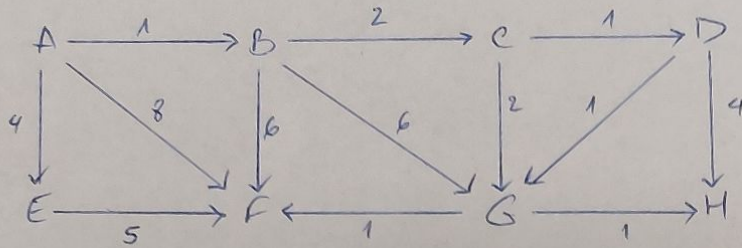


CS 216

Quan Nguyen

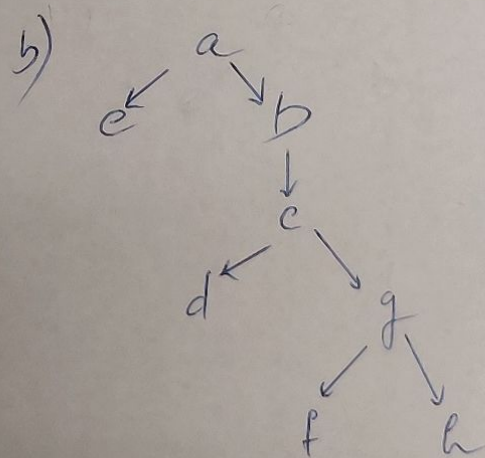
HW 7: graph paths

4.1



	Dist						Prev					
a	0						∅					
b	∅ → 1						∅ → a					
c	∅ → 3						∅ → b					
d	∅ → 4						∅ → c					
e	∅ → 4						∅ → a					
f	∅ → 8 → 6						∅ → d → b → g					
g	∅ → 7 → 5						∅ → b → c					
h	∅ → 8 → 6						∅ → d → g					
Loop	0	1	2	3	4	5	0	1	2	3	4	5

Loop	Priority Queue
0	a b c d e f g h
1	b c e f c d g h
2	c e f g d h
3	d e g f h
4	e g f h
5	f g h
6	g h
7	h
8	∅



4.3

Method():

HashTable <key; Boolean> u_neigh, v_neigh

for u in V:

u_neigh.put(u, True if it is u's neighbor else False)
count = 0

for v in V:

v_neigh.put(v, True if it is v's neighbors else False)

u_neigh.put(v, False)

v_neigh.put(u, False)

for (Key k: V):

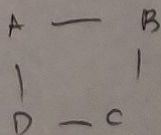
if u_neigh.get(k) & v_neigh.get(k):
count++

if count == 2:

return True

return False

/* How it works.



Randomly choose A and C ($O(|V|^2)$)

so that they share exactly 2 neighbors (B, D)

(loop through takes $O(|V|)$)

$$\Rightarrow \Sigma = O(|V|^3)$$

*/

4.5

// using Dijkstra

Method (s, t) : from s to t

for u in V :

$$\text{dist}(u) = \infty$$
$$\text{prev}(u) = \emptyset$$

HashTable < Key, Integer> num-path

for u in V :

```
num_path.put(u, 0)
```

$$\text{dist}(8) = 0$$
$$PQ = pq(v)$$

while ($PQ \neq \emptyset$):

$u = PQ$. delete Min

for $(u, v) \in E$:

if $\text{dist}(v) > \text{dist}(u) + \text{len}(u, v)$: // find shorter path

$$\text{dist}(v) = \text{dist}(u) + \text{len}(u, v)$$

$\text{dist}(v) = \text{dist}(u) + \text{len}(u, v)$
 $\text{num_path.put}(v, \text{num_path.get}(u))$ // num-path stays same

elif $\text{dist}(v) == \text{dist}(u) + \text{len}(u, v)$

$\text{dist}(v) = \text{dist}(u) + \text{len}(u, v)$
 $\text{num_path.put}(v, \text{num_path.get}(v) + 1)$ // add num-path
 when find another
 path

PO. decrease key (v , $\text{dist}(v)$)

return num-path.get(t)

Method (s, t) // edge e connects s and t

if t is not s's neighbor:

return

tmpGraph = copy(this Graph)

tmpGraph.vertices.get(s).remove(t) // remove edge e

// use Dijkstra

V' = vertices of tmpGraph

E' = edges of tmpGraph

for u in V' :

dist(u) = ∞

prev(u) = \emptyset

dist(s) = 0

PQ = pq(V')

While PQ $\neq \emptyset$:

u = PQ.deleteMin

for (u, v) in E' :

if dist(v) > dist(u) + len(u, v)

dist(v) = dist(u) + len(u, v)

prev(v) = u

PQ.decreaseKey(v, dist(v)) $\rightarrow O(1)$

return dist(t) + len(s, t)

$$\Sigma = O(|V|^2) + O(|V| + |E|)$$

worst case: $|E| \approx |V|^2$

$$\Rightarrow \Sigma = O(|V|^2)$$

$$\rightarrow O(|V|) \approx O(|V'|)$$

$$\rightarrow O(|V|)$$

$$\rightarrow O(|E'_u|) \approx O(|E'_u|)$$