Answer Set Programming¹

Abdallah Saffidine

COMP4418

¹Slides designed by Christoph Schwering





 $\forall x (\operatorname{Car}(x) \to \neg \operatorname{Entry}(x))$



$$\forall x (\operatorname{Car}(x) \to \neg \operatorname{Entry}(x)) \forall x (\operatorname{Car}(x) \land \operatorname{Auth}(x) \to \operatorname{Entry}(x))$$



$$\forall x (\operatorname{Car}(x) \to \neg \operatorname{Entry}(x)) \forall x (\operatorname{Car}(x) \land \operatorname{Auth}(x) \to \operatorname{Entry}(x))$$

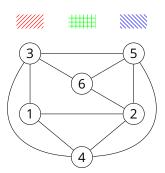
ASP at a Glance

- ASP = Answer Set Programming
 - ightharpoonup ASP \neq Microsoft's Active Server Pages
- ASP belongs to logic programming
 - ▶ Like Prolog: $Head \leftarrow Body$ or Head : Body.
 - Like Prolog: negation as failure
 - ▶ Unlike Prolog: Head may be empty \Rightarrow constraints
- Declarative programming
 - Unlike Prolog: no procedural control
 - Order has no impact on semantics
- ASP programs compute models
 - Unlike Prolog: not query-oriented, no resolution
 - Unlike Prolog: not Turing-complete
 - Tool for problems in NP and NP^{NP}

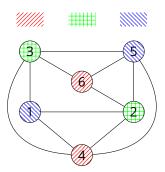
Motivation for ASP and this Lecture

- Very useful in practice!
 - Declarative problem solving
 - Very fast to write
 - Very fast to run
 - Few experts
- Interesting case study
 - Small, simple core language
 - Great expressivity by reduction to core language
- Knowing the theory is essential

Definition: graph colouring problem



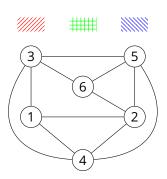
Definition: graph colouring problem



Definition: graph colouring problem

- Graph Coulouring is NP-complete
 - ▶ NP: guess solution, verify in polynomial time
 - NP-complete: among hardest in NP
- Many applications:
 - Mapping (neighbouring countries to different colors)
 - Compilers (register allocation)
 - Scheduling (e.g., conflicting jobs to different time slots)
 - Allocation problems, Sudoku, ...

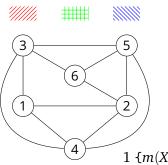
Definition: graph colouring problem



$$\begin{array}{l} c(r) \cdot c(g) \cdot c(b) \cdot \\ \nu(1) \cdot \nu(2) \cdot \nu(3) \cdot \nu(4) \cdot \nu(5) \cdot \nu(6) \cdot \\ e(1,2) \cdot e(1,3) \cdot e(1,4) \cdot \\ e(2,4) \cdot e(2,5) \cdot e(2,6) \cdot \\ e(3,4) \cdot e(3,5) \cdot e(3,6) \cdot \\ e(4,5) \cdot \\ e(5,6) \cdot \end{array}$$

Definition: graph colouring problem

Input: graph with vertices V and edges $E \subseteq V \times V$, set of colors C. Is there a mapping $m: V \to C$ with $m(x) \neq m(y)$ for all $(x,y) \in E$?

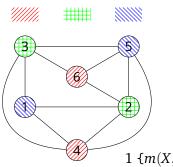


$$c(r) \cdot c(g) \cdot c(b) \cdot v(1) \cdot v(2) \cdot v(3) \cdot v(4) \cdot v(5) \cdot v(6) \cdot e(1,2) \cdot e(1,3) \cdot e(1,4) \cdot e(2,4) \cdot e(2,5) \cdot e(2,6) \cdot e(3,4) \cdot e(3,5) \cdot e(3,6) \cdot e(4,5) \cdot e(5,6) \cdot e$$

 $1 \{m(X,C) : c(C)\} \ 1 := v(X)$. guess mapping m: = e(X,Y), m(X,C), m(Y,C). verify $m(X) \neq m(Y)$

Definition: graph colouring problem

Input: graph with vertices V and edges $E \subseteq V \times V$, set of colors C. Is there a mapping $m: V \to C$ with $m(x) \neq m(y)$ for all $(x,y) \in E$?



$$c(r) \cdot c(g) \cdot c(b) \cdot v(1) \cdot v(2) \cdot v(3) \cdot v(4) \cdot v(5) \cdot v(6) \cdot e(1,2) \cdot e(1,3) \cdot e(1,4) \cdot e(2,4) \cdot e(2,5) \cdot e(2,6) \cdot e(3,4) \cdot e(3,5) \cdot e(3,6) \cdot e(4,5) \cdot e(5,6) \cdot e$$

 $1 \{m(X,C) : c(C)\} \ 1 := v(X)$. guess mapping m: = e(X,Y), m(X,C), m(Y,C). verify $m(X) \neq m(Y)$

Applications of ASP

- Automated product configuration
- Linux package manager
- Decision-support system for space shuttle
- Bioinformatics (diagnosis, inconsistency detection)
- General game playing
- Several implementations are available
- For this lecture: Clingo
 - https://potassco.org/
 - https://github.com/potassco/clingo/releases/
 - https://potassco.org/clingo/run/