KU | Fall 2018 | Drew Davidson

# EECS 665 COMPILER CONSTRUCTION

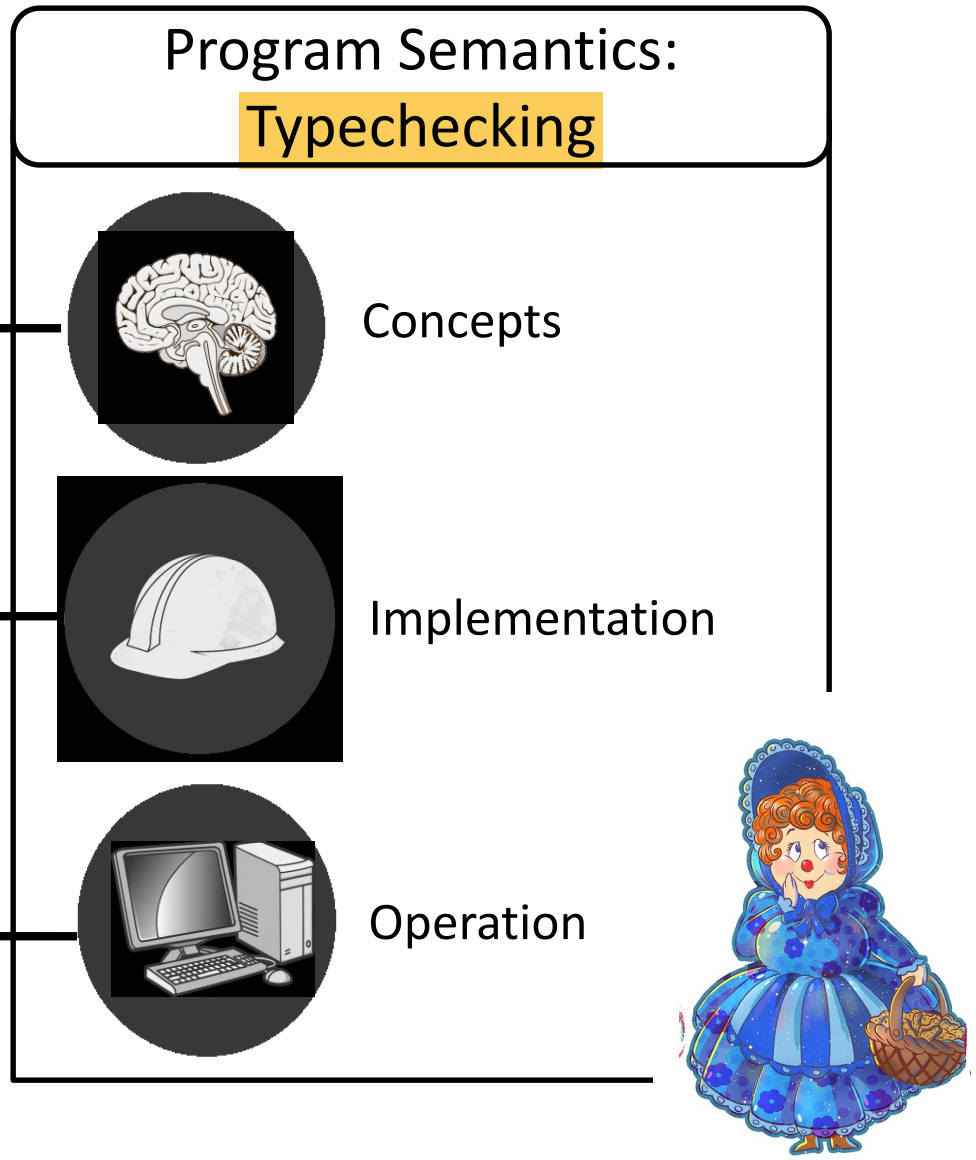
## 22 – AST Transformation

# Last Lecture

- What it is
- Why its done

- Translation rules on node type

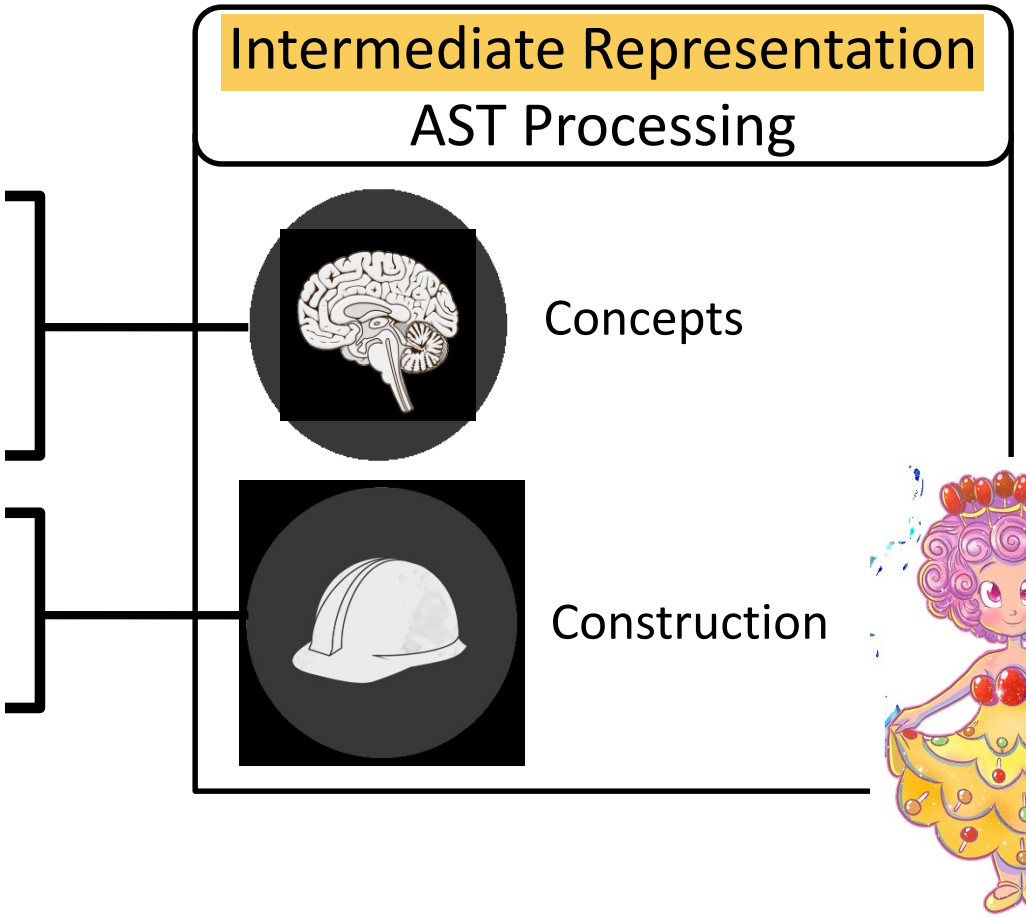
- Bottom-up / Inside-out AST traversal



# This Lecture

- What IRs are
- Why IRs are used
- Some IR forms

- IR Rewriting



# What IRs are

*Intermediate Representations*

Src->IR->Target

- Such a big, basic concept that its hard to get a precise definition
  - "An encoding of a program"
  - "What is output by the frontend of the compiler and processed by the backend of the compiler"

Good Defintion

  - "A representation of what a compiler knows about a program"
  - "A simpler language to which the source language is mapped"

# Why use Intermediate Representations?

Basic answer: help the compiler achieve its goals

… but why use a compiler?

