



Exceptional Control Flow: Signals and Nonlocal Jumps

These slides adapted from materials provided by the textbook authors.

Signals and Nonlocal Jumps

- Review
- Shells
- Signals
 - Sending Signals
 - Receiving Signals
- Nonlocal jumps

Receiving Signals

 Suppose kernel is returning from an exception handler and is ready to pass control to process p

- Kernel computes pnb = pending & ~blocked
 - The set of pending nonblocked signals for process p
- If (pnb == 0)
 - Pass control to next instruction in the logical flow for p
- Else
 - Choose least nonzero bit k in pnb and force process p to receive signal k
 - The receipt of the signal triggers some action by p
 - Repeat for all nonzero k in pnb
 - Pass control to next instruction in logical flow for p

Default Actions

- Each signal type has a predefined default action, which is one of:
 - The process terminates
 - The process stops until restarted by a SIGCONT signal
 - The process ignores the signal

Installing Signal Handlers

- The signal function modifies the default action associated with the receipt of signal signum:
 - handler_t *signal(int signum, handler_t *handler)

Different values for handler:

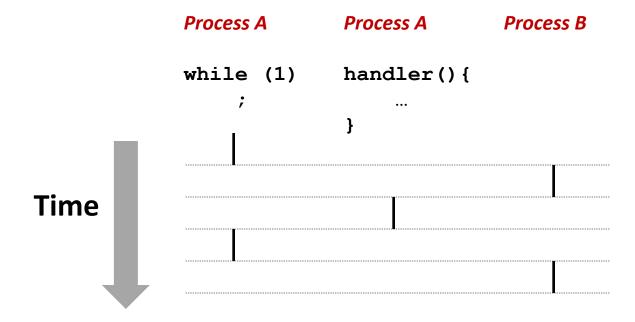
- SIG_IGN: ignore signals of type signum
- SIG_DFL: revert to the default action on receipt of signals of type signum
- Otherwise, handler is the address of a user-level signal handler
 - Called when process receives signal of type signum
 - Referred to as "installing" the handler
 - Executing handler is called "catching" or "handling" the signal
 - When the handler executes its return statement, control passes back to instruction in the control flow of the process that was interrupted by receipt of the signal

Signal Handling Example

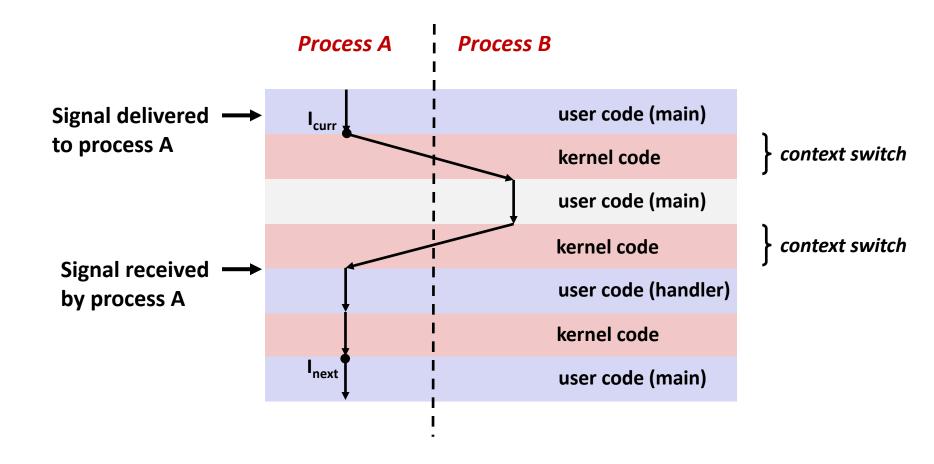
```
void sigint_handler(int sig) /* SIGINT handler */
{
    printf("So you think you can stop the bomb with ctrl-c, do you?\n");
    sleep(2);
    printf("Well...");
    fflush(stdout);
    sleep(1);
    printf("OK. :-)\n");
    exit(0);
int main()
{
    /* Install the SIGINT handler */
    if (signal(SIGINT, sigint_handler) == SIG_ERR)
        unix_error("signal error");
    /* Wait for the receipt of a signal */
    pause();
    return 0;
                                                                     sigint.c
```

