Outline of this lecture

Communication Modes

- Standard Send
- Synchronous Send
- Buffered send

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General concepts

Blocking call

- A call is blocking if the memory space used for the communication can be reused immediately after the exit of the call
- The data sent can be modified after the call.
- The data received can be read after the call.

Point-to-Point Send Modes

Mode	Blocking	Non-blocking
Standard Send	MPI_SEND()	MPI_ISEND()
Synchronous Send	MPI_SSEND()	MPI_ISSEND()
Buffered send	MPI_BSEND()	<pre>MPI_IBSEND()</pre>
Receive	MPI_RECV()	MPI_IRECV()

Table: Point-to-Point Send Modes

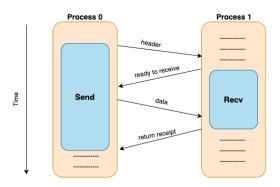
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Synchronous sends

A synchronous send involves a synchronization between the involved processes. A send cannot start until its receive is posted. There can be no communication before the two processes are ready to communicate.

Rendezvous Protocol

The rendezvous protocol is generally the protocol used for synchronous sends (implementation-dependent). The return receipt is optional.



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Interface of : MPI_SSEND()

```
1 MPI_SSEND(values, count, msgtype, dest, tag, comm, code)
2
3 type(*), intent(in) :: values
4 integer, intent(in) :: count, msgtype, dest, tag, comm
5
6 integer, intent(out) :: code
```

- Advantages of synchronous mode
 - ☐ Low resource consumption (no buffer)
 - Rapid if the receiver is ready (no copying in a buffer)
 - ☐ Knowledge of receipt through synchronization
- Disadvantages of synchronous mode
 - □ Waiting time if the receiver is not there/not ready
 - Risk of deadlocks

```
COMM.Ssend(self, data, int dest, int tag=0)
#or
COMM.ssend(self, data, int dest, int tag=0)
```

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Deadlock example

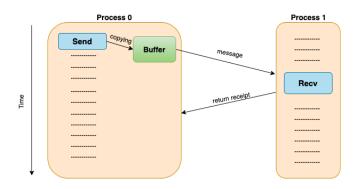
8

```
from mpi4py import MPI
    COMM = MPI.COMM_WORLD
    RANK = COMM.Get_rank()
    tag = 99
    sendbuf = 1000
    COMM.ssend(sendbuf, dest=1, tag=tag)
10
11
    recvbuf = COMM.recv(source=1, tag=tag)
12
    print("I, process 1, I received ",recvbuf," from the process 2.")
```

mpirun -n 2 python ssendrecv.py

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- Buffered sends
 A buffered send implies the copying of data into an intermediate memory space.
- Protocol with user buffer on the sender side In this approach, the buffer is on the sender side and is managed explicitly by the application. A buffer managed by MPI can exist on the receiver side. Many variants are possible. The return receipt is optional.



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Interface of : MPI_BSEND()

Buffered sends

The buffers have to be managed manually (with calls to MPI_BUFFER_ATTACH() and MPI_BUFFER_DETACH()). Message header size needs to be taken into account when allocating buffers (by adding the constant MPI_BSEND_OVERHEAD() for each message occurrence).

Interfaces

```
1 MPI_BUFFER_ATTACH (buf, typesize, code)
2 MPI_BSEND (values, count, msgtype, dest, tag, comm, code)
3 MPI_BUFFER_DETACH (buf, typesize, code)
4
5 TYPE(*), intent(in) :: values
6 integer, intent(in) :: count, msgtype, dest, tag, comm
7 integer, intent(out) :: code
8 TYPE(*) :: buf
9
10 integer :: typesize
```

```
COMM.Bsend(self, data, int dest, int tag=0)
#or
COMM.bsend(self, data, int dest, int tag=0)
```

