

Robust Principal Component Analysis



Yuxin Chen

Princeton University, Spring 2017

Disentangling sparse and low-rank matrices

Suppose we are given a matrix

$$M = \underbrace{L}_{\text{low-rank}} + \underbrace{S}_{\text{sparse}} \in \mathbb{R}^{n \times n}$$

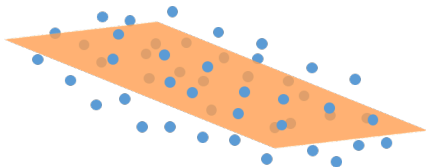
Question: Can we hope to recover both L and S from M ?

Principal component analysis (PCA)

- N samples $\mathbf{X} = [\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_N] \in \mathbb{R}^{n \times N}$ that are centered
- PCA: seeks r directions that explain most variance of data

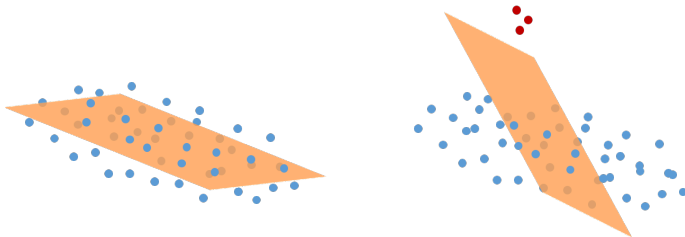
$$\text{minimize}_{\mathbf{L}: \text{rank}(\mathbf{L})=r} \quad \|\mathbf{X} - \mathbf{L}\|_F$$

- best rank- r approximation of \mathbf{X}



Sensitivity to corruptions / outliers

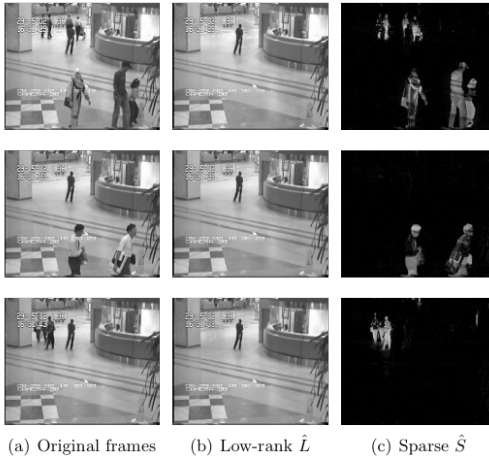
What if some samples are corrupted (e.g. due to sensor errors / attacks)?



Classical PCA fails even with a few outliers

Video surveillance

Separation of background (low-rank) and foreground (sparse)



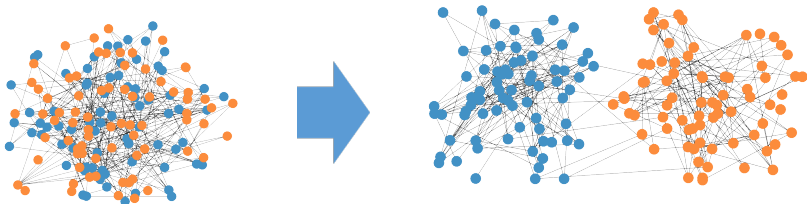
Candes, Li, Ma, Wright '11

Graph clustering / community recovery


- n nodes, 2 (or more) clusters
- A friendship graph \mathcal{G} : for any pair (i, j) ,

$$M_{i,j} = \begin{cases} 1, & \text{if } (i, j) \in \mathcal{G} \\ 0, & \text{else} \end{cases}$$

- Edge density **within** clusters $>$ edge density **across** clusters
- **Goal:** recover cluster structure



Graph clustering / community recovery


$$M = \underbrace{L}_{\text{low-rank}} + \underbrace{M - L}_{\text{sparse}}$$

- An equivalent goal: recover ground truth matrix

$$L_{i,j} = \begin{cases} 1, & \text{if } i \text{ and } j \text{ are in same community} \\ 0, & \text{else} \end{cases}$$

- Clustering \iff robust PCA

Gaussian graphical models

Fact 9.1

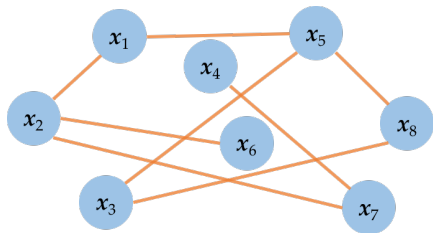
Consider a Gaussian vector $\mathbf{x} \sim \mathcal{N}(\mathbf{0}, \Sigma)$. For any u and v ,

$$x_u \perp\!\!\!\perp x_v \mid \mathbf{x}_{\mathcal{V} \setminus \{u,v\}}$$

iff $\Theta_{u,v} = 0$, where $\Theta = \Sigma^{-1}$ is inverse covariance matrix.

