

# Quantum Imaginary-Time Evolution

Prepare ITE state with polynomial resources in time



Preprint  
(v3 updated soon)



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# Imaginary-time evolution (ITE)

**Problem** For Hamiltonian  $H$ , given initial state  $|\phi(0)\rangle$ , prepare the normalized solution of imaginary-time Schrödinger equation

$$\begin{aligned}\partial_\tau |\phi(\tau)\rangle &= -H |\phi(\tau)\rangle \text{ with } |\phi\rangle = |\phi(0)\rangle \\ \Rightarrow |\phi(\tau)\rangle &= e^{-\tau H} |\phi\rangle / \|e^{-\tau H} |\phi\rangle\|\end{aligned}$$

$\tau$

imaginary evolution time

$|\phi(\tau)\rangle$

normalized imaginary-time evolved state (ITE state)

$|\phi\rangle \rightarrow |\phi(\tau)\rangle$

imaginary-time evolution operator (ITE operator)

# Imaginary-time evolution

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## Application

- Ground state preparation  $|\psi_0\rangle \propto \lim_{\tau \rightarrow \infty} \mathbf{e}^{-\tau H} |\phi(0)\rangle$  if  $|\langle\psi_0|\phi(0)\rangle| > 0$
- Thermal state preparation [1]  $\rho = \mathbf{e}^{-\tau H} / \text{Tr}(e^{-\tau H})$
- Open-system simulation [2]  $e^{\mathcal{L}t}[\rho] = V\rho V^\dagger, V = \lim_{N \rightarrow \infty} (e^{-itH_1/N} \mathbf{e}^{-tH_2/N})^N$
- Others:
  - computing Hamiltonian-related properties [3]
  - data classification [4], optimization [5]
  - quantum mechanics [6], quantum field theory [7]

[1] Motta, Mario, et al. "Determining eigenstates and thermal states on a quantum computer using quantum imaginary time evolution." *Nature Physics* 16.2 (2020): 205-210.

[2] Kamakari, Hirsh, et al. "Digital quantum simulation of open quantum systems using quantum imaginary-time evolution." *PRX quantum* 3.1 (2022): 010320.

[3] Wang, Xiaoyang, et al. "Computing \$n\$-Time Correlation Functions without Ancilla Qubits." *Physical Review Letters* 135 (2025): 230602.

[4] Ye, Qi, et al. "Quantum automated learning with provable and explainable trainability." *arXiv preprint arXiv:2502.05264* (2025).

[5] Wang, Xiaoyang, et al. "Imaginary Hamiltonian variational Ansatz for combinatorial optimization problems." *Physical Review A* 111.3 (2025): 032612.

[6] Wick, Gian-Carlo. "Properties of Bethe-Salpeter wave functions." *Physical Review* 96.4 (1954): 1124.

[7] Lancaster, Tom, and Stephen J. Blundell. *Quantum field theory for the gifted amateur*. OUP Oxford, 2014.

# Existing works

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## Routine 1

Fix one  $|\phi\rangle$ . Find a physical unitary  $U$  such that  $U|\phi\rangle = |\phi(\tau)\rangle$ .

- ① (PQC scheme) find  $\theta^*$  such that  $U(\theta^*) \approx U$  [1]
- ② (Manifold scheme) find  $V \in \mathrm{SU}(2^n)$  such that  $V \approx U$  [2]
- ③ (Trotter scheme) compute  $\{(t_j, A_j)\}$  such that  $\prod_j e^{-iA_j t_j} \approx U$  [3]

- ✓ polynomial resource complexity in system size
- ✓ execute without failure
- ✗ one solution cannot be converted to the other that has unknown initial state
- ✗ unclear relation between  $\tau$  and precision, possibly an exponential dependency

[1] McArdle, Sam, et al. "Variational ansatz-based quantum simulation of imaginary time evolution." *npj Quantum Information* 5.1 (2019): 75.

[2] Gluza, Marek, et al. "Double-bracket quantum algorithms for quantum imaginary-time evolution." *arXiv preprint arXiv:2412.04554* (2024).

[3] Motta, Mario, et al. "Determining eigenstates and thermal states on a quantum computer using quantum imaginary time evolution." *Nature Physics* 16.2 (2020): 205-210.

# Existing works

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## Routine 2

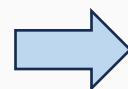
Unknown  $|\phi\rangle$ . Find a physical operation  $\mathcal{E}$  such that  $\mathcal{E}(\phi) = \phi(\tau)$  or FAILURE.

- ✓ polynomial resource complexity in system size
- ✓ theoretically analyzed; work for any initial state
- ✗ success probability decrease exponentially with  $\tau$

exponential decay problem

## Short Conclusion

- For long imaginary evolution, algorithms are either
  - heuristic or
  - theoretically infeasible



NEED ITE algorithm for  $\tau \gg 0$

# Transformation of quantum data

- A unitary  $U$  applies rotations with respect to its eigenspaces.

