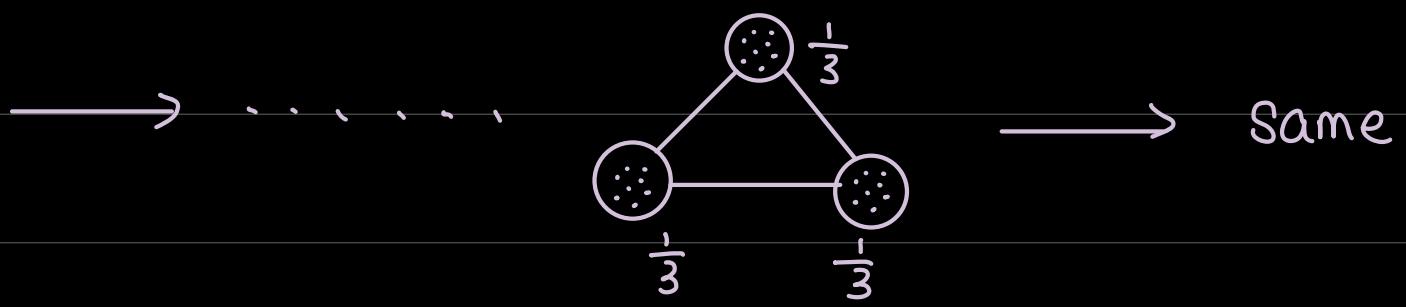
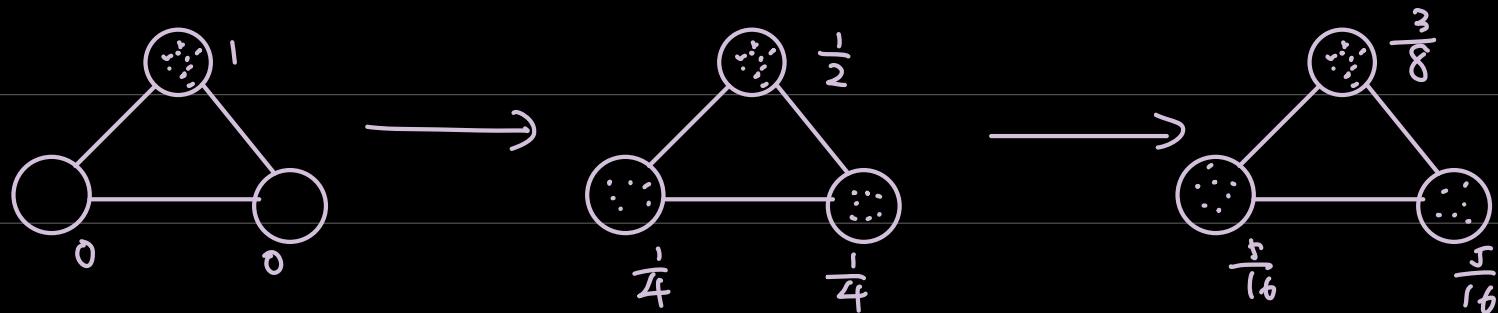


# Markov Chain

random walk in state space

1 million people



Population migration

probability  $\frac{1}{2}$  , stay. probability  $\frac{1}{4}$  to one other state .

Arrow of time , coverage to stationary distribution

Entropy (randomness) increases

Notation :

$X_t$  : State at time  $t$

$P^{(t)}$  : distribution of  $X_t$  (population distribution)

$P^{(t)}$	$P_1^{(t)}$	$P_2^{(t)}$	$P_3^{(t)}$	
$t=0$	1	0	0	
$t=1$	$\frac{1}{2}$	$\frac{1}{4}$	$\frac{1}{4}$	$P^{(t)} \xrightarrow[t]{\infty} \pi$
$t=2$	$\frac{3}{8}$	$\frac{5}{16}$	$\frac{5}{16}$	"arrow of time"
$\vdots$	$\vdots$	$\vdots$	$\vdots$	
$t \rightarrow \infty \pi$	$\frac{1}{3}\pi_1$	$\frac{1}{3}\pi_2$	$\frac{1}{3}\pi_3$	

Key / driving force :

Transition      Probability

$$k_{ij} = P(X_{t+1} = j \mid X_t = i)$$

$$\text{or } [k(x,y) = P(X_{i+1} = y \mid X_i = x)]$$

transition matrix:

	$j$	1	2	3
$i$		$\frac{1}{2}$	$\frac{1}{4}$	$\frac{1}{4}$
1		$\frac{1}{4}$	$\frac{1}{2}$	$\frac{1}{4}$
2		$\frac{1}{4}$	$\frac{1}{4}$	$\frac{1}{2}$

# Master Equation

$$\begin{aligned}
 P_j^{(t+1)} &= P(X_{t+1} = j) \xrightarrow{\text{# of ppl inj at } t-1} \\
 &= \sum_i P(X_{t+1} = j \text{ and } X_t = i) \xrightarrow{\text{# inj at } t+1 \text{ and in } i \text{ at } t} \\
 &= \sum_i P(X_t = i) \xrightarrow{\text{# in } i \text{ at } t} P(X_{t+1} = j | X_t = i) \xrightarrow{\substack{\downarrow \\ \text{fraction of these in } i}} \\
 &= \sum_i P_i^{(t)} K_{ij} \xrightarrow{\substack{\text{who will go to } j}}
 \end{aligned}$$

So the master equation is:

$$P_j^{(t+1)} = \sum_i P_i^{(t)} K_{ij} \quad \text{eigen-analysis}$$

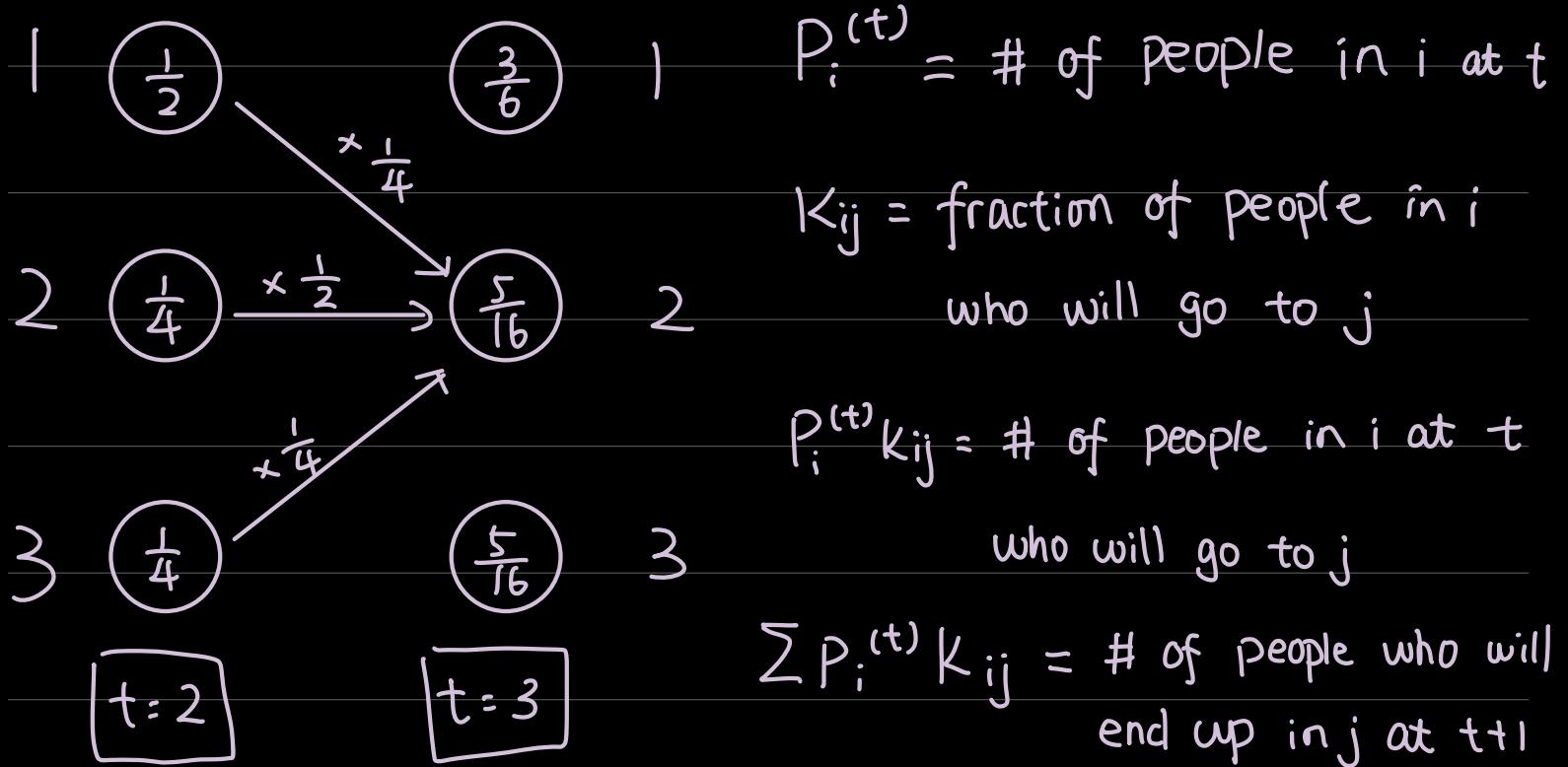
$$P^{(t)} \xrightarrow{K} P^{(t+1)}$$

$$P^{(t+1)} = P^{(t)} K \Rightarrow P^{(t)} = P^{(0)} K^t$$

$$\left[ \begin{array}{c|c|c} \frac{3}{8} & \frac{5}{16} & \frac{5}{16} \end{array} \right] = \left[ \begin{array}{c|c|c} \frac{1}{2} & \frac{1}{4} & \frac{1}{4} \end{array} \right]$$