# Yeah, why DO we need a Compiler?

- The obvious answer: PL implementation
  - Avoid dealing with target language directly
- By strict definition of “compiler”, other options may exist

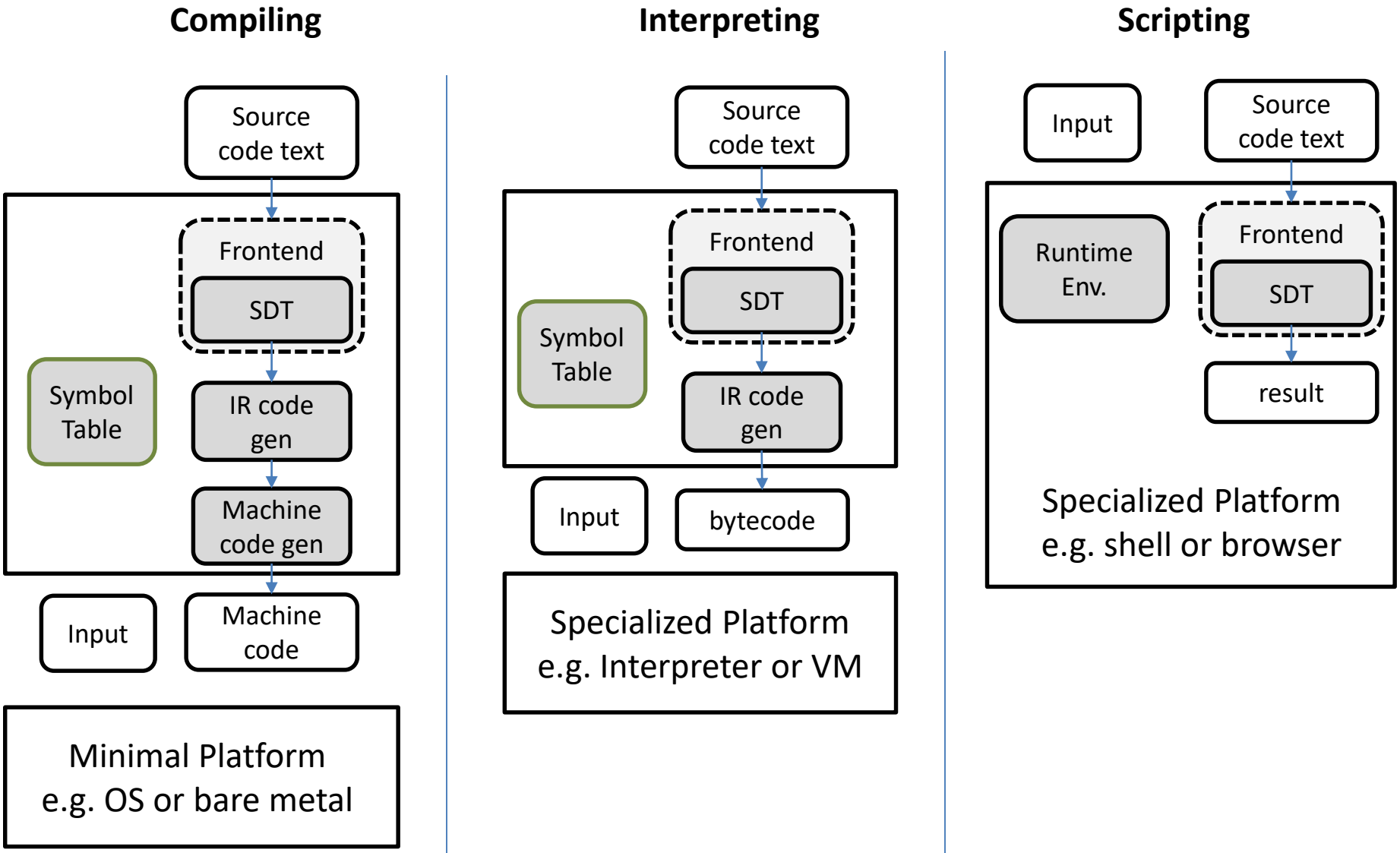
com·pil·er

/kəmˈpīlər/

*noun*

1. a person who produces a list or book by assembling information or written material collected from other sources.
   "this passage was revised in different ways by later compilers"
2. COMPUTING
   a program that converts instructions into a machine-code or lower-level form so that they can be read and executed by a computer.
   "conversion would require more than just running it through a different compiler"