eigenphase: angle of the rotation

$$U = \sum_j e^{i\lambda_j} |\psi_j\rangle\langle\psi_j|$$

eigenvalue: amplitude of the rotation

eigenbasis: basis of the rotation

# Transformation of quantum data

- Transformation of  $U$  in terms of function  $f: \mathbb{R} \rightarrow \mathbb{C}$  can be described as

eigenphase: angle of the rotation

$$f(U) = \sum_j f(\lambda_j) |\psi_j\rangle\langle\psi_j|$$

eigenvalue: amplitude of the rotation

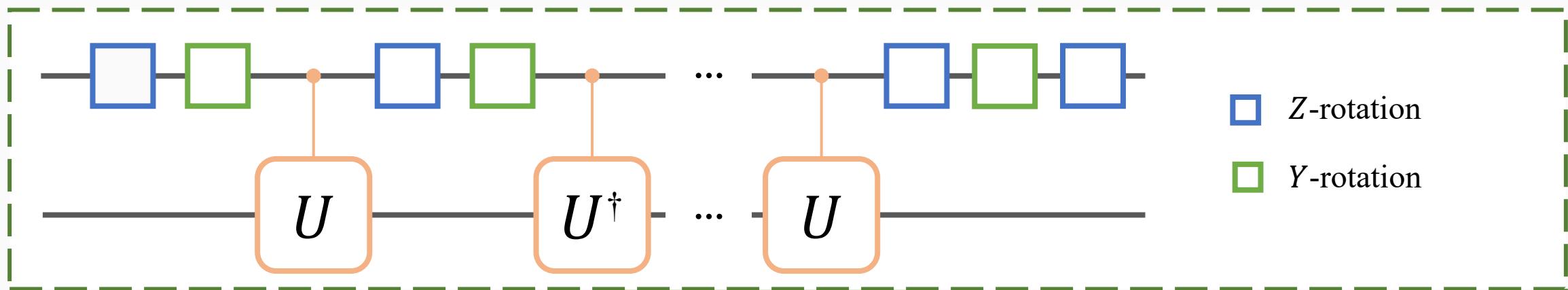
eigenbasis: basis of the rotation

- Quantum algorithms using unitaries deal with eigen-information procession.
  - Extract the target eigenphase/value – Shor, HHL, QAE
  - Amplify the target eigenphase/value – Hamiltonian simulation, Grover

# Transformation of quantum data

## Core Idea 1

Use quantum phase processing (QPP) [1], a QSP-based structure to simulate transformation



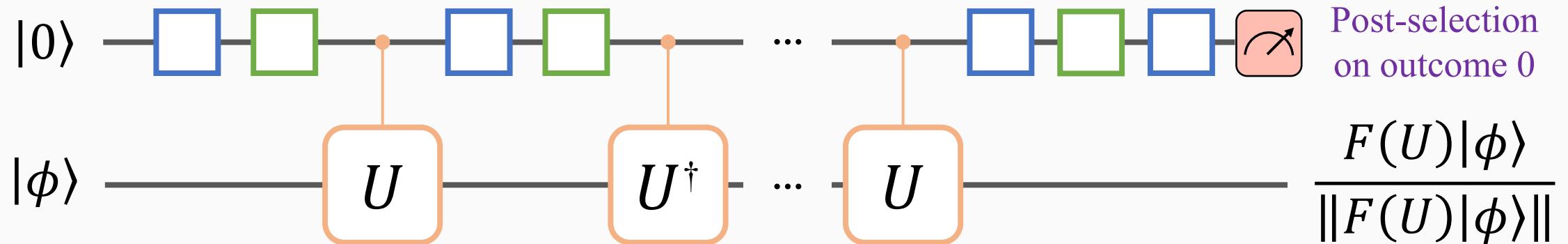
$$\begin{pmatrix} P(U) & -Q(U) \\ Q^*(U) & P^*(U) \end{pmatrix}$$

$P, Q$  satisfy  $\begin{cases} P, Q \in \mathbb{C}[e^{-ix/2}, e^{ix/2}], \\ \deg P = \deg Q \\ P, Q \text{ has the same parity} \end{cases}$

# Transformation of quantum data

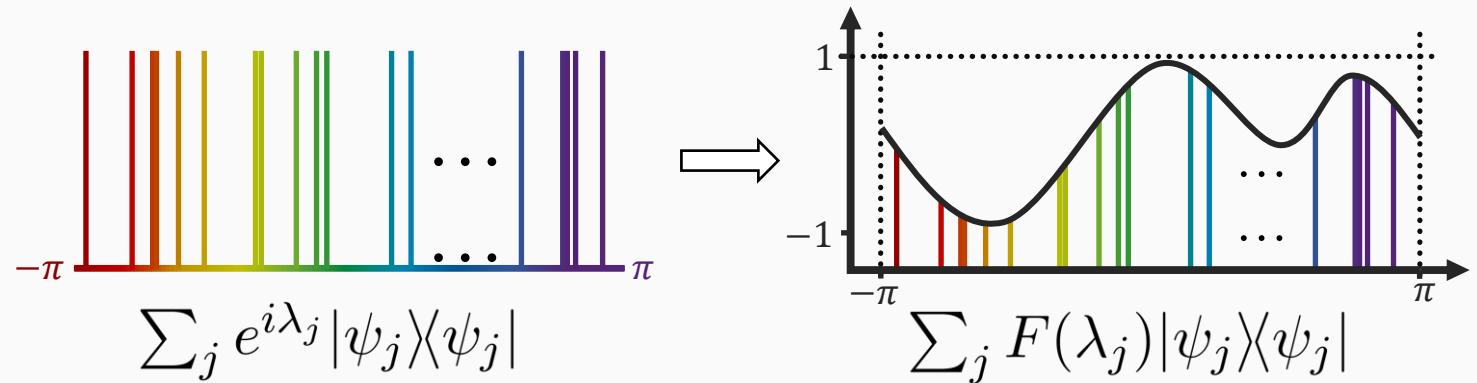
## Core Idea 1

Use quantum phase processing (QPP), a QSP-based structure to simulate transformation  $F$



How to simulate  $f(U)$ ?