$$\begin{matrix} \frac{1}{2} & \frac{1}{4} & \frac{1}{4} \\ \frac{1}{4} & \frac{1}{2} & \frac{1}{4} \\ \frac{1}{4} & \frac{1}{4} & \frac{1}{2} \end{matrix}$$

Interpret the Master equation



$\approx$  equilibrium

## Stationary $\pi$

$$\pi = \pi K$$

① condition

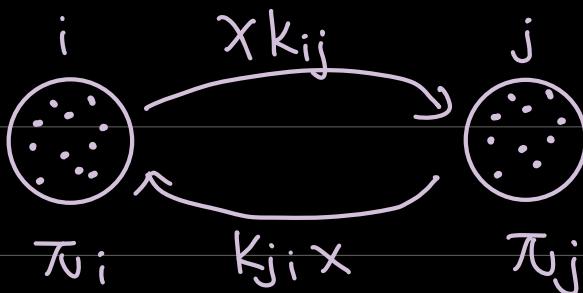
$$\pi_j = \sum_i \pi_i K_{ij} \rightarrow \text{Overall balance}$$

②

stronger condition :

Detailed balance  $\pi_i K_{ij} = \pi_j K_{ji}$  satisfies overall balance when it's satisfied.

$$\pi_i K_{ij} = \pi_j K_{ji}$$



Page Rank : Know  $K$ , want  $\pi$

Start from  $P^{(0)}$  (e.g. Uniform)

$$P^{(0)} \xrightarrow{K} P^{(1)} \xrightarrow{K} P^{(2)} \xrightarrow{K} \dots \xrightarrow{K} P^{(T)} \approx \pi$$

$K$  is a stochastic matrix, the sum of each row = 1

Why recursive : more stable, accurate measure of popularity

MCMC : know  $\pi$ , want  $K$

(target distribution)  $\uparrow$  (iterative algorithm)  $\downarrow$   
so that  $\pi = \pi K$  (design principle) to implement

Start from  $X_0 \sim P^{(i)}$

$$X_0 \rightarrow X_1 \rightarrow X_2 \rightarrow \dots \rightarrow X_t \rightarrow X_{t+1} \rightarrow \dots \rightarrow X_T$$