Signals Handlers as Concurrent Flows

 A signal handler is a separate logical flow (not process) that runs concurrently with the main program

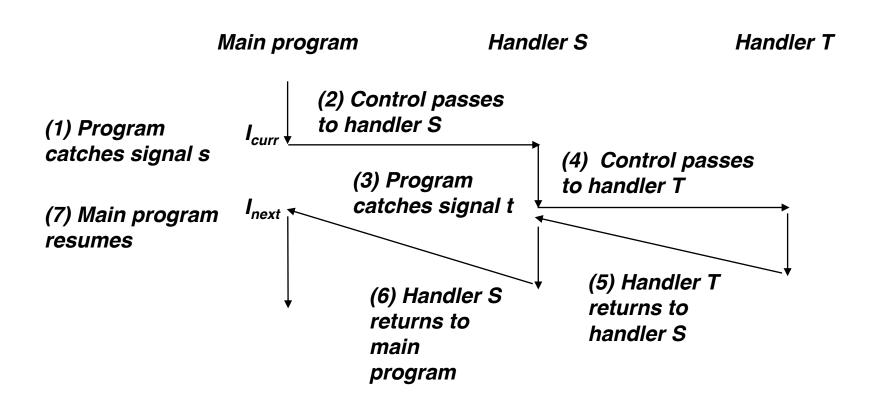


Another View of Signal Handlers as Concurrent Flows



Nested Signal Handlers

Handlers can be interrupted by other handlers



Blocking and Unblocking Signals

Implicit blocking mechanism

- Kernel blocks any pending signals of type currently being handled.
- E.g., A SIGINT handler can't be interrupted by another SIGINT

Explicit blocking and unblocking mechanism

sigprocmask function

Supporting functions

- sigemptyset Create empty set
- sigfillset Add every signal number to set
- sigaddset Add signal number to set
- sigdelset Delete signal number from set

Temporarily Blocking Signals

```
sigset_t mask, prev_mask;
Sigemptyset(&mask);
Sigaddset(&mask, SIGINT);

/* Block SIGINT and save previous blocked set */
Sigprocmask(SIG_BLOCK, &mask, &prev_mask);

/* Code region that will not be interrupted by SIGINT */

/* Restore previous blocked set, unblocking SIGINT */
Sigprocmask(SIG_SETMASK, &prev_mask, NULL);
```

Safe Signal Handling

- Handlers are tricky because they are concurrent with main program and share the same global data structures.
 - Shared data structures can become corrupted.
- We'll explore concurrency issues later in the term.
- For now here are some guidelines to help you avoid trouble.

Guidelines for Writing Safe Handlers

- G0: Keep your handlers as simple as possible
 - e.g., Set a global flag and return
- G1: Call only async-signal-safe functions in your handlers
 - printf, sprintf, malloc, and exit are not safe!
- G2: Save and restore errno on entry and exit
 - So that other handlers don't overwrite your value of errno
- G3: Protect accesses to shared data structures by temporarily blocking all signals.
 - To prevent possible corruption
- G4: Declare global variables as volatile
 - To prevent compiler from storing them in a register
- G5: Declare global flags as volatile sig atomic t
 - flag: variable that is only read or written (e.g. flag = 1, not flag++)
 - Flag declared this way does not need to be protected like other globals

Async-Signal-Safety

- Function is async-signal-safe if either reentrant (e.g., all variables stored on stack frame, CS:APP3e 12.7.2) or non-interruptible by signals.
- Posix guarantees 117 functions to be async-signal-safe
 - Source: "man 7 signal"
 - Popular functions on the list:
 - _exit, write, wait, waitpid, sleep, kill
 - Popular functions that are **not** on the list:
 - printf, sprintf, malloc, exit
 - Unfortunate fact: write is the only async-signal-safe output function