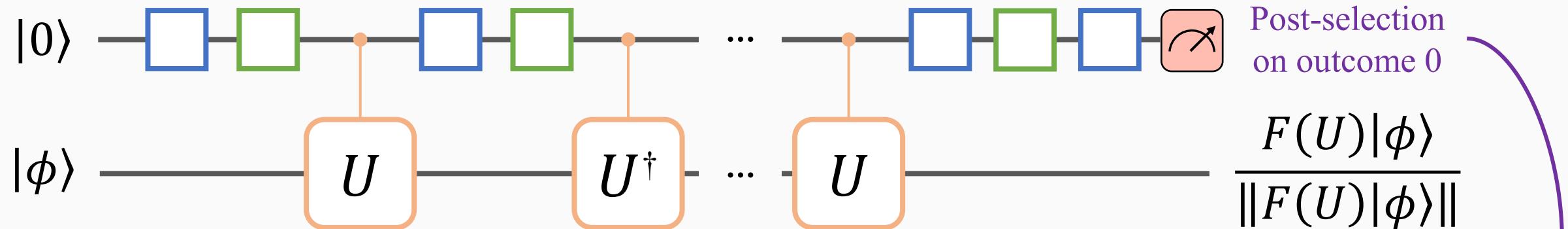
- ① Assume access to control –  $U$
- ② Compute  $f$ 's Fourier approximation  $F$
- ③ Compute rotation angles for  $\square$   $\square$
- ④ Construct the QPP circuit
- ⑤ Post-select the ancilla qubit



# Transformation of quantum data

## Core Idea 1

Use quantum phase processing (QPP) to simulate transformation



For ITE problem, choose  $U = e^{-iH}$

$$F(x) \approx e^{\tau x} / e^\tau, \text{ then}$$

output state  $\approx |\phi(\tau)\rangle$  with success probability  $\mathcal{O}(e^{-2\tau})$

Naïve choice raises exponential decay problem

# Adaptive normalization

## Core Idea 2

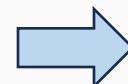
Choose an exponential transformation with adaptive normalization factor  $\lambda$ . Consider function

$$\boxed{\cancel{F(x) \approx e^{\tau x}}} \Rightarrow F(x) \approx \begin{cases} \alpha e^{\tau(x-\lambda)}, & x \in [-\pi, \lambda] \\ \xi_{\tau, \lambda}(x), & x \in (\lambda, \pi] \end{cases}$$

where  $\alpha, \xi_{\tau, \lambda}$  ensure approximation error decays *super-polynomially* with the degree of approximation.

**Lemma 1** Let  $C \geq \tau(\lambda - |\lambda_0|) \geq 0$ . Under Assumptions (i,iv,v), the output state  $|\tilde{\phi}(\tau)\rangle$  from the ITE circuit  $V_{f_{\tau, \lambda}}^\epsilon(U_H)$  is obtained with success probability lower bounded by  $\alpha^2 \gamma^2 e^{-2C} - \epsilon$ . Moreover, the state fidelity between the output state and the ITE state is approximately lower bounded as

$$|\langle \phi(\tau) | \tilde{\phi}(\tau) \rangle| \gtrsim 1 - \mathcal{O}(\alpha^{-1} \epsilon \cdot e^C). \quad (5)$$



Finding  $\lambda \in [|\lambda_0|, |\lambda_0| + \tau^{-1}]$  outputs an approximated ITE state with success probability nearly lower bounded by a CONSTANT  $\alpha^2 e^{-2\gamma^2}$

