

## 1. INTRODUCTION

In [2, Section 1], Erdős proposed the following problem (see also Problem #614 in [1]):

**Problem 1.1.** *Let  $f(n, k)$  be minimal such that there is a graph with  $n$  vertices and  $f(n, k)$  edges where every set of  $k + 2$  vertices induces a subgraph with maximum degree at least  $k$ . Determine  $f(n, k)$ .*

Now, we denote  $f(n, k)$  for the minimum number of edges of an  $n$ -vertex graph  $G$  with the property that every  $(k + 2)$ -subset of vertices induces a subgraph whose maximum degree is at least  $k$ . Equivalently, if  $H = \overline{G}$  denotes the complement of  $G$ , then

every induced subgraph of  $H$  on  $k + 2$  vertices has minimum degree at most 1. (★)

Thus  $f(n, k) = \binom{n}{2} - \max\{e(H) : H \text{ satisfies } (★)\}$  and our task is to maximize the number of edges under the constraint (★).

## 2. WHEN $k = 1$

We first dispose of the easy base  $k = 1$ .

**Theorem 2.1.** *For all  $n \geq 2$ , we have*

$$f(n, 1) = \binom{n}{2} - \left\lfloor \frac{n^2}{4} \right\rfloor.$$

*Proof.* For  $k = 1$ , property (★) says that  $\overline{G}$  is triangle-free. Hence  $e(\overline{G}) \leq \lfloor n^2/4 \rfloor$  by Mantel's theorem, with equality attained by the balanced complete bipartite graph. The formula for  $f(n, 1)$  follows. □

## 3. WHEN $k = 2$

**Theorem 3.1.** *For all  $n \geq 4$ , we have*

$$f(n, 2) = \binom{n}{2} - \text{ex}(n, C_4).$$

*Proof.* Let  $k = 2$ . First assume that  $H$  satisfies (★). If  $H$  contains a copy of  $C_4$  on vertex set  $S$  with  $|S| = 4$ , then the induced subgraph  $H[S]$  contains that cycle and hence has minimum degree at least 2, contradicting (★). Thus  $H$  is  $C_4$ -free.

Conversely, assume that  $H$  is  $C_4$ -free and let  $S \subseteq V(H)$  with  $|S| = 4$ . If  $\delta(H[S]) \geq 2$ , then  $H[S]$  is a graph on 4 vertices with minimum degree at least 2, hence by Dirac's theorem it is Hamiltonian. In particular,  $H[S]$  contains a spanning cycle of length 4, contradicting that  $H$  is  $C_4$ -free. Therefore  $\delta(H[S]) \leq 1$  for every 4-set  $S$ , so  $H$  satisfies (★).

Finally,

$$f(n, 2) = \binom{n}{2} - \max\{e(H) : H \text{ satisfies } (★)\} = \binom{n}{2} - \text{ex}(n, C_4),$$

since the admissible  $H$  are exactly the  $C_4$ -free graphs. □

## 4. A GENERAL TURÁN REDUCTION

**Definition 4.1.** Fix an integer  $k \geq 2$  and set  $m := k + 2$ . We recall that an  $n$ -vertex graph  $H$  satisfies (★) if

$$\forall S \subseteq V(H) \text{ with } |S| = m, \quad \delta(H[S]) \leq 1.$$

Let  $\mathcal{F}_k$  be the (finite) family of all graphs  $F$  on exactly  $m$  vertices with

$$\delta(F) \geq 2.$$

Let  $\mathcal{F}_k^{\min} \subseteq \mathcal{F}_k$  be the subfamily consisting of those  $F \in \mathcal{F}_k$  that are *edge-minimal* with  $\delta(F) \geq 2$ , i.e.  $\delta(F) \geq 2$  but for every edge  $e \in E(F)$  we have

$$\delta(F - e) \leq 1.$$

**Lemma 4.2.** *A graph  $H$  satisfies  $(\star)$  if and only if  $H$  is  $\mathcal{F}_k$ -free as a (not necessarily induced) subgraph.*

*Proof.*  $(\Rightarrow)$  Suppose  $H$  satisfies  $(\star)$ . If  $H$  contained some  $F \in \mathcal{F}_k$  as a subgraph on a vertex set  $S$  with  $|S| = m$ , then in the induced graph  $H[S]$  every vertex would have degree at least its degree in  $F$ , hence  $\delta(H[S]) \geq \delta(F) \geq 2$ , contradicting  $(\star)$ .