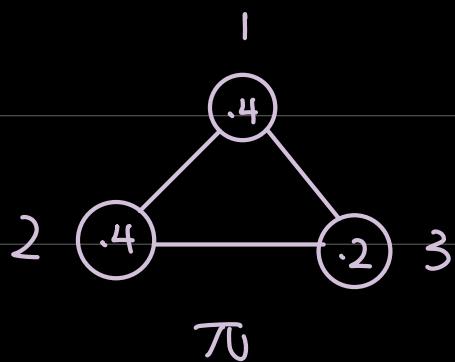
Approx:  $X_T \sim \pi$

# Metropolis algorithm

1950s

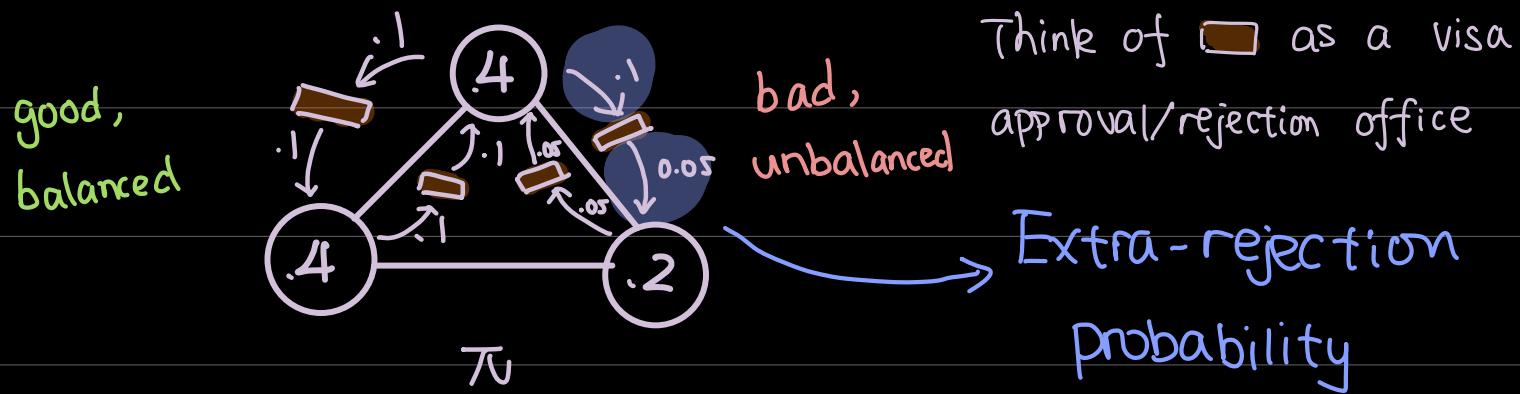
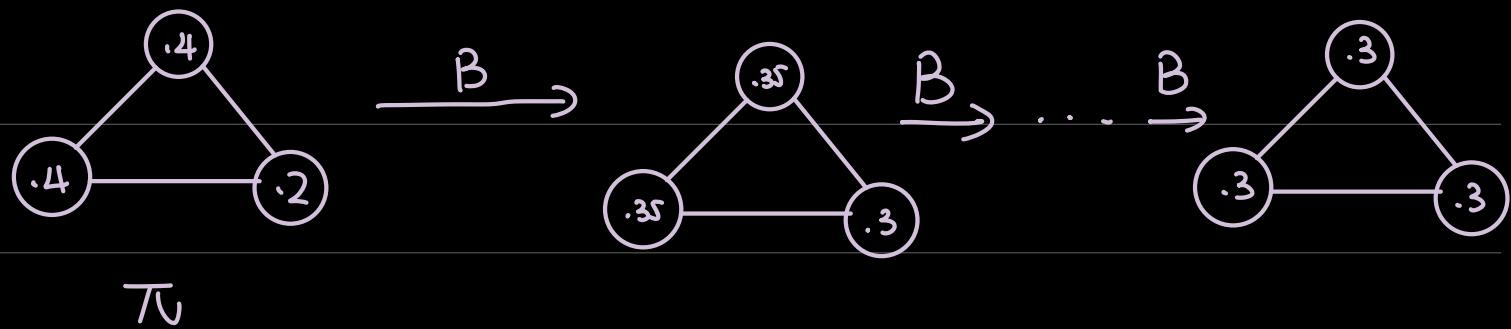
Metropolis , Roseblatt<sup>2</sup> , Teller<sup>2</sup>

Basic idea

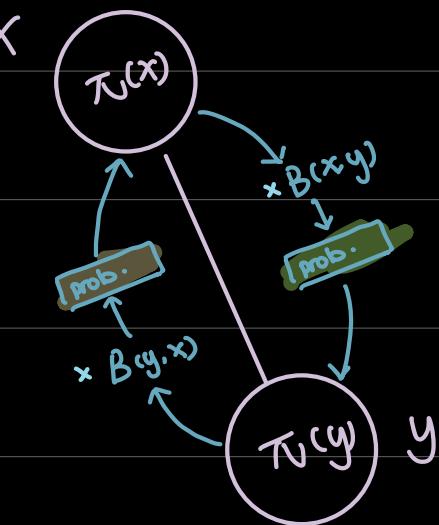


Base Chain : B

	1	2	3
1	$\frac{1}{2}$	$\frac{1}{4}$	$\frac{1}{4}$
2	$\frac{1}{4}$	$\frac{1}{2}$	$\frac{1}{4}$
3	$\frac{1}{4}$	$\frac{1}{4}$	$\frac{1}{2}$