Avoids the exponential decay problem

# Main results

## Assumptions without loss of generality

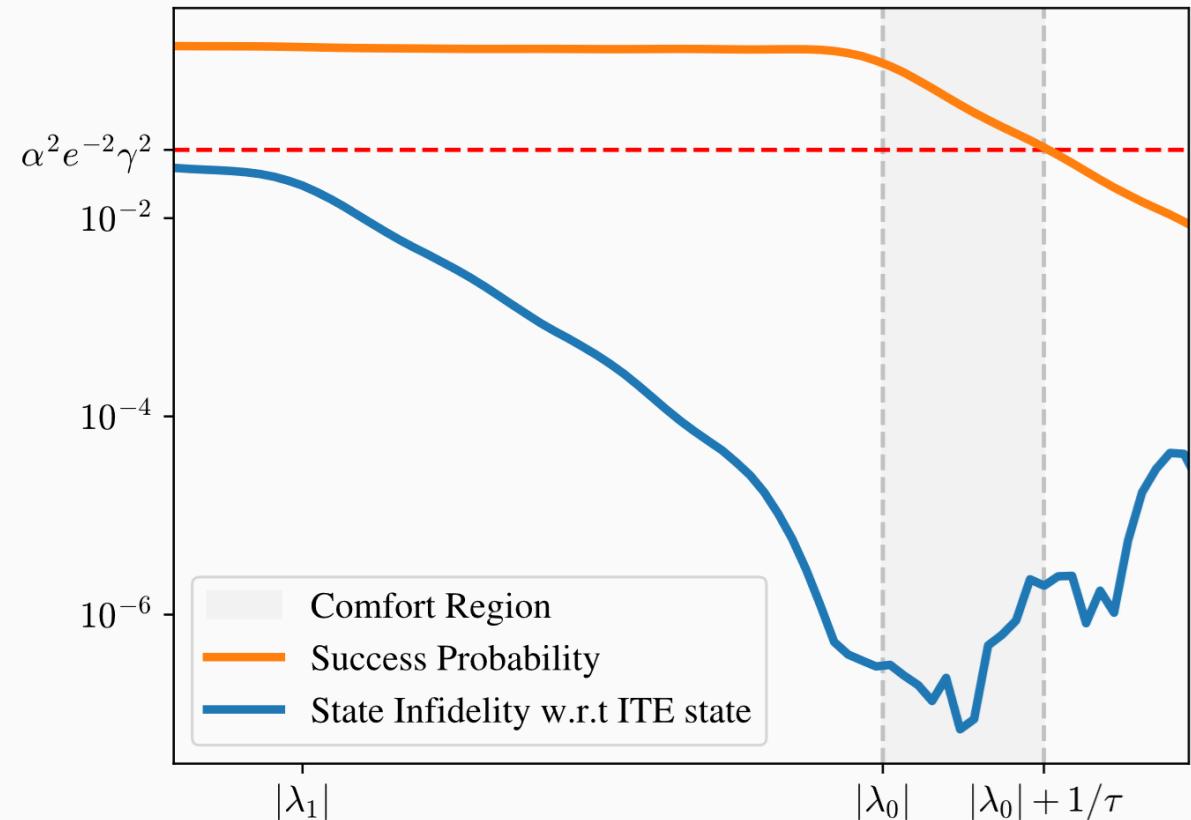
- i) eigenvalues of  $H$  is within  $[-1, 1]$ , and ground-state energy  $\lambda_0$  is negative
- iv) imaginary-time evolution is long i.e.,  $\tau \gg 0$
- v) the overlap between input state  $|\phi(0)\rangle$  and the ground state is non-zero

**Lemma 1** Let  $C \geq \tau(\lambda - |\lambda_0|) \geq 0$ . Under Assumptions (i,iv,v), the output state  $|\tilde{\phi}(\tau)\rangle$  from the ITE circuit  $V_{f_{\tau,\lambda}}^{\epsilon}(U_H)$  is obtained with success probability lower bounded by  $\alpha^2\gamma^2e^{-2C} - \epsilon$ . Moreover, the state fidelity between the output state and the ITE state is approximately lower bounded as

$$|\langle\phi(\tau)|\tilde{\phi}(\tau)\rangle| \gtrsim 1 - \mathcal{O}(\alpha^{-1}\epsilon \cdot e^C). \quad (5)$$

$$\lambda \in [|\lambda_0|, |\lambda_0| + \tau^{-1}]$$

Error of ITE state preparation:  $\mathcal{O}(10^{-5})$   
 Success probability  $\gtrsim \alpha^2 e^{-2\gamma^2}$



# Main results

Assumptions that are specific to this problem

- vii) the overlap between input state  $|\phi(0)\rangle$  and the ground state is at least  $\mathcal{O}(\text{poly}(n^{-1}))$
- viii) the gap  $\Delta$  between the ground-state energy  $\lambda_0$  and the first-excited-state energy  $\lambda_1$  is not zero
- ix)  $\Delta$  satisfies  $e^{\tau\Delta}$  is at least  $\mathcal{O}(\text{poly}(\tau))$



**Theorem 3** Under Assumptions (i, iii, iv, v, vi, vii, viii, ix), one can prepare the ITE state  $|\phi(\tau)\rangle$  up to fidelity  $1 - \mathcal{O}(L^2\Lambda^2 \text{poly}(\tau^{-1}))$ , using the following cost:

- $\tilde{\mathcal{O}}(L \text{poly}(n\tau))$  queries to controlled Pauli rotations,
- $\mathcal{O}(\text{poly}(n))$  copies of  $|\phi\rangle$ ,
- $\tilde{\mathcal{O}}(L \text{poly}(\tau))$  maximal query depth, and
- one ancilla qubit initialized in the zero state,

where  $L$  is the number of Pauli terms and  $\Lambda = \max_j |h_j|$ .

**Main Result** Approximates the target state to **polynomially small errors in inverse imaginary time** using polynomially many elementary quantum gates and a single ancilla qubit.

