

MH1300 Foundations of Mathematics – Solutions

Final Examination, Academic Year 2021/2022, Semester 1

Compiled and typeset by QRS from the original handwritten solution

November 25, 2025

Overview of the 2021/2022 Semester 1 Paper

This typeset version is based on the original handwritten solutions from AY24/25 Sem 1, with minor corrections for clarity, consistent notation, and formatting. Alternative solution methods have been added for selected sub-questions where helpful (e.g. additional methods for Q1, Q3(b), Q4(a), Q4(b), and others). All mathematical content remains faithful to the original intent; any expanded explanations or supplementary methods are clearly marked as QRS notes.

Question 1**(18 marks)**

- (a) Let a and b be integers. Prove that $a - b$ and $a^2 + b^2$ have the same parity. (8 marks)
- (b) Is the following statement form a tautology, a contradiction, or neither?

$$(p \rightarrow (q \wedge \neg r)) \rightarrow (\neg q \rightarrow \neg p)$$

Justify your answer.

(6 marks)

- (c) Write down a tautology using only the statement variables s and t , and the connective \rightarrow . All three symbols must be used at least once in your statement form, and you cannot use any other variables or connectives. Prove that your statement form is a tautology. (4 marks)

Solution

- (a) Suppose a, b are arbitrary integers.

Method 1 (factorisation and cases).

We divide into two cases, according to the parity of $a - b$.

Case 1: $a - b$ is odd. Then

$$a + b = (a - b) + 2b = \text{odd} + \text{even} = \text{odd}.$$

Therefore

$$a^2 - b^2 = (a + b)(a - b) = \text{odd} \times \text{odd} = \text{odd}.$$

Hence

$$a^2 + b^2 = (a^2 - b^2) + 2b^2 = \text{odd} + \text{even} = \text{odd}.$$

So when $a - b$ is odd, $a^2 + b^2$ is also odd, and thus they have the same parity.

Case 2: $a - b$ is even. Then

$$a + b = (a - b) + 2b = \text{even} + \text{even} = \text{even},$$

so

$$a^2 - b^2 = (a + b)(a - b) = \text{even} \times \text{even} = \text{even},$$

and hence

$$a^2 + b^2 = (a^2 - b^2) + 2b^2 = \text{even} + \text{even} = \text{even}.$$

So when $a - b$ is even, $a^2 + b^2$ is also even.

In both cases, $a - b$ and $a^2 + b^2$ have the same parity.

Method 2 (four parity cases).

We divide into four cases for the parities of a and b :

- Case 1: a even, b even. Then $a - b$ is even, and $a^2 + b^2$ is even+even=even.
- Case 2: a even, b odd. Then $a - b$ is odd, and $a^2 + b^2$ is even+odd=odd.
- Case 3: a odd, b even. Then $a - b$ is odd, and $a^2 + b^2$ is odd+even=odd.
- Case 4: a odd, b odd. Then $a - b$ is even, and $a^2 + b^2$ is odd+odd=even.

In each case $a - b$ and $a^2 + b^2$ have the same parity, so the statement holds.

Method 3 (parity of squares).

Note that

$$\text{even}^2 = \text{even}, \quad \text{odd}^2 = \text{odd}.$$

Thus a and a^2 have the same parity, and $-b$ and b^2 have the same parity.

Hence $a - b$ (which has the same parity as $a + (-b)$) and $a^2 + b^2$ also have the same parity.

Therefore, in all methods, $a - b$ and $a^2 + b^2$ have the same parity.

Remarks. You may also argue using parity rules such as

$$\text{odd} \times \text{odd} = \text{odd}, \quad \text{odd} \times \text{even} = \text{even}, \quad \text{even} + \text{even} = \text{even}, \text{ etc.},$$

instead of always expanding from the definitions of even and odd.

(b) We show that the given statement form

$$(p \rightarrow (q \wedge \neg r)) \rightarrow (\neg q \rightarrow \neg p)$$

is a tautology.