General:



Algorithm :

(1) If  $\pi(x) B(x,y) > \pi(y) B(y,x)$

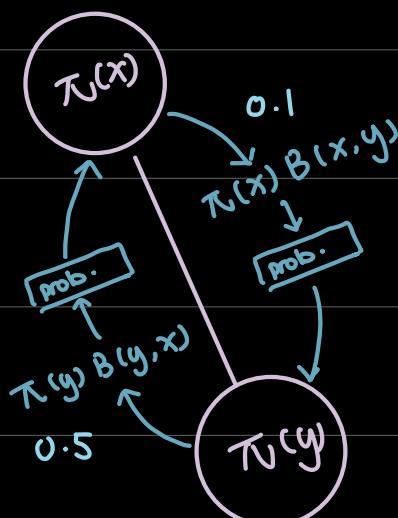
prob( accept  $x \rightarrow y$ ) =  $\frac{\pi(y) B(y,x)}{\pi(x) B(x,y)}$

(2) If  $\pi(x) B(x,y) \leq \pi(y) B(y,x)$

prob( accept  $x \rightarrow y$ ) = 1

(3) Combine (1), (2) :

prob( accept  $x \rightarrow y$ ) = min ( 1,  $\frac{\pi(y) B(y,x)}{\pi(x) B(x,y)}$  )



$$X_t = x$$

$$X_{\text{propose}} \sim B(x, y) = P(X_{\text{propose}} = y \mid X_t = x)$$

(let  $X_{t+1} = \begin{cases} X_{\text{propose}} & \text{with } P = \min(1, \frac{\pi(y)B(y,x)}{\pi(x)B(x,y)}) \\ X_t = x & \text{with } 1-P \end{cases}$ )

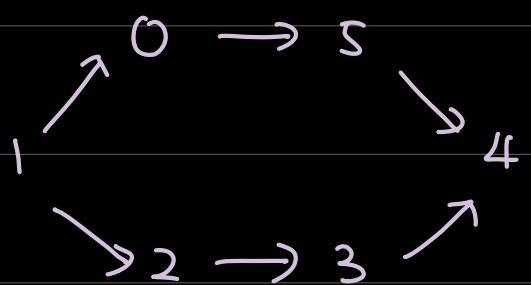
Special Case :

Metropolis algorithm :  $B(x, y) = B(y, x)$

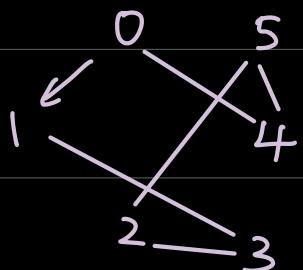
$$P = \min(1, \frac{\pi(y)}{\pi(x)})$$

Hastings Extension : general  $B$

Traveling Salesman Problem



$X$ : path  $0 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow 0$



another  $X$  :  $0 \rightarrow 1 \rightarrow 3 \rightarrow 2 \rightarrow 4 \rightarrow 5 \rightarrow 0$

$\mathcal{E}(x)$  = total length of path  $x$ .

