



College of Engineering & Applied Sciences

CSPB 2824

Discrete Structures

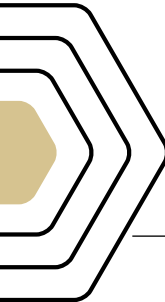
Proof Techniques Mastery Workbook

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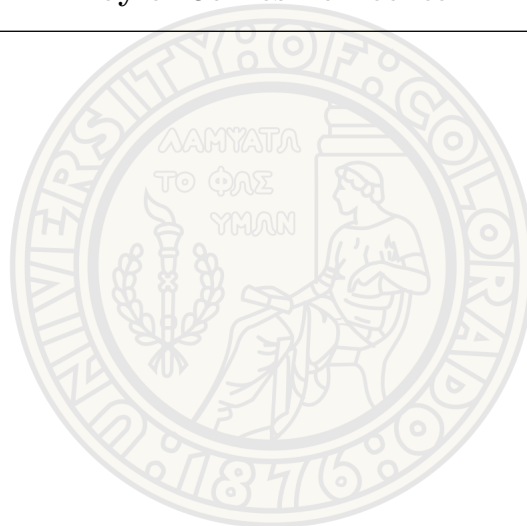
Mastery Workbook 3



Proof Techniques Workbook

I have neither given nor received unauthorized assistance.

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Problem 0 - Warm Up THMs

Problem Statement

- (a) Prove: For all positive integers, A, B , if A is even and B is odd, $A + B = C$ is odd.
- (b) Prove: For all positive integers, A, B , if A is odd and B is odd, $A + B = C$ is even.
- (c) Prove: For all positive integers, A, B , if A and B are even, $A \cdot B = C$ is even.
- (d) Prove: For all positive integers, A, B , if A is even B is odd, $A \cdot B = C$ is even.

Solution - Part (a)

Theorem: For all positive integers, A, B , if A is even and B is odd, $A + B = C$ is odd.

Direct Proof:

There is k such that $A = 2k$	(By definition of even)
There is l such that $B = 2l + 1$	(By definition of odd)
$C = A + B$	(Premise)
$= 2k + 2l + 1$	(Substitution)
$= 2(k + l) + 1$	(Factoring)
Let α integer $= k + l$	(Closure of integers renaming $k + l$)
$= 2\alpha + 1$	(Substitution)

Consequently, $C = 2\alpha + 1$ with α is an integer, and thus by the **definition of odd** C is odd. \square

Solution - Part (b)

Theorem: For all positive integers, A, B , if A is odd and B is odd, $A + B = C$ is even.

Direct Proof:

There is k such that $A = 2k + 1$	(By definition of odd)
There is l such that $B = 2l + 1$	(By definition of odd)
$C = A + B$	(Premise)
$= 2k + 1 + 2l + 1$	(Substitution)
$= 2k + 2l + 1 + 1$	(Commutativity)
$= 2k + 2l + 2$	(Simplification by addition)
$= 2(k + l) + 2$	(Factoring)
Let α integer $= k + l$	(Closure of integers renaming $k + l$)
$= 2\alpha + 2$	(Substitution)
$= 2(\alpha + 1)$	(Factoring)
Let β integer $= \alpha + 1$	(Closure of integers renaming $\alpha + 1$)
$= 2\beta$	(Substitution)

Consequently, $C = 2\beta$ with β is an integer, and thus by the **definition of even** C is even. \square

Solution - Part (c)

Theorem: For all positive integers, A, B , if A and B are even, $A \cdot B = C$ is even.

Direct Proof:

There is k such that $A = 2k$	(By definition of even)
There is l such that $B = 2l$	(By definition of even)
$C = A \cdot B$	(Premise)
$= (2k)(2l)$	(Substitution)
$= (2)(2)(k)(l)$	(Commutativity)
$= 4(k)(l)$	(Simplification by multiplication)
$= 4kl$	(Associativity)
$= 2(2kl)$	(Factoring)
Let α integer $= 2kl$	(Closure of integers renaming $2kl$)
$= 2\alpha$	(Substitution)

Consequently, $C = 2\alpha$ with α is an integer, and thus by the **definition of even** C is even. \square

Solution - Part (d)

Theorem: For all positive integers, A, B , if A is even B is odd, $A \cdot B = C$ is even.

