

# MH1300 Foundations of Mathematics – Solutions

Final Examination, Academic Year 2021/2022, Semester 1

*Compiled and typeset by QRS from the original handwritten solution*

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## Overview of the 2021/2022 Semester 1 Paper

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This typeset version is based on the original handwritten solutions from AY24/25 Sem 1, with minor corrections for clarity, consistent notation, and formatting. Alternative solution methods have been added for selected sub-questions where helpful (e.g. additional methods for Q1, Q3(b), Q4(a), Q4(b), and others). All mathematical content remains faithful to the original intent; any expanded explanations or supplementary methods are clearly marked as QRS notes.

**Question 1** (18 marks)

- (a) Let  $a$  and  $b$  be integers. Prove that  $a - b$  and  $a^2 + b^2$  have the same parity. (8 marks)
- (b) Is the following statement form a tautology, a contradiction, or neither?

$$(p \rightarrow (q \wedge \neg r)) \rightarrow (\neg q \rightarrow \neg p)$$

Justify your answer. (6 marks)

- (c) Write down a tautology using only the statement variables  $s$  and  $t$ , and the connective  $\rightarrow$ . All three symbols must be used at least once in your statement form, and you cannot use any other variables or connectives. Prove that your statement form is a tautology. (4 marks)

**Solution**

- (a) Suppose  $a, b$  are arbitrary integers.

**Method 1 (factorisation and cases).**

We divide into two cases, according to the parity of  $a - b$ .

*Case 1:  $a - b$  is odd.* Then

$$a + b = (a - b) + 2b = \text{odd} + \text{even} = \text{odd}.$$

Therefore

$$a^2 - b^2 = (a + b)(a - b) = \text{odd} \times \text{odd} = \text{odd}.$$

Hence

$$a^2 + b^2 = (a^2 - b^2) + 2b^2 = \text{odd} + \text{even} = \text{odd}.$$

So when  $a - b$  is odd,  $a^2 + b^2$  is also odd, and thus they have the same parity.

*Case 2:  $a - b$  is even.* Then

$$a + b = (a - b) + 2b = \text{even} + \text{even} = \text{even},$$

so

$$a^2 - b^2 = (a + b)(a - b) = \text{even} \times \text{even} = \text{even},$$

and hence

$$a^2 + b^2 = (a^2 - b^2) + 2b^2 = \text{even} + \text{even} = \text{even}.$$

So when  $a - b$  is even,  $a^2 + b^2$  is also even.

In both cases,  $a - b$  and  $a^2 + b^2$  have the same parity.

**Method 2 (four parity cases).**

We divide into four cases for the parities of  $a$  and  $b$ :

- Case 1:  $a$  even,  $b$  even. Then  $a - b$  is even, and  $a^2 + b^2$  is even+even=even.
- Case 2:  $a$  even,  $b$  odd. Then  $a - b$  is odd, and  $a^2 + b^2$  is even+odd=odd.
- Case 3:  $a$  odd,  $b$  even. Then  $a - b$  is odd, and  $a^2 + b^2$  is odd+even=odd.
- Case 4:  $a$  odd,  $b$  odd. Then  $a - b$  is even, and  $a^2 + b^2$  is odd+odd=even.

In each case  $a - b$  and  $a^2 + b^2$  have the same parity, so the statement holds.

### Method 3 (parity of squares).

Note that

$$\text{even}^2 = \text{even}, \quad \text{odd}^2 = \text{odd}.$$

Thus  $a$  and  $a^2$  have the same parity, and  $-b$  and  $b^2$  have the same parity.

Hence  $a - b$  (which has the same parity as  $a + (-b)$ ) and  $a^2 + b^2$  also have the same parity.

Therefore, in all methods,  $a - b$  and  $a^2 + b^2$  have the same parity.

**Remarks.** You may also argue using parity rules such as

$$\text{odd} \times \text{odd} = \text{odd}, \quad \text{odd} \times \text{even} = \text{even}, \quad \text{even} + \text{even} = \text{even}, \text{ etc.},$$

instead of always expanding from the definitions of even and odd.

(b) We show that the given statement form

$$(p \rightarrow (q \wedge \neg r)) \rightarrow (\neg q \rightarrow \neg p)$$

is a tautology.

