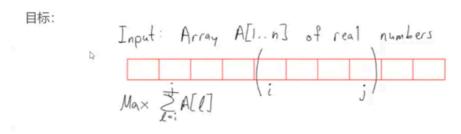
Dynamic programming

- 1、问题目标
- 2、状态的定义: opt[n]
- 3、状态转移方程: opt[n] = best_of(opt[n-1], opt[n-2], ...)

1. 最大子序和



子问题:

$$M(j)$$
: max sum over all windows ending at j .
 $M(j) = \max \left\{ M(j-1) + A[j], A[j] \right\}$

```
def maxSubArray(self, nums):
    :type nums: List[int]
    :rtype: int
    if len(nums)==1:
        return nums[0]
    max_ret = nums[0]
    cur_max = last_max = nums[0]
    for i in range(1, len(nums)):
        if last_max + nums[i] < nums[i]:
            cur_max = nums[i]
        else:
            cur_max = last_max + nums[i]
        if cur_max>max_ret:
            max_ret = cur_max
        last_max = cur_max
        return max_ret
```

2. 最长上升子序列

目标:

Input: Sequence A, ... An
Goal: find a longest strictly increasing
subsequence (not necessarily contiguous).

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子问题:

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3. 零钱兑换

目标:

Input: n denominations of coins
of values 1=V, < V2 < V3 < ... < Vn
Goal: make change for amount of money C.
Use as few coins as possible.

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子问题:

$$M(j)$$
: minimum # coins required to make change for amount of money j.

 $M(j) = \min\{M(j-v_i)\} + 1$

有一些零钱,给定一个数值。用已有的硬币,去组成这个数值。用最少数量的硬币。

建立一个 memory, 一直完善这个 memory。

4. 背包问题

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目标:

Input: n items, with integer sizes since sizes six and values violate of capacity C.

F问题:

M(i,j): optimal value for filling exactly

M(i,j): optimal value for filling exactly
a capacity j knapsack with some
subset of items 1....

M(i,j) = max{M(i-1,j), M(i-1,j-si)+vi}
ith item not
ith item used)

价值数组v = {8, 10, 6, 3, 7, 2},

重量数组w = {4, 6, 2, 2, 5, 1},

背包容量C = 12时对应的m[i][j]数组。

0	1	2	3	4	5	6	7	8	9	10	11	12
1	0	0	0	8	8	8	8	8	8	8	8	8
2	0	0	0	8	8	10	10	10	10	18	18	18
3	0	6	6	8	8	14	14	16	16	18	18	24
4	0	6	6	9	9	14	14	17	17	19	19	24
5	0	6	6	9	9	14	14	17	17	19	21	24
6	2	6	8	9	11	14	16	17	19	19	21	24

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