$(\Leftarrow)$  Conversely, suppose  $H$  violates  $(\star)$ . Then there exists  $S \subseteq V(H)$  with  $|S| = m$  such that  $\delta(H[S]) \geq 2$ . Taking  $F := H[S]$ , we have  $F \in \mathcal{F}_k$  and  $F$  appears as a subgraph of  $H$ .  $\square$

**Lemma 4.3.** *A graph  $H$  is  $\mathcal{F}_k$ -free if and only if it is  $\mathcal{F}_k^{\min}$ -free.*

*Proof.* The forward implication is trivial since  $\mathcal{F}_k^{\min} \subseteq \mathcal{F}_k$ .

For the reverse implication, suppose  $H$  contains a copy of some  $F \in \mathcal{F}_k$  on a fixed vertex set  $S$  of size  $m$ . Starting from this copy, delete edges (within  $S$ ) one by one as long as the minimum degree remains at least 2. Since  $S$  is finite, this process terminates at a subgraph  $F' \subseteq F$  on the same vertex set  $S$  such that  $\delta(F') \geq 2$  but deleting any further edge would force the minimum degree to drop to at most 1. Thus  $F' \in \mathcal{F}_k^{\min}$  and  $H$  contains  $F'$ .  $\square$

**Proposition 4.4.** *Let  $k \geq 2$  be fixed. Then*

$$\max\{e(H) : |V(H)| = n, H \text{ satisfies } (\star)\} = \text{ex}(n, \mathcal{F}_k^{\min}).$$

*Proof.* By Lemma 4.2, graphs satisfying  $(\star)$  are precisely the  $\mathcal{F}_k$ -free graphs, so the maximum number of edges among them is  $\text{ex}(n, \mathcal{F}_k)$  by definition of  $\text{ex}(\cdot, \cdot)$ . By Lemma 4.3,  $\mathcal{F}_k$ -freeness is equivalent to  $\mathcal{F}_k^{\min}$ -freeness, hence  $\text{ex}(n, \mathcal{F}_k) = \text{ex}(n, \mathcal{F}_k^{\min})$ .  $\square$

Let  $f(n, k)$  be as in Problem 1.1 and set  $m := k + 2$ . Passing to the complement  $H = \overline{G}$ , the defining property for  $G$  is equivalent to the condition that  $H$  satisfies  $(\star)$ . Consequently, we obtain the following Turán-type expression for  $f(n, k)$ .

**Theorem 4.5.**

$$f(n, k) = \binom{n}{2} - \text{ex}(n, \mathcal{F}_k^{\min}).$$

For  $3 \leq k \leq 6$ , we computed the family  $\mathcal{F}_k^{\min}$  by an exhaustive Mathematica search; see Figures 1–5.

In particular, let  $K_1 \vee 2K_2$  denote the join of a single vertex with two disjoint edges. Our computation in Figure 2 yields

$$\mathcal{F}_3^{\min} = \{C_5, K_{2,3}, K_1 \vee 2K_2\},$$

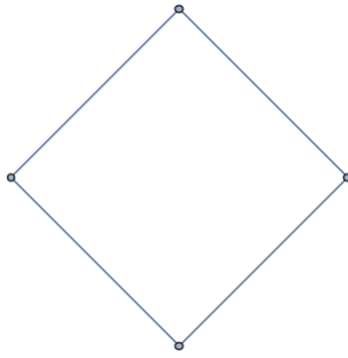
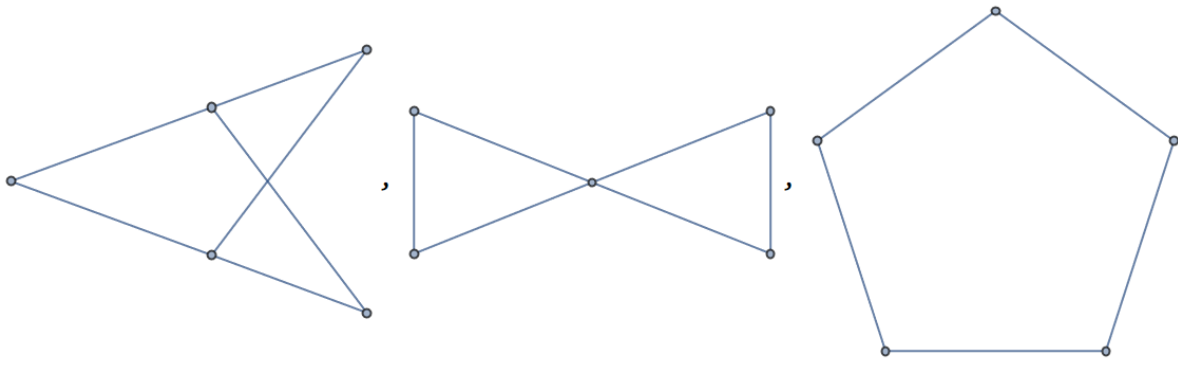
and hence we obtain an explicit Turán formulation for  $f(n, 3)$ .

