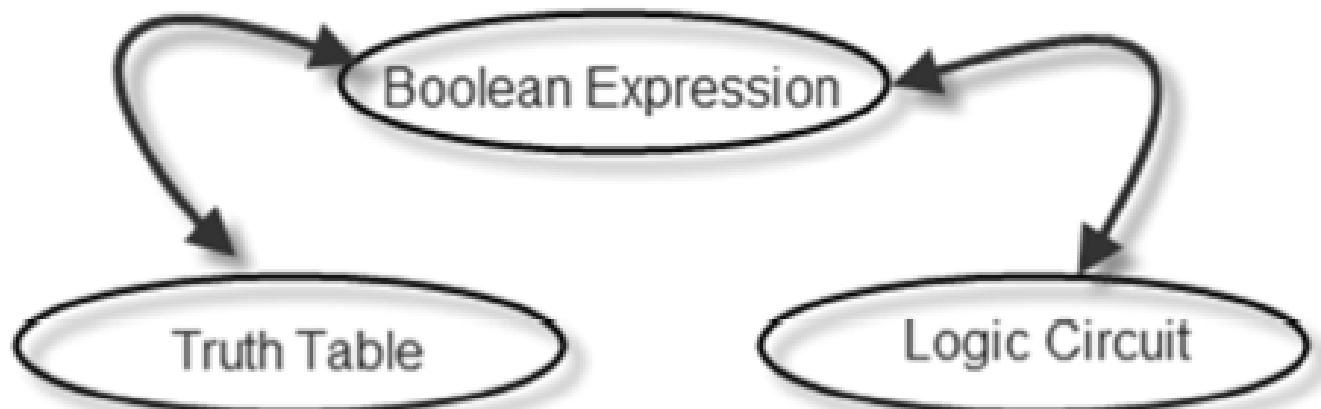
Exercise 4a.2:

Apply DeMorgan's theorems to each of the following expressions:

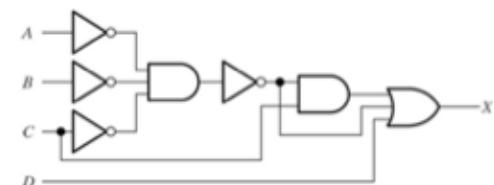
- (a) $\overline{(A + B + C)D}$
- (b) $\overline{ABC + DEF}$
- (c) $\overline{\overline{AB} + \overline{CD} + EF}$

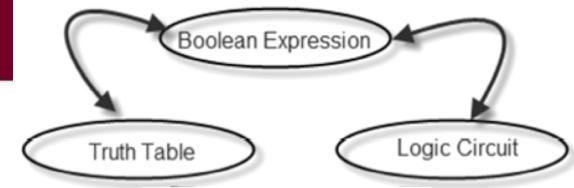
Combinational Logic Representation

$$\overline{(AB + C)(A + BC)}$$



| A | B | AB |
|---|---|----|
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