### Method 1 (truth table).

Construct the truth table with columns for  $p, q, r$ , then  $q \wedge \neg r$ ,  $p \rightarrow (q \wedge \neg r)$ ,  $\neg q \rightarrow \neg p$ , and finally the whole implication. One finds that in all eight rows the final column is  $T$ . Hence the statement is a tautology.

### Method 2 (logical equivalences).

$$\begin{aligned}
 (p \rightarrow (q \wedge \neg r)) \rightarrow (\neg q \rightarrow \neg p) &\equiv \neg(p \rightarrow (q \wedge \neg r)) \vee (\neg q \rightarrow \neg p) && \text{(implication law)} \\
 &\equiv \neg(\neg p \vee (q \wedge \neg r)) \vee (\neg q \rightarrow \neg p) && \text{(implication law)} \\
 &\equiv (\neg \neg p \wedge \neg(q \wedge \neg r)) \vee (\neg q \rightarrow \neg p) && \text{(De Morgan)} \\
 &\equiv (p \wedge (\neg q \vee r)) \vee (\neg q \rightarrow \neg p) \\
 &\equiv (p \wedge (\neg q \vee r)) \vee (\neg \neg q \vee \neg p) && \text{(implication law)} \\
 &\equiv (p \wedge (\neg q \vee r)) \vee (q \vee \neg p) \\
 &\equiv (p \wedge (\neg q \vee r)) \vee ((q \vee \neg p) \vee (q \vee \neg p)) \\
 &\equiv (p \wedge (\neg q \vee r)) \vee T && \text{(negation / univ. bound)} \\
 &\equiv T.
 \end{aligned}$$

Thus the statement form is a tautology.

- (c) One possible choice is

$$s \rightarrow (t \rightarrow t).$$

**Proof that this is a tautology.**

$$\begin{aligned} s \rightarrow (t \rightarrow t) &\equiv \neg s \vee (t \rightarrow t) && \text{(implication law)} \\ &\equiv \neg s \vee (\neg t \vee t) && \text{(implication law)} \\ &\equiv \neg s \vee T && \text{(negation law)} \\ &\equiv T && \text{(universal bound law).} \end{aligned}$$

Hence  $s \rightarrow (t \rightarrow t)$  is a tautology.

Other correct choices are also acceptable, for example  $s \rightarrow (t \rightarrow s)$ ,  $(s \rightarrow s) \rightarrow (t \rightarrow t)$ , or  $(s \rightarrow t) \rightarrow (s \rightarrow t)$ , together with valid proofs that they are tautologies.

### **Mark Scheme**

- (a) 8 marks.

- Correct interpretation/restatement of the claim in terms of parity. [1]
- Coherent proof strategy (e.g. case analysis on parity of  $a, b$  or  $a - b$ , or observation that squares preserve parity). [2]
- Correct parity reasoning in all required cases, including identifying the parity of  $a - b$  and  $a^2 + b^2$  in each case. [4]
- Clear concluding sentence explicitly stating that  $a - b$  and  $a^2 + b^2$  always have the same parity. [1]

- (b) 6 marks.

- Correctly classifies the formula as a *tautology*. [1]
- Sets up an appropriate method: complete truth table *or* a correct chain of logical equivalences (using  $\rightarrow, \neg, \wedge$ , De Morgan, etc.). [2]
- Correctly carries the chosen method through to the end, showing that the final column is all  $T$  or that the formula is equivalent to  $\top$ . [3]

- (c) 4 marks.

- Constructs a valid propositional formula using only  $s, t$  and  $\rightarrow$ , each appearing at least once, which is indeed a tautology. [2]
- Provides a correct justification (truth table or equivalence steps) that the chosen formula is a tautology. [2]

## Question 2 (14 marks)

Determine if each of the following is true or false. Justify your answers.