Method 1 (truth table).

Construct the truth table with columns for p, q, r , then $q \wedge \neg r$, $p \rightarrow (q \wedge \neg r)$, $\neg q \rightarrow \neg p$, and finally the whole implication. One finds that in all eight rows the final column is T . Hence the statement is a tautology.

Method 2 (logical equivalences).

$$\begin{aligned}
 (p \rightarrow (q \wedge \neg r)) \rightarrow (\neg q \rightarrow \neg p) &\equiv \neg(p \rightarrow (q \wedge \neg r)) \vee (\neg q \rightarrow \neg p) && \text{(implication law)} \\
 &\equiv \neg(\neg p \vee (q \wedge \neg r)) \vee (\neg q \rightarrow \neg p) && \text{(implication law)} \\
 &\equiv (\neg\neg p \wedge \neg(q \wedge \neg r)) \vee (\neg q \rightarrow \neg p) && \text{(De Morgan)} \\
 &\equiv (p \wedge (\neg q \vee r)) \vee (\neg q \rightarrow \neg p) \\
 &\equiv (p \wedge (\neg q \vee r)) \vee (\neg\neg q \vee \neg p) && \text{(implication law)} \\
 &\equiv (p \wedge (\neg q \vee r)) \vee (q \vee \neg p) \\
 &\equiv (p \wedge (\neg q \vee r)) \vee ((q \vee \neg p) \vee (q \vee \neg p)) \\
 &\equiv (p \wedge (\neg q \vee r)) \vee T && \text{(negation / univ. bound)} \\
 &\equiv T.
 \end{aligned}$$

Thus the statement form is a tautology.

(c) One possible choice is

$$s \rightarrow (t \rightarrow t).$$

Proof that this is a tautology.

$$\begin{aligned} s \rightarrow (t \rightarrow t) &\equiv \neg s \vee (t \rightarrow t) && \text{(implication law)} \\ &\equiv \neg s \vee (\neg t \vee t) && \text{(implication law)} \\ &\equiv \neg s \vee T && \text{(negation law)} \\ &\equiv T && \text{(universal bound law).} \end{aligned}$$

Hence $s \rightarrow (t \rightarrow t)$ is a tautology.

Other correct choices are also acceptable, for example $s \rightarrow (t \rightarrow s)$, $(s \rightarrow s) \rightarrow (t \rightarrow t)$, or $(s \rightarrow t) \rightarrow (s \rightarrow t)$, together with valid proofs that they are tautologies.

Mark Scheme

(a) 8 marks.

- Correct interpretation/restatement of the claim in terms of parity. [1]
- Coherent proof strategy (e.g. case analysis on parity of a, b or $a - b$, or observation that squares preserve parity). [2]
- Correct parity reasoning in all required cases, including identifying the parity of $a - b$ and $a^2 + b^2$ in each case. [4]
- Clear concluding sentence explicitly stating that $a - b$ and $a^2 + b^2$ always have the same parity. [1]

(b) 6 marks.

- Correctly classifies the formula as a *tautology*. [1]
- Sets up an appropriate method: complete truth table *or* a correct chain of logical equivalences (using \rightarrow , \neg , \wedge , De Morgan, etc.). [2]
- Correctly carries the chosen method through to the end, showing that the final column is all T or that the formula is equivalent to \top . [3]

(c) 4 marks.

- Constructs a valid propositional formula using only s , t and \rightarrow , each appearing at least once, which is indeed a tautology. [2]
- Provides a correct justification (truth table or equivalence steps) that the chosen formula is a tautology. [2]

Question 2**(14 marks)**

Determine if each of the following is true or false. Justify your answers.