conditional independence \iff sparsity

Gaussian graphical models



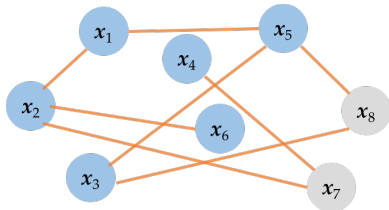
$$\underbrace{\begin{bmatrix} * & * & 0 & 0 & * & 0 & 0 & 0 \\ * & * & 0 & 0 & 0 & * & * & 0 \\ 0 & 0 & * & 0 & * & 0 & 0 & * \\ 0 & 0 & 0 & * & 0 & 0 & * & 0 \\ * & 0 & * & 0 & * & 0 & 0 & * \\ 0 & * & 0 & 0 & 0 & * & 0 & 0 \\ 0 & * & 0 & * & 0 & 0 & * & 0 \\ 0 & 0 & * & 0 & * & 0 & 0 & * \end{bmatrix}}_{\Theta}$$

Inverse covariance matrix Θ is often sparse

Graphical models with latent factors

What if one only observes a subset of variables?

$$\begin{bmatrix} \mathbf{x}_o \\ \mathbf{x}_h \end{bmatrix} \quad \begin{array}{l} \text{(observed variables)} \\ \text{(hidden variables)} \end{array}$$



$$\mathbf{x}_o = [x_1, \dots, x_6]^\top, \mathbf{x}_h = [x_7, x_8]^\top$$

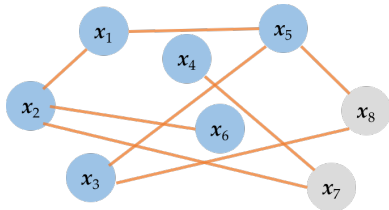
Covariance and precision matrices can be partitioned as

$$\Sigma = \begin{bmatrix} \overbrace{\Sigma_o}^{\text{observed part}} & \Sigma_{o,h} \\ \Sigma_{o,h}^\top & \Sigma_h \end{bmatrix} = \begin{bmatrix} \Theta_o & \Theta_{o,h} \\ \Theta_{o,h}^\top & \Theta_h \end{bmatrix}^{-1}$$

Graphical models with latent factors

What if one only observes a subset of variables?

$$\begin{bmatrix} \mathbf{x}_o \\ \mathbf{x}_h \end{bmatrix} \quad \begin{array}{l} \text{(observed variables)} \\ \text{(hidden variables)} \end{array}$$



$$\mathbf{x}_o = [x_1, \dots, x_6]^\top, \mathbf{x}_h = [x_7, x_8]^\top$$

$$\underbrace{\Sigma_o^{-1}}_{\text{observed}} = \underbrace{\Theta_o}_{\text{sparse}} - \underbrace{\Theta_{o,h} \Theta_h^{-1} \Theta_{h,o}}_{\text{low-rank if } \# \text{ latent vars is small}}$$

sparse + low-rank decomposition

When is decomposition possible?

Identifiability issues: a matrix might be simultaneously low-rank and sparse!

$$\underbrace{\begin{bmatrix} 1 & 0 & 0 & \cdots & 0 \\ 0 & 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & 0 \end{bmatrix}}_{\text{sparse and low-rank}} \quad \text{vs.} \quad \underbrace{\begin{bmatrix} 1 & 0 & 1 & \cdots & 1 \\ 0 & 1 & 0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & 0 & 0 & \cdots & 1 \end{bmatrix}}_{\text{sparse but not low-rank}}$$

Nonzero entries of sparse component need to be spread out
— This lecture: assume locations of nonzero entries are random

When is decomposition possible?

Identifiability issues: a matrix might be simultaneously low-rank and sparse!

$$\underbrace{\begin{bmatrix} 1 & 1 & 1 & \cdots & 1 \\ 1 & 1 & 1 & \cdots & 1 \\ \vdots & \vdots & \vdots & & \\ 1 & 1 & 1 & \cdots & 1 \end{bmatrix}}_{\text{low-rank and dense}} \quad \text{vs.} \quad \underbrace{\begin{bmatrix} 1 & 1 & 1 & \cdots & 1 \\ 0 & 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \vdots & & \\ 0 & 0 & 0 & \cdots & 0 \end{bmatrix}}_{\text{low-rank but sparse}}$$

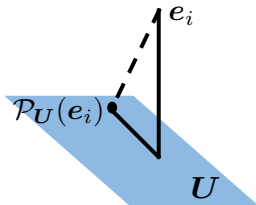
Low-rank component needs to be incoherent.