- (a) If  $a$  and  $b$  are composite numbers, then  $a + b$  is composite. (3 marks)
- (b) For all positive integers  $c, d, e$ , if  $c \mid e$  and  $d \mid e$ , then either  $c = e$ ,  $d = e$  or  $cd \mid e$ . (4 marks)
- (c) Let  $A$  be a subset of a set  $B$ . Then  $A \times A \subseteq B \times B$ . (3 marks)
- (d) If  $C, D$  and  $E$  are sets, then  $(C \cup D) \cap E = C \cup (D \cap E)$ . (4 marks)

### Solution

- (a) **False.** Take  $a = 4$  and  $b = 9$ . Then  $a$  and  $b$  are composite, since  $a = 2 \cdot 2$  and  $b = 3 \cdot 3$ . But  $a + b = 13$  is prime, so the statement is false.
- (b) **False.** Take  $c = d = 4$  and  $e = 8$ . Then  $c, d, e$  are positive integers and
 
$$c \mid e \quad \text{and} \quad d \mid e,$$
 since  $4 \cdot 2 = 8$ . However  $c \neq e$ ,  $d \neq e$ , and  $cd = 16$  does not divide  $8 = e$ . Hence the statement is false.
- (c) **True.** Assume  $A \subseteq B$  and let  $(a, b) \in A \times A$ . Then  $a \in A$  and  $b \in A$ . Since  $A \subseteq B$ , we have  $a \in B$  and  $b \in B$ , so  $(a, b) \in B \times B$ . Thus  $A \times A \subseteq B \times B$ .
- (d) **False.** Take  $C = \{0\}$  and  $D = E = \emptyset$ . Then
 
$$(C \cup D) \cap E = (\{0\} \cup \emptyset) \cap \emptyset = \{0\} \cap \emptyset = \emptyset,$$
 while
 
$$C \cup (D \cap E) = \{0\} \cup (\emptyset \cap \emptyset) = \{0\} \cup \emptyset = \{0\}.$$
 Hence  $(C \cup D) \cap E \neq C \cup (D \cap E)$ , so the statement is false.

### Mark Scheme

- (a) 3 marks.
  - Chooses valid composite numbers  $a, b$  (e.g. 4, 9) and notes they are composite. [1]
  - Correctly computes  $a + b$  and identifies it as prime/non-composite. [1]
  - Explicitly states that this contradicts the claim, so the statement is false. [1]
- (b) 4 marks.
  - Chooses positive integers  $c, d, e$  satisfying  $c \mid e$  and  $d \mid e$  (e.g.  $c = d = 4, e = 8$ ). [1]

- Checks  $c \neq e$  and  $d \neq e$ . [1]
- Computes  $cd$  and checks that  $cd \nmid e$ . [1]
- Concludes that the implication fails and the statement is false. [1]

(c) 3 marks.

- Correctly uses  $A \subseteq B$  to deduce: if  $a \in A$  then  $a \in B$  (and same for  $b$ ). [1]
- Shows that any  $(a, b) \in A \times A$  is therefore in  $B \times B$ . [1]
- States the inclusion  $A \times A \subseteq B \times B$  as the final conclusion. [1]

(d) 4 marks.

- Chooses explicit sets  $C, D, E$  (such as  $C = \{0\}$ ,  $D = E = \emptyset$ ). [1]
- Correctly computes  $(C \cup D) \cap E$ . [1]
- Correctly computes  $C \cup (D \cap E)$ . [1]
- Observes the two results differ, hence the equality is false. [1]

**Question 3** (18 marks)(a) Let  $F : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$  be a function satisfying

$$F(a, 0) = a, \quad F(a, b + 1) = F(a, b) + 1$$

for each  $a, b \in \mathbb{N}$ . Using the definition of  $F$ , prove that for every  $a, b, c \in \mathbb{N}$ ,

$$F(F(a, b), c) = F(a, F(b, c)).$$

[8 marks]

**Solution**Fix  $a, b \in \mathbb{N}$  and prove the statement by induction on  $c$ .Let  $P(c)$  be the statement

$$F(F(a, b), c) = F(a, F(b, c)).$$

**Base case ( $c = 0$ )**.

$$F(F(a, b), 0) = F(a, b) \quad (\text{by the rule with second argument } 0),$$

and

$$F(a, F(b, 0)) = F(a, b) \quad (\text{since } F(b, 0) = b).$$

So  $P(0)$  holds.**Inductive step.** Let  $k \geq 0$  and suppose  $P(k)$  holds, i.e.