- (a) If a and b are composite numbers, then $a + b$ is composite. (3 marks)
- (b) For all positive integers c, d, e , if $c \mid e$ and $d \mid e$, then either $c = e$, $d = e$ or $cd \mid e$. (4 marks)
- (c) Let A be a subset of a set B . Then $A \times A \subseteq B \times B$. (3 marks)
- (d) If C, D and E are sets, then $(C \cup D) \cap E = C \cup (D \cap E)$. (4 marks)

Solution

- (a) **False.** Take $a = 4$ and $b = 9$. Then a and b are composite, since $a = 2 \cdot 2$ and $b = 3 \cdot 3$. But $a + b = 13$ is prime, so the statement is false.
- (b) **False.** Take $c = d = 4$ and $e = 8$. Then c, d, e are positive integers and

$$c \mid e \quad \text{and} \quad d \mid e,$$

since $4 \cdot 2 = 8$. However $c \neq e$, $d \neq e$, and $cd = 16$ does not divide $8 = e$. Hence the statement is false.

- (c) **True.** Assume $A \subseteq B$ and let $(a, b) \in A \times A$. Then $a \in A$ and $b \in A$. Since $A \subseteq B$, we have $a \in B$ and $b \in B$, so $(a, b) \in B \times B$. Thus $A \times A \subseteq B \times B$.
- (d) **False.** Take $C = \{0\}$ and $D = E = \emptyset$. Then

$$(C \cup D) \cap E = (\{0\} \cup \emptyset) \cap \emptyset = \{0\} \cap \emptyset = \emptyset,$$

while

$$C \cup (D \cap E) = \{0\} \cup (\emptyset \cap \emptyset) = \{0\} \cup \emptyset = \{0\}.$$

Hence $(C \cup D) \cap E \neq C \cup (D \cap E)$, so the statement is false.

Mark Scheme

- (a) 3 marks.
- Chooses valid composite numbers a, b (e.g. 4, 9) and notes they are composite. [1]
 - Correctly computes $a + b$ and identifies it as prime/non-composite. [1]
 - Explicitly states that this contradicts the claim, so the statement is false. [1]
- (b) 4 marks.
- Chooses positive integers c, d, e satisfying $c \mid e$ and $d \mid e$ (e.g. $c = d = 4$, $e = 8$). [1]

- Checks $c \neq e$ and $d \neq e$. [1]
- Computes cd and checks that $cd \nmid e$. [1]
- Concludes that the implication fails and the statement is false. [1]

(c) 3 marks.

- Correctly uses $A \subseteq B$ to deduce: if $a \in A$ then $a \in B$ (and same for b). [1]
- Shows that any $(a, b) \in A \times A$ is therefore in $B \times B$. [1]
- States the inclusion $A \times A \subseteq B \times B$ as the final conclusion. [1]

(d) 4 marks.

- Chooses explicit sets C, D, E (such as $C = \{0\}$, $D = E = \emptyset$). [1]
- Correctly computes $(C \cup D) \cap E$. [1]
- Correctly computes $C \cup (D \cap E)$. [1]
- Observes the two results differ, hence the equality is false. [1]

Question 3**(18 marks)**

(a) Let $F : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$ be a function satisfying

$$F(a, 0) = a, \quad F(a, b + 1) = F(a, b) + 1$$

for each $a, b \in \mathbb{N}$. Using the definition of F , prove that for every $a, b, c \in \mathbb{N}$,

$$F(F(a, b), c) = F(a, F(b, c)).$$

[8 marks]

Solution

Fix $a, b \in \mathbb{N}$ and prove the statement by induction on c .

Let $P(c)$ be the statement

$$F(F(a, b), c) = F(a, F(b, c)).$$

Base case ($c = 0$).

$$F(F(a, b), 0) = F(a, b) \quad (\text{by the rule with second argument } 0),$$

and

$$F(a, F(b, 0)) = F(a, b) \quad (\text{since } F(b, 0) = b).$$

So $P(0)$ holds.