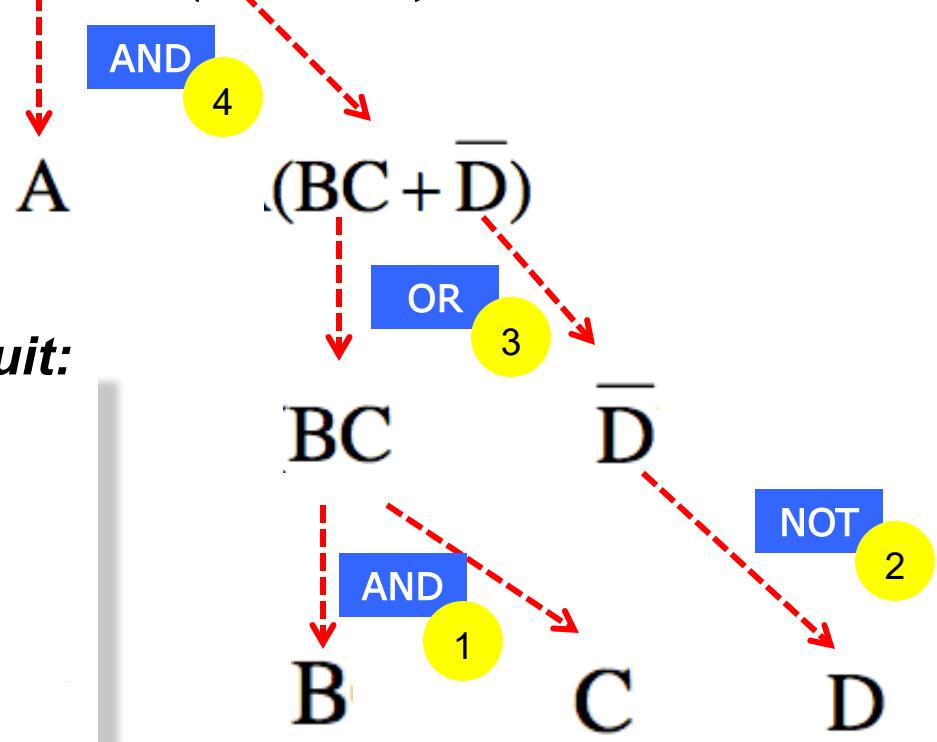




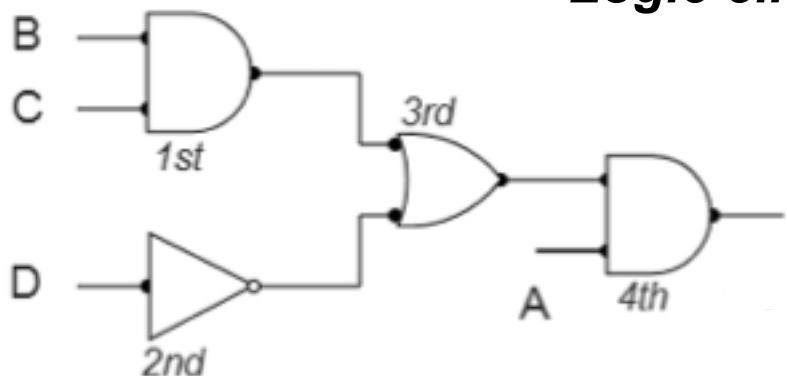
From Boolean Expression to Logic Circuit

- Treat a Boolean expression as if you try to solve a normal mathematical equation
- Just like a normal mathematic we have the precedence of operation
 - start with () then multiplication (AND) then addition (OR)
 - Therefore to draw the circuit following the same sequence

Example: $X = A(BC + \bar{D}) = A \cdot (BC + \bar{D})$



Logic circuit:



1 → 2 → 3 → 4



Exercise 4a.3:

Draw the logic circuit represented by each expression:

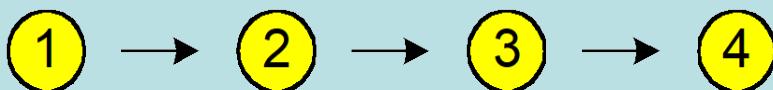
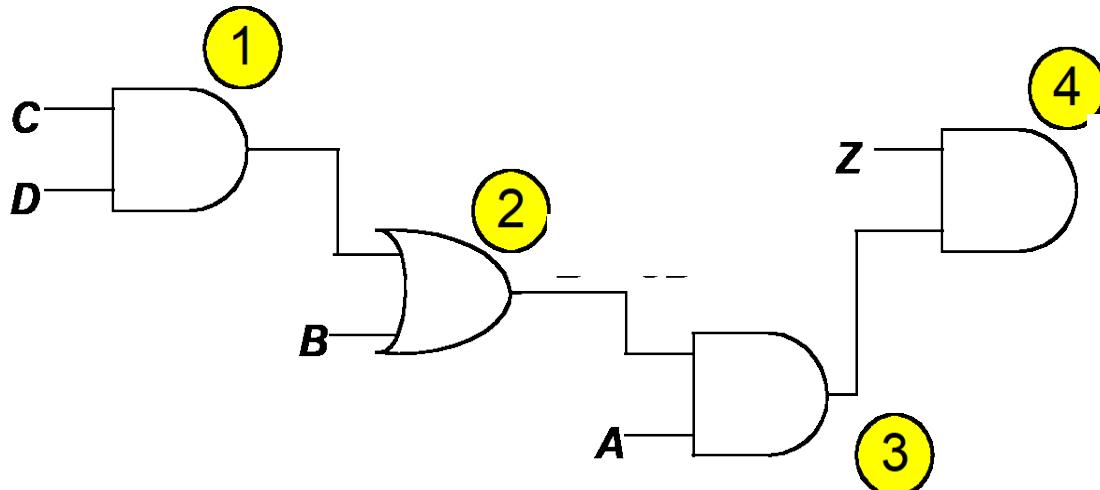
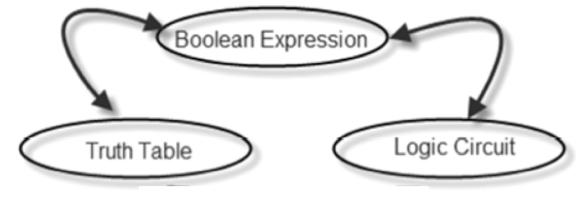
(i) $A\bar{B} + \bar{A}B$