$$F(F(a, b), k) = F(a, F(b, k)).$$

Then

$$\begin{aligned} F(F(a, b), k + 1) &= F(F(a, b), k) + 1 && (\text{definition of } F) \\ &= F(a, F(b, k)) + 1 && (\text{induction hypothesis}) \\ &= F(a, F(b, k) + 1) && (\text{since } F(a, d + 1) = F(a, d) + 1) \\ &= F(a, F(b, k + 1)) && (\text{definition of } F \text{ for } F(b, k + 1)). \end{aligned}$$

Thus  $P(k + 1)$  holds.By mathematical induction,  $P(c)$  is true for all  $c \in \mathbb{N}$ , i.e.

$$F(F(a, b), c) = F(a, F(b, c))$$

for all  $a, b, c \in \mathbb{N}$ .(b) Prove that for every positive integer  $m$ ,

$$4^{m+1} + 5^{2m-1}$$

is divisible by 21.

[10 marks]

## Solution

**Method 1 (mathematical induction).**

Let  $P(m)$  be the statement

$$4^{m+1} + 5^{2m-1} \text{ is divisible by } 21, \quad m \geq 1.$$

**Base case ( $m = 1$ ).**

$$4^{1+1} + 5^{2 \cdot 1 - 1} = 4^2 + 5^1 = 16 + 5 = 21,$$

which is divisible by 21. So  $P(1)$  is true.

**Inductive step.** Let  $k \geq 1$  and assume  $P(k)$  holds. Then there exists an integer  $x$  such that

$$4^{k+1} + 5^{2k-1} = 21x.$$

Consider

$$\begin{aligned} 4^{(k+1)+1} + 5^{2(k+1)-1} &= 4^{k+2} + 5^{2k+1} \\ &= 4 \cdot 4^{k+1} + 25 \cdot 5^{2k-1} \\ &= 4 \cdot 4^{k+1} + 4 \cdot 5^{2k-1} + 21 \cdot 5^{2k-1} \\ &= 4(4^{k+1} + 5^{2k-1}) + 21 \cdot 5^{2k-1}. \end{aligned}$$

Using the induction hypothesis,

$$4^{k+1} + 5^{2k-1} = 21x,$$

so

$$\begin{aligned} 4^{k+2} + 5^{2k+1} &= 4 \cdot 21x + 21 \cdot 5^{2k-1} \\ &= 21(4x + 5^{2k-1}), \end{aligned}$$

which is divisible by 21. Hence  $P(k+1)$  holds.

By mathematical induction,  $4^{m+1} + 5^{2m-1}$  is divisible by 21 for all positive integers  $m$ .

**Method 2 (modular arithmetic).**

We show  $4^{m+1} + 5^{2m-1}$  is divisible by both 3 and 7.

*Modulo 3.* Since  $4 \equiv 1 \pmod{3}$ , we have  $4^{m+1} \equiv 1 \pmod{3}$ . Also  $5 \equiv -1 \pmod{3}$ , so

$$5^{2m-1} \equiv (-1)^{2m-1} = -1 \pmod{3}.$$

Thus

$$4^{m+1} + 5^{2m-1} \equiv 1 - 1 \equiv 0 \pmod{3}.$$

*Modulo 7.* Note that  $25 \equiv 4 \pmod{7}$ , so

$$5^{2m-1} = 5 \cdot 25^{m-1} \equiv 5 \cdot 4^{m-1} \pmod{7}.$$

Hence

$$\begin{aligned} 4^{m+1} + 5^{2m-1} &\equiv 4^{m+1} + 5 \cdot 4^{m-1} \\ &= 4^{m-1}(4^2 + 5) = 4^{m-1}(16 + 5) = 4^{m-1} \cdot 21 \equiv 0 \pmod{7}. \end{aligned}$$

Therefore  $4^{m+1} + 5^{2m-1}$  is divisible by both 3 and 7. Since  $\gcd(3, 7) = 1$ , it is divisible by 21.

### **Mark Scheme**

(a) 8 marks.

- Correctly fixes  $a, b$  and sets up induction on  $c$  (statement  $P(c)$  written clearly). [2]
- Correct base case  $c = 0$  with both sides computed using  $F(a, 0) = a$ . [2]
- Inductive step: correctly assumes  $P(k)$ , applies the recursive definition  $F(a, b+1) = F(a, b) + 1$  on both sides, and derives  $P(k+1)$ . [4]

(b) 10 marks.