Inductive step. Let $k \geq 0$ and suppose $P(k)$ holds, i.e.

$$F(F(a, b), k) = F(a, F(b, k)).$$

Then

$$\begin{aligned} F(F(a, b), k + 1) &= F(F(a, b), k) + 1 && (\text{definition of } F) \\ &= F(a, F(b, k)) + 1 && (\text{induction hypothesis}) \\ &= F(a, F(b, k) + 1) && (\text{since } F(a, d + 1) = F(a, d) + 1) \\ &= F(a, F(b, k + 1)) && (\text{definition of } F \text{ for } F(b, k + 1)). \end{aligned}$$

Thus $P(k + 1)$ holds.

By mathematical induction, $P(c)$ is true for all $c \in \mathbb{N}$, i.e.

$$F(F(a, b), c) = F(a, F(b, c))$$

for all $a, b, c \in \mathbb{N}$.

(b) Prove that for every positive integer m ,

$$4^{m+1} + 5^{2m-1}$$

is divisible by 21.

[10 marks]

Solution

Method 1 (mathematical induction).

Let $P(m)$ be the statement

$$4^{m+1} + 5^{2m-1} \text{ is divisible by } 21, \quad m \geq 1.$$

Base case ($m = 1$).

$$4^{1+1} + 5^{2 \cdot 1 - 1} = 4^2 + 5^1 = 16 + 5 = 21,$$

which is divisible by 21. So $P(1)$ is true.

Inductive step. Let $k \geq 1$ and assume $P(k)$ holds. Then there exists an integer x such that

$$4^{k+1} + 5^{2k-1} = 21x.$$

Consider

$$\begin{aligned} 4^{(k+1)+1} + 5^{2(k+1)-1} &= 4^{k+2} + 5^{2k+1} \\ &= 4 \cdot 4^{k+1} + 25 \cdot 5^{2k-1} \\ &= 4 \cdot 4^{k+1} + 4 \cdot 5^{2k-1} + 21 \cdot 5^{2k-1} \\ &= 4(4^{k+1} + 5^{2k-1}) + 21 \cdot 5^{2k-1}. \end{aligned}$$

Using the induction hypothesis,

$$4^{k+1} + 5^{2k-1} = 21x,$$

so

$$\begin{aligned} 4^{k+2} + 5^{2k+1} &= 4 \cdot 21x + 21 \cdot 5^{2k-1} \\ &= 21(4x + 5^{2k-1}), \end{aligned}$$

which is divisible by 21. Hence $P(k+1)$ holds.

By mathematical induction, $4^{m+1} + 5^{2m-1}$ is divisible by 21 for all positive integers m .

Method 2 (modular arithmetic).

We show $4^{m+1} + 5^{2m-1}$ is divisible by both 3 and 7.

Modulo 3. Since $4 \equiv 1 \pmod{3}$, we have $4^{m+1} \equiv 1 \pmod{3}$. Also $5 \equiv -1 \pmod{3}$, so

$$5^{2m-1} \equiv (-1)^{2m-1} = -1 \pmod{3}.$$

Thus

$$4^{m+1} + 5^{2m-1} \equiv 1 - 1 \equiv 0 \pmod{3}.$$

Modulo 7. Note that $25 \equiv 4 \pmod{7}$, so

$$5^{2m-1} = 5 \cdot 25^{m-1} \equiv 5 \cdot 4^{m-1} \pmod{7}.$$

Hence

$$\begin{aligned} 4^{m+1} + 5^{2m-1} &\equiv 4^{m+1} + 5 \cdot 4^{m-1} \\ &= 4^{m-1}(4^2 + 5) = 4^{m-1}(16 + 5) = 4^{m-1} \cdot 21 \equiv 0 \pmod{7}. \end{aligned}$$

Therefore $4^{m+1} + 5^{2m-1}$ is divisible by both 3 and 7. Since $\gcd(3, 7) = 1$, it is divisible by 21.

Mark Scheme

(a) 8 marks.

- Correctly fixes a, b and sets up induction on c (statement $P(c)$ written clearly). [2]
- Correct base case $c = 0$ with both sides computed using $F(a, 0) = a$. [2]
- Inductive step: correctly assumes $P(k)$, applies the recursive definition $F(a, b + 1) = F(a, b) + 1$ on both sides, and derives $P(k + 1)$. [4]

(b) 10 marks.