# “Alternatives” to “Compiling”

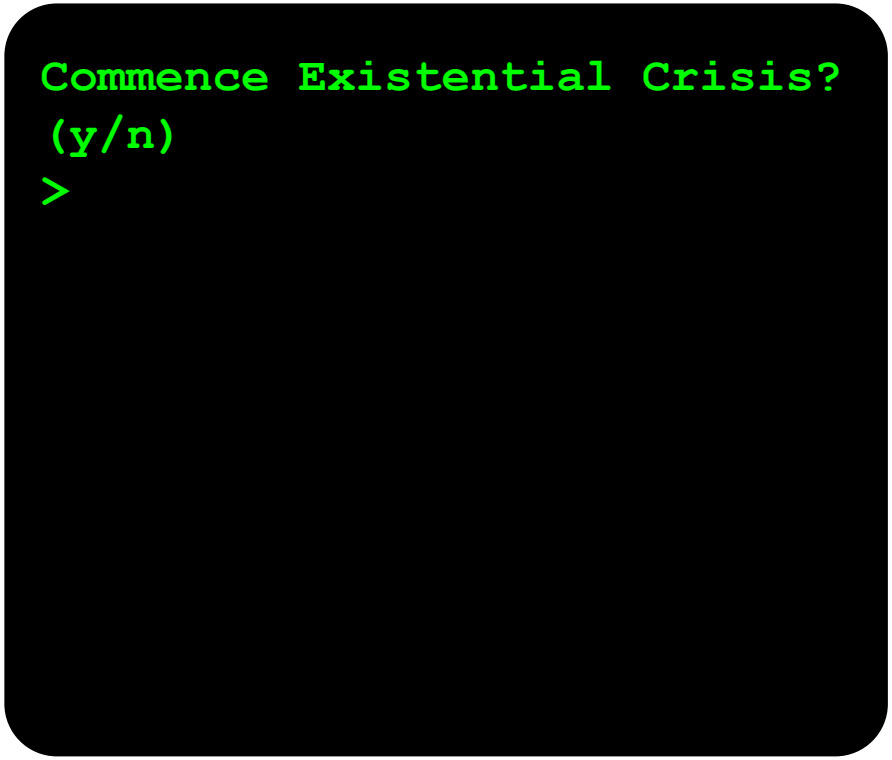


# So Compiling is Optional???

- In some contexts, no
  - If you have no runtime support, you're compiling (or coding ASM)
- In some contexts, yes
  - Interpreted languages might forgo machine code
    - May evaluate IR directly
  - Scripts might forgo source translation entirely
    - Evaluate source code directly

# Then Why Compile at All?!?!

- Abstraction
  - Don't have to deal with the target language
- Analysis
  - Error checking: Predict bugs before they strike
  - Optimization: Generate more efficient code

