

Grover's algorithm

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Searching problem

Searching problem:

$$f : \{1, 2, \dots, N\} \rightarrow \{0, 1\}$$

Find $x \in \{1, 2, \dots, N\}$ satisfying $f(x) = 1$.

How many times, do we have to evaluate $f(x)$?

Obviously, $O(N)$.

Quantum searching problem

Unitary oracle

$$U_f |x\rangle |y\rangle = |x\rangle |y \oplus f(x)\rangle .$$

Find $x \in \{1, 2, \dots, N\}$ satisfying $f(x) = 1$.

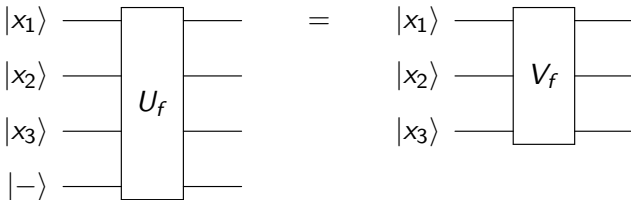
How many times, do we have to evaluate U_f ?

$O(\sqrt{N})$ by Grover's algorithm.

Unitary matrix for Grover's algorithm

Another unitary

$$V_f |x\rangle = (-1)^{f(x)} |x\rangle .$$



$$|x\rangle |-\rangle \mapsto U_f |x\rangle |-\rangle = (-1)^{f(x)} |x\rangle |-\rangle .$$

Grover's algorithm

$$|\psi\rangle := \frac{1}{\sqrt{N}} \sum_{x=1}^N |x\rangle$$

$$V_f = I - 2 \sum_{x:f(x)=1} |x\rangle \langle x|$$

$$W := I - 2 |\psi\rangle \langle \psi|.$$

Then, $G := -WV_f$ is called the Grover's operator.

The Grover's algorithm just measures $G^k |\psi\rangle$ by the computational basis $\{|x\rangle\}_x$ for some **appropriately chosen** k .

The two dimensional subspace

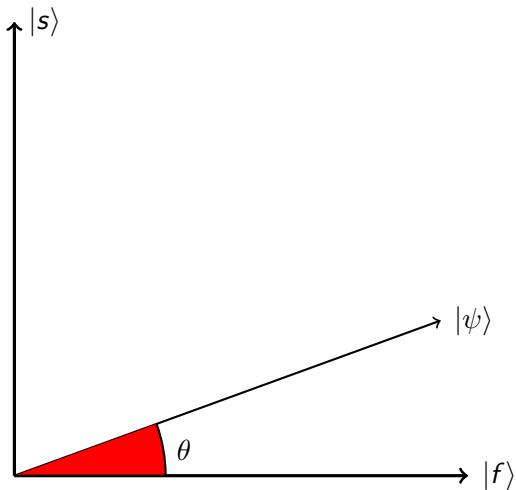
$$|s\rangle := \frac{1}{\sqrt{M}} \sum_{x:f(x)=1} |x\rangle$$
$$|f\rangle := \frac{1}{\sqrt{N-M}} \sum_{x:f(x)=0} |x\rangle.$$

Then,

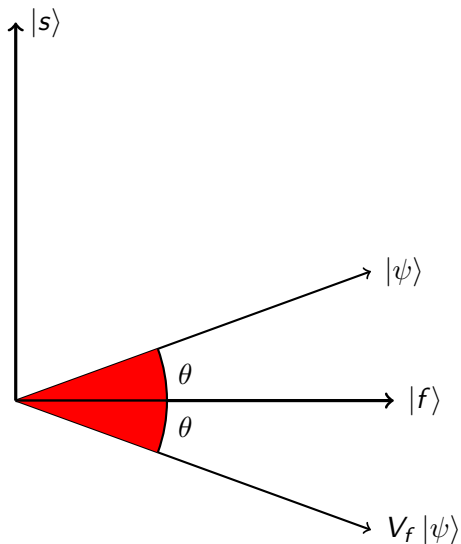
$$\begin{aligned} |\psi\rangle &= \frac{1}{\sqrt{N}} \sum_{x=1}^N |x\rangle = \sqrt{\frac{M}{N}} |s\rangle + \sqrt{\frac{N-M}{N}} |f\rangle \\ &= \sin \theta |s\rangle + \cos \theta |f\rangle \end{aligned}$$

where $\theta := \arcsin \sqrt{\frac{M}{N}}$.