$$\pi(x) = \frac{1}{Z} e^{-\mathcal{E}(x)}$$

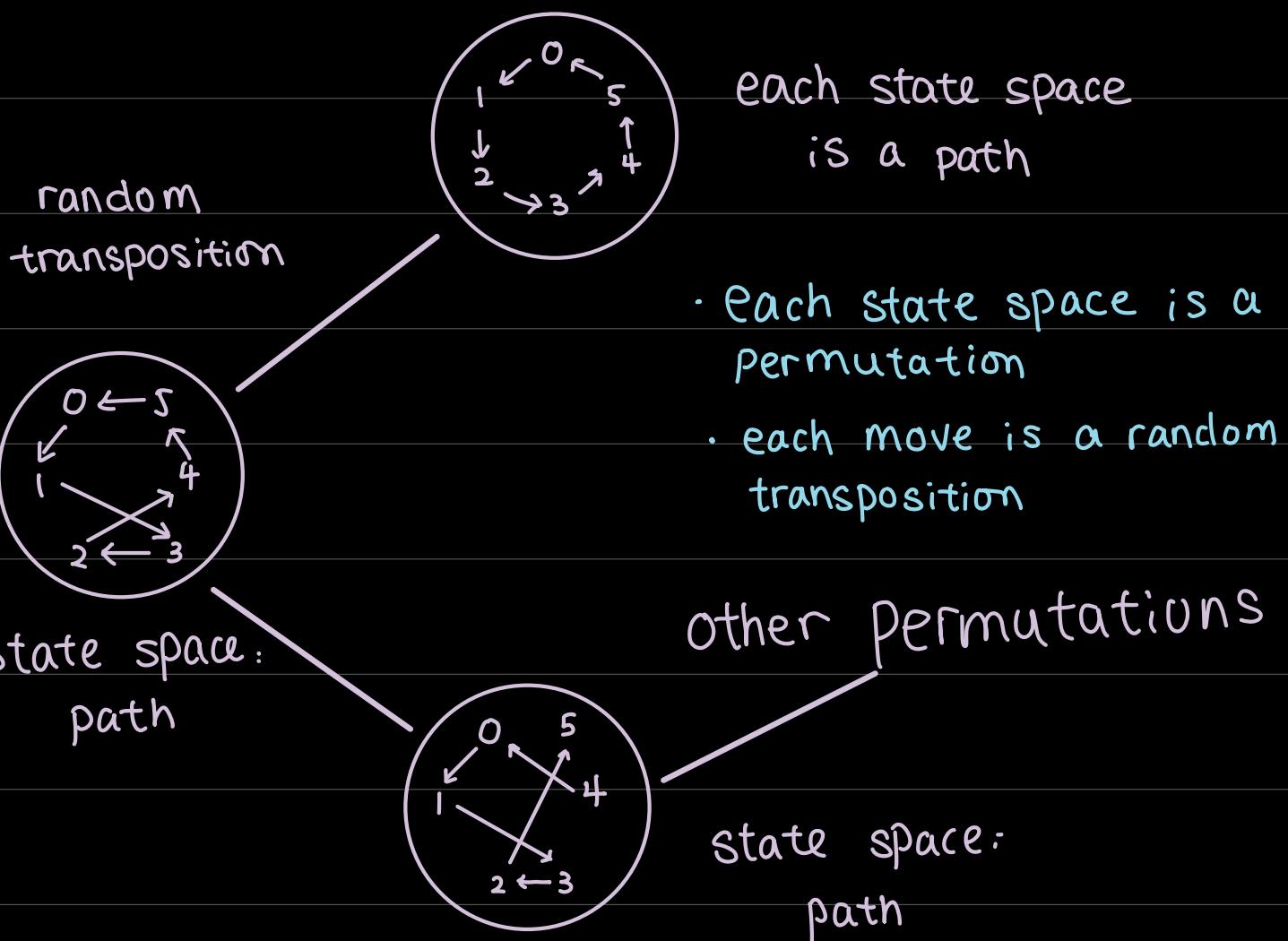
$$Z = \sum_x e^{-\mathcal{E}(x)}$$

Gibbs distribution / Boltzman distribution

Searching thru such huge space is extremely difficult

Sampling and optimization are related here

State Space (n factorial)



## Base Chain

$x \quad 1 \quad 2 \quad 3 \quad 4 \quad 5$



$y \quad 1 \quad 4 \quad 3 \quad 2 \quad 5$

$1 \quad 4 \quad 2 \quad 3 \quad 5$

.

:

:

Randomly pick 2 cities

& exchange their positions,

Propose change by random transposition.

accept change with

$$P = \min(1, \frac{\pi(y)}{\pi(x)})$$

$$= \min\left(1, e^{\frac{\epsilon(x) - \epsilon(y)}{k}}\right)$$

recall

here  $\frac{1}{k}$  gets cancelled

$$\pi(\cdot) = \frac{1}{Z} e^{-\epsilon(\cdot)}$$

$X_0 \rightarrow X_1 \rightarrow X_2 \rightarrow X_3 \rightarrow \dots$

(12345)                    (14325)                    (14235)                    (14235)

$\rightarrow X_{1000}$  we end up w/ a good permutation

Approx:  $X_{1000} \sim \pi(x)$