# Comparison with theoretical works

Methods	Circuit depth	Expected circuit runs	Ancilla	$H$ in Pauli form?
Manifold-based, $k$ steps [23]	$\mathcal{O}(3^k n)$	1	0	No
Grover-based [25]	$\mathcal{O}(\tau)$	\	$\mathcal{O}(n)$	No
TE-based [26]	$\mathcal{O}(\tau)$	\	1	No
QSP-based [27, 28]	$\tilde{\mathcal{O}}(\tau)$	\	1	No
QSP-based with AN (Theorem 2)	$\tilde{\mathcal{O}}(\tau)$	$\mathcal{O}(\gamma^{-2})$	1	No
QSP-based with AN (Theorem 3)	$\tilde{\mathcal{O}}(L \text{poly}(\tau))$	$\mathcal{O}(\text{poly}(n))$	1	Yes

adaptive normalization

holds for reasonable ground-state overlap

proves ITE state can be prepared in polynomial resource  
in terms of evolution time

[23] Gluza, Marek, et al. "Double-bracket quantum algorithms for quantum imaginary-time evolution." *arXiv preprint arXiv:2412.04554* (2024).

[25] Liu, Tong, Jin-Guo Liu, and Heng Fan. "Probabilistic nonunitary gate in imaginary time evolution." *Quantum Information Processing* 20.6 (2021): 204.

[26] Kosugi, Taichi, et al. "Imaginary-time evolution using forward and backward real-time evolution with a single ancilla: First-quantized eigensolver algorithm for quantum chemistry." *Physical Review Research* 4.3 (2022): 033121.

[27] Silva, Thais L., et al. "Fragmented imaginary-time evolution for early-stage quantum signal processors." *Scientific Reports* 13.1 (2023): 18258.

[28] Chan, Hans Hon Sang, David Muñoz Ramo, and Nathan Fitzpatrick. "Simulating non-unitary dynamics using quantum signal processing with unitary block encoding." *arXiv preprint arXiv:2303.06161* (2023).

\* This table is not shown on arXiv and will be updated in the version 3 soon

# Summary for the ITE part

Ground-state energy estimation  
of precision  $\mathcal{O}(\tau^{-1})$



ITE state preparation  
of precision  $\mathcal{O}(\text{poly}(\tau^{-1}))$

## Why ITE state can be prepared in polynomial resource?

- ① Multi-qubit QSP circuits are independent of system size
- ② Exponential transformation can be approximated via super-polynomial Fourier convergence
- ③ Assumption on reasonable ground state overlap (otherwise may not be scalable in system size)  
 not scalable for directly applying to Gibbs state preparation ☺

# Application

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- Ground state preparation and ground-state energy estimation

# Application to Ground-state problems

- Given a non-degenerate Hamiltonian

$$H = \sum_j \lambda_j |\psi_j\rangle\langle\psi_j|, \text{ with } \lambda_0 < \lambda_1 \leq \dots$$

find

$\lambda_0 \rightarrow$  Ground-state energy estimation

$|\psi_0\rangle \rightarrow$  Ground-state preparation

- Common query model  $U(t) = e^{iHt}$

# Application to Ground-state problems

**Lemma 4** *Under Assumptions (iii, vi),*

$$|\langle \psi_0 | \phi(\tau) \rangle| \geq \gamma / \sqrt{e^{-2\tau\Delta} + \gamma^2}.$$

*Moreover, the lower bound is tight for some Hamiltonians.*

- ITE state converges to the ground state as  $\tau \rightarrow \infty$
- The rate of convergence is at least polynomial if  $e^{\tau\Delta} = \Omega(\text{poly}(\tau))$

If  $\tau$  satisfies  $e^{\tau\Delta} = \Omega(\text{poly}(\tau))$ :

Ground-state  
energy estimation  
of precision  $\mathcal{O}(\tau^{-1})$

ITE state preparation  
of precision  $\mathcal{O}(\text{poly}(\tau^{-1}))$

Ground-state preparation  
of precision  $\mathcal{O}(\text{poly}(\tau^{-1}))$

How to find such  $\tau$  ?

# Application to Ground-state problems

**Lemma 4** *Under Assumptions (iii, vi),*

$$|\langle \psi_0 | \phi(\tau) \rangle| \geq \gamma / \sqrt{e^{-2\tau\Delta} + \gamma^2}.$$

