

— Predicates Q (ambiguous / unspecified)

↳ Statement (unambiguous / specified) by specifying the values of the variables $Q(x)$

↳ Statement (quantifier \exists, \forall) $\exists x. Q(x)$ $\forall x. Q(x)$

Statements \Rightarrow Compound Statements

1. Negation

2. Conjunction

3. Disjunction

4. Condition

↳ truth table

$p \rightarrow q$	$p \wedge q$	$p \vee q$
T	T	T
F	F	T
T	F	F
F	F	F



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— Some statement can be difficult to interpret

p : We are going to eat McDonalds

q : We are going to eat Pizza Hut

$\neg(p \vee q)$: It is not true that we are going to eat
M or P today.

$\neg p \wedge \neg q$: We are not going to eat M, nor P today.

— Different ways of presenting a statement must have the same meaning.

P	q	$P \vee q$	$\neg(P \vee q)$	$\neg P$	$\neg q$	$\neg P \wedge \neg q$
T	T	T	F	F	F	F
T	F	T	F	F	T	F
F	T	T	F	T	F	F
F	F	F	T	T	T	T

"For any combination of values of the elementary statements, the values for the compound statements expressed in different ways are identical."

Logical Equivalent " \equiv "
 $\neg(P \vee q) \equiv \neg P \wedge \neg q$

1. Commutative Law $p \vee q \equiv q \vee p$; $p \wedge q \equiv q \wedge p$

2. Associative Law $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$; $(p \vee q) \vee r \equiv p \vee (q \vee r)$

3. Distributive Law

4. Identity Law

5. Negation Law

6. Double negation law

7. Universal bound law

8. Idempotent law.

9. de Morgan's Law

10. Absorption law $p \vee (p \wedge q) \equiv p$; $p \wedge (p \vee q) \equiv p$

11. Negations of tautology and contradiction.



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- Statements:

Predicates Q
(parameters/
quantifiers)

Statements
 $Q(x)$

?

Argument form: a sequence of Statement form.

\approx

Statement form $Q \wedge P$

"de-parameterize"

Compound Statements
 $Q(x) \wedge P(x)$

$\neg, \vee, \wedge, \rightarrow$

Premise: all Statement/Statement forms
in argument except the last one.

Conclusion: The last Statement/Statement form
in an argument.

$$\frac{p \rightarrow q, q \rightarrow r, p, \text{there } q, q, \text{therefor } r.}{p \quad p \quad p \quad p \quad p \quad c}$$



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- Showing an argument form is correct (or not)
Using Truth Table.

Example: Show that $p \rightarrow q, p$, then q is a correct argument form

$p \quad q$		$p \rightarrow q \quad p$		q
elementary Statements		premises		Conclusion
T	T	T	T	T
T	F	F	T	F
F	T	T	F	T
F	F	T	F	F

"If, and only if, no matter what values are substituted for the statement variables in the premises, if all premises are true, then the conclusion is also true."



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- Showing an argument form is valid.

Example: Show that " $(p \vee q) \rightarrow r, (p \vee q), \text{ then } r$ " is valid

P	q	r	$p \vee q$	$(p \vee q) \rightarrow r$	$(p \vee q)$	r
E			P			C
T	T	T	T	T	T	T
T	F	T	T	T	T	T
F	T	T	T	T	T	T
F	F	T	F	T	F	T
T	T	F	T	F	T	F
T	F	F	T	F	T	F
F	T	F	F	T	F	F
F	F	F	F	T	F	F

Example: Show that $p \rightarrow q, q, \text{ then } p$ is invalid.

P	q	$p \rightarrow q$	q	p
E			P	C
T	T	T	T	T
T	F	F	F	T
F	T	T	T	F
F	F	T	F	F



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1. Modus Ponens
2. Modus Tollens
3. Generalization
4. Specialization
5. Elimination
6. Transitivity
7. Proof by division into cases.

Sound argument:

if all premises are true
and the argument form is valid

Premises \checkmark Argument form \times

\checkmark P: Tom is a vegetarian.

\times Q: Tom did not eat meat this Friday.

\checkmark $P \rightarrow Q$: if Tom is a vegetarian, he did not eat meat this Friday.

\times if $P \rightarrow Q, Q, P$.

Argument form \checkmark

Premises \times

\checkmark P: "John is a rock star"

\times Q: "John has red hair"

\times $P \rightarrow Q$

\times if $P \rightarrow Q, P, Q$.

" If John is a rock star, John has red hair. John is a rock star, so

John has red hair."