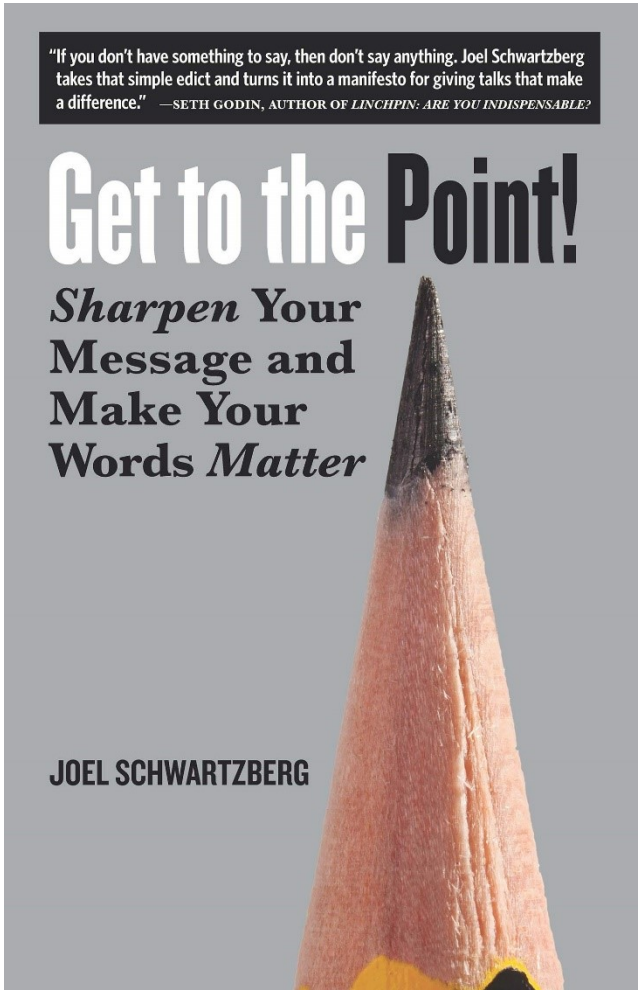


# Cool… Why did we detour?

- Minor: contextualize interpreted and scripting languages
- Major: The goals of the IR serve the goals of the compiler