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- ITE state converges to the ground state as  $\tau \rightarrow \infty$
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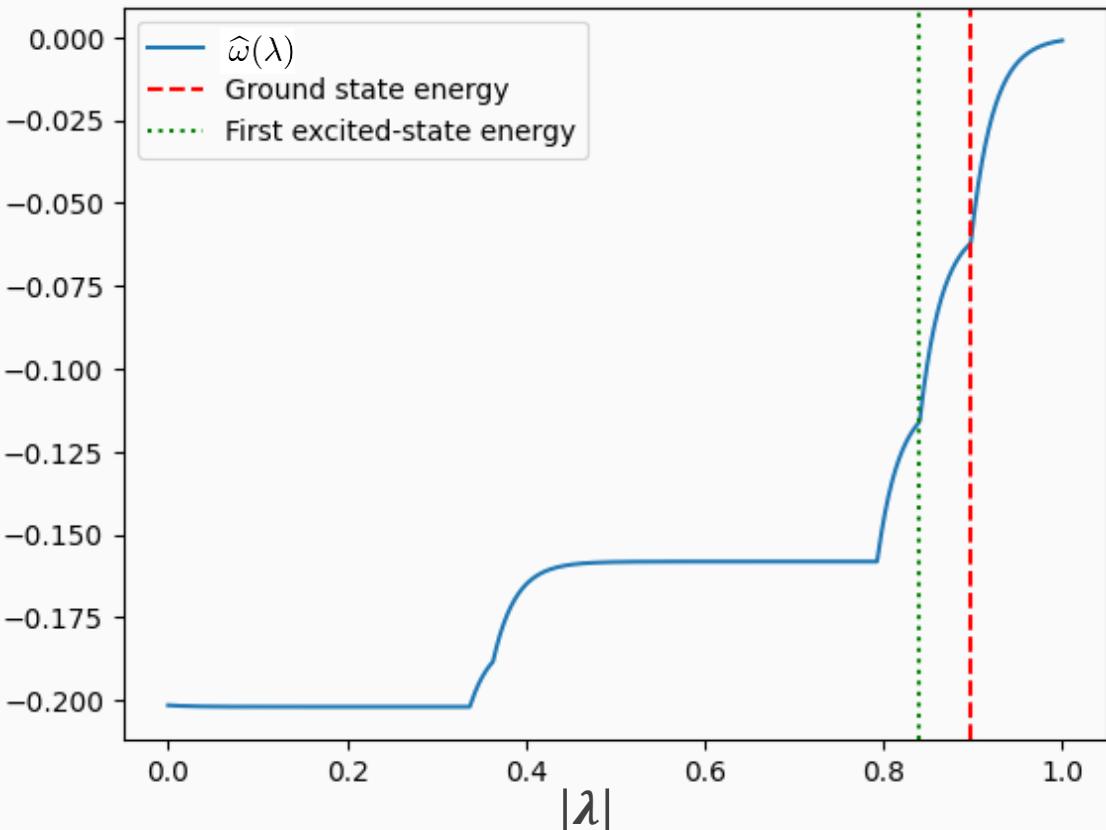
**Problem to be solved** For Hamiltonian  $H$ , find

1.  $\tau \gg 0$  s.t.  $e^{\tau\Delta} = \Omega(\text{poly}(\tau))$  (**ITE state  $\approx$  ground state**)
2.  $\lambda \in [|\lambda_0|, |\lambda_0| + \tau^{-1}]$  (**efficient ITE state preparation**)
3.  $E$ , an estimation of ITE state's expectation value w.r.t.  $H$   
**(ground-state energy estimation)**

# Our work

## Expectation value of circuit wrt. $H$

$$\hat{\omega}(\lambda) = \langle \phi | f_{\tau, \lambda}(U_H)^\dagger H f_{\tau, \lambda}(U_H) | \phi \rangle$$



## Algorithm Sketch

1. Start with a guess time  $t$ , and a search interval  $[\lambda_l, \lambda_r]$
2. Measure QPP circuits that prepare ITE states, which will give an estimation  $E_i$  and the interval where  $\lambda_0$  lies.
3. If search interval is small enough and the estimation sequence  $\{E_i\}_i$  converges, return  $t, \lambda_r, E_{-1}$
4. Otherwise, update guess time, search interval, and go back to Step 1.

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**Algorithm 1:** Ground state preparation and energy estimation via ITE

---

**Input :** Hamiltonian  $H$ , initial state  $|\phi\rangle$ , step size  $\Delta t$ , lower bound  $B$ , a boolean function  $\mathcal{X}$  for testing convergence

**Output:**  $\tau, \lambda, E$  in Problem 1

```

1 Guess  $t \gg 0$ ;
2  $E_0 \leftarrow 0, i \leftarrow 0$ ;
3  $\lambda_l \leftarrow 0, \lambda_r \xleftarrow{\sim} \max \{ \lambda : |\omega(\lambda)| > B \}$ ;
4 while  $\lambda_r - \lambda_l > t^{-1}$  or  $\mathcal{X}(\{E_i\}_i) = \text{False}$  do
5   Measurement shots #  $\leftarrow 8L\Lambda^{\frac{1}{2}}t^3B^{-2}$ ;
6    $\delta \leftarrow (\lambda_r - \lambda_l)/3, \lambda_{lm} \leftarrow \lambda_l + \delta, \lambda_{rm} \leftarrow \lambda_r - \delta$ ;
7   Estimate  $\omega(\lambda_{lm}), \omega(\lambda_r)$ ;
8    $r \leftarrow (\omega(\lambda_{lm}) - \omega(\lambda_r)) / \omega(\lambda_r)$ ;
9   if  $|r - (e^{4\tau\delta} - 1)| > \tau^{-1}(e^{4\tau\delta} + 1)$  then
10     $E_i \leftarrow$  selected samples that estimate  $\omega(\lambda_r)$ ;
11     $[\lambda_l, \lambda_r] \leftarrow [\lambda_{lm}, \lambda_r]$ ;
12  else
13     $E_i \leftarrow$  selected samples that estimate
         $\omega(\lambda_{lm}), \omega(\lambda_r)$ ;
14     $[\lambda_l, \lambda_r] \leftarrow [\lambda_l, \lambda_{rm}]$ ;
15   $t \leftarrow t + \Delta t, i \leftarrow i + 1$ ;
16 return  $\tau \leftarrow t, \lambda_r, E_i$ ;

```

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# Our work

**Theorem 6** Suppose Assumptions (i,ii,v,vi,vii,viii,x) hold. Algorithm 1 returns a time  $\tau$  that satisfies Assumption (ix), an estimate  $\lambda \in [|\lambda_0|, |\lambda_0| + \tau^{-1}]$ , and an estimate of  $\lambda_0$  within precision  $\mathcal{O}(B\gamma^{-1}\tau^{-1})$ , with failure probability  $\mathcal{O}(e^{-\tau} \log \tau)$ . Moreover, there are at most  $\mathcal{O}(L \log \tau)$  distinct circuit constructed in Algorithm 1, and each circuit takes at most:

- $\mathcal{O}(\tau)$  queries to controlled- $U_H$  and its inverse,
- $\mathcal{O}(\tau)$  query depth of  $U_H$ ,
- 1 ancilla qubit, and
- $\mathcal{O}(L\Lambda^2 B^{-2}\tau^3)$  measurement shots,

where  $L$  is the number of Pauli terms and  $\Lambda = \max_j |h_j|$ .

