Instruction: You must show all your work clearly for credit. Partial credit will only be given to meaningful answers.

1. Use the definition of big-O to prove that $n^4 - 8n^3 + 16n^2 - 3n + 560 \neq O(n^3)$.

Proof: If true, we must have constants k and n_0 such that $n^4 - 8n^3 + 16n^2 - 3n + 560 \le n^3$,

$$\forall n \ge n_0$$
. Hence, $\frac{n^4 - 8n^3 + 16n^2 - 3n + 560}{n^3} \le k$. As $n \to \infty$, we have $\infty \le k = \text{constant}$. A

contradiction is reached, implying that $n^4 - 8n^3 + 16n^2 - 3n + 560 \neq O(n^3)$.

2. For any two functions f(n) and g(n), by definition, f(n) = o(g(n)) iff f(n) = O(g(n)) and $f(n) \neq \Theta(g(n))$. Use the definition of little-o to prove or disprove that $n^4 - 8n^3 + 16n^2 - 3n + 560 = o(n^5)$.

Proof: (i) Prove that $n^4 - 8n^3 + 16n^2 - 3n + 560 = O(n^5)$.

$$n^4 - 8n^3 + 16n^2 - 3n + 560$$

$$\leq n^4 + 16n^4 + 560n^4, \forall n \geq 1$$

$$\leq 577n^4, \forall n \geq 1$$

$$\leq 577n^5, \forall n \geq 1.$$

Hence,
$$n^4 - 8n^3 + 16n^2 - 3n + 560 = O(n^5)$$
.

(ii) Prove that $n^4 - 8n^3 + 16n^2 - 3n + 560 \neq \Theta(n^5)$.

If true, we must have $n^4 - 8n^3 + 16n^2 - 3n + 560 = \Omega(n^5)$, implying that there exist constants k and n_0 such that $n^4 - 8n^3 + 16n^2 - 3n + 560 \ge kn^5$, $\forall n \ge n_0$.

Hence,
$$\frac{n^4 - 8n^3 + 16n^2 - 3n + 560}{n^5} \ge k, \forall n \ge n_0$$
.

As $n \to \infty$, we have $0 \ge k > 0$, which is a contradiction.

Hence,
$$n^4 - 8n^3 + 16n^2 - 3n + 560 \neq \Theta(n^5)$$
.

Conclusion:
$$n^4 - 8n^3 + 16n^2 - 3n + 560 = o(n^5)$$
.

3. Let f(n) and g(n) be any two positive functions. Prove or disprove the statement that if f(n) = O(g(n)), then $2^{f(n)} = O(2^{g(n)})$.

Solution:

If
$$f(n) = O(g(n))$$
, then $2^{f(n)} = O(2^{g(n)})$ may or may not be true.

Take
$$f(n) = 2n$$
, $g(n) = n$.

$$2^{2n} = 4^n \neq O(2^n).$$

4. By assuming that all basic operations require the same constant cost K, compute the cost of the resource function $R_w(n)$ in closed-form for the following program segment using the simplified approach as discussed in class. *Remark*: You must first set up a summation equation for $R_w(n)$ and then evaluate the sum(s) clearly for credit.

$$x = 210;$$

 $y = 560;$
 $for i = 1 \text{ to } n*n \text{ do}$
 $for j = i \text{ to } n \text{ do}$
 $y = x * y / 660 + 388;$
 $endfor;$
 $endfor;$

Wrong Solution:

$$T(n) = \sum_{i=1}^{n^2} \sum_{j=i}^{n} K$$

$$= K \sum_{i=1}^{n^2} (n - i + 1)$$

$$= K [n^3 - \frac{n^2 (n^2 + 1)}{2} + n^2]$$
< 0. (Q: What is wrong?)