Aid Abstraction

Aid Analysis



# Why Use Intermediate Representations?

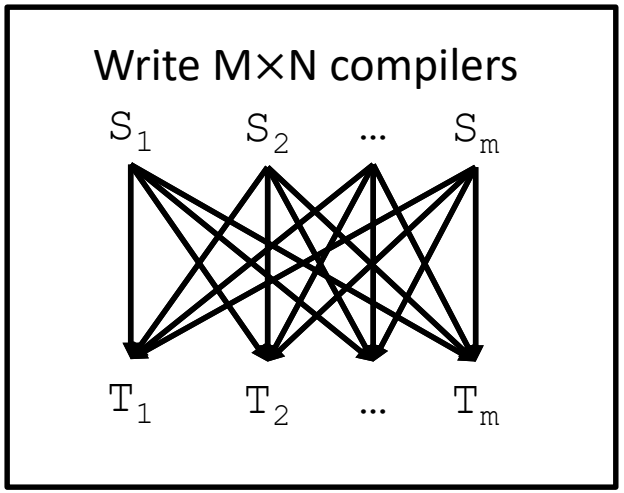
- Abstraction

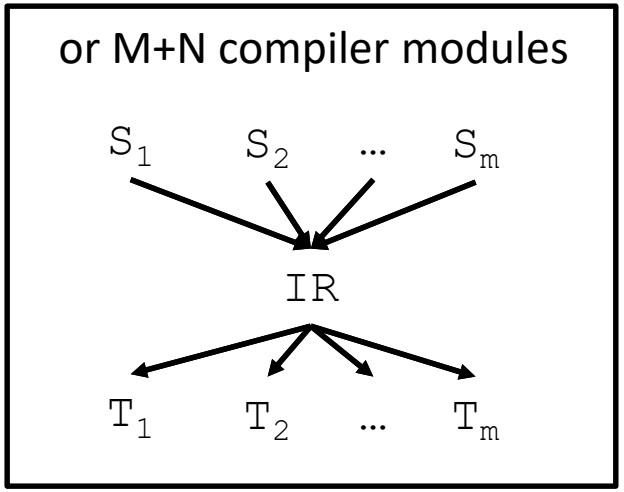
- Analysis

# Why Use Intermediate Representations?

- Abstraction
  - Decouple compiler frontend from backend
- Analysis

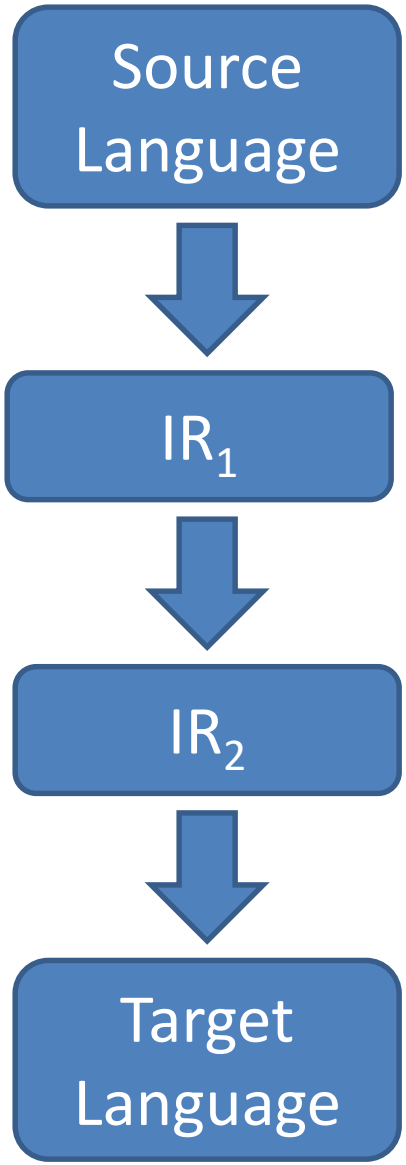
M source languages
N target languages



# Why Use Intermediate Representations?

- Abstraction
  - Decouple compiler frontend from backend
  - Break down source language constructs into target language over several steps
- Analysis



# Why Use Intermediate Representations?

- Abstraction
  - Decouple compiler frontend from backend
  - Break down source language constructs into target language over several steps
- Analysis
  - Optimize programs

Improve…

- Runtime
- Memory usage
- Power usage
- Security

# Why Use Intermediate Representations?

- Abstraction
  - Decouple compiler frontend from backend
  - Break down source language constructs into target language over several steps
- Analysis
  - Optimize programs
  - Predict faults

For example…

- typechecking

But isn't this an analysis on the AST?

# Why Use Intermediate Representations?

- Abstraction
  - Deco
    front
  - Brea
    langu
    targe
    seve
- Analys
  - Opti
  - Pred

For example

sis

**ASTs are an example of an IR!!**

We've been talking about IR for a week!!


*M. Night Shyamalan, famous for (ill-considered) plot twists in movies he writes/directs*

# Summary

Intermediate representations (like the AST) are a compiler's "internal formats" for the program, aid in abstraction and analysis