- Base case $m = 1$ computed correctly and checked for divisibility by 21. [2]
- Inductive step set up correctly: clear statement of the induction hypothesis and what needs to be shown for $m = k + 1$. [2]
- Algebraic manipulation to express $4^{k+2} + 5^{2k+1}$ in terms of $4^{k+1} + 5^{2k-1}$ (e.g. factoring out 4 and 25). [3]
- Correct use of the induction hypothesis to factor out 21, and conclusion that $4^{k+2} + 5^{2k+1}$ is divisible by 21. [3]
- * A fully correct and clearly explained modular arithmetic solution (showing divisibility by 3 and by 7) can receive up to full credit if the induction step is incomplete or omitted.

Question 4**(10 marks)**

(a) Let a, b, c, d be integers such that $d \mid a$, $d \mid b$ and $d \mid c$. Prove that

$$d^2 \mid ab + ac + bc.$$

[4 marks]

Solution**Method 1 (direct substitution).**

Since $d \mid a$, $d \mid b$ and $d \mid c$, there exist integers x, y, z such that

$$a = dx, \quad b = dy, \quad c = dz.$$

Then

$$\begin{aligned} ab + ac + bc &= (dx)(dy) + (dx)(dz) + (dy)(dz) \\ &= d^2xy + d^2xz + d^2yz \\ &= d^2(xy + xz + yz). \end{aligned}$$

Because $xy + xz + yz \in \mathbb{Z}$, we conclude that $d^2 \mid ab + ac + bc$.

Method 2 (using basic divisibility facts).

First note the lemma: if $d \mid r$ and $d \mid s$, then $d^2 \mid rs$. Indeed, write $r = dr_1$, $s = ds_1$ with $r_1, s_1 \in \mathbb{Z}$. Then

$$rs = (dr_1)(ds_1) = d^2r_1s_1,$$

so $d^2 \mid rs$.

Using this:

- From $d \mid a$ and $d \mid b$ we get $d^2 \mid ab$.
- From $d \mid a$ and $d \mid c$ we get $d^2 \mid ac$.
- From $d \mid b$ and $d \mid c$ we get $d^2 \mid bc$.

Now use the fact that if $d^2 \mid u$ and $d^2 \mid v$, then $d^2 \mid (u + v)$ (since $u = d^2u_1$, $v = d^2v_1 \Rightarrow u + v = d^2(u_1 + v_1)$).

First, $d^2 \mid ab$ and $d^2 \mid ac$ imply $d^2 \mid (ab + ac)$. Then $d^2 \mid (ab + ac)$ and $d^2 \mid bc$ imply

$$d^2 \mid (ab + ac) + bc = ab + ac + bc.$$

Thus $d^2 \mid ab + ac + bc$.

- (b) Prove that every non-zero rational number is the product of two irrational numbers. You may use the fact that the product of a non-zero rational number with an irrational number is irrational. [6 marks]

Solution

Let $r \in \mathbb{Q}$ with $r \neq 0$.

Method 1 (direct construction).

Set

$$x = r\sqrt{2}, \quad y = \frac{1}{\sqrt{2}}.$$

Then x is irrational, because it is the product of the non-zero rational number r and the irrational number $\sqrt{2}$. Similarly,

$$y = \frac{1}{\sqrt{2}} = \frac{\sqrt{2}}{2} = \left(\frac{1}{2}\right)\sqrt{2}$$

is the product of a non-zero rational and an irrational, so y is irrational.

Moreover,

$$xy = r\sqrt{2} \cdot \frac{1}{\sqrt{2}} = r.$$

Thus every non-zero rational r can be written as a product of two irrational numbers.

Method 2 (proof by contradiction).

Suppose, for contradiction, that there exists a non-zero rational number r that cannot be expressed as the product of two irrational numbers.