- Base case  $m = 1$  computed correctly and checked for divisibility by 21. [2]
- Inductive step set up correctly: clear statement of the induction hypothesis and what needs to be shown for  $m = k+1$ . [2]
- Algebraic manipulation to express  $4^{k+2} + 5^{2k+1}$  in terms of  $4^{k+1} + 5^{2k-1}$  (e.g. factoring out 4 and 25). [3]
- Correct use of the induction hypothesis to factor out 21, and conclusion that  $4^{k+2} + 5^{2k+1}$  is divisible by 21. [3]
- \* A fully correct and clearly explained modular arithmetic solution (showing divisibility by 3 and by 7) can receive up to full credit if the induction step is incomplete or omitted.

**Question 4 (10 marks)**

- (a) Let  $a, b, c, d$  be integers such that  $d \mid a$ ,  $d \mid b$  and  $d \mid c$ . Prove that

$$d^2 \mid ab + ac + bc.$$

[4 marks]

**Solution****Method 1 (direct substitution).**

Since  $d \mid a$ ,  $d \mid b$  and  $d \mid c$ , there exist integers  $x, y, z$  such that

$$a = dx, \quad b = dy, \quad c = dz.$$

Then

$$\begin{aligned} ab + ac + bc &= (dx)(dy) + (dx)(dz) + (dy)(dz) \\ &= d^2xy + d^2xz + d^2yz \\ &= d^2(xy + xz + yz). \end{aligned}$$

Because  $xy + xz + yz \in \mathbb{Z}$ , we conclude that  $d^2 \mid ab + ac + bc$ .

**Method 2 (using basic divisibility facts).**

First note the lemma: if  $d \mid r$  and  $d \mid s$ , then  $d^2 \mid rs$ . Indeed, write  $r = dr_1$ ,  $s = ds_1$  with  $r_1, s_1 \in \mathbb{Z}$ . Then

$$rs = (dr_1)(ds_1) = d^2r_1s_1,$$

so  $d^2 \mid rs$ .

Using this:

- From  $d \mid a$  and  $d \mid b$  we get  $d^2 \mid ab$ .
- From  $d \mid a$  and  $d \mid c$  we get  $d^2 \mid ac$ .
- From  $d \mid b$  and  $d \mid c$  we get  $d^2 \mid bc$ .

Now use the fact that if  $d^2 \mid u$  and  $d^2 \mid v$ , then  $d^2 \mid (u + v)$  (since  $u = d^2u_1$ ,  $v = d^2v_1 \Rightarrow u + v = d^2(u_1 + v_1)$ ).

First,  $d^2 \mid ab$  and  $d^2 \mid ac$  imply  $d^2 \mid (ab + ac)$ . Then  $d^2 \mid (ab + ac)$  and  $d^2 \mid bc$  imply

$$d^2 \mid (ab + ac) + bc = ab + ac + bc.$$

Thus  $d^2 \mid ab + ac + bc$ .

- (b) Prove that every non-zero rational number is the product of two irrational numbers. You may use the fact that the product of a non-zero rational number with an irrational number is irrational. [6 marks]

## Solution

Let  $r \in \mathbb{Q}$  with  $r \neq 0$ .

### Method 1 (direct construction).

Set

$$x = r\sqrt{2}, \quad y = \frac{1}{\sqrt{2}}.$$

Then  $x$  is irrational, because it is the product of the non-zero rational number  $r$  and the irrational number  $\sqrt{2}$ . Similarly,

$$y = \frac{1}{\sqrt{2}} = \frac{\sqrt{2}}{2} = \left(\frac{1}{2}\right)\sqrt{2}$$

is the product of a non-zero rational and an irrational, so  $y$  is irrational.

Moreover,

$$xy = r\sqrt{2} \cdot \frac{1}{\sqrt{2}} = r.$$

Thus every non-zero rational  $r$  can be written as a product of two irrational numbers.

### Method 2 (proof by contradiction).

Suppose, for contradiction, that there exists a non-zero rational number  $r$  that cannot be expressed as the product of two irrational numbers.

Consider again

$$x = r\sqrt{2}, \quad y = \frac{1}{\sqrt{2}}.$$

By the given fact, both  $x$  and  $y$  are irrational (each is the product of a non-zero rational and an irrational). But then

$$xy = r\sqrt{2} \cdot \frac{1}{\sqrt{2}} = r,$$

so  $r$  is the product of two irrational numbers, contradicting our assumption.