Direct Proof:

There is k such that $A = 2k$	(By definition of even)
There is l such that $B = 2l + 1$	(By definition of odd)
$C = A \cdot B$	(Premise)
$= (2k)(2l + 1)$	(Substitution)
$= (2k)(2l) + (2k)(1)$	(Distributivity)
$= (2)(2)(k)(l) + (2k)(1)$	(Commutativity)
$= 4(k)(l) + 2k$	(Simplification by multiplication)
$= 4kl + 2k$	(Associativity)
$= 2(kl + k)$	(Factoring)
Let α integer $= kl + k$	(Closure of integers renaming $kl + k$)
$= 2\alpha$	(Substitution)

Consequently, $C = 2\alpha$ with α is an integer, and thus by the **definition of even** C is even. \square

Problem 0 Summary

Procedure

- (a) Part (a)
 - Use the definitions of odd and even and make a definition with the use of closure
 - Use the respective properties, make substitutions, and arrive at the correct conclusion
- (b) Part (b)
 - Same as part (a)
- (c) Part (c)
 - Same as parts (a) and (b)
- (d) Part (d)
 - Same as parts (a)-(c)

Key Concepts

- These proofs are done with the method of the direct proof format
- These proofs use the definition of even ($n = 2k$) and the definition of odd ($n = 2k + 1$)
- These proofs use simplification by addition, subtraction, and multiplication
- These proofs use properties such as associativity, commutativity, and distributivity, and closure properties

Variations

- We could be asked to prove a different theorem with different initial values
 - In this case we would use the same procedure but with a different initial theorem

Problem 1

Problem Statement

Prove: For all positive integers, A , B , if A and B are odd, $A \cdot B = C$ is odd.

Solution

Theorem: For all positive integers, A , B , if A and B are odd, $A \cdot B = C$ is odd.

Direct Proof:

There is k such that $A = 2k + 1$	(By definition of odd)
There is l such that $B = 2l + 1$	(By definition of odd)
$C = A \cdot B$	(Premise)
$= (2k + 1)(2l + 1)$	(Substitution)
$= (2k + 1)(2l) + (2k + 1)(1)$	(Distributivity)
$= (2l)(2k + 1) + (1)(2k + 1)$	(Commutativity)
$= (2l)(2k) + (2k)(1) + (1)(2k) + (1)(1)$	(Distributivity)
$= (2)(2)(k)(l) + (2k)(1) + (2k)(1) + (1)(1)$	(Commutativity)
$= (4)(k)(l) + 2k + 2k + 1$	(Simplification by multiplication)
$= 4kl + 2k + 2k + 1$	(Associativity)
$= 4kl + 4k + 1$	(Simplification by addition)
$= 4(kl + k) + 1$	(Factoring)
Let α integer $= kl + k$	(Closure by renaming $kl + k$)
$= 4\alpha + 1$	(Substitution)
$= 2(2\alpha) + 1$	(Factoring)
Let β integer $= 2\alpha$	(Closure by renaming 2α)
$= 2\beta + 1$	(Substitution)

Consequently, C (the quantity $A \cdot B$) with β is an integer, and thus by the **definition of odd** C is **odd**. \square

Problem 1 Summary

Procedure

- Use the direct proof format
- Use the definitions of odd and even along with field axioms of algebra to reach the conclusion

Key Concepts

- This problem uses the direct proof format to prove that a quantity is odd based off of the hypothesis' being odd

Variations

- The hypothesis' could change
 - In this case we would have to work through the problem as normal with the new hypothesis'
- The conclusion could change
 - In this case we would have to show if the conclusion is true or false



Problem 2

Problem Statement

Prove that if x is odd, then $5x + 3$ is even two different ways.

- (a) A direct proof by **applying the Warm Up THMs** - this means just USE the RESULTS of what you have posted above.
- (b) Using the *techniques* of the warm up THM and the formal definitions of even and odd. (This will look a lot like your WUT proofs)

Solution - Part (a)

Theorem: If x is odd, then $5x + 3$ is even.

Direct Proof: Using the axioms from the warm up theorems.

