

# Answer Set Programming

(4) ASP as modelling language

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COMP4418

# Overview of the Lecture

- Semantics of ASP programs
- Extensions of ASP programs
- Handling of variables in ASP
- **ASP as modelling language**

# ASP Modelling

$c(r). c(g). c(b).$   
 $v(1). \dots v(6).$   
 $e(1, 2). e(1, 3). e(1, 4).$   
 $e(2, 4). e(2, 5). e(2, 6).$   
 $e(3, 1). e(3, 4). e(3, 5).$   
 $e(4, 1). e(4, 2).$   
 $e(5, 3). e(5, 4). e(5, 6).$   
 $e(6, 2). e(6, 3). e(6, 5).$

Typical ASP structure:

- Problem **instance**: a set of facts
- Problem **class**: a set of rules

- ▶ Generator rules: often choice rules  $1 \{m(X, C) : c(C)\} 1 :- v(X).$
- ▶ Test rules: often integrity constraints  $:- e(X, Y), m(X, C), m(Y, C).$

Ideal modeling is **uniform**: problem class encoding fits all instances

Semantically equivalent encodings may differ immensely in performance!

## Example: Non-monotonic Reasoning

Tweety the penguin:

- (Normal) Birds fly.
- Penguins are abnormal.
- Tweety is a bird. So Tweety flies.
- Tweety is a penguin. So Tweety doesn't fly.

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$$U = \{f(X) \leftarrow b(X), \text{not } a(X). \quad a(X) \leftarrow p(X). \quad b(t).\}$$

$$P = \{f(t) \leftarrow b(t), \text{not } a(t). \quad a(t) \leftarrow p(t). \quad b(t).\}$$

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$$S_1 = \{b(t), f(t)\} \Rightarrow P^{S_1} = \{f(t) \leftarrow b(t), \text{not } a(t). \quad a(t) \leftarrow p(t). \quad b(t).\} \checkmark$$

$$S_2 = \{a(t), b(t), p(t)\} \Rightarrow P^{S_2} = \{f(t) \leftarrow b(t), \text{not } a(t). \quad a(t) \leftarrow p(t). \quad b(t).\} \times$$

Tweety flies!

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Tweety the penguin:

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$$U = \{f(X) \leftarrow b(X), \text{not } a(X). \quad a(X) \leftarrow p(X). \quad b(t).\}$$

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Tweety flies!

$$S_1 = \{b(t), f(t)\} \Rightarrow (P \cup \{p(t).\})^{S_1} = P_2^{S_1} \cup \{p(t).\} \quad \times$$

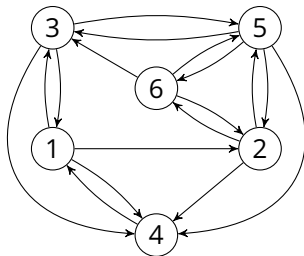
$$S_2 = \{a(t), b(t), p(t)\} \Rightarrow (P \cup \{p(t).\})^{S_2} = P_2^{S_1} \cup \{p(t).\} \quad \checkmark$$

Tweety doesn't fly.

# Example: Hamilton Cycle

## Definition: Hamilton cycle problem

Input: graph with vertex set  $V$  and edges  $E \subseteq V \times V$ .  
Is there a cycle that visits every vertex exactly once?

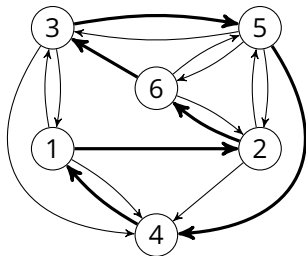




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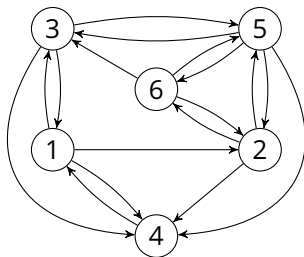
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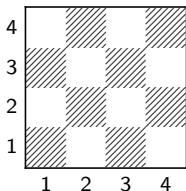
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$$\begin{aligned} \{p(X, Y)\} &\leftarrow e(X, Y). \\ r(X) &\leftarrow p(1, X). \\ r(Y) &\leftarrow r(X), p(X, Y). \\ &\leftarrow 2 \{p(X, Y)\}, v(X). \\ &\leftarrow 2 \{p(X, Y)\}, v(Y). \\ &\leftarrow \text{not } r(X), v(X). \end{aligned}$$

## Example: $N$ -Queens

### Definition: $N$ -queens problem

Place  $N$  queens on a  $N \times N$  chessboard so that they do not attack each other, i.e., share no row, column, or diagonal.

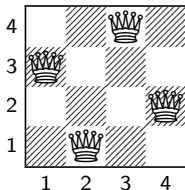


Program on paper

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Program on paper