Consider again

$$x = r\sqrt{2}, \quad y = \frac{1}{\sqrt{2}}.$$

By the given fact, both x and y are irrational (each is the product of a non-zero rational and an irrational). But then

$$xy = r\sqrt{2} \cdot \frac{1}{\sqrt{2}} = r,$$

so r is the product of two irrational numbers, contradicting our assumption.

Hence our assumption was false, and every non-zero rational number is indeed the product of two irrational numbers.

Mark Scheme

(a) 4 marks.

- Uses $d \mid a, d \mid b, d \mid c$ to introduce integers x, y, z with $a = dx, b = dy, c = dz$ (or equivalent lemma $d \mid r, d \mid s \Rightarrow d^2 \mid rs$). [2]
- Correct substitution and simplification showing $ab + ac + bc = d^2(\dots)$ and concluding $d^2 \mid ab + ac + bc$. [2]

(b) 6 marks.

- Chooses explicit expressions $x = r\sqrt{2}$ and $y = 1/\sqrt{2}$ (or another valid pair of irrationals depending on r). [2]
- Correctly argues that x is irrational using the given fact about (non-zero rational) \times (irrational). [2]
- Correctly argues that y is also irrational (e.g. $y = \frac{1}{2}\sqrt{2}$). [1]
- Computes xy and shows $xy = r$, then concludes that every non-zero rational is a product of two irrationals. [1]

Question 5**(12 marks)**

- (a) Let A and B be sets and $f : A \rightarrow B$ be a function. Let $X \subseteq A$ and $Y \subseteq B$. Prove that

$$f(f^{-1}(Y)) \subseteq Y \quad \text{and} \quad X \subseteq f^{-1}(f(X)).$$

[6 marks]

- (b) Let $g : \mathbb{Z} \rightarrow \mathbb{Z}$ be the function $g(n) = 2n \bmod 3$.

- (i) Is g injective?
- (ii) Is g surjective?
- (iii) What is the range of g ?

Justify your answers.

[6 marks]

Solution

- (a) We first show that $f(f^{-1}(Y)) \subseteq Y$.

Let $y \in f(f^{-1}(Y))$. Then, by definition of $f(f^{-1}(Y))$, there exists some $x \in f^{-1}(Y)$ such that $y = f(x)$. Since $x \in f^{-1}(Y)$, by definition of preimage we know $f(x) \in Y$. Hence $y = f(x) \in Y$. Therefore $f(f^{-1}(Y)) \subseteq Y$.

Next, we show that $X \subseteq f^{-1}(f(X))$.

Let $x \in X$. Then $f(x) \in f(X)$ (since $f(X) = \{f(a) \in B \mid a \in X\}$). Hence, by the definition of preimage, $x \in f^{-1}(f(X))$. Therefore $X \subseteq f^{-1}(f(X))$.

- (b) For $g : \mathbb{Z} \rightarrow \mathbb{Z}$ defined by $g(n) = 2n \bmod 3$ we first understand its values.

For any $n \in \mathbb{Z}$, we can write

$$2n = 3(2n \operatorname{div} 3) + (2n \bmod 3), \quad 0 \leq 2n \bmod 3 < 3.$$

Thus $g(n) = 2n \bmod 3$ is always one of 0, 1, 2, so

$$\operatorname{range}(g) \subseteq \{0, 1, 2\}.$$

Now we show the reverse inclusion:

$$g(0) = 0 \bmod 3 = 0, \quad g(1) = 2 \bmod 3 = 2, \quad g(2) = 4 \bmod 3 = 1.$$

Hence $0, 1, 2 \in \operatorname{range}(g)$, so

$$\operatorname{range}(g) = \{0, 1, 2\}.$$

- (i) g is *not injective*, since

$$0 \neq 3, \quad g(0) = 0 \bmod 3 = 0, \quad g(3) = 6 \bmod 3 = 0,$$

so $g(0) = g(3)$ with $0 \neq 3$.