## Algorithm Sketch

1. Start with a guess time  $\tau$ , and a search interval  $[\lambda_l, \lambda_r]$
2. Measure QPP circuits that prepare ITE states, which will give an estimation  $E_i$  and the location where  $\lambda_0$  lies.
3. If search interval is small enough and the estimation sequence  $\{E_i\}_i$  converges, return  $\tau, \lambda_r, E_{-1}$
4. Otherwise, update guess time, search interval, and go back to Step 1.

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**Algorithm 1:** Ground state preparation and energy estimation via ITE

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**Input :** Hamiltonian  $H$ , initial state  $|\phi\rangle$ , step size  $\Delta t$ , lower bound  $B$ , a boolean function  $\mathcal{X}$  for testing convergence

**Output:**  $\tau, \lambda, E$  in Problem 1

```

1 Guess  $t \gg 0$ ;
2  $E_0 \leftarrow 0, i \leftarrow 0$ ;
3  $\lambda_l \leftarrow 0, \lambda_r \overset{\sim}{\leftarrow} \max \{ \lambda : |\omega(\lambda)| > B \}$ ;
4 while  $\lambda_r - \lambda_l > t^{-1}$  or  $\mathcal{X}(\{E_i\}_i) = \text{False}$  do
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8    $r \leftarrow (\omega(\lambda_{lm}) - \omega(\lambda_r)) / \omega(\lambda_r)$ ;
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13     $E_i \leftarrow$  selected samples that estimate
         $\omega(\lambda_{lm}), \omega(\lambda_r)$ ;
14     $[\lambda_l, \lambda_r] \leftarrow [\lambda_l, \lambda_{rm}]$ ;
15     $t \leftarrow t + \Delta t, i \leftarrow i + 1$ ;
16 return  $\tau \leftarrow t, \lambda_r, E_i$ ;

```

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# Comparison with existing works

Methods	Query depth	Expected circuit runs	Ancilla
Conventional QPE [31]	$\Omega(\gamma^{-1}\tau)$	$\mathcal{O}(\gamma^{-2}\tau)$	$\Omega(\log(\gamma^{-1}\tau))$
Semi-classical QPE [32, 33]	$\Omega(\gamma^{-3}\tau)$	$\mathcal{O}(\gamma^{-2}\tau)$	1
QET-based [34] (state preparation)	$\Omega(\gamma^{-1}\tau)$	$\tilde{\mathcal{O}}(\gamma^{-2}\tau)$	3
QET-based [34] (energy estimation)	$\Omega(\gamma^{-1}\tau)$	$\tilde{\mathcal{O}}(\gamma^{-2}\tau)$	1
HT-based [35] (energy estimation)	$\Omega((1 - \gamma^2)^{1/2}\gamma^{-1}\tau)$	$\tilde{\mathcal{O}}(\gamma^{-4}\tau)$	1
ITE-based [3, 7]	\	\	0
ITE-based (Theorem 6)	$\mathcal{O}(\tau)$	$\Omega(\gamma^{-4}\tau^3)$	1

Obtain shorter circuit depth as a trade-off of more queries to  $U_H$

[3] Motta, Mario, et al. "Determining eigenstates and thermal states on a quantum computer using quantum imaginary time evolution." *Nature Physics* 16.2 (2020): 205-210.

[7] McArdle, Sam, et al. "Variational ansatz-based quantum simulation of imaginary time evolution." *npj Quantum Information* 5.1 (2019): 75.

[31] Kitaev, A. Yu. "Quantum measurements and the Abelian stabilizer problem." *arXiv preprint quant-ph/9511026* (1995).

[32] Griffiths, Robert B., and Chi-Sheng Niu. "Semiclassical Fourier transform for quantum computation." *Physical Review Letters* 76.17 (1996): 3228.