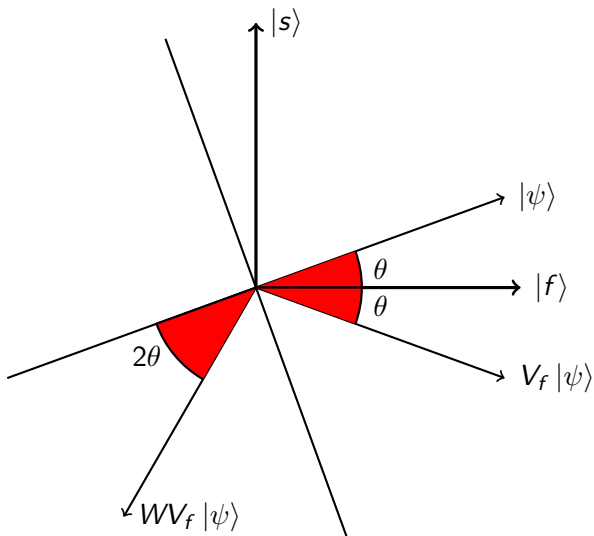
Analysis of Grover's algorithm



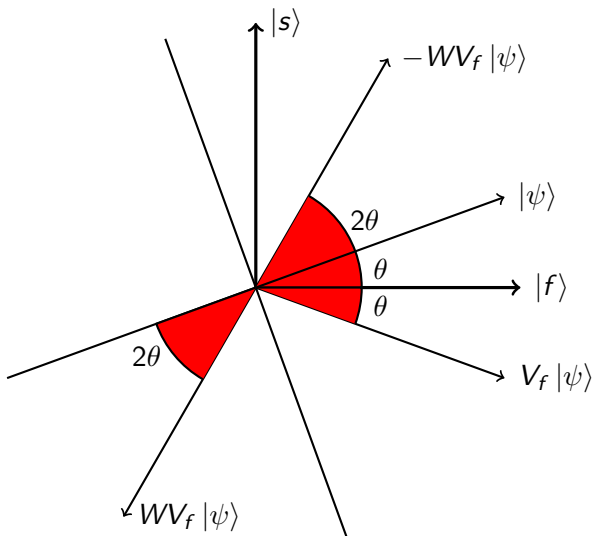
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Analysis of Grover's algorithm

$$(-WV_f)^k |\psi\rangle = \sin((2k+1)\theta) |s\rangle + \cos((2k+1)\theta) |f\rangle$$

The probability of success is $\sin^2((2k+1)\theta)$.

Choose k satisfying

$$(2k+1)\theta \approx \frac{\pi}{2} \iff k \approx \frac{\pi}{4\theta}$$

Here, $\sin \theta = \sqrt{\frac{M}{N}} \iff \theta \approx \sqrt{\frac{M}{N}}$. Hence, $k \approx \frac{\pi}{4} \sqrt{\frac{N}{M}}$.

[Grover 1996]

Grover's search is exactly **optimal** for $M = 1$ [Zalka 1999], and general M [Ito and Mori 2021].

Grover's algorithm

[Boyer, Brassard, Høyer, and Tapp 1998]

- 1 Initialize $m = 1$ and set $\lambda = 8/7$.
- 2 Choose an integer j uniformly from $0, 1, \dots, m$.
- 3 Apply Grover's algorithm with j iterations.
- 4 If solution is not found, set $m \leftarrow \min(\lambda m, \sqrt{N})$ and go back to step 2.

This algorithm solves the “OR problem”
with $O(\sqrt{N/M})$ query for U_f .

Applications of Grover's algorithm

- $O^*(2^{n/2})$ algorithm for SAT.
- $O^*(2^{n/3})$ algorithm for the subset sum [Brassard et al. 1997].
- $O(1.728^n)$ algorithm for the travelling salesman problem [Ambainis et al. 2019].
- $O(1.914^n)$ algorithm for the graph coloring problem [Shimizu and Mori 2020].

Summary

- Grover's search solves the quantum searching problem in time $O(\sqrt{N})$.
- Grover's search is exactly **optimal** if $M = 1$ [Zalka 1999].
- For general M , Grover's search is exactly **optimal** [Ito and Mori 2021].

Assignments

- 1 Show two distinct eigenvalues and corresponding eigenvectors of the Grover operator $-WV_f$.