Safely Generating Formatted Output

Use the reentrant SIO (Safe I/O library) from csapp.c in your handlers.

```
ssize_t sio_puts(char s[]) /* Put string */
ssize_t sio_putl(long v) /* Put long */
void sio_error(char s[]) /* Put msg & exit */
```

```
void sigint_handler(int sig) /* Safe SIGINT handler */
{
    Sio_puts("So you think you can stop the bomb with ctrl-
c, do you?\n");
    sleep(2);
    Sio_puts("Well...");
    sleep(1);
    Sio_puts("OK. :-)\n");
    _exit(0);
}
```

Correct Signal Handling

```
int ccount = 0;
void child_handler(int sig) {
    int olderrno = errno;
    pid_t pid;
    if ((pid = wait(NULL)) < 0)</pre>
        Sio_error("wait error");
    ccount--:
    Sio_puts("Handler reaped child ");
    Sio_putl((long)pid);
    Sio_puts(" \n");
    sleep(1);
    errno = olderrno;
void fork14() {
    pid_t pid[N];
    int i;
    ccount = N:
    Signal(SIGCHLD, child_handler);
    for (i = 0; i < N; i++) {
        if ((pid[i] = Fork()) == 0) {
            Sleep(1);
            exit(0); /* Child exits */
        }
    while (ccount > 0) /* Parent spins */
```

```
> ./forks 14
Handler reaped child 23240
Handler reaped child 23241
```

```
int ccount = 0;
void child_handler(int sig) {
    int olderrno = errno:
    pid_t pid;
    if ((pid = wait(NULL)) < 0)</pre>
        Sio_error("wait error");
    ccount--:
    Sio_puts("Handler reaped child ");
    Sio_putl((long)pid);
    Sio_puts(" \n");
    sleep(1);
    errno = olderrno;
void fork14() {
    pid_t pid[N];
    int i;
    ccount = N;
    Signal(SIGCHLD, child_handler);
    for (i = 0; i < N; i++) {
        if ((pid[i] = Fork()) == 0) {
            Sleep(1);
            exit(0); /* Child exits */
        }
    while (ccount > 0) /* Parent spins */
```

Signal Handling

- Pending signals are not queued
 - For each signal type, one bit indicates whether or not signal is pending...
 - ...thus at most one pending signal of any particular type.
- You can't use signals to count events, such as children terminating.

```
> ./forks 14
Handler reaped child 23240
Handler reaped child 23241
```

Correct Signal Handling

- Must wait for all terminated child processes
 - Put wait in a loop to reap all terminated children

```
void child_handler2(int sig)
{
    pid_t pid;
    while ((pid = wait(NULL)) > 0) {
         ccount--;
                                > ./forks 15
                                Handler reaped child 23246
                                Handler reaped child 23247
                                Handler reaped child 23248
                                Handler reaped child 23249
                                Handler reaped child 23250
                                >
```

Synchronizing Flows to Avoid Races

```
int main(int argc, char **argv)
{
    int pid:
    sigset_t mask_all, prev_all;
    Sigfillset(&mask_all);
    Signal(SIGCHLD, handler);
    initjobs(); /* Initialize the job list */
    while (1) {
        if ((pid = Fork()) == 0) { /* Child */
            Execve("/bin/date", argv, NULL);
        Sigprocmask(SIG_BLOCK, &mask_all, &prev_all); /* Parent */
        addjob(pid); /* Add the child to the job list */
        Sigprocmask(SIG_SETMASK, &prev_all, NULL);
    exit(0);
                                                          procmask1.c
```

Synchronizing Flows to Avoid Races

Simple shell with a subtle synchronization error because it assumes parent runs before child.

```
int main(int argc, char **argv)
{
    int pid;
    sigset_t mask_all, prev_all;
    Sigfillset(&mask_all);
    Signal(SIGCHLD, handler);
    initjobs(); /* Initialize the job list */
    while (1) {
        if ((pid = Fork()) == 0) { /* Child */
            Execve("/bin/date", argv, NULL);
        Sigprocmask(SIG_BLOCK, &mask_all, &prev_all); /* Parent */
        addjob(pid); /* Add the child to the job list */
        Sigprocmask(SIG_SETMASK, &prev_all, NULL);
    exit(0);
                                                          procmask1.c
```