(ii) $AB + \bar{A}\bar{B} + \bar{A}BC$

(iii) $\bar{A}B(C + \bar{D})$

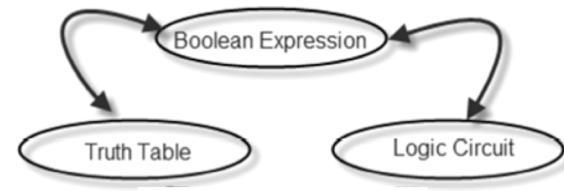
Logic circuit to Boolean expression

- To derive the Boolean expression for a given circuit, follow **left-2-right** rule.
 - Begin from the left-most inputs and work towards the last.



Boolean expression to truth table

- A truth table shows the output for all possible input values.
- From a Boolean expression, a truth table can be developed.



x = number of input variables

Possible combinations of values, $n = 2^x$

- Example:

$$F = A + B \rightarrow x = 2; n = 2^2 = 4$$



Steps in construction a truth table

- **Step 1:** Identify **x** and **n** from the Boolean exp.
- **Step 2:** Find the values of the variables that make the expression equal to **1**.
(Hint: use the rules for Boolean addition and multiplication)
- **Step 3:** List in a table
 - all the **n** combinations of 1s and 0s (**inputs**)
 - the values of variables from step 2 (**outputs=1**)
 - all the **other output** values will be **0**



Example: $F = A + B$

1. $x = ?$

$n = 2^x ?$

2. Combination of inputs $\rightarrow F = 1$

$A + B = 1$

3. Fill in the table:

$x=2, n = 4$

| INPUT | | OUTPUT |
|-------|---|--------|
| A | B | F |
| 0 | 0 | |
| 0 | 1 | |
| 1 | 0 | |
| 1 | 1 | |



Example: $A(B + CD)$

Step 1: $x = 4$; $n = 2^x = 2^4 = 16$

Step 2:

$$A(B + CD) = 1 \cdot 1 = 1 \rightarrow A = 1$$

What makes $B + CD = 1$?

$$B + CD = 1 + 0 = 1$$

$$B + CD = 0 + 1 = 1$$

$$B + CD = 1 + 1 = 1 \rightarrow B = 1 \text{ or } 0$$

What makes $CD = 1$?

$$CD = 1 \cdot 1 = 1 \rightarrow C = 1; D = 1$$

Therefore, the output of $A(B + CD)$ will be 1 if

$$A = 1, B = 1, C = 0/1, D = 1/0$$

$$A = 1, B = 0, C = 1, D = 1$$

Step 3 : Fill in the table with results from Step 2

| INPUTS | | | | OUTPUT |
|--------|---|---|---|-------------|
| A | B | C | D | $A(B + CD)$ |
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 |
| 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 |



- Tips on ‘**table-making**’ :
 - For n possible combinations, the input part of the table will register the binary value of 0 to $n-1$.
 - (e.g. $n=16$; 0 to 15)
 - Remember the sequence

| 2^4 | 2^2 | 2^1 | 2^0 | |
|-------|-------|-------|-------|-----|
| 8 | 4 | 2 | 1 | |
| 0 | 0 | 0 | 0 | (0) |
| 0 | 0 | 0 | 1 | (1) |
| 0 | 0 | 1 | 0 | (2) |
| 0 | 0 | 1 | 1 | (3) |