Hence our assumption was false, and every non-zero rational number is indeed the product of two irrational numbers.

### Mark Scheme

(a) 4 marks.

- Uses  $d \mid a, d \mid b, d \mid c$  to introduce integers  $x, y, z$  with  $a = dx, b = dy, c = dz$  (or equivalent lemma  $d \mid r, d \mid s \Rightarrow d^2 \mid rs$ ). [2]
- Correct substitution and simplification showing  $ab + ac + bc = d^2(\dots)$  and concluding  $d^2 \mid ab + ac + bc$ . [2]

(b) 6 marks.

- Chooses explicit expressions  $x = r\sqrt{2}$  and  $y = 1/\sqrt{2}$  (or another valid pair of irrationals depending on  $r$ ). [2]
- Correctly argues that  $x$  is irrational using the given fact about (non-zero rational)  $\times$  (irrational). [2]
- Correctly argues that  $y$  is also irrational (e.g.  $y = \frac{1}{2}\sqrt{2}$ ). [1]
- Computes  $xy$  and shows  $xy = r$ , then concludes that every non-zero rational is a product of two irrationals. [1]

## Question 5 (12 marks)

- (a) Let  $A$  and  $B$  be sets and  $f : A \rightarrow B$  be a function. Let  $X \subseteq A$  and  $Y \subseteq B$ . Prove that

$$f(f^{-1}(Y)) \subseteq Y \quad \text{and} \quad X \subseteq f^{-1}(f(X)).$$

[6 marks]

- (b) Let  $g : \mathbb{Z} \rightarrow \mathbb{Z}$  be the function  $g(n) = 2n \bmod 3$ .

- (i) Is  $g$  injective?
- (ii) Is  $g$  surjective?
- (iii) What is the range of  $g$ ?

Justify your answers.

[6 marks]

### Solution

- (a) We first show that  $f(f^{-1}(Y)) \subseteq Y$ .

Let  $y \in f(f^{-1}(Y))$ . Then, by definition of  $f(f^{-1}(Y))$ , there exists some  $x \in f^{-1}(Y)$  such that  $y = f(x)$ . Since  $x \in f^{-1}(Y)$ , by definition of preimage we know  $f(x) \in Y$ . Hence  $y = f(x) \in Y$ . Therefore  $f(f^{-1}(Y)) \subseteq Y$ .

Next, we show that  $X \subseteq f^{-1}(f(X))$ .

Let  $x \in X$ . Then  $f(x) \in f(X)$  (since  $f(X) = \{f(a) \in B \mid a \in X\}$ ). Hence, by the definition of preimage,  $x \in f^{-1}(f(X))$ . Therefore  $X \subseteq f^{-1}(f(X))$ .

- (b) For  $g : \mathbb{Z} \rightarrow \mathbb{Z}$  defined by  $g(n) = 2n \bmod 3$  we first understand its values.

For any  $n \in \mathbb{Z}$ , we can write

$$2n = 3(2n \bmod 3) + (2n \bmod 3), \quad 0 \leq 2n \bmod 3 < 3.$$

Thus  $g(n) = 2n \bmod 3$  is always one of 0, 1, 2, so

$$\text{range}(g) \subseteq \{0, 1, 2\}.$$

Now we show the reverse inclusion:

$$g(0) = 0 \bmod 3 = 0, \quad g(1) = 2 \bmod 3 = 2, \quad g(2) = 4 \bmod 3 = 1.$$

Hence  $0, 1, 2 \in \text{range}(g)$ , so

$$\text{range}(g) = \{0, 1, 2\}.$$

- (i)  $g$  is *not injective*, since

$$0 \neq 3, \quad g(0) = 0 \bmod 3 = 0, \quad g(3) = 6 \bmod 3 = 0,$$

so  $g(0) = g(3)$  with  $0 \neq 3$ .

- (ii)  $g$  is *not surjective*, because for example  $3 \notin \{0, 1, 2\} = \text{range}(g)$ .  
(iii) The range of  $g$  is  $\{0, 1, 2\}$ , as shown above.