Corrected Shell Program without Race

```
int main(int argc, char **argv)
{
    int pid;
    sigset_t mask_all, mask_one, prev_one;
    Sigfillset(&mask_all);
    Sigemptyset(&mask_one);
    Sigaddset(&mask_one, SIGCHLD);
    Signal(SIGCHLD, handler);
    initjobs(); /* Initialize the job list */
    while (1) {
        Sigprocmask(SIG_BLOCK, &mask_one, &prev_one); /* Block SIGCHLD */
        if ((pid = Fork()) == 0) { /* Child process */
            Sigprocmask(SIG_SETMASK, &prev_one, NULL); /* Unblock SIGCHLD */
            Execve("/bin/date", argv, NULL);
        Sigprocmask(SIG_BLOCK, &mask_all, NULL); /* Parent process */
        addjob(pid); /* Add the child to the job list */
        Sigprocmask(SIG_SETMASK, &prev_one, NULL); /* Unblock SIGCHLD */
    exit(0);
                                                                   procmask2.c
```

Explicitly Waiting for Signals

Handlers for program explicitly waiting for SIGCHLD to arrive.

```
volatile sig_atomic_t pid;

void sigchld_handler(int s)
{
    int olderrno = errno;
    pid = Waitpid(-1, NULL, 0); /* Main is waiting for nonzero pid */
    errno = olderrno;
}

void sigint_handler(int s)
{
}

waitforsignal.c
```

Explicitly Waiting for Signals

```
int main(int argc, char **argv) {
                                                   Similar to a shell waiting
    sigset_t mask, prev;
                                                   for a foreground job to
    Signal(SIGCHLD, sigchld_handler);
                                                   terminate.
    Signal(SIGINT, sigint_handler);
    Sigemptyset(&mask);
    Sigaddset(&mask, SIGCHLD);
    while (1) {
        Sigprocmask(SIG_BLOCK, &mask, &prev); /* Block SIGCHLD */
        if (Fork() == 0) /* Child */
                         exit(0):
        /* Parent */
        pid = 0;
        Sigprocmask(SIG_SETMASK, &prev, NULL); /* Unblock SIGCHLD */
        /* Wait for SIGCHLD to be received (wasteful!) */
        while (!pid);
        /* Do some work after receiving SIGCHLD */
        printf(".");
    exit(0);
                                                           waitforsignal.c
}
```

Explicitly Waiting for Signals

- Program is correct, but very wasteful
- Other options:

```
while (!pid) /* Race! */
  pause();
```

```
while (!pid) /* Too slow! */
    sleep(1);
```

Solution: sigsuspend

Waiting for Signals with sigsuspend

- int sigsuspend(const sigset_t *mask)
- Equivalent to atomic (uninterruptable) version of:

```
sigprocmask(SIG_BLOCK, &mask, &prev);
pause();
sigprocmask(SIG_SETMASK, &prev, NULL);
```

Waiting for Signals with sigsuspend

```
int main(int argc, char **argv) {
    sigset t mask, prev;
    Signal(SIGCHLD, sigchld handler);
    Signal(SIGINT, sigint handler);
    Sigemptyset(&mask);
    Sigaddset(&mask, SIGCHLD);
   while (1) {
        Sigprocmask(SIG BLOCK, &mask, &prev); /* Block SIGCHLD */
        if (Fork() == 0) /* Child */
            exit(0);
       /* Wait for SIGCHLD to be received */
       pid = 0;
        while (!pid)
            Sigsuspend(&prev);
       /* Optionally unblock SIGCHLD */
        Sigprocmask(SIG SETMASK, &prev, NULL);
        /* Do some work after receiving SIGCHLD */
       printf(".");
    exit(0);
```