There exists an integer k such that $5 = 2k + 1$	(By definition of odd)
There exists an integer l such that $x = 2l + 1$	(By definition of odd)
There exists an integer m such that $3 = 2m + 1$	(By definition of odd)
Let $A = 5$ and $B = x$ integer	(Closure by renaming 5 and x)
Let $C = A \cdot B$ integer	(Closure by renaming $A \cdot B$)
Then C is odd	(By WUT # 6)
Let $D = 3$	(Closure by renaming 3)
Let $E = C + D$ integer	(Closure by renaming $C + D$)
Then E is even	(By WUT # 3)

Consequently, E (the quantity $5x + 3$) an integer, and thus by the **WUT's** E is **even**. \square

Solution - Part (b)

Theorem: If x is odd, then $5x + 3$ is even.

Direct Proof:

There exists an integer k such that $x = 2k + 1$	(By definition of odd)
Let $C = 5x + 3$ integer	(Closure by renaming $5x + 3$)
$C = 5x + 3$	(Premise)
$= 5(2k + 1) + 3$	(Substitution)
$= (5)(2k) + (5)(1) + 3$	(Distributivity)
$= (5)(2)k + 5(1) + 3$	(Associativity)
$= 10k + 5 + 3$	(Simplification by multiplication)
$= 10k + 8$	(Simplification by addition)
$= 2(5k + 4)$	(Factoring)
Let $\alpha = 5k + 4$ integer	(Closure by renaming $5k + 4$)
$= 2\alpha$	(Substitution)

Consequently, C (the quantity $5x + 3$) with α is an integer, and thus by the **definition of even** C is **even**. \square

Problem 2 Summary

Procedure

(a) Part (a)

- Use the results from the warm up theorems to show that the theorem is true

(b) Part (b)

- Use the same procedure as the warm up theorems with the field axioms to show that the theorem is true

Key Concepts

- This problem uses the direct proof format and field axioms to prove that a theorem is true

Variations

- We could be given a different theorem to prove
 - We would use the same procedure as the warm up theorems to show that the theorem is true
- The hypothesis could change
 - In this case we would still use the same procedure as the warm up theorems to prove or disprove the theorem



Problem 3

Problem Statement

Prove using a proof by contrapositive. (Be sure to use the right format)

If $3n + 7$ is odd, then n is even (n is an integer). Why does this work, when a direct proof did not?

Solution

Theorem: If $3n + 7$ is odd, then n is even (n is an integer)

Contrapositive Theorem: If n is odd, then $3n + 7$ is even.

Contrapositive Proof:

There exists an integer k such that $n = 2k + 1$	(By definition of odd)
Let $C = 3n + 7$ integer	(Closure by renaming $3n + 7$)
$C = 3n + 7$	(Premise)
$= 3(2k + 1) + 7$	(Substitution)
$= (3)(2k) + (3)(1) + 7$	(Distributivity)
$= (3)(2)k + 3(1) + 7$	(Associativity)
$= 6k + 3 + 7$	(Simplification by multiplication)
$= 6k + 10$	(Simplification by addition)
$= 2(k + 5)$	(Factoring)
Let $\alpha = k + 5$ integer	(Closure by renaming $k + 5$)
$= 2\alpha$	(Substitution)

Consequently, C (the quantity $3n + 7$) with α is an integer, and thus by the **definition of even** C is **even**.
□

Insights

Proving this theorem by using the *contrapositive* statement is easier than doing a direct proof because of the algebra that will happen at some point while attempting a direct proof. There will come a time in the direct proof process where we are not currently equipped with the tools to proceed. Since the *contrapositive* statement is equivalent in truth value to the original statement, we know that proving the *contrapositive* will suffice.

Problem 3 Summary

Procedure

- Because the theorem is written in a difficult way, we would need to use a contrapositive statement to simplify the statement
- State the contrapositive of the statement
- Use the same procedure from here as the warm up theorems to prove the contrapositive statement

Key Concepts

- The theorem is written in a way that makes proving it difficult with a direct proof
- To prove this theorem, we use a contrapositive proof by first stating the contrapositive
- Once the contrapositive statement has been stated, we then apply a direct proof on the contrapositive statement
- Contrapositive statements have the same truth value as the original statement
- Because the contrapositive statement has the same truth value as the original statement, we consequently prove the original statement by proving the contrapositive

Variations

- We could be given a different theorem that requires a contrapositive proof
 - In this case we would use the same procedure but with a different theorem

Problem 4

Problem Statement

Prove for positive integers, if n^2 is even, then n is even.