- (ii) g is *not surjective*, because for example $3 \notin \{0, 1, 2\} = \text{range}(g)$.
- (iii) The range of g is $\{0, 1, 2\}$, as shown above.

Mark Scheme

- (a) Correct proof that $f(f^{-1}(Y)) \subseteq Y$ using element chase. (3 marks)
Correct proof that $X \subseteq f^{-1}(f(X))$. (3 marks)
- (b) – Correct identification that g is not injective with a valid counterexample. (2 marks)
 - Correct identification that g is not surjective with justification (e.g. 3 is not in the range). (2 marks)
 - Correct determination and justification that $\text{range}(g) = \{0, 1, 2\}$. (2 marks)

Remarks

The definitions used are:

$$f(X) = \{ f(a) \in B \mid a \in X \}, \quad f^{-1}(Y) = \{ a \in A \mid f(a) \in Y \}.$$

Question 6**(12 marks)**

- (a) Find all sixth roots of unity. That is, find all complex numbers z satisfying $z^6 - 1 = 0$.
Leave your answer in terms of $re^{i\theta}$. [6 marks]
- (b) Determine if the following argument is valid. State all rules of inference used.

$$\begin{aligned} &\neg p \vee \neg q \\ &\neg r \rightarrow (p \wedge q) \\ &\neg r \vee s \\ &s \rightarrow (t \wedge u) \\ &\therefore u \end{aligned}$$

[6 marks]

Solution

- (a) The sixth roots of unity are the solutions to $z^6 = 1$.
Write 1 in polar form as $1 = e^{i \cdot 0}$. Using the general n th root formula

$$z = r^{1/n} e^{i(\theta + 2k\pi)/n}, \quad k = 0, 1, \dots, n-1,$$

with $r = 1$, $\theta = 0$ and $n = 6$, we obtain

$$z = e^{i(0+2k\pi)/6} = e^{i\frac{2k\pi}{6}}, \quad k = 0, 1, 2, 3, 4, 5.$$

Thus the six roots are

$$z = e^{i \cdot 0} = 1, \quad e^{i\pi/3}, \quad e^{i2\pi/3}, \quad e^{i\pi}, \quad e^{i4\pi/3}, \quad e^{i5\pi/3},$$

each having modulus $r = 1$ and argument $\theta = \frac{2k\pi}{6}$ for $k = 0, \dots, 5$.

- (b) We show that the argument is valid by deriving u from the premises.

- | | | |
|------|-----------------------------------|--|
| (1) | $\neg p \vee \neg q$ | (Premise #1) |
| (2) | $\neg(p \wedge q)$ | (from (1) by De Morgan's Law) |
| (3) | $\neg r \rightarrow (p \wedge q)$ | (Premise #2) |
| (4) | $\neg \neg r$ | (from (2) and (3) by Modus Tollens) |
| (5) | r | (from (4) by Double Negation) |
| (6) | $\neg r \vee s$ | (Premise #3) |
| (7) | s | (from (5) and (6) by <i>Disjunctive Elimination</i> / Elimination) |
| (8) | $s \rightarrow (t \wedge u)$ | (Premise #4) |
| (9) | $t \wedge u$ | (from (7) and (8) by Modus Ponens) |
| (10) | u | (from (9) by <i>Conjunction Elimination</i> / Specialisation) |

Since u has been derived from the premises using valid rules of inference, the argument is valid.

Mark Scheme

- (a) – Correct use of polar form and the general n th root formula. (2 marks)
 - Expression $z = e^{i2k\pi/6}$ with correct range of k . (2 marks)
 - Correct explicit list of the six roots in the requested form. (2 marks)
- (b) – Correct transformation of $\neg p \vee \neg q$ to $\neg(p \wedge q)$ by De Morgan. (1 mark)
 - Correct use of Modus Tollens (or equivalent reasoning) to obtain r . (2 marks)
 - Correct use of $\neg r \vee s$ with r to deduce s . (1 mark)
 - Correct use of $s \rightarrow (t \wedge u)$ to get $t \wedge u$, and then u . (2 marks)

Question 7**(16 marks)**

- (a) For the relation
- R
- on
- \mathbb{R}^2
- defined by

$$(x_1, x_2) R (x_3, x_4) \iff \exists i \neq j \text{ with } x_i = x_j, \quad i, j \in \{1, 2, 3, 4\},$$

determine whether R is reflexive, symmetric and transitive. Justify your answers.

