Instruction: You must show all your work clearly for credit. Partial credit will only be given to meaningful answers.

- 1. When implementing an ADT for a set of records S, $|S| = 2^6$, it is determined that a find operation, find(x,S), will requires 0.5ms (10^{-3} s) to execute. If the complexity of the find operation is given by the following closed-form expressions T(n), compute the time required to execute this operation when $|S| = 2^{16}$.
 - (a) T(n) = 560.
 (b) T(n) = nlgn.
 (c) T(n) = n²lgn.
 - (d) $T(n) = n^3$.
- 2. If an algorithm requires 0.5ms to solve a problem with input size of 100, how large a problem it can solve in 1 min if the complexity of the algorithm is given by the following function T(n) in closed-form?
 - (a) T(n) = n. (b) $T(n) = n^2$.
- 3. Given the following algorithm for finding the two largest integers in an array A[1..n] of n distinct positive integers. Base on the number of comparisons between elements in A, compute T_b(n) and T_w(n). You must justify your answer and show your work clearly for credit.

```
if A[1] > A[2]
                                      // Initialization
  then largest = A[1];
         s\_largest = A[2]
  else largest = A[2];
         s\_largest = A[1]
endif:
for i = 3 to n do
                                      // Checking A[3], ..., A[n]
                                      // A[i] is one of the two largest integers
  if A[i] > s\_largest
                                      // A[i] is the current largest integer
     then if A[i] > largest
                then s\_largest = largest;
                      largest = A[i]
                     s \ largest = A[i]
               else
           endif
  endif
endfor;
```

4. Assuming that all basic operations require the same constant cost C, by concentrating on the dominating step(s), compute the cost of the resource function R(n) for the following program segment in closed-form.

```
x = 2;

y = 10;

for i = 1 to n do

for j = i to n do

y = x * y / 2;

endfor;

for k = 1 to n do

x = x + y - 10;

endfor;

endfor;
```

5. By concentrating on the dominating step and by assuming that all basic operations require the same constant cost C, compute T_W(n) in closed-form for the following program segment as discussed in class.

Remark: You must first set up the equation for $T_W(n)$ and then evaluate its sums for credits. Do not simplify the final expression.

```
x = 2;
y = 10;
k = 1;
while k \le n do
     x = x + x*y + 210;
     y = y - x + 560;
     k = k+1;
endwhile;
for i = 1 to n do
 for j = i to n do
      y = x * y / 2;
      for k = i to n do
          x = x + y - 10;
      endfor;
  endfor;
endfor;
```

- 6. Let A_1 and A_2 be two algorithms with closed-form complexity $T_1(n) = 10n^2$ and $T_2(n) = 499n + 50$. Find smallest integer n_0 such that for all $n > n_0$, algorithm A_2 will always be more efficient than algorithm A_1 .
- 7. Use the definition of big-O to prove or disprove that $2^{2n} = O(3^n)$.
- 8. Prove or disprove that if $T_1(n) = O(f(n))$ and $T_2(n) = O(f(n))$, the $T_1(n) + T_2(n) = O(f(n))$.

- 9. Prove or disprove that if $T_1(n) = O(f(n))$ and $T_2(n) = O(f(n))$, then $\frac{T_1(n)}{T_2(n)} = O(1)$.
- 10. Using the definition of big-O to prove that $\frac{n^4 n^3 2n^2 + 4}{2n^2 2n 27} = \Omega(n^2).$
- 11. Use the definition of big- Θ to prove that $\frac{2n^4 n^3 5n^2 + 4}{n^2 6n + 7} = \Theta(n^2).$