Solution

Theorem: For positive integers, if n^2 is even, then n is even.

Contrapositive Theorem: For positive integers, if n is odd, then n^2 is odd.

Contrapositive Proof:

There exists an integer k such that $n = 2k + 1$	(By definition of odd)
Let $C = n^2$ integer	(Closure by renaming n^2)
$C = n^2$	(Premise)
$= (2k + 1)^2$	(Substitution)
$= (2k + 1)(2k + 1)$	(By definition of exponents)
$= (2k + 1)(2k) + (2k + 1)(1)$	(Distributivity)
$= (2k)(2k) + (1)(2k) + (2k)(1) + (1)(1)$	(Distributivity)
$= (2)(2)(k)(k) + (2k)(1) + (2k)(1) + (1)(1)$	(Commutativity)
$= 4(k)(k) + 2k + 2k + 1$	(Simplification by multiplication)
$= 4(k)(k) + 4k + 1$	(Simplification by addition)
$= 4k^2 + 4k + 1$	(By definition of exponents)
$= 2(2k^2 + 2k) + 1$	(Factoring)
Let $\alpha = 2k^2 + 2k$ integer	(Closure by renaming $2k^2 + 2k$)
$= 2\alpha + 1$	(Substitution)

Consequently, C (the quantity n^2) with α is an integer, and thus by the **definition of odd** C is **odd**. \square

This proof employs a contrapositive strategy to demonstrate that if n is an odd positive integer, then n^2 is also odd. It begins with the assumption that n is odd, represented as $n = 2k + 1$ for some integer k . The goal is to prove that n^2 is odd. By renaming variables (letting $C = n^2$ and $\alpha = 2k^2 + 2k$), the proof simplifies the algebraic expressions. The critical step is to show that $C = 2\alpha + 1$, indicating that C is an odd integer by definition. In conclusion, this contrapositive proof establishes the original theorem's validity, affirming that when n is an odd positive integer, n^2 is also odd.

Problem 4 Summary

Procedure

- Employ a contrapositive proof for this theorem
- Use the field axioms and the warm up theorem procedure for proving the contrapositive statement

Key Concepts

- Similar to problem (3), this theorem is stated in a way that cannot be proven easily with a direct proof
- We use a contrapositive proof to prove the theorem because it is easier than a direct proof
- We use field axioms and other properties to prove the theorem

Variations

- We could be given a different theorem of the same format
 - This would require the same procedure as the original theorem but with a new theorem



Problem 5

Problem Statement

Prove that $13n + 3$ is even if and only if n is odd. n is an integer. (Must prove both directions, why?)

Solution

Theorem: $13n + 3$ is even if and only if n is odd. n is an integer.

Contrapositive Theorem: If n is even, then $13n + 3$ is odd.

Contrapositive Proof:

There exists an integer k such that $n = 2k$	(By definition of even)
Let $C = 13n + 3$ integer	(Closure by renaming $13n + 3$)
$C = 13n + 3$	(Premise)
$= 13(2k) + 3$	(Substitution)
$= (13)(2)k + 3$	(Associativity)
$= 26k + 3$	(Simplification by multiplication)
$= 26k + 2 + 1$	(Addition)
$= 2(13k + 1) + 1$	(Factoring)
Let $\alpha = 13k + 1$ integer	(Closure by renaming $13k + 1$)
$= 2\alpha + 1$	(Substitution)

Consequently, C (the quantity $13n + 3$) with α is an integer, and thus by the **definition of odd** C is **odd**.
□

Theorem: If n is odd then $13n + 3$ is even.

Direct Proof:

There exists an integer k such that $n = 2k + 1$	(By definition of even)
Let $C = 13n + 3$ integer	(Closure by renaming $13n + 3$)
$C = 13n + 3$	(Premise)
$= 13(2k + 1) + 3$	(Substitution)
$= (13)(2)(k) + (13)(1) + 3$	(Distributivity)
$= (13)(2)k + (13)(1) + 3$	(Associativity)
$= 26k + 13 + 3$	(Simplification by multiplication)
$= 26k + 16$	(Simplification by addition)
$= 2(k + 13)$	(Factoring)
Let $\alpha = k + 13$ integer	(Closure by renaming $k + 13$)
$= 2\alpha$	(Substitution)

Consequently, C (the quantity $13n + 3$) with α is an integer, and thus by **definition of even** C is **even**. □

This proof must be done both ways because the theorem that is presented in the problem statement is a bi-conditional statement.