**Subtopics**

Alternatives to compilation | IR Goals Overview | Abstraction | Analysis

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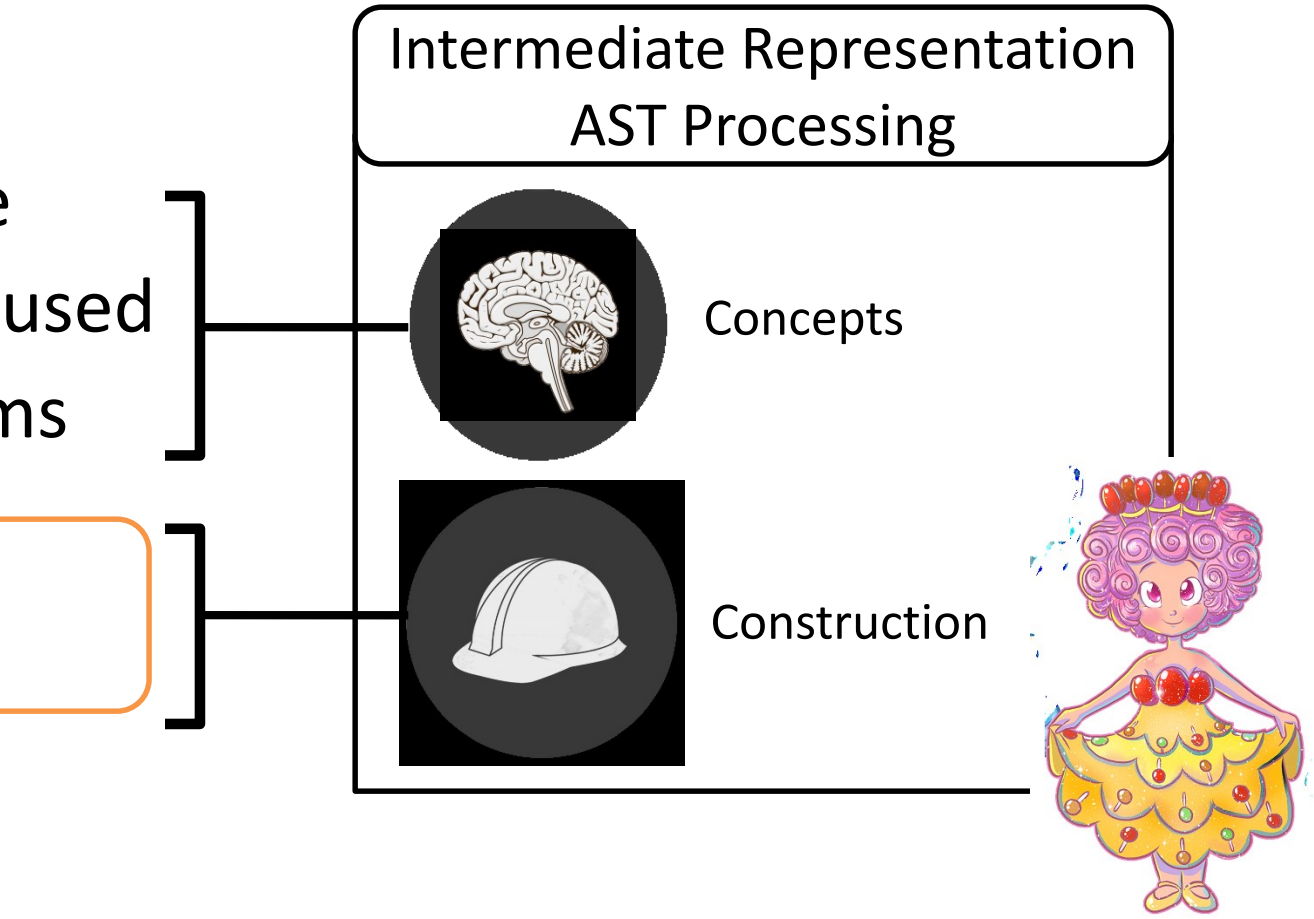

# AST Transformations

What Optimizations can we do on the AST?

- Function inlining
  - Replace the function call with the entire body of the function
- Constant folding
  - Replace computations with their results, if statically available

```
def jabber(int a){
  print a + 2;
}
def jibber(){
  jabber(2);
}
```

```
def jabber(int a){
  print a + 2;
}
def jibber(){
  print 2 + 2;
}
```

# Function Inlining: Why/Why not?

Why

- *Might* simplify code
- Function calls are (relatively) expensive

Why Not

- Might *not* simplify code
- Likely to increase overall code size

# Function Inlining: How

pros and cons of function inling

- Place the function body into the caller's context
- Disambiguate name clashes
- Replace parameter passing with assignments
- Replace function return with assignments
- Remove function definition ?

```
int g(int b){
int a;
print(a+b)
}
int f(){
int a =4;
g(a)
}
```

function inline clashes

```
int f(){
int a =4;
int a;
print(a+b);
}
```

# AST Transformations

What Optimizations can we do on the AST?

- Function inlining
  - Replace the function call with the entire body of the function
- Constant folding
  - Replace computations with their results, if statically available`

```
def doIt(){
  print 2 + 2;
}
```

```
def doIt (){
  print 4;
}
```

# Summary

The AST can be a good target for analysis such as function inlining, typechecking, and constant folding