Correct Solution:

$$T(n) = \left(\left(\sum_{i=1}^{n} + \sum_{i=n+1}^{n^{2}}\right)\sum_{j=i}^{n}\right)K$$

$$= \sum_{i=1}^{n} \sum_{j=i}^{n} K + \sum_{i=n+1}^{n^{2}} K$$

$$= K \sum_{i=1}^{n} (n-i+1) + (n^{2}-n)K$$

$$= K \left[2n^{2} - \frac{n(n+1)}{2}\right]$$

$$= K \left(\frac{3n^{2}-n}{2}\right)$$

$$= \Theta(n^{2}).$$

- 5. Given a set of records with 7 keys $S = \{35, 28, 43, 17, 39, 3, 46\}$.
 - (a) By using the hash function $h(x) = x \mod m$ and chaining with singly linked list in constructing an open hash table H with m = 11 buckets, insert the records in S, in the given order, into H. You must show your computations for locations and illustrate the final structure of your hash table H clearly for credit. Remark: Insertion must be done at the beginning of the list.
 - (b) By using the hash function $h(x) = x \mod m$ and quadratic probing in constructing a closed hash table H with m = 11 buckets, insert the records in S, in the given order, into H. You must show your computations for locations and illustrate the final structure of your hash table H clearly for credit.
 - (c) Given two hash functions $h(x) = x \mod m$ and $h^+(x) = p x \mod p$. By using open addressing with $f_i = i * h^+(x)$ and double hashing in constructing a closed hash table H with m = 11 buckets and p = 5, insert the records in S, in the given order, into H. You must show your computations for locations and illustrate the final structure of your hash table H clearly for credit.

Solution:

(a) Address computations:

35 % 11 = 2,

28 % 11 = 6,

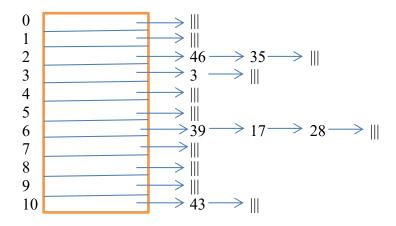
43 % 11 = 10,

17 % 11 = 6,

39 % 11 = 6,

3% 11 = 3,

46 % 11 = 2.



Open hash table with chaining

(b) Address computations:

$$35 \% 11 = 2$$
,
 $28 \% 11 = 6$,
 $43 \% 11 = 10$,
 $17 \% 11 = 6 \rightarrow 7$,
 $39 \% 11 = 6 \rightarrow 7 \rightarrow 10 \rightarrow 15 \% 11 = 4$,
 $3 \% 11 = 3$,
 $46 \% 11 = 2 \rightarrow 3 \rightarrow 6 \rightarrow 11 \% 11 = 0$.

0	46
1	
2	35
2 3 4 5 6 7	35 3 39
4	39
5	
6	28
	17
8	
9	
10	43

Closed hash table with quadratic probing

(c) Address computations:

$$35 \% 11 = 2$$
,

$$28 \% 11 = 6$$
,

$$17\% 11 = 6$$
,

$$h^+(x) = p - (x \mod p) = 5 - (17 \mod 5) = 3,$$

$$h_1(x) = (h(x) + 1h^+(x)) \mod 11 = (6+3) \mod 11 = 9.$$

$$39 \% 11 = 6$$
,

$$h^+(x) = p - (x \mod p) = 5 - (39 \mod 5) = 1,$$

$$h_1(x) = (h(x) + 1h^+(x)) \mod 11 = (6+1) \mod 11 = 7.$$

$$3\% 11 = 3$$
,

$$46 \% 11 = 2$$
,

$$h^+(x) = p - (x \mod p) = 5 - (46 \mod 5) = 4,$$

$$h_1(x) = (h(x) + 1h^+(x)) \mod 11 = (2+4) \mod 11 = 6,$$

$$h_2(x) = (h(x) + 2h^+(x)) \mod 11 = (2 + 8) \mod 11 = 1.$$

0	
1	46
2	35 3
2 3 4 5 6 7	3
4	
5	
6	28
7	39
9	
9	17
10	43

Closed hash table with double hashing

- 6. If a set of 4090 records is being stored using a binary tree T with 4090 nodes (one record per node), answer the following questions with *integer solution* if possible.
 - (a) What is the min height of T?
 - (b) What is the max height of T?
 - (c) What is the min number of leaves in T?
 - (d) What is the max number of leaves in T?
 - (e) If T is being implemented using the sequential array data structure, what is the size of an array A in order to store T?