www

From Truth Table to Logic Circuit

Step 1: Get the product term from HIGH outputs **1**

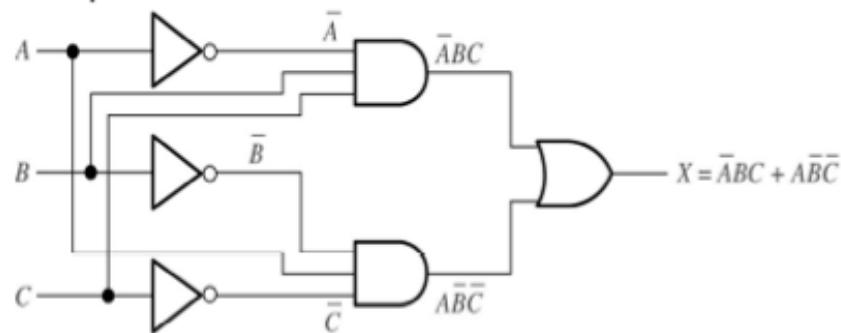
Step 2: From the product term get the expression

- This is done by OR-ing the product terms

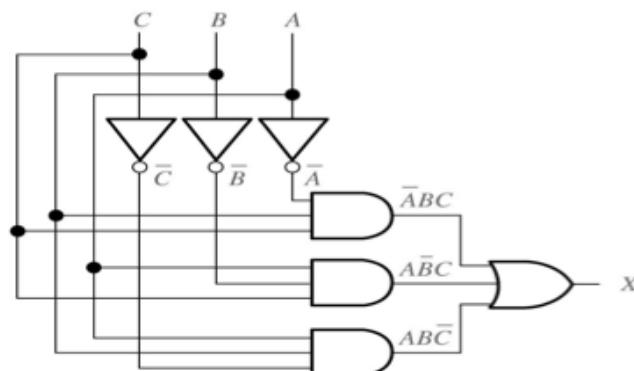
$$X = \overline{A}BC + A\overline{B}\overline{C}$$

| INPUTS | | | PRODUCT TERM | |
|--------|---|---|--------------|-----------------------------|
| A | B | C | X | |
| 0 | 0 | 0 | 0 | |
| 0 | 0 | 1 | 0 | |
| 0 | 1 | 0 | 0 | |
| 0 | 1 | 1 | 1 | $\overline{A}BC$ |
| 1 | 0 | 0 | 1 | $A\overline{B}\overline{C}$ |
| 1 | 0 | 1 | 0 | |
| 1 | 1 | 0 | 0 | |
| 1 | 1 | 1 | 0 | |

Step 3: Implement the circuit



| INPUTS | | | |
|--------|---|---|---|
| A | B | C | X |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 0 |



Example:

Given the truth table below, produce the logic circuit (Module v5: page 116).

| INPUTS | | | |
|--------|---|---|---|
| A | B | C | X |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 0 |

←————— $X = \overline{ABC} + \overline{A}\overline{B}C + A\overline{B}\overline{C}$



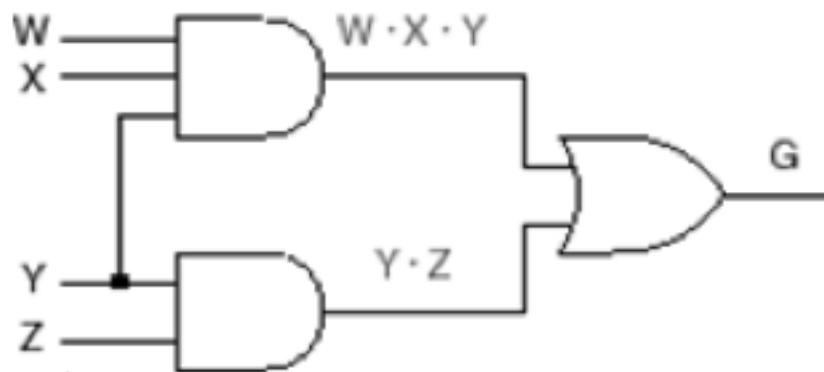
Logic Circuit to Truth Table

- Requires 2 steps
 1. From circuit derive the Boolean expression of the output
 2. From Boolean expression produce a Truth Table

Both step has been mentioned



Example:



Produce a truth table from the above circuit.

Step 1: derive the Boolean expression of the output , $G = WXY + YZ$

Step 2: from the expression produce the truth table

G will be 1 if either term = 1

term 1 which is WXY will only be 1 if $W = 1$, $X = 1$ and $Y = 1$

term2 which is YZ will only be 1 if $Y = 1$ and $Z = 1$

| w | x | y | z | g |
|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 1 | 1 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 |
| 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 0 | 0 | 0 |
| 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 |