Low-rank component: coherence

Definition 9.2

Coherence parameter μ_1 of $M = U\Sigma V^\top$ is smallest quantity s.t.

$$\max_i \|U^\top e_i\|^2 \leq \frac{\mu_1 r}{n} \quad \text{and} \quad \max_i \|V^\top e_i\|^2 \leq \frac{\mu_1 r}{n}$$



Low-rank component: joint coherence

Definition 9.3 (Joint coherence)

Joint coherence parameter μ_2 of $M = U\Sigma V^\top$ is smallest quantity s.t.

$$\|UV^\top\|_\infty \leq \sqrt{\frac{\mu_2 r}{n^2}}$$

This prevents UV^\top from being too peaky.

- $\mu_1 \leq \mu_2 \leq \mu_1^2 r$, since

$$|(UV^\top)_{ij}| = |e_i^\top UV^\top e_j| \leq \|e_i^\top U\| \cdot \|V^\top e_j\| \leq \frac{\mu_1 r}{n}$$

$$\|UV^\top\|_\infty^2 \geq \frac{\|UV^\top e_j\|_F^2}{n} = \frac{\|V^\top e_j\|^2}{n} = \frac{\mu_1 r}{n^2} \text{ (suppose } \|V^\top e_j\|^2 = \frac{\mu_1 r}{n})$$

Convex relaxation

$$\text{minimize}_{\mathbf{L}, \mathbf{S}} \quad \text{rank}(\mathbf{L}) + \lambda \|\mathbf{S}\|_0, \quad \text{s.t.} \quad \mathbf{M} = \mathbf{L} + \mathbf{S} \quad (9.1)$$

\Downarrow

$$\text{minimize}_{\mathbf{L}, \mathbf{S}} \quad \|\mathbf{L}\|_* + \lambda \|\mathbf{S}\|_1, \quad \text{s.t.} \quad \mathbf{M} = \mathbf{L} + \mathbf{S} \quad (9.2)$$

- $\|\cdot\|_*$ is nuclear norm; $\|\cdot\|_1$ is entry-wise ℓ_1 norm
- $\lambda > 0$: regularization parameter that balances two terms

Theoretical guarantee

Theorem 9.4 (Candes, Li, Ma, Wright '11)

- $\text{rank}(\mathbf{L}) \lesssim \frac{n}{\max\{\mu_1, \mu_2\} \log^2 n}$;
- *Nonzero entries of \mathbf{S} are randomly located, and $\|\mathbf{S}\|_0 \leq \rho_s n^2$ for some constant $\rho_s > 0$ (e.g. $\rho_s = 0.2$).*

Then (9.2) with $\lambda = 1/\sqrt{n}$ is exact with high prob.

- $\text{rank}(\mathbf{L})$ can be quite high (up to $n/\text{polylog}(n)$)
- Parameter free: $\lambda = 1/\sqrt{n}$
- Ability to correct gross error: $\|\mathbf{S}\|_0 \asymp n^2$
- Sparse component \mathbf{S} can have arbitrary magnitudes / signs!

Geometry

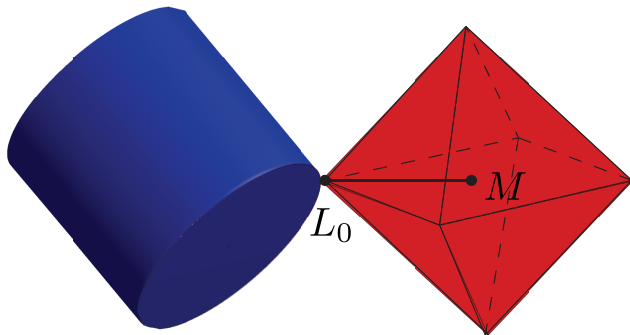
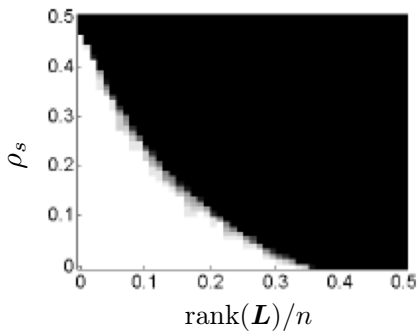


Fig. credit: Candes '14

Empirical success rate



$$n = 400$$

Fig. credit: Candes, Li, Ma, Wright '11

Dense error correction

Theorem 9.5 (Ganesh, Wright, Li, Candes, Ma '10, Chen, Jalali, Sanghavi, Caramanis '13)

- $\text{rank}(\mathbf{L}) \lesssim \frac{n}{\max\{\mu_1, \mu_2\} \log^2 n}$;
- Nonzero entries of \mathbf{S} are randomly located, have *random sign*, and $\|\mathbf{S}\|_0 = \rho_s n^2$.