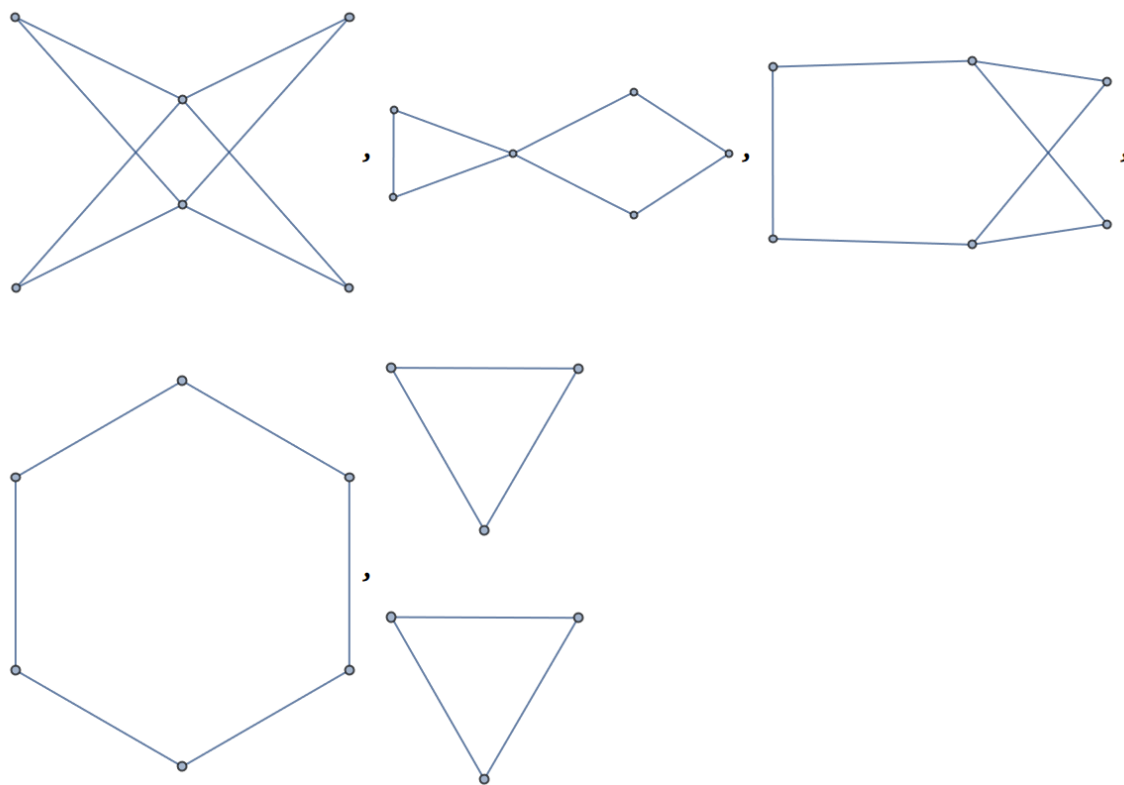
**Corollary 4.6.** *For all  $n \geq 5$ ,*

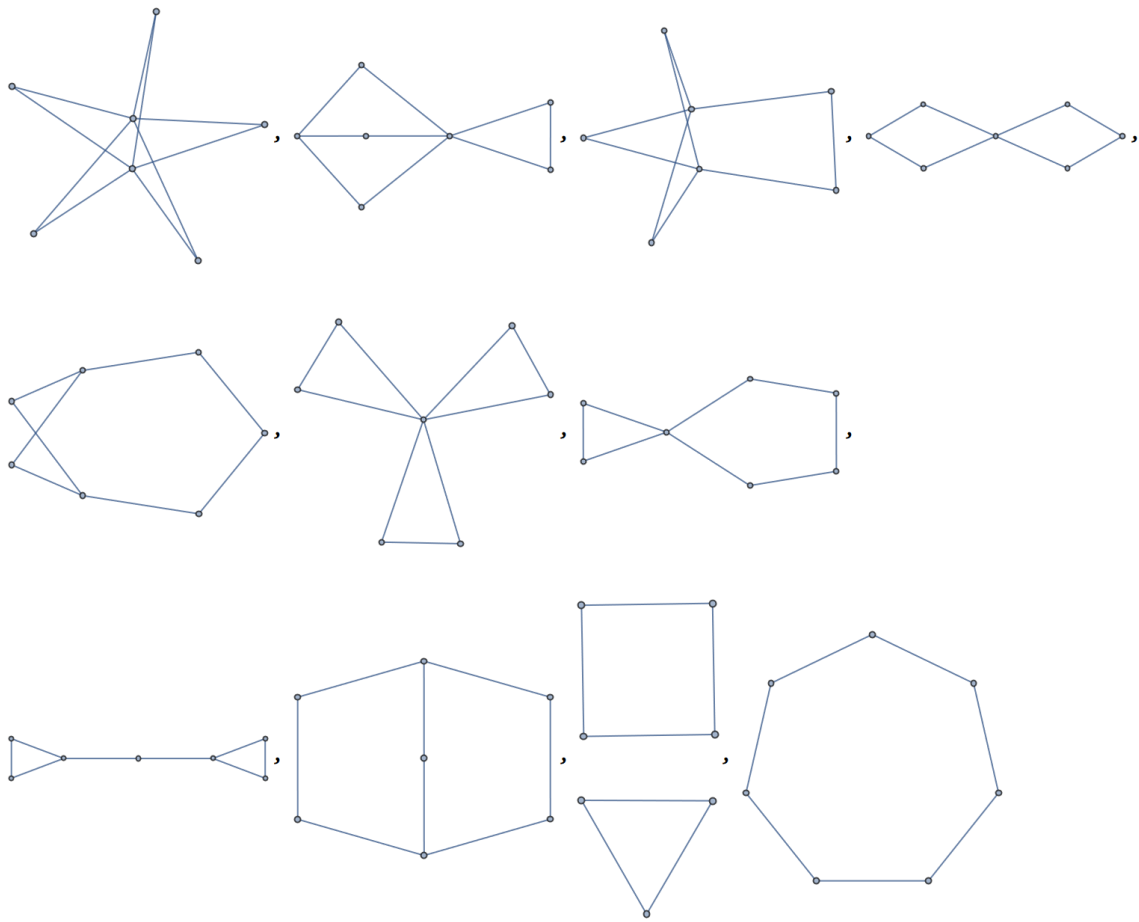
$$f(n, 3) = \binom{n}{2} - \text{ex}(n, \{C_5, K_{2,3}, K_1 \vee 2K_2\}).$$

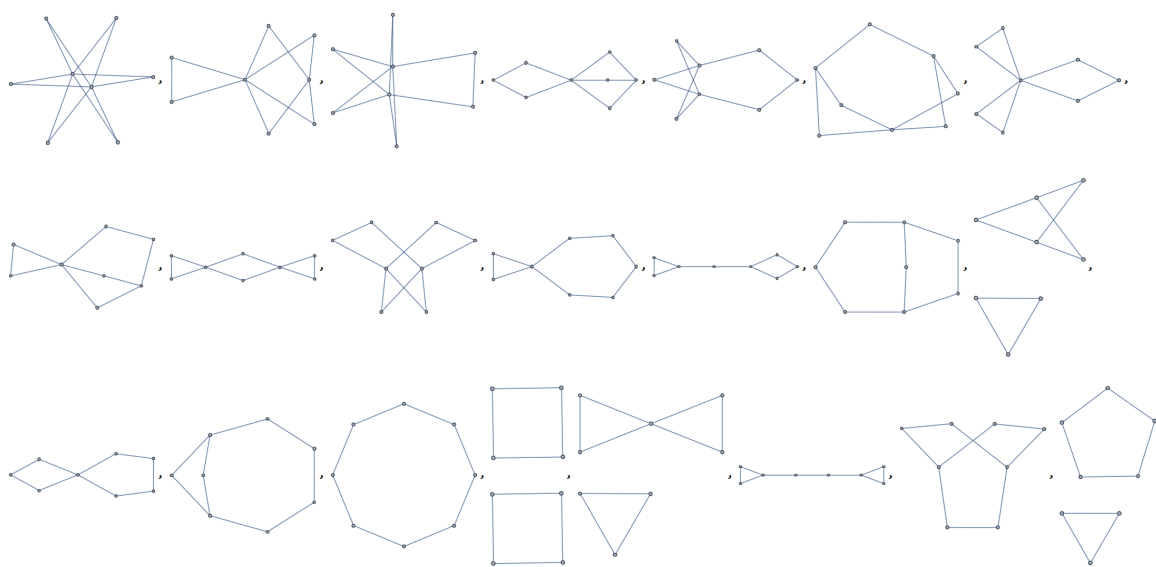
## REFERENCES

- [1] T. F. Bloom, Erdős Problem #614, <https://www.erdosproblems.com/614>, accessed 2025-10-26.
- [2] Faudree, R. J. and Rousseau, C. C. and Schelp, R. H., Problems in graph theory from Memphis. The mathematics of Paul Erdős, II (1997), 7–26.

FIGURE 1.  $k = 2$ FIGURE 2.  $k = 3$

FIGURE 3.  $k = 4$

FIGURE 4.  $k = 5$

FIGURE 5.  $k = 6$