Solution:

$$4090 = 2^{12} - 6$$
.

(a) What is the min height of T?

$$\lfloor \operatorname{lgn} \rfloor = 11$$

(b) What is the max height of T?

$$n - 1 = 4089$$
.

(C) What is the min number of leaf in T?

(d) What is the max number of leaf in T?

Height h = 11 implies tree with $2^{11+1} - 1 = 4095$ nodes and $2^{11} = 2048$ leaves.

Remove 5 leaves at level h and add back 2 new leaves at level h-1:

$$2048 - 5 + 2 = 2045$$
 leaves.

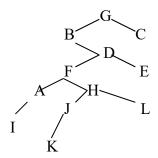
(e) If T is being implemented using the sequential array implementation, what is the size of array A in order to store T?

$$2^{4089+1} - 1 = 2^{4090} - 1$$

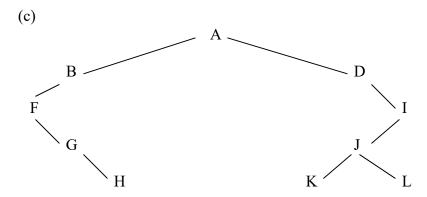
- 7. Construct the (unique) binary tree corresponding to the given pair of tree traversals if possible. **Remark:** You must show all your steps clearly as illustrated in class for credit. If no such a tree is possible, you must justify your answer.
 - (a) Preorder: GBDFAIHJKLEC Inorder: BIAFKJHLDEGC
 - (b) Postorder: HIBCAKGEJDF Inorder: IHCBKGJEFDA
 - (c) Postorder: H G F B K L J I D A Inorder: F G H B A D K J L I

Solution:

(a)



(b) No tree.



8. Given a set S of 4 records with keys $\{x_1, x_2, x_3, x_4\}$, $x_1 < x_2 < x_3 < x_4$. Construct all possible binary search trees (BST) that can be used to store S. Remark: You must illustrate all your BSTs clearly for credit.

Solution:

Number of BST = $C_{n+1} = \frac{1}{n+1} {2n \choose n}$, which is the (n+1)th-term Catalan number.

For
$$n = 4$$
, $C_5 = \frac{1}{5} {8 \choose 4} = 14$.

Distinct BST's for 4 records:

1:



2.



3:



1.



5:



6:



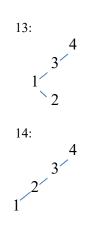


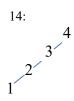












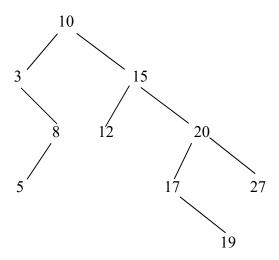
- 9. Given a set of 10 records with priorities $S = \{10, 15, 3, 8, 20, 5, 17, 27, 19, 12\}.$
 - (a) Construct a BST T for S by inserting the records, in the given order, into an initially empty binary search tree. When done, delete 10, and then 20 from the tree.
 - (b) Construct a BST T for S by inserting the records, in the reverse given order, into an initially empty binary search tree. When done, delete 5, and then 19 from the tree.

Remark: You must show your BST after each insertion/deletion for credits. For deletion, you must use deleteMin operations as discussed in class.

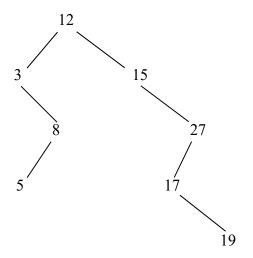
Solution:

(a)

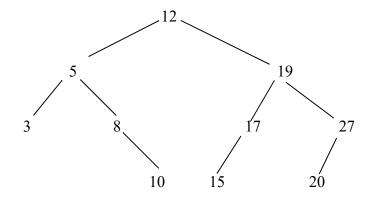
After insertions:



After deletions:



(b) After insertions:



After deletions:

