NOTES ON ERDŐS PROBLEM #451: DENSITY ESTIMATES AND HEURISTIC EVIDENCE

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ABSTRACT. We consider Erdős Problem #451 from the Erdős Problems website: estimate n_k , the smallest integer > 2k such that

$$\prod_{1 \le i \le k} (n_k - i)$$

has no prime factor in the interval (k, 2k). For this problem it is natural to introduce the quantity

$$D_k := \prod_{\substack{k$$

which represents the density of admissible integers n (those for which the block [n-k, n-1] contains no multiple of a prime p with k). We rigorously establish the asymptotic behavior

$$\log \frac{1}{D_k} = (\log 4 + o(1)) \frac{k}{\log k},$$

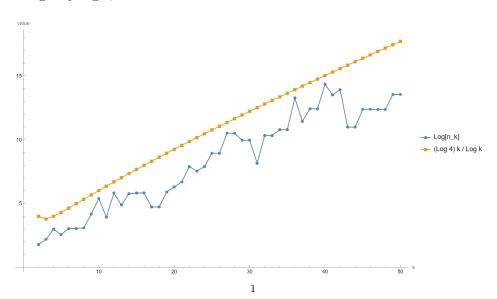
and conjecture that this density asymptotic also governs the growth of the minimal solution n_k .

1. Introduction and Motivation

On Bloom's Erdős Problems website, Problem #451 asks:

Problem 1.1. Estimate n_k , the smallest integer > 2k such that $\prod_{1 \le i \le k} (n_k - i)$ has no prime factor in the open interval (k, 2k).

Numerical experiments for $1 \le k \le 50$ show that $\log n_k$ grows roughly like $k/\log k$. In particular, the plot of $\log n_k$ against $(\log 4) k/\log k$ exhibits very similar growth. Moreover, if we replace $\log 4$ by $\log 3$, then the two curves intersect:



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These data suggest the asymptotic order

Conjecture 1.2.
$$\log n_k = \Theta\left(\frac{k}{\log k}\right)$$
,

and more precisely make it natural to conjecture the upper bound

Conjecture 1.3.
$$\log n_k \leq (\log 4 + o(1)) \frac{k}{\log k}$$
.

Why does the constant $\log 4$ appear? The point is that one can pass from the "no prime factor in (k, 2k)" constraint to a clean density calculation. Let

$$\mathcal{P}_k := \{p: p \text{ is prime and } k$$

Requiring that the whole block [n-k, n-1] contain no multiple of any $p \in \mathcal{P}_k$ is equivalent to requiring that the k integers $n-1, n-2, \ldots, n-k$ are all coprime to M_k . Fix a prime $p \in \mathcal{P}_k$. Modulo p, the interval [n-k, n-1] contains a multiple of p if and only if

$$n \equiv 1, 2, \dots, k \pmod{p}$$
.

Hence the allowed residue classes modulo p are exactly the p-k classes

$$\{0, k+1, k+2, \ldots, p-1\}.$$

Taking all primes $p \in \mathcal{P}_k$ together and applying the Chinese Remainder Theorem, the number of allowed residue classes modulo M_k is

$$A_k = \prod_{\substack{k$$

so the "exact density" of admissible n in the integers is

$$D_k := \frac{A_k}{M_k} = \prod_{\substack{k$$

Heuristically, the "typical gap" between admissible n_k is $\approx D_k^{-1}$, so one expects the minimal solution n_k also to be of order D_k^{-1} . (However, I do not yet have a proof that n_k is indeed of order D_k^{-1} ; this is precisely where things are currently stuck.)

In the next section we will explain how to obtain the asymptotic for D_k^{-1} , namely

$$\log \frac{1}{D_k} = (\log 4 + o(1)) \frac{k}{\log k},$$

which is the key input toward the conjectured bound for $\log n_k$.

2. On the asymptotics of D_k^{-1}

Let

$$\log \frac{1}{D_k} = \sum_{\substack{k$$

We aim to show that this sum has the following asymptotic behavior.

Theorem 2.1. As $k \to \infty$,

$$\log \frac{1}{D_k} = (\log 4) \frac{k}{\log k} + O\left(\frac{k}{(\log k)^2}\right). \tag{2.1}$$

Before proving Theorem 2.1, we first establish some lemmas.

2.1. Auxiliary lemmas. Write $\pi(x)$ for the prime counting function and

$$\operatorname{Li}(x) = \int_2^x \frac{dt}{\log t}.$$

The following lemma is standard, see [1, 2].

Lemma 2.2. For any function F of bounded variation on [k, 2k],

$$\sum_{\substack{k$$

Lemma 2.3. Let $F_k(t) := \log \frac{t}{t-k}$ for t > k and

$$I_k := \int_k^{2k} \frac{F_k(t)}{\log t} \, dt.$$

Then

$$I_k = (\log 4) \frac{k}{\log k} + O\left(\frac{k}{(\log k)^2}\right). \tag{2.3}$$

Proof. Substitute $t = k(1 + u), u \in (0, 1]$:

$$I_k = k \int_0^1 \frac{\log\left(\frac{1+u}{u}\right)}{\log k + \log(1+u)} du.$$

Since $0 \le \log(1+u) \le \log 2$, uniformly in u we have

$$\frac{1}{\log k + \log(1+u)} = \frac{1}{\log k} + O\left(\frac{1}{(\log k)^2}\right).$$

Moreover $\int_0^1 \left| \log \left(\frac{1+u}{u} \right) \right| du < \infty$, so we may integrate termwise to get

$$I_k = \frac{k}{\log k} \int_0^1 \log \frac{1+u}{u} du + O\left(\frac{k}{(\log k)^2}\right).$$

A direct computation gives $\int_0^1 \log \frac{1+u}{u} du = 2 \log 2 = \log 4$. This completes the proof.

Lemma 2.4. Let $\pi(x) = \text{Li}(x) + E(x)$, where the prime number theorem holds with de la Vallée Poussin-type error

$$E(x) \ = \ O\left(xe^{-c\sqrt{\log x}}\right) \qquad (x \to \infty)$$

for some constant c > 0. Then

$$\int_{(k,2k]} \log \frac{t}{t-k} dE(t) = O\left(\frac{k}{(\log k)^2}\right) \qquad (k \to \infty).$$
 (2.4)

Proof. Write

$$F(t) := \log \frac{t}{t - k}$$
 $(t > k),$ $I_k := \int_{(k,2k]} F(t) dE(t).$

Split the interval:

$$I_k = \underbrace{\int_{(k,k+1]} F(t) dE(t)}_{=:I_0} + \underbrace{\int_{(k+1,2k]} F(t) dE(t)}_{=:I_1}.$$

Since (k, k + 1] contains at most one prime,

$$I_0 = \sum_{k$$

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hence $|I_0| \le \log(k+1) + O(1) \ll k/(\log k)^2$.

On [k+1,2k] the function F is bounded and C^1 , with

$$F(2k) = \log 2,$$
 $F(k+1) = \log(k+1),$ $F'(t) = \frac{1}{t} - \frac{1}{t-k} = -\frac{k}{t(t-k)}.$

Integration by parts gives

$$I_1 = F(2k)E(2k) - F(k+1)E(k+1) - \int_{k+1}^{2k} E(t)F'(t) dt.$$

Put

$$H_k := \sup_{t \in [k,2k]} t e^{-c\sqrt{\log t}} \approx k e^{-c\sqrt{\log k}},$$

then $|E(t)| \ll H_k$ uniformly on [k, 2k]. Therefore

$$|F(2k)E(2k)| \ll H_k$$
, $|F(k+1)E(k+1)| \ll \log(k+1) H_k \ll H_k \log k$,

and

$$\int_{k+1}^{2k} |E(t)F'(t)| dt \ll H_k \int_{k+1}^{2k} \frac{k}{t(t-k)} dt = H_k \left[\log \frac{t-k}{t} \right]_{k+1}^{2k} \times H_k \log k.$$

Hence

$$|I_1| \ll H_k \log k \approx k e^{-c\sqrt{\log k}} \cdot \log k.$$

Hence $|I_1| \ll k/(\log k)^2$.

Combining the bounds for I_0 and I_1 yields (2.4).

We are now ready to provide the following.

2.2. **Proof of Theorem 2.1.** By Lemma 2.2 with $F_k(t) = \log \frac{t}{t-k}$,

$$\log \frac{1}{D_k} = \sum_{\substack{k$$

The first integral equals $(\log 4) \frac{k}{\log k} + O\left(\frac{k}{(\log k)^2}\right)$ by Lemma 2.3, and the second is $O\left(\frac{k}{(\log k)^2}\right)$ by Lemma 2.4. This proves (2.1).

3. Concluding remarks

We have now established the asymptotic density estimate

$$\log \frac{1}{D_k} = (\log 4 + o(1)) \frac{k}{\log k}.$$

The remaining problem is to understand the connection between n_k and $\log \frac{1}{D_k}$. Intuitively, the average spacing of admissible integers n_k is $\approx D_k^{-1}$, so one expects that the minimal solution n_k should also be of order D_k^{-1} . However, making this heuristic precise is delicate. At present, we do not know how to establish such a result rigorously. Bridging this gap between the density asymptotic and the actual minimal solution n_k is what remains to be done.

Based on the numerical figure in Section 1, we also computed n_k and D_k for $1 \le k \le 50$ and obtained the ratios

$$R_k := n_k D_k$$
.

In this range the smallest and largest values of R_k occur at

$$\mathbf{Min} \ R_k \ \text{at} \ k = 46: \quad n_{46} = 241{,}251, \quad D_{46} = 5.5915609400 \times 10^{-7}, \quad R_{46} = 0.1348969668,$$

$$\mathbf{Max}\ R_k\ \mathrm{at}\ k=23:\quad n_{23}=1884,\quad D_{23}=0.0041253238,\quad R_{23}=7.7721101000.$$

References

- [1] H. Iwaniec and E. Kowalski, Analytic Number Theory.
- [2] H. L. Montgomery and R. C. Vaughan, Multiplicative Number Theory I. Classical Theory.

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