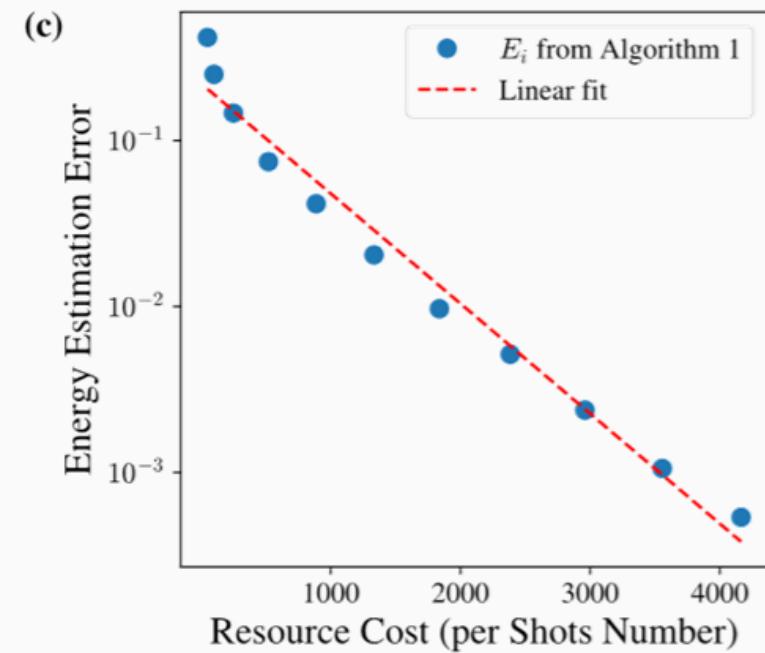
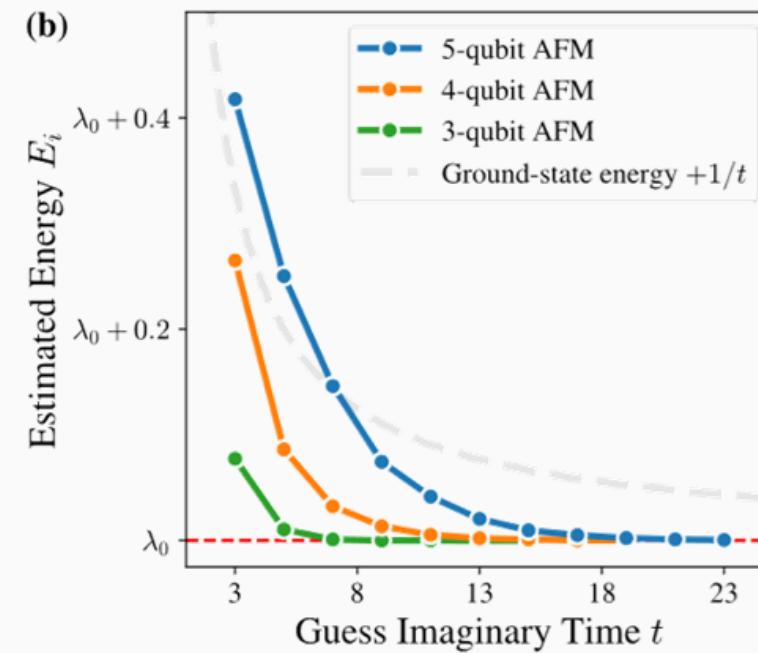
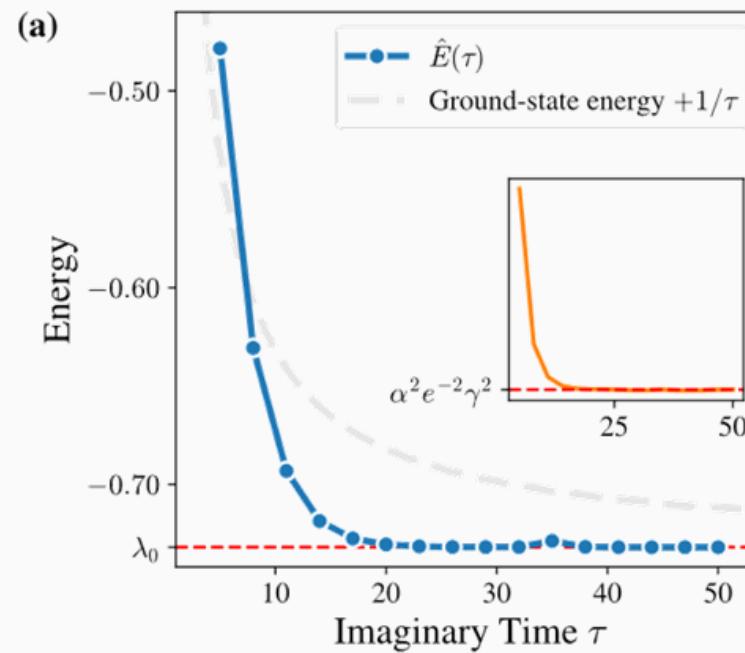
[34] Dong, Yulong, Lin Lin, and Yu Tong. "Ground-state preparation and energy estimation on early fault-tolerant quantum computers via quantum eigenvalue transformation of unitary matrices." *PRX Quantum* 3.4 (2022): 040305.

[35] Ding, Zhiyan, and Lin Lin. "Even shorter quantum circuit for phase estimation on early fault-tolerant quantum computers with applications to ground-state energy estimation." *PRX Quantum* 4.2 (2023): 020331.

\* This table is not shown on arXiv and will be updated in the version 3 soon

# Experiments

Hamiltonian Type Antiferromagnetic Heisenberg Model (AFM)  $H \propto \sum_{j=0}^n (X_j X_{j+1} + Y_j Y_{j+1} + Z_j Z_{j+1} - I)$



Complete tasks with imaginary-time evolution time up to 50

# Summary

- The first ITE algorithm proved to have polynomial resource dependence in evolution time
- Applying ITE to ground-state problems can reduce circuit depth
- Potential caveats
  - Does not consider classical machine error

# Thank you

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Research Group



Preprint  
(v3 updated soon)



Code Repo