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Ad	vantages of buffered mode
	No need to wait for the receiver (copying in a buffer)
	No risk of deadlocks
Dis	sadvantages of buffered mode
	Uses more resources (memory use by buffers with saturation risk)
	The send buffers in the MPI_BSEND() or MPI_IBSEND() calls have to be
	managed manually (often difficult to choose a suitable size)
	Slightly slower than the synchronous sends if the receiver is ready
	No knowledge of receipt (send-receive decoupling)
	Risk of wasted memory space if buffers are too oversized
	Application crashes if buffer is too small
	There are often hidden buffers managed by the MPI implementation on the sender
	side and/or on the receiver side (and consuming memory resources)

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Interface of : MPI_SEND()

Standard sends

A standard send is made by calling the MPI_SEND() subroutine. In most implementations, the mode is buffered (eager) for small messages but is synchronous for larger messages.

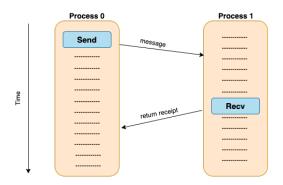
Interfaces

```
COMM.Send(self, data, int dest, int tag=0)
#or
COMM.send(self, data, int dest, int tag=0)
```

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The eager protocol

The eager protocol is often used for standard sends of small-size messages. It can also be used for sends with MPI_BSEND() for small messages (implementation-dependent) and by bypassing the user buffer on the sender side. In this approach, the buffer is on the receiver side. The return receipt is optional.



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- Advantages of standard mode
 - Often the most efficient (because the constructor chose the best parameters and algorithms)
 - ☐ The most portable for performance
- Disadvantages of standard mode
 - Little control over the mode actually used (often accessible via environment variables)
 - Risk of deadlocks depending on the mode used
 - Behavior can vary according to the architecture and problem size

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■ Example :

```
if RANK == 2:
    sendbuf = 1000
    COMM.send(sendbuf, dest=1, tag=tag)

if RANK == 5
    recvbuf = COMM.recv(source=2, tag=tag)
    print("I, process 5, I received ",recvbuf," from the process 2.")
```

```
mpirun -n 6 python bsendrecv.py
I, process 5, I received 1000 from the process 2.
```

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Non blocking communication

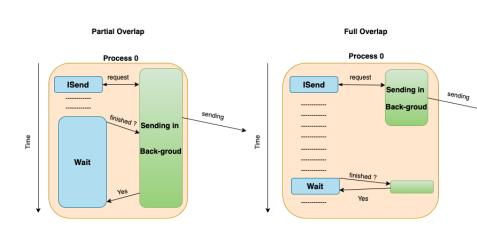
Presentation

The overlap of communications by computations is a method which allows executing
communications operations in the background while the program continues to operate
On Ada, the latency of a communication internode is 1.5 μs , or 4000 processor cycles
It is thus possible, if the hardware and software architecture allows it, to hide all of part of communications costs.
 The computation-communication overlap can be seen as an additional level of parallelism.
 This approach is used in MPI by using nonblocking subroutines (i.e.MPI_ISEND() MPI_IRECV() and MPI_WAIT()).

■ Non blocking communication

A nonblocking call returns very quickly but it does not authorize the immediate re-use of the memory space which was used in the communication. It is necessary to make sure that the communication is fully completed (with MPI_WAIT(), for example) before using it again.

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Ad	vantages of non blocking mode
	Possibility of hiding all or part of communications costs (if the architecture allows it)
	No risk of deadlock
Dis	advantages of non blocking mode
	Greater additional costs (several calls for one single send or receive, request management)
	Higher complexity and more complicated maintenance
	Less efficient on some machines (for example with transfer starting only at the MPI_WAIT() call)
	Risk of performance loss on the computational kernels (for example, differentiated management between the area near the border of a domain and the interior area, resulting in less efficient use of memory caches)
	Limited to point-to-point communications (it is extended to collective communications in MPI 3.0)