Problem 5 Summary

Procedure

- State the contrapositive of the original statement
- Prove the contrapositive statement for both directions to prove the parity of the statement

Key Concepts

- This is a biconditional statement which means that it must be proven in both ways
- Biconditional statements that are written in this way must be proven in both directions
- This particular statement uses a contrapositive proof that makes the statement easier to prove

Variations

- We could be given a different theorem
 - In this case we would use the same procedure but with the new theorem



Problem 6

Problem Statement

Assume the domain of positive integers. Prove if $n = ab$, then $a \leq \sqrt{n}$ or $b \leq \sqrt{n}$. Prove by contradiction (using required steps).

Solution

Theorem: If $n = ab$, then $a \leq \sqrt{n}$ or $b \leq \sqrt{n}$

Negation: There exists a positive integer a and there exists a positive integer b such that $n = ab$ and $a > \sqrt{n}$ and $b > \sqrt{n}$.

Proof By Contradiction:

$aa = a^2$	(Definition of exponents)	(1)
$bb = b^2$	(Definition of exponents)	(2)
$(n^{\frac{1}{2}})^2 = n$	(Law of exponents)	(3)
$a > \sqrt{n}$	(Premise)	(4)
$a^2 > n$	(Squaring both sides)	(5)
$a^2 > ab$	(Substitution)	(6)
$a > b$	(Simplification by division)	(7)
$b > \sqrt{n}$	(Premise)	(8)
$b^2 > n$	(Squaring both sides)	(9)
$b^2 > ab$	(Substitution)	(10)
$b > a$	(Simplification by division)	(11)

From lines (7) and (11) we have a contradiction of $a > b$ and $b > a$, therefore if $n = ab$, then $a \leq \sqrt{n}$ or $b \leq \sqrt{n}$. \square

Problem 6 Summary

Procedure

- State the negation of the original statement
- Perform a proof by contradiction by attempting a direct proof
- Arrive at a contradiction and state the contradiction

Key Concepts

- Proofs by contradiction require us to state the negation
- We then attempt to prove the negation with a direct proof and arrive at a contradiction
- The fact that the negation is false implicates that the original statement is true

Variations

- We could be given a different theorem and be asked to prove it by contradiction
 - We would then have to go through the same machinery of how to employ a proof by contradiction



Problem 7

Problem Statement

We define an exponent b of a , (a is an element of the positive real numbers, and b is an element of the natural numbers) as:

a^b is a multiplied by itself b times. Example $a^4 = aaaa$
 $a^0 = 1$.

We will use only the basic principles of algebra (associativity, commutativity, etc), and the definition of exponents given above to prove the rules of exponents.

Assume a and b are elements of the positive real numbers. You must explain each step.

Notice the format of this example proof and follow this structure for #8 and #9. Do not ‘work both sides of the equation.’

Notice how this proof does not use the rule it is trying to prove. **For #7, just fill in the blanks with the correct justification of each step.**

Prove that $a^3 \cdot a^5 = a^{3+5}$

$a^3 \cdot a^5 = (aaa) \cdot (aaaaa)$	Def. of exponents
$= (aaaaaaaa)$	‘ ’ law
$= a^8$	Definition of ‘ ’
$= a^{(3+5)}$	Substitution using $8 = 3 + 5$

Solution

Theorem: $a^3 \cdot a^5 = a^{3+5}$

Direct Proof:

$a^3 \cdot a^5 = (aaa) \cdot (aaaaa)$	Def. of exponents
$= (aaaaaaaa)$	Associativity law of multiplication
$= a^8$	Def. of exponents
$= a^{(3+5)}$	Substitution using $8 = 3 + 5$

Consequently, the quantity $a^3 \cdot a^5$ is equal to that of a^{3+5} . \square

Problem 7 Summary

Procedure

- Fill in the blanks with their respective properties

Key Concepts

- This direct proof uses rules of exponents and the field axioms to prove a theorem

Variations

- We could be given a different proof to fill in the blanks
 - We would then fill in the correct properties for each blank



Problem 8

Problem Statement

Prove that $(a^{(3)})^5 = a^{3 \cdot 5}$.