[10 marks]

QRS Note: The original question stated “for some $i \neq j$ ”. Based on the official solution, this condition is interpreted as an existential condition (“there exists such i, j ”), not a universal one. Thus the notation has been rewritten using \exists for clarity.

SolutionWrite (x_1, x_2) and (x_3, x_4) for arbitrary elements of \mathbb{R}^2 .

Reflexive: Let $(x, y) \in \mathbb{R}^2$. Then, when we consider $(x_1, x_2, x_3, x_4) = (x, y, x, y)$, we have

$$x_1 = x_3 \quad (\text{and also } x_2 = x_4).$$

Thus there exist $i \neq j$ (for example $i = 1, j = 3$) with $x_i = x_j$. Hence $(x, y)R(x, y)$, so R is reflexive.

Symmetric: Suppose $(x_1, x_2)R(x_3, x_4)$. Then, by definition, there exist $i \neq j$ with $x_i = x_j$, where $i, j \in \{1, 2, 3, 4\}$. But the condition “some coordinates among x_1, x_2, x_3, x_4 are equal” is unchanged if we interchange the first ordered pair with the second. Therefore we also have $(x_3, x_4)R(x_1, x_2)$, and R is symmetric.

Not transitive: We provide a counterexample. Take

$$(1, 2)R(1, 3), \quad (1, 3)R(4, 3).$$

Indeed,

$$(1, 2, 1, 3) \text{ has } x_1 = x_3 = 1, \quad (1, 3, 4, 3) \text{ has } x_2 = x_4 = 3,$$

so in both cases the defining condition of R holds. However, for $(1, 2)$ and $(4, 3)$ we have

$$(1, 2, 4, 3),$$

whose four coordinates 1, 2, 4, 3 are pairwise distinct, so there are no $i \neq j$ with $x_i = x_j$. Thus $(1, 2) \not R (4, 3)$ and R is *not* transitive.

- (b) Suppose that T is a reflexive relation on a set A such that for all $x, y, z \in A$, if xTy and xTz then $y = z$. Show that T is an equivalence relation, and describe the equivalence classes of T . [6 marks]

Solution

To prove that T is an equivalence relation, we must show that T is symmetric and transitive (it is given to be reflexive).

T is symmetric: Let $x, y \in A$ with xTy . Since T is reflexive, xTx also holds. By the given property, from xTy and xTx we obtain

$$y = x.$$

Hence $x = y$, and so yTx (because T is reflexive). Therefore T is symmetric.

T is transitive: Let $x, y, z \in A$ with xTy and yTz . Again, by reflexivity, xTx and yTy hold. From xTy and xTx we obtain $y = x$. From yTy and yTz we obtain $z = y$. Thus $x = y = z$, and so xTz (since T is reflexive). Hence T is transitive.

Because T is reflexive, symmetric and transitive, T is an equivalence relation on A .

Equivalence classes: Fix $x \in A$. If $y \in [x]_T$, then xTy . Using the argument above with xTy and xTx , we get $y = x$. Thus

$$[x]_T = \{x\}$$

for every $x \in A$. So all equivalence classes of T are singletons.

Mark Scheme

- (a) Correctly identify reflexive, symmetric, not transitive (3). Clear justification of reflexivity and symmetry using the definition (3). Counterexample showing failure of transitivity, with explanation (4).
- (b) Use the given condition and reflexivity to prove symmetry (3). Use the condition again to prove transitivity (2). Correct description of equivalence classes as singletons $\{x\}$ (1).