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Interface of : MPI_ISEND(), MPI_ISSEND(), MPI_IBSEND(), MPI_IRECV()

```
1 MPI_ISEND(values, count, datatype, dest, tag, comm, req, code)
 2 MPI_ISSEND(values, count, datatype, dest, tag, comm, req, code)
 3 MPI IBSEND(values, count, datatype, dest, tag, comm, reg, code)
 5 TYPE(*), intent(in) :: values
 7 integer, intent(in) :: count, datatype, dest, tag, comm
 8 integer, intent(out) :: req, code
 1 MPI IRECV(values, count, msgtvpe, source, tag, comm, reg, code)
 3 TYPE(*), intent(in) :: values
 4 integer, intent(in) :: count, msgtype, source, tag, comm
 6 integer, intent(out) :: req, code
   COMM. Isend (self. data. int dest. int tag=0)
1
   COMM. Issend (self. data. int dest. int tag=0)
3
   COMM. Ibsend (self, data, int dest, int tag=0)
5
   data = COMM. Irecv(self, source, int tag=0)
   #or
```

Comm. IRecv (self. buf. int source. int tag=0. Status status=None)

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Interface of : MPI_WAIT()

MPI_WAIT() wait for the end of a communication, MPI_TEST() is the nonblocking version.

```
1 MPI_WAIT(req, statut, code)
2 MPI_TEST(req, flag, statut, code)
3
4 integer, intent(inout) :: req
5 integer, dimension(MPI_STATUS_SIZE), intent(out) :: statut
6 integer, intent(out) :: code
7
8 logical, intent(out) :: flag
```

```
req = COMM.Isend(self, data, int dest, int tag=0)
req.Wait(self, Status status=None)
req.Test(self, Status status=None)
```

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Interface of : MPI_WAIT()

MPI_WAITALL() (MPI_TESTALL()) await the end of all communications.

```
1 MPI_WAITALL(taille, reqs, statuts, code)
2 MPI_TESTALL(taille, reqs, statuts, flag, code)
3
4 integer, intent(in) :: count
5 integer, dimension(count) :: reqs
6 integer, dimension(MPI_STATUS_SIZE,count), intent(out) :: statuts
7 integer, intent(out) :: code
8
9 logical, intent(out) :: flag
```

```
req = COMM.Isend(self, data, int dest, int tag=0)
req.Waitall(type cls, requests, statuses=None)
req.Test(type cls, requests, statuses=None)
```

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Request management

- After a call to a blocking wait function (MPI_WAIT(), MPI_WAITALL(),...), the request argument is set to MPI_REQUEST_NULL.
- ☐ The same for a nonblocking wait when the flag is set to true.
- ☐ A wait call with a MPI_REQUEST_NULL request does nothing.
- Example :

```
1  data = np.ones(1)
2
3  if RANK == 2:
4   data = np.random.random_sample(1)
5   print("I, process 2, I send", data[0], "to process 5")
6   req = COMM.Isend(data, dest=5)
7   req.Wait()
8   req.Test()
9
10  if RANK == 5:
   req = COMM.Irecv(data, source=2)
12   req.Wait()
13  print("I, process 5, I received ",data[0]," from the process 2.")
```

```
mpirun -n 6 python isendirecv.py

I, process 2, I send 0.7493708552194022 to process 5
I, process 5, I received 0.7493708552194022 from the process 2.
```

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Number of received elements

```
1 MPI_RECV(buf,count,datatype,source,tag,comm,msgstatus,code)
2
3 <type >:: buf
4 integer :: count, datatype
integer :: source, tag, comm, code
6
7 integer, dimension(MPI_STATUS_SIZE) :: msgstatus
```

- In MPI_RECV() or MPI_IRECV() call, the count argument in the standard is the number of elements in the buffer buf.
- This number must be greater than the number of elements to be received.
- When it is possible, for increased clarity, it is adviced to put the number of elements to be received
- We can obtain the number of elements received with MPI_GET_COUNT() and the msgstatus argument returned by the MPI_RECV() or MPI_WAIT() call.

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