Solution

Theorem: $(a^3)^5 = a^{3 \cdot 5}$

Direct Proof:

$(a^{(3)})^5 = (aaa)^5$	(Definition of exponents)
Let $\alpha = aaa$	(Closure by renaming aaa)
$= (\alpha)^5$	(Substitution)
$= \alpha \cdot \alpha \cdot \alpha \cdot \alpha \cdot \alpha$	(Definition of exponents)
$= (aaa) \cdot (aaa) \cdot (aaa) \cdot (aaa) \cdot (aaa)$	(Substitution)
$= aaaaaaaaaaaaaaa$	(Associativity)
$= a^{15}$	(Definition of exponents)
$= a^{3 \cdot 5}$	(Substitution using $15 = 3 \cdot 5$)

Consequently, the quantity $(a^{(3)})^5$ is equal to that of $a^{3 \cdot 5}$. \square



Problem 8 Summary

Procedure

- Use the definition of exponents to simplify the premise
- Use a direct proof
- Use closure, field axioms, and definition of exponents to prove the theorem

Key Concepts

- This problem uses the rules of exponents to prove an exponent property
- Exponent rules are a consequence of the field axioms

Variations

- We could be given a different theorem with using exponent rules
 - We would then use exponent rules and field axioms to prove the new theorem



Problem 9

Problem Statement

Prove that $a^5b^5 = (ab)^5$.

Solution

Theorem: $a^5b^5 = (ab)^5$

Direct Proof:

$$\begin{aligned}
 a^5b^5 &= (aaaaa)(bbbbb) && \text{(Definition of exponents)} \\
 &= aaaaabbbbb && \text{(Associativity)} \\
 &= abaaaaabbbb && \text{(Commutativity)} \\
 &= ababaaaabbb && \text{(Commutativity)} \\
 &= abababaabb && \text{(Commutativity)} \\
 &= ababababab && \text{(Commutativity)} \\
 \text{Let } \alpha &= ab && \text{(Closure by renaming } ab) \\
 &= \alpha\alpha\alpha\alpha\alpha && \text{(Substitution)} \\
 &= \alpha^5 && \text{(Definition of exponents)} \\
 &= (ab)^5 && \text{(Substitution)}
 \end{aligned}$$

Consequently, the quantity a^5b^5 is equal to that of $(ab)^5$. \square



Problem 9 Summary

Procedure

- Use the definition of exponents to simplify the premise
- Use a direct proof
- Use closure, field axioms, and definition of exponents to prove the theorem

Key Concepts

- This problem uses the rules of exponents to prove an exponent property
- Exponent rules are a consequence of the field axioms

Variations

- We could be given a different theorem with using exponent rules
 - We would then use exponent rules and field axioms to prove the new theorem



Problem 10

Problem Statement

Considering the above, when we add the definition of exponents to the algebra, are the “rules” of exponents really a new idea or just a consequence of the existing structure of algebra? Give a thoughtful and complete answer.

Solution

The rules of exponents are just a consequence of the existing structure of algebra. Since, in some senses, math is a pattern, these rules emerge from finding a more concise way of representing repeated multiplication. The rules of exponents also obey the field axioms, this can be seen below

$$\begin{array}{ll}
 \alpha = a^n & \text{(Closure)} \\
 a^0 = 1 & \text{(Identity)} \\
 a^{-n} = \frac{1}{a^n} & \text{(Inverse)} \\
 (a^m)^n = a^{mn} & \text{(Associativity)} \\
 a^m \cdot a^n = a^n \cdot a^m & \text{(Commutativity)} \\
 a^m \cdot (a^n + b^n) = a^m \cdot a^n + a^m \cdot b^n & \text{(Distributivity).}
 \end{array}$$

The rules of exponents adhere to the foundation of mathematical structure and adhere to fundamental properties such as the field axioms. The rules of exponents, like stated previously, are a mere consequence of recognizing patterns in mathematics to then formally state equivalences in scenarios. One scenario is that of repeated multiplication, this pattern is then recognized and hence the rules of exponents have been born.



Problem 10 Summary

Procedure

- Answer the prompt

Key Concepts

- Exponent rules are a consequence of the field axioms

Variations

- We could be given a different prompt to respond to
 - We then would answer the prompt