**Subtopics**

Pos/Neg/:-O comments  Lab Sucks  Q&A issues  Me  Suggestions/Thoughts

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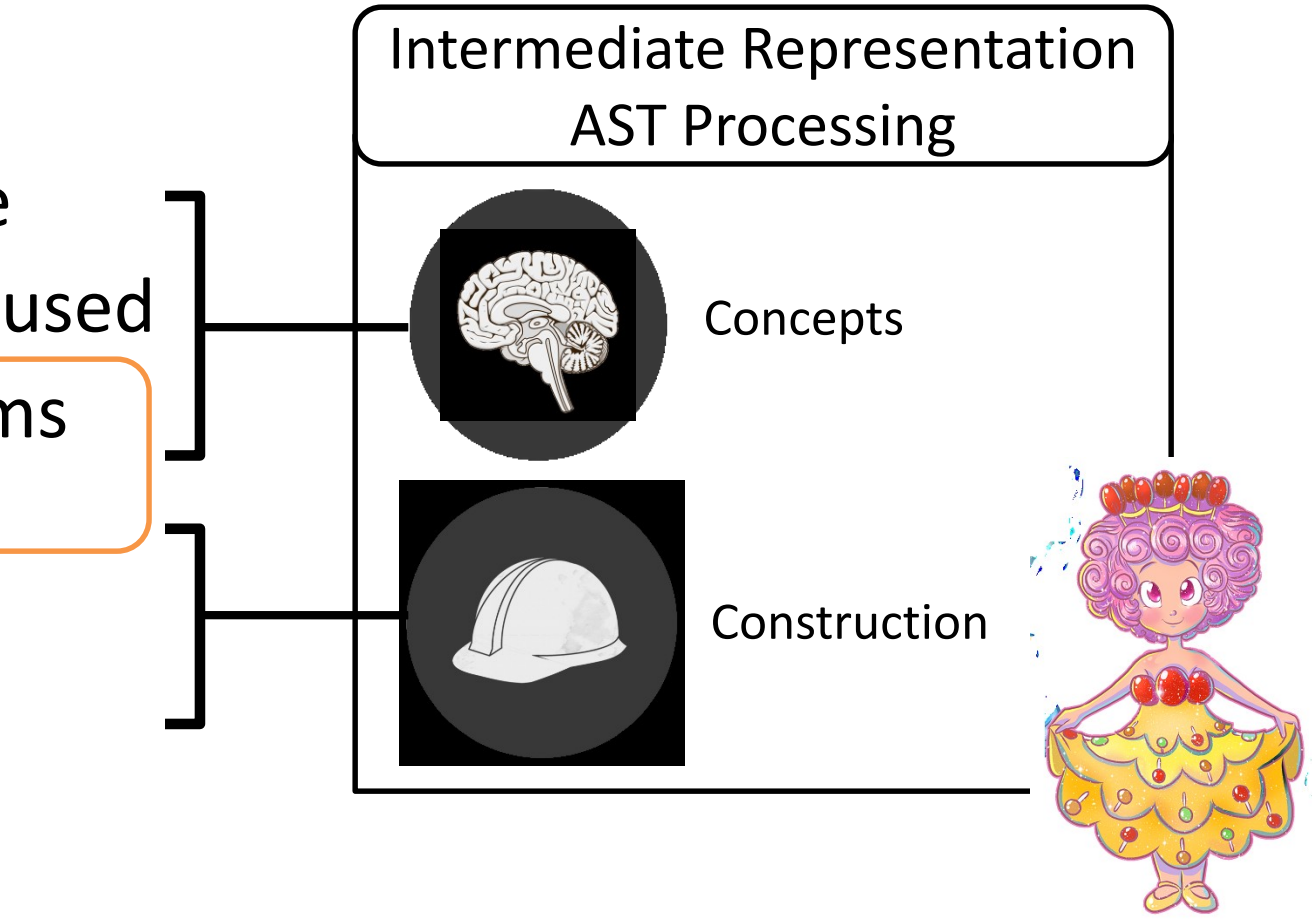
Other



# Other IR Forms: Why?

- AST is great for some things, but not everything
  - Doesn't represent control flow very well
- (NB: compilers could go directly from AST to machine code)

# Other IR Forms: Examples

- Structural
  - Abstract-Syntax Tree (AST)
  - Abstract Syntax DAG
- Linear
  - Three-Address Code (3AC)
  - Stack machine code
- Hybrid
  - Control-Flow Graph

Next time

Next next time