Then (9.2) with $\lambda \asymp \sqrt{\frac{1-\rho_s}{\rho_s n}}$ succeeds with high prob., provided that

$$\underbrace{1 - \rho_s}_{\text{non-corruption rate}} \gtrsim \sqrt{\frac{\max\{\mu_1, \mu_2\} r \text{polylog}(n)}{n}}$$

- When additive corruptions have random signs, (9.2) works even when *a dominant fraction* of entries are corrupted

Is joint coherence needed?

- Matrix completion: does not need μ_2
- Robust PCA: so far we need μ_2

Question: is μ_2 needed? can we recover \mathbf{L} with rank up to $\frac{n}{\mu_1 \text{polylog}(n)}$ (rather than $\frac{n}{\max\{\mu_1, \mu_2\} \text{polylog}(n)}$)?

Answer: no (example: planted clique)

Planted clique problem

Setup: a graph \mathcal{G} of n nodes generated as follows

1. connect each pair of nodes independently with prob. 0.5
2. pick n_0 nodes and make them a clique (fully connected)

Goal: find hidden clique from \mathcal{G}

Information theoretically, one can recover a clique if $n_0 > 2 \log_2 n$

Conjecture on computational barrier

Conjecture: \forall constant $\epsilon > 0$, if $n_0 \leq n^{0.5-\epsilon}$, then no tractable algorithm can find the clique from \mathcal{G} with prob. $1 - o(1)$

— often used as hardness assumption

Lemma 9.6

If there is an algorithm that allows recovery of any \mathbf{L} from \mathbf{M} with $\text{rank}(\mathbf{L}) \leq \frac{n}{\mu_1 \text{polylog}(n)}$, then the above conjecture is violated

Proof of Lemma 9.6

Suppose L is true adjacency matrix,

$$L_{i,j} = \begin{cases} 1, & \text{if } i, j \text{ are both in the clique} \\ 0, & \text{else} \end{cases}$$

Let A be adjacency matrix of \mathcal{G} , and generate M s.t.

$$M_{i,j} = \begin{cases} A_{i,j}, & \text{with prob. } 2/3 \\ 0, & \text{else} \end{cases}$$

Therefore, one can write

$$M = L + \underbrace{\quad M - L \quad}_{\text{each entry is nonzero w.p. } 1/3}$$

Proof of Lemma 9.6

Note that

$$\mu_1 = \frac{n}{n_0} \quad \text{and} \quad \mu_2 = \frac{n^2}{n_0^2}$$

If there is an algorithm that can recover any \mathbf{L} of rank $\frac{n}{\mu_1 \text{polylog}(n)}$ from \mathbf{M} , then

$$\text{rank}(\mathbf{L}) = 1 \leq \frac{n}{\mu_1 \text{polylog}(n)} \iff n_0 \geq \text{polylog}(n)$$

But this contradicts the conjecture (which claims computational infeasibility to recover \mathbf{L} unless $n_0 \geq n^{0.5-o(1)}$)

Matrix completion with corruptions

What if we have missing data + corruptions?

- Observed entries

$$M_{ij} = L_{ij} + S_{ij}, \quad (i, j) \in \Omega$$

for some observation set Ω , where $\mathbf{S} = (S_{ij})$ is sparse

- A natural extension of RPCA

$$\text{minimize}_{\mathbf{L}, \mathbf{S}} \quad \|\mathbf{L}\|_* + \lambda \|\mathbf{S}\|_1 \quad \text{s.t. } \mathcal{P}_\Omega(\mathbf{M}) = \mathcal{P}_\Omega(\mathbf{L} + \mathbf{S})$$

- Theorems 9.4 - 9.5 easily extend to this setting

Efficient algorithm

In the presence of noise, one needs to solve

$$\text{minimize}_{\mathbf{L}, \mathbf{S}} \quad \|\mathbf{L}\|_* + \lambda \|\mathbf{S}\|_1 + \frac{\mu}{2} \|\mathbf{M} - \mathbf{L} - \mathbf{S}\|_F^2$$

which can be solved efficiently

Algorithm 9.1 Iterative soft-thresholding

for $t = 0, 1, \dots$:

$$\begin{aligned}\mathbf{L}^{t+1} &= \mathcal{T}_{1/\mu}(\mathbf{M} - \mathbf{S}^t) \\ \mathbf{S}^{t+1} &= \psi_{\lambda/\mu}(\mathbf{M} - \mathbf{L}^{t+1})\end{aligned}$$

where \mathcal{T} is singular-value thresholding operator, and ψ is soft thresholding operator

Reference

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Reference

- [6] "*Dense error correction for low-rank matrices via principal component pursuit*," A. Ganesh, J. Wright, X. Li, E. Candes, Y. Ma, *ISIT*, 2010.
- [7] "*Low-rank matrix recovery from errors and erasures*," Y. Chen, A. Jalali, S. Sanghavi, C. Caramanis, *IEEE Transactions on Information Theory*, 2013.