*Mark Scheme*

- (a) Correct proof that  $f(f^{-1}(Y)) \subseteq Y$  using element chase. (3 marks)  
Correct proof that  $X \subseteq f^{-1}(f(X))$ . (3 marks)
- (b) – Correct identification that  $g$  is not injective with a valid counterexample. (2 marks)  
– Correct identification that  $g$  is not surjective with justification (e.g. 3 is not in the range). (2 marks)  
– Correct determination and justification that  $\text{range}(g) = \{0, 1, 2\}$ . (2 marks)

**Remarks**

The definitions used are:

$$f(X) = \{ f(a) \in B \mid a \in X \}, \quad f^{-1}(Y) = \{ a \in A \mid f(a) \in Y \}.$$

## Question 6 (12 marks)

- (a) Find all sixth roots of unity. That is, find all complex numbers  $z$  satisfying  $z^6 - 1 = 0$ . Leave your answer in terms of  $re^{i\theta}$ . [6 marks]

- (b) Determine if the following argument is valid. State all rules of inference used.

$$\begin{aligned} & \neg p \vee \neg q \\ & \neg r \rightarrow (p \wedge q) \\ & \neg r \vee s \\ & s \rightarrow (t \wedge u) \\ & \therefore u \end{aligned}$$

[6 marks]

### Solution

- (a) The sixth roots of unity are the solutions to  $z^6 = 1$ .

Write 1 in polar form as  $1 = e^{i \cdot 0}$ . Using the general  $n$ th root formula

$$z = r^{1/n} e^{i(\theta+2k\pi)/n}, \quad k = 0, 1, \dots, n-1,$$

with  $r = 1$ ,  $\theta = 0$  and  $n = 6$ , we obtain

$$z = e^{i(0+2k\pi)/6} = e^{i\frac{2k\pi}{6}}, \quad k = 0, 1, 2, 3, 4, 5.$$

Thus the six roots are

$$z = e^{i \cdot 0} = 1, \quad e^{i\pi/3}, \quad e^{i2\pi/3}, \quad e^{i\pi}, \quad e^{i4\pi/3}, \quad e^{i5\pi/3},$$

each having modulus  $r = 1$  and argument  $\theta = \frac{2k\pi}{6}$  for  $k = 0, \dots, 5$ .

- (b) We show that the argument is valid by deriving  $u$  from the premises.

- |      |                                   |   |
|------|-----------------------------------|---|
| (1)  | $\neg p \vee \neg q$              | (Premise #1)  |
| (2)  | $\neg(p \wedge q)$                | (from (1) by De Morgan's Law)                                       |
| (3)  | $\neg r \rightarrow (p \wedge q)$ | (Premise #2)  |
| (4)  | $\neg\neg r$                      | (from (2) and (3) by Modus Tollens)                                 |
| (5)  | $r$                               | (from (4) by Double Negation)                                       |
| (6)  | $\neg r \vee s$                   | (Premise #3)  |
| (7)  | $s$                               | (from (5) and (6) by <i>Disjunctive Elimination / Elimination</i> ) |
| (8)  | $s \rightarrow (t \wedge u)$      | (Premise #4)  |
| (9)  | $t \wedge u$                      | (from (7) and (8) by Modus Ponens)                                  |
| (10) | $u$                               | (from (9) by <i>Conjunction Elimination / Specialisation</i> )      |

Since  $u$  has been derived from the premises using valid rules of inference, the argument is valid.

*Mark Scheme*

- (a) – Correct use of polar form and the general  $n$ th root formula. (2 marks)
  - Expression  $z = e^{i2k\pi/6}$  with correct range of  $k$ . (2 marks)
  - Correct explicit list of the six roots in the requested form. (2 marks)
- (b) – Correct transformation of  $\neg p \vee \neg q$  to  $\neg(p \wedge q)$  by De Morgan. (1 mark)
  - Correct use of Modus Tollens (or equivalent reasoning) to obtain  $r$ . (2 marks)
  - Correct use of  $\neg r \vee s$  with  $r$  to deduce  $s$ . (1 mark)
  - Correct use of  $s \rightarrow (t \wedge u)$  to get  $t \wedge u$ , and then  $u$ . (2 marks)

**Question 7****(16 marks)**

- (a) For the relation  $R$  on  $\mathbb{R}^2$  defined by

$$(x_1, x_2) R (x_3, x_4) \iff \exists i \neq j \text{ with } x_i = x_j, \quad i, j \in \{1, 2, 3, 4\},$$

determine whether  $R$  is reflexive, symmetric and transitive. Justify your answers.

[10 marks]

*QRS Note: The original question stated “for some  $i \neq j$ ”. Based on the official solution, this condition is interpreted as an existential condition (“there exists such  $i, j$ ”), not a universal one. Thus the notation has been rewritten using  $\exists$  for clarity.*

**Solution**

Write  $(x_1, x_2)$  and  $(x_3, x_4)$  for arbitrary elements of  $\mathbb{R}^2$ .

**Reflexive:** Let  $(x, y) \in \mathbb{R}^2$ . Then, when we consider  $(x_1, x_2, x_3, x_4) = (x, y, x, y)$ , we have

$$x_1 = x_3 \quad (\text{and also } x_2 = x_4).$$

Thus there exist  $i \neq j$  (for example  $i = 1, j = 3$ ) with  $x_i = x_j$ . Hence  $(x, y) R (x, y)$ , so  $R$  is reflexive.

**Symmetric:** Suppose  $(x_1, x_2) R (x_3, x_4)$ . Then, by definition, there exist  $i \neq j$  with  $x_i = x_j$ , where  $i, j \in \{1, 2, 3, 4\}$ . But the condition “some coordinates among  $x_1, x_2, x_3, x_4$  are equal” is unchanged if we interchange the first ordered pair with the second. Therefore we also have  $(x_3, x_4) R (x_1, x_2)$ , and  $R$  is symmetric.

**Not transitive:** We provide a counterexample. Take

$$(1, 2) R (1, 3), \quad (1, 3) R (4, 3).$$

Indeed,

$$(1, 2, 1, 3) \text{ has } x_1 = x_3 = 1, \quad (1, 3, 4, 3) \text{ has } x_2 = x_4 = 3,$$

so in both cases the defining condition of  $R$  holds. However, for  $(1, 2)$  and  $(4, 3)$  we have

$$(1, 2, 4, 3),$$

whose four coordinates 1, 2, 4, 3 are pairwise distinct, so there are no  $i \neq j$  with  $x_i = x_j$ . Thus  $(1, 2) \not R (4, 3)$  and  $R$  is *not* transitive.

- (b) Suppose that  $T$  is a reflexive relation on a set  $A$  such that for all  $x, y, z \in A$ , if  $xTy$  and  $xTz$  then  $y = z$ . Show that  $T$  is an equivalence relation, and describe the equivalence classes of  $T$ . [6 marks]

## Solution

To prove that  $T$  is an equivalence relation, we must show that  $T$  is symmetric and transitive (it is given to be reflexive).

**$T$  is symmetric:** Let  $x, y \in A$  with  $xTy$ . Since  $T$  is reflexive,  $xTx$  also holds. By the given property, from  $xTy$  and  $xTx$  we obtain

$$y = x.$$

Hence  $x = y$ , and so  $yTx$  (because  $T$  is reflexive). Therefore  $T$  is symmetric.

**$T$  is transitive:** Let  $x, y, z \in A$  with  $xTy$  and  $yTz$ . Again, by reflexivity,  $xTx$  and  $yTy$  hold. From  $xTy$  and  $xTx$  we obtain  $y = x$ . From  $yTy$  and  $yTz$  we obtain  $z = y$ . Thus  $x = y = z$ , and so  $xTz$  (since  $T$  is reflexive). Hence  $T$  is transitive.

Because  $T$  is reflexive, symmetric and transitive,  $T$  is an equivalence relation on  $A$ .

**Equivalence classes:** Fix  $x \in A$ . If  $y \in [x]_T$ , then  $xTy$ . Using the argument above with  $xTy$  and  $xTx$ , we get  $y = x$ . Thus

$$[x]_T = \{x\}$$

for every  $x \in A$ . So all equivalence classes of  $T$  are singletons.

### Mark Scheme

- (a) Correctly identify reflexive, symmetric, not transitive (3). Clear justification of reflexivity and symmetry using the definition (3). Counterexample showing failure of transitivity, with explanation (4).
- (b) Use the given condition and reflexivity to prove symmetry (3). Use the condition again to prove transitivity (2). Correct description of equivalence classes as